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On Infinitary Rational Relations and Borel Sets

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Abstract. We prove in this paper that there exists some infinitary rational relations which are Σ_3^0 -complete Borel sets and some others which are Π_3^0 -complete. These results give additional answers to questions of Simonnet [Sim92] and of Lescow and Thomas [Tho90] [LT94].

Keywords: infinitary rational relations; topological properties; Borel sets.

1 Introduction

Acceptance of infinite words by finite automata was firstly considered in the sixties by Büchi in order to study decidability of the monadic second order theory of one successor over the integers [Büc62]. Then the so called ω -regular languages have been intensively studied and many applications have been found, see [Tho90] [Sta97] [PP01] for many results and references.

Since then many extensions of ω -regular languages have been investigated as the classes of ω -languages accepted by pushdown automata, Petri nets, Turing machines, see [Tho90] [EH93] [Sta97] for a survey of this work.

On the other side rational relations on finite words were studied in the sixties and played a fundamental role in the study of families of context free languages [Ber79]. Investigations on their extension to rational relations on infinite words were carried out or mentioned in the books [BT70] [LS77]. Gire and Nivat studied infinitary rational relations in [Gir81] [GN84]. Infinitary rational relations are subsets of $\Sigma_1^\omega \times \Sigma_2^\omega$, where Σ_1 and Σ_2 are finite alphabets, which are recognized by Büchi transducers or by 2-tape finite Büchi automata with asynchronous reading heads. So the class of infinitary rational relations extends both the class of finitary rational relations **and** the class of ω -regular languages.

They have been much studied, in particular in connection with the rational functions they may define, see for example [CG99] [BCPS00] [Sim92] [Sta97] [Pri00] for many results and references. Notice that a rational relation $R \subseteq \Sigma_1^\omega \times \Sigma_2^\omega$ may be seen as an ω -language over the alphabet $\Sigma_1 \times \Sigma_2$.

A way to study the complexity of languages of infinite words accepted by finite machines is to study their topological complexity and firstly to locate them with regard to the Borel and the projective hierarchies. This work is analysed for

example in [Sta86] [Tho90] [EH93] [LT94] [Sta97]. It is well known that every ω -language accepted by a Turing machine with a Büchi or Muller acceptance condition is an analytic set and that ω -regular languages are boolean combinations of Π_2^0 -sets hence Δ_3^0 -sets, [Sta97] [PP01].

The question of the topological complexity of relations on infinite words also naturally arises and is asked by Simonnet in [Sim92]. It is also posed in a more general form by Lescow and Thomas in [LT94] (for infinite labelled partial orders) and in [Tho89] where Thomas suggested to study reducibility notions and associated completeness results.

Every infinitary rational relation is an analytic set. We showed in [Fin01b] that there exist some infinitary rational relations which are analytic but non Borel sets. Considering Borel infinitary rational relations we prove in this paper that there exist some infinitary rational relations which are Σ_3^0 -complete Borel sets and some others which are Π_3^0 -complete. This implies that there exist also some infinitary rational relations which are Δ_4^0 -sets but not $(\Sigma_3^0 \cup \Pi_3^0)$ -sets.

These results may be compared with examples of Σ_3^0 -complete ω -languages accepted by deterministic pushdown automata with the acceptance condition: “some stack content appears infinitely often during an infinite run”, given by Cachat, Duparc, and Thomas in [CDT02] or with examples of Σ_n^0 -complete and Π_n^0 -complete ω -languages, $n \geq 1$, accepted by non-deterministic pushdown automata with Büchi acceptance condition given in [Fin01a].

The paper is organized as follows. In section 2 we introduce the notion of transducers and of infinitary rational relations. In section 3 we recall definitions of Borel sets, and we prove our main results in sections 4 and 5.

2 Infinitary rational relations

Let Σ be a finite alphabet whose elements are called letters. A non-empty finite word over Σ is a finite sequence of letters: $x = a_1 a_2 \dots a_n$ where for all i in $[1; n]$, $a_i \in \Sigma$. We shall denote $x(i) = a_i$ the i^{th} letter of x and $x[i] = x(1) \dots x(i)$ for $i \leq n$. The length of x is $|x| = n$. The empty word will be denoted by λ and has 0 letter. Its length is 0. The set of finite words over Σ is denoted Σ^* . A (finitary) language L over Σ is a subset of Σ^* . The usual concatenation product of u and v will be denoted by $u.v$ or just uv . For $V \subseteq \Sigma^*$, we denote $V^* = \{v_1 \dots v_n \mid n \in \mathbb{N} \text{ and } \forall i \in [1; n] \ v_i \in V\}$.

The first infinite ordinal is ω . An ω -word over Σ is an ω -sequence $a_1 a_2 \dots a_n \dots$, where for all $i \geq 1$, $a_i \in \Sigma$. When σ is an ω -word over Σ , we write $\sigma = \sigma(1)\sigma(2)\dots\sigma(n)\dots$ and $\sigma[n] = \sigma(1)\sigma(2)\dots\sigma(n)$ the finite word of length n , prefix of σ . The set of ω -words over the alphabet Σ is denoted by Σ^ω . An ω -language over an alphabet Σ is a subset of Σ^ω . For $V \subseteq \Sigma^*$, $V^\omega = \{\sigma = u_1 \dots u_n \dots \in \Sigma^\omega \mid \forall i \geq 1 \ u_i \in V\}$ is the ω -power of V . The concatenation product is extended to the product of a finite word u and an ω -word v : the infinite word $u.v$ is then the ω -word such that: $(u.v)(k) = u(k)$ if $k \leq |u|$, and $(u.v)(k) = v(k - |u|)$ if $k > |u|$.

If A is a subset of B we shall denote $A^- = B - A$ the complement of A (in B).

We assume the reader to be familiar with the theory of formal languages and of ω -regular languages. We recall that ω -regular languages form the class of ω -languages accepted by finite automata with a Büchi acceptance condition and this class is the omega Kleene closure of the class of regular finitary languages. We are going now to introduce the notion of infinitary rational relation via definition of Büchi transducers:

Definition 1. A Büchi transducer is a sextuple $\mathcal{T} = (K, \Sigma, \Gamma, \Delta, q_0, F)$, where K is a finite set of states, Σ and Γ are finite sets called the input and the output alphabets, Δ is a finite subset of $K \times \Sigma^* \times \Gamma^* \times K$ called the set of transitions, q_0 is the initial state, and $F \subseteq K$ is the set of accepting states.

A computation \mathcal{C} of the transducer \mathcal{T} is an infinite sequence of transitions

$$(q_0, u_1, v_1, q_1), (q_1, u_2, v_2, q_2), \dots, (q_{i-1}, u_i, v_i, q_i), (q_i, u_{i+1}, v_{i+1}, q_{i+1}), \dots$$

The computation is said to be successful iff there exists a final state $q_f \in F$ and infinitely many integers $i \geq 0$ such that $q_i = q_f$.

The input word of the computation is $u = u_1.u_2.u_3 \dots$

The output word of the computation is $v = v_1.v_2.v_3 \dots$

Then the input and the output words may be finite or infinite.

The infinitary rational relation $R(\mathcal{T}) \subseteq \Sigma^\omega \times \Gamma^\omega$ recognized by the Büchi transducer \mathcal{T} is the set of pairs $(u, v) \in \Sigma^\omega \times \Gamma^\omega$ such that u and v are the input and the output words of some successful computation \mathcal{C} of \mathcal{T} .

The set of infinitary rational relations will be denoted RAT.

Remark 2. An infinitary rational relation is a subset of $\Sigma^\omega \times \Gamma^\omega$ for two finite alphