

Time modeling in MARTE

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Abstract

This article introduces the Time Model subprofile of MARTE, a new OMG UML Profile dedicated to Modeling and Analysis of Real-Time and Embedded systems. After a brief presentation of former time modeling elements present in SPT and UML2, we introduce the Time meta-model of MARTE. It defines physical and logical time, timed model elements and their associated properties. We present both the time domain view and the UML representation of the most important concepts. Various time bases (called clocks in the profile) can be correlated using clock relations and constraints, built from a core set predefined in the profile. Constraints are usually collected from scheduling and partitioning decisions taken in the course of design flow for embedded systems. We illustrate this on two simple examples.

I. Introduction

Modeling of Time should certainly be a central concern in Model-Driven Engineering dedicated to Real-Time Embedded Systems (RTES). Timed extensions should allow to support a rich design flow, that can encompass both established and emerging new techniques for model-based optimization, transformation and analysis of systems. Indeed, embedded system models very often consists of a predefined set of *application* functions, and of *execution platforms* on which to *allocate* these functions. Application elements are increasingly componentized, may coexist and possibly cooperate concurrently. Execution platforms increasingly comprise parallel resources for both communications and computations.

The design challenge in embedded system modeling is then to provide model-level compilation techniques that provide support for both spatial distribution and temporal scheduling of applications onto platforms (collectively called allocation). This approach is therefore akin to system level design techniques such as advocated in SysML [1]. But SysML, just like UML, hardly formalizes

its real-time aspects.

This global issue triggered over the years a number of proposals for specific representation formalisms, and their associated particular design techniques. These models then can be, and often have been, represented inside the scope of the UML [2], [3], [4], [5], [6]. But this is typically done, a) mostly at the metamodel level, and b) without any clear interpretation of any kind of the time annotations in the framework of the UML. This raises the risk of mismatch between the private interpretation inherited from the formalism and the existing UML semantics [7].

The primary objective of the Time subprofile in MARTE [8] was to provide basic and advanced time modeling concepts, with interpretation *inside* the UML modeling level, not *outside*. These time-related concepts could then be used to build various Models of Computation and Communication (MoCC). Importantly, the profile should subsume the former SPT [9] and the UML2 Simple Time models [10], while extending them towards the possibility of modeling much richer MoCC-based design approaches [11], [12], [13], [14].

Time as considered in design can be of physical or logical nature. Physical time is continuous, but can usually be discretized into *chronometric* clocks under appropriate assumptions. Logical time is less often recognized in itself as an explicit modeling concept. Processings and execution steps performed at the rate of a processor cycle (which may vary according to power consumption management), or triggered by successive occurrences of an external event (such as completion of an engine revolution) are simple example of that. Often the allocation process may be perceived as this : asynchronous concurrent application components are each considered as being governed by their own (local) logical clock, connected to appropriate events; then the allocation itself consists in *fitting* these various clock threads onto a single (or at least more correlated) synchronous clock, subject to constraints of various sources abstracted from physical time properties and requirements. The transformation and analysis steps involved in the proper mapping are (at least implicitly) dealing with scheduling objects that are

relations between logical and physical clocks attached to the various processings. MARTE Time profile is meant exactly to represent that.

In the sequel, we shall describe the profile in greater details and its position with respect to other parts of MARTE. We start with the domain view and give an overview of the UML representation. Two examples illustrate the usage of the profile.

II. Time domain view

This section describes the MARTE *Time domain view*, i.e., the main concepts related to time and their relationships. In an MDE approach, this is done through meta-modeling. Before reviewing the MARTE Time models, we will have a look at the former UML profile for Time: SPT

A. SPT

The UML Profile for Schedulability, Performance, and Time (SPT) aimed at filling the lacks of UML 1.4 in some key areas that are of particular concern to real-time system designers and developers. SPT introduces a *quantifiable* notion of time and resources. It annotates model elements with quantitative information related to time, information used for timeliness, performance, and schedulability analyses.

SPT only considers (*chrono*)metric time, which makes implicit reference to physical time. It provides time-related concepts: concepts of instant and duration, concepts for modeling events in time and time-related stimuli. SPT also addresses modeling of timing mechanisms (clocks, timers), and timing services. SPT, which relies on UML 1.4, had to be aligned with UML 2.1. This is one of the objectives of the MARTE profile, presented below.

UML 2 has included a new package called SimpleTime, part of the CommonBehavior package. The main new meta-classes are TimeEvent and Observation (Time and Duration). The model is kept very simple and is intended to be extended in specialized profiles. This is the specific agenda of MARTE Time model.

B. MARTE time model

As a successor of SPT, MARTE has to support a *metric* time with implicit reference to physical time. However, MARTE goes beyond this quantitative model of time and adopts more general time models suitable for system design. In MARTE, Time can be *physical*, and considered as *dense* or *discretized*, but it can also be *logical*, and related to user-defined clocks. Time may even be *multiform*, allowing different times to progress in a non-uniform fashion, and possibly independently to any (direct) reference to physical time.

1) *Concept of time structure*: The building element in a time structure is the *TimeBase* (Fig. 1). A time base is a totally ordered set of instants. A set of instants can be

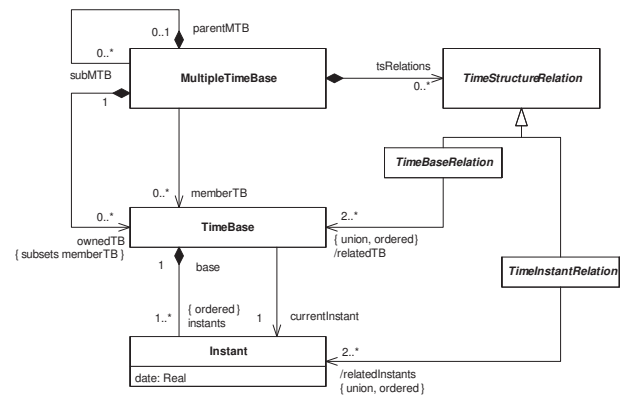


Fig. 1. Time structure (Domain view).

discrete or *dense*. The linear vision of time represented by a single time base is not sufficient for most of the applications, especially in the case of multithreaded or distributed applications. Multiple time bases are then used. A *MultipleTimeBase* consists of one or many time bases. A time structure contains a tree of multiple time bases.

Time bases are *a priori* independent. They become dependent when instants from different time bases are linked by relationships (coincidence or precedence). The abstract class *TimeInstantRelation* in Fig. 1 has *CoincidenceRelation* and *PrecedenceRelation* as concrete subclasses. Instead of imposing local dependencies between instants, dependencies can be directly imposed between time bases. A *TimeBaseRelation* (or more precisely one of its concrete subclasses) specifies many (possibly an infinity of) individual time instant relations. This will be illustrated later on some time base relations. *TimeBaseRelation* and *TimeInstantRelation* have a common generalization: the abstract class *TimeStructureRelation*. As a result of adding time structure relations to multiple time bases, time bases are no longer independent and the instants are partially ordered. This partial ordering of instants characterizes the *time structure* of the application.

This model of time is sufficient to check the logical correctness of the application. Quantitative information, attached to the instants, can be added to this structure when quantitative analyses become necessary (see below).

2) *Access to time*: In real world technical systems, special devices, called clocks, are used to measure the progress of physical time. In MARTE, we adopt a more general point of view: a clock is a *model element* giving access to the time structure. Time may be logical or physical or both.

A *Clock* makes reference to a *TimeBase* (Fig. 2), and thus indirectly to the instants of this timebase. The nature attribute indicates whether the accessed time is dense or discrete. A *Clock* accepts units (*acceptedUnits* property). Unit is a concept introduced in the MARTE NFP (Non-functional property) domain view. One of these accepted units is the *defaultUnit*. A clock may own an event

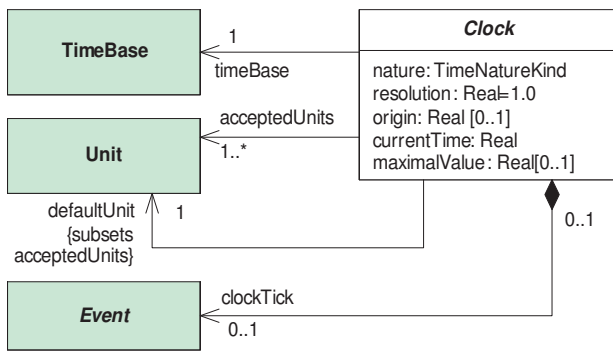


Fig. 2. Clock

(clockTick). This event occurs at each change of the current time of the clock. Other attributes characterize quantitative information that can be attached to a clock. The resolution property specifies the readout granularity of the clock, expressed in defaultUnit unit. Its default value is 1.0. The optional attribute origin specifies the possible offset in the clock reading. The optional attribute maximalValue expresses the limited capability of usual clocks to represent arbitrary large instant values: the clock “rolls over” when the currentTime value gets at the maximal value. For instance, for a discrete periodic clock, the time value attached to the k^{th} instant is given by $(origin + (k - 1) * resolution) \bmod maximalValue$.

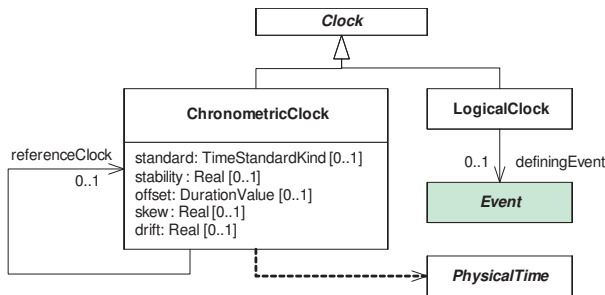


Fig. 3. Logical and chronometric clock

Clock is an abstract concept. There exist two concrete specializations of Clock: LogicalClock and ChronometricClock. A *chronometric* clock is a clock making reference to physical time. A special attention is then put on the quantitative information attached to this model. Non functional properties like stability, skew, etc can be defined for these clocks with respect to some referenceClock. On the other hand, a *logical* clock can be defined by any event (definingEvent property); in this case, the clock ticks at each occurrence of the defining event. Logical time is usually counted in the number of ticks. So, tick is a predefined unit often used as the default unit for a logical clock.

3) *Time value specifications*: An application may use time in two ways: either as a reference to a time instant or as a time span. So, MARTE introduces two distinct concepts: InstantValue and DurationValue, specializations of the abstract concept of TimeValue. Since the access

to time is done through clocks, a TimeValue refers to a Clock (the onClock property). A time value also have an associated unit. When optional property unit is given, it must be used instead of default unit of the associated clock. The attribute nature specifies whether the time values associated with the clock take their values in a dense or discrete domain.

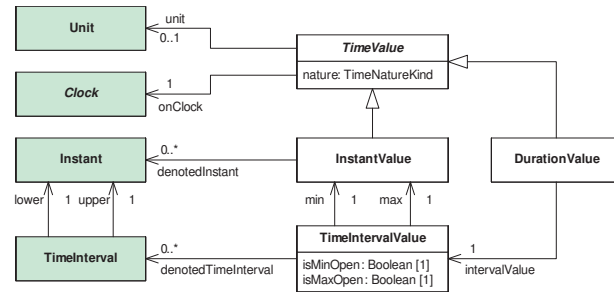


Fig. 4. Time values

4) *Time-related concepts*: A timed element is a most general concept. *TimedElement* is an abstract class generalization of all other timed concepts. It associates a non empty set of clocks with a model element. The semantics of the association with clocks depends on the kind of timed element.

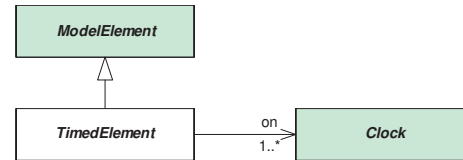


Fig. 5. Timed element

Events and behaviors can be directly bound to time. The occurrences of a (timed) event refer to points of time (instants). The executions of a (timed) behavior refer to points of time (start and finish instants) or to segments of time (duration of the execution).

TimedEvent (TimedProcessing, resp.) is a concept representing an event (a processing, resp.) *explicitly* bound to time through a clock. In this way, time is not a mere annotation: it changes the semantics of the timed model elements.

Other timed elements—not detailed in this presentation—are also defined in the MARTE Time domain: timed observations, and timed constraints. As timed elements they explicitly bind observations or constraints to clocks.

III. UML view of Time in MARTE

A. The Time sub-profile

The time structure presented above constitutes the semantic domain of our time model. The UML view is defined in the “MARTE Time profile”. This profile introduces a limited number of powerful stereotypes. We have striven to avoid the multiplication of too specialized

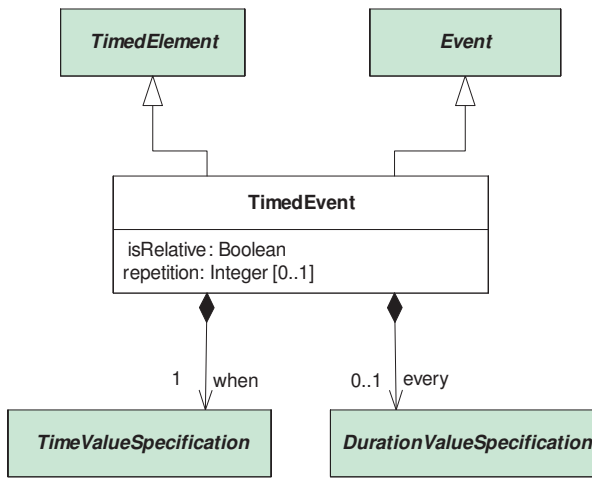


Fig. 6. Timed event

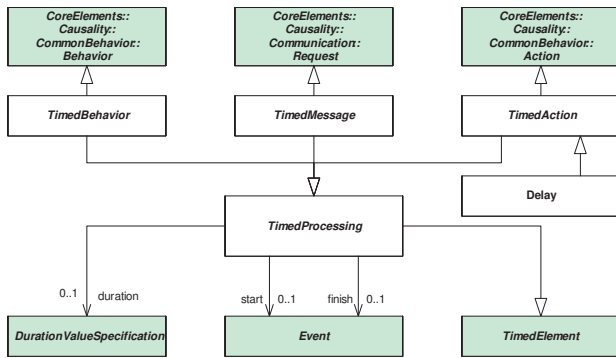


Fig. 7. Timed processing

stereotypes. Thanks to the sound semantic grounds of our stereotypes, modeling environments may propose patterns for more specific uses.

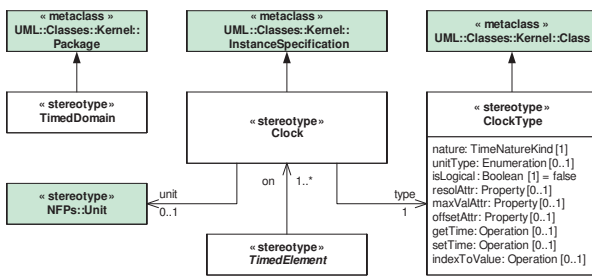


Fig. 8. MARTE Time profile: Clock.

1) *ClockType and Clock*: The main stereotypes are presented in figure 8. ClockType is a stereotype of the UML Class. Its properties specifies the kind (chronometric or logical) of clock, the nature (dense or discrete) of the represented time, a set of clock properties (e.g., resolution, maximal value...), and a set of accepted time units. Clock is a stereotype of InstanceSpecification. An OCL rule imposes to apply the Clock stereotype only to instance specifications of a class stereotyped by ClockType. The unit of the clock is given when the stereotype

is applied. Unit is defined in the Non Functional Property modeling (NFPs) subprofile of MARTE; it extends EnumerationLiteral. It is very convenient because a unit can be used like any user-defined enumeration literal, and conversion factors between units can be specified (e.g., $1ms = 10^{-3}s$). TimedElement is an abstract stereotype with no defined metaclass. It stands for model elements which reference clocks. All other *timed* stereotypes specialize TimedElement.

2) *Clock constraints*: ClockConstraint is a stereotype of the UML Constraint. The clock constraints are used to specify the time structure relations of a time domain.

The *context* of the constraint must be a TimedDomain. The *constrained elements* are clocks of this timed domain and possibly other objects. The *specification* of a clock constraint is a set of declarative statements. This raises the question of choosing a language for expressing the clock constraints. A natural language is not sufficiently precise to be a good candidate. UML encourages the use of OCL. However, our clocks usually deal with infinite sets of instants, the relations may use many *mathematical* quantifiers, which are not supported by OCL. Additionally, OCL [15] is made to be evaluable, while our constraints often have to be processed altogether to get a set of possible solutions. So, we have chosen to define a simple constraint expression language endowed with a mathematical semantics. The specification of a clock constraint is a UML::OpaqueExpression that makes use of *pre-defined* (clock) relations, the meaning of which is given in mathematical terms, outside the UML. Our *Clock Constraint Specification Language* is not normative. Other languages can be used, so long as the semantics of clocks and clock constraints is respected.

3) *TimedEvent and TimedProcessing*: In UML, an Event describes a set of possible occurrences; an occurrence may trigger effects in the system. A UML2 TimeEvent is an Event that defines a point in time (instant) when the event occurs. The MARTE stereotype TimedEvent extends TimeEvent. Its instant specification *explicitly* refers to a clock. If the event is recurrent, a repetition period—duration between two successive occurrences of the event—and the number of repetitions may be specified.

In UML, a Behavior describes a set of possible executions; an execution is the performance of an algorithm according to a set of rules. MARTE associates a duration, an instant of start, an instant of termination with an execution, these times being read on a clock. The stereotype TimedProcessing extends the metaclasses Behavior, Action, and also Message. The latter extension assimilates a message transfer to a *communication* action.

Note that, StateMachine, Activity, Interaction being Behavior, they can be stereotyped by TimedProcessing, and thus, can be bound to clocks.

B. The Time model library

The TimeLibrary (Fig 9) is a user's model library that provides enumerations related to time and facilities for using the ideal chronometric time (*i.e.*, the time referenced in physical laws). TimeUnitKind contains the main chronometric time units. s (second) is an SI unit. Other units are derived units. All the enumeration literals are stereotyped by Unit. LogicalTimeUnitKind is a special enumeration which contains one enumeration literal only. This literal is tick. The IdealClock and its instance idealClk model the abstract and ideal time which is used in physical laws. It is a dense time. idealClk should be imported in models that refer to chronometric time. TimedValueType is a templated data type. The template parameter is an enumeration which contains time units.

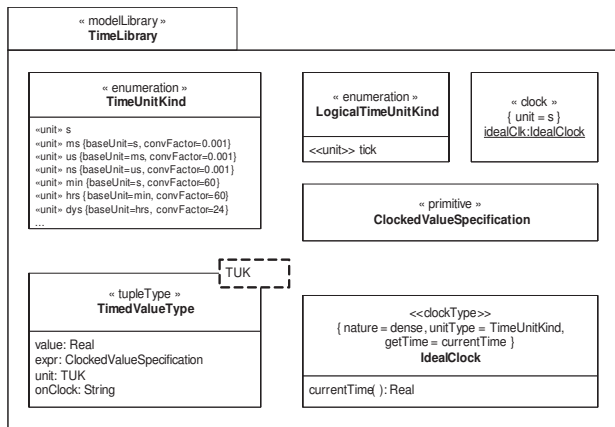


Fig. 9. TimeLibrary: a user library for time

Besides libraries, other facilities are offered to MARTE users: concrete languages dedicated to value expressions (Value Specification Language—VSL) and to clock constraint expressions (Clock Constraint Specification Language—CCSL) [16, Annexes B & C]. The latter defines a core set of constraints that can be extended to express desired relation patterns between timed elements.

Examples of such clock constraints are described in a technical report [17]. For lack of room we cannot describe them here. Some constraint relations state that variations in rate or jitter between two clocks are somehow bounded, that clocks are related up to some drift, or that a clock has to be a subclock of some other (with the subclocking mechanism following possibly some pattern). Some relations are more imperative in that they denote the single solution to some clock transformations (as for instance in `b isPeriodicOn a ofPeriod n`).

IV. Examples

The first step for the designer is to construct its own, not always perfect, clocks. The first subsection shows how to create chronometric clocks and the second subsection illustrates the use of logical clocks.

A. Chronometric clock specifications

The MARTE TimeLibrary provides a model for the *ideal time* used in physical laws: idealClk, which is an instance of the class IdealClock, stereotyped by ClockType (Fig. 10, upper part). idealClk is a dense time clock, its unit is the SI time unit s.

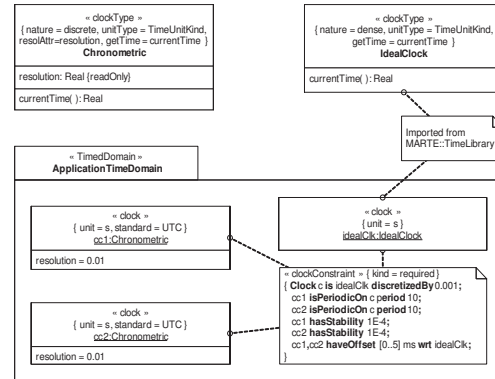


Fig. 10. Chronometric clocks.

After importing the library, new user-defined chronometric clocks can be defined. For instance, Fig. 10 defines the class Chronometric with an attribute resolution of type Real and an operation currentTime that allows for reading the current time. The stereotype ClockType is applied and the tagged values characterize the nature of clocks represented by this class. Here, the clock is a discrete clock and accepts the units defined by the predefined Enumeration TimeUnitKind (see Fig. 9). By default, the clock types are chronometric, not logical.

Actual clocks belong to timed domains, *i.e.*, a package stereotyped by TimedDomain (Fig. 10, lower part). Here, a single time domain is considered. It owns three clocks. idealClk is imported from the library. Two instances of the class Chronometric, cc1 and cc2, are defined. They both use s (second) as a time unit and their resolution is 0.01 s. The three clocks are *a priori* independent. A clock constraint specifies relationships among them. According to the given constraints, cc1 and cc2 are two 100 Hz clocks, the stability of which is 10^{-4} , and with an offset less than 5ms.

B. Logical clock specifications

Fig. 11 illustrates the definition of logical clocks. The discrete, logical, clock type AngleClock is defined. It has three attributes, resolution, offset and maximalValue. The label function angle associates a real value, the clock reading, with each instant of the clock. Each instant is uniquely identified by a natural number, its index, inferred from the linear order defined on the clock instants.

The enumeration AngleUnitKind defines the units, enumeration literals stereotyped by NFP::Unit, that can be used by the clocks.

Two clocks, instances of AngleClock, are then created. crkClk represents the crankshaft revolutions, its unit is

$^{\circ}CRK$ (degree crank), its resolution is $1^{\circ}CRK$ and its maximal value is $720^{\circ}CRK$, *i.e.*, two revolutions of the crank shaft. $camClk$ represents the camshaft resolutions, its unit is $^{\circ}CAM$ (degree cam), its resolution is $1^{\circ}CAM$ and its maximal value is $360^{\circ}CAM$, one rotation of the camshaft. One revolution of the camshaft (mechanically) implies two revolutions of the crankshaft, hence $1^{\circ}CAM = 2^{\circ}CRK$.

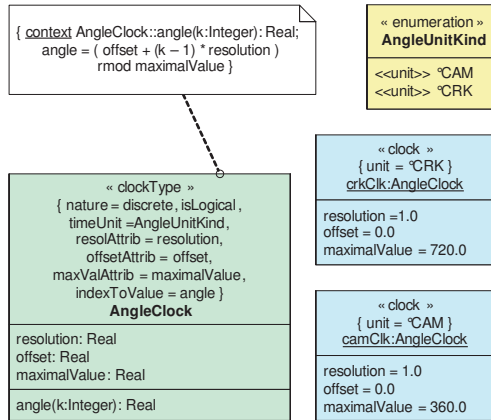


Fig. 11. Logical Clocks.

Fig. 12 uses the clock $camClk$ to represent a four-stroke engine cycle. This state machine is stereotyped by $TimedProcessing$. The on attribute identifies the clock used and therefore the unit ($^{\circ}CAM$).

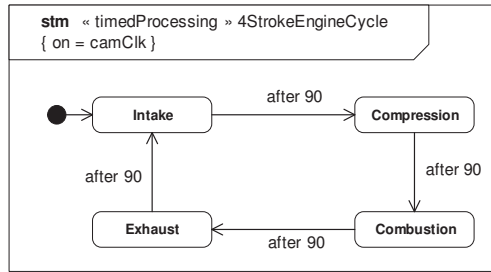


Fig. 12. State machine of a 4-stroke engine cycle.

This example is developed in a previous paper [18] dedicated to the use of multiform time, and the modeling of the “knock control” problem with MARTE.

V. Conclusion

The MARTE Time subprofile provides a limited number of time concepts from which to build Models of Computations with timed interpretations (as exemplified by our modeling of AADL aspects in the same proceedings [19]). These concepts rely on the existence of various time threads (or *logical clocks*) that may drive application elements. Constraint relations may exist between clocks, from the loose asynchronous compositions to stricter simultaneous coincidence. More constraints are raised as result of scheduling decisions, or by abstraction

of timely requirements demanded by the view or imposed by the execution platform. Solving constraints and committing progressively to particular schedule results from the intended flow of design promoted by MARTE in model-based engineering of embedded systems.

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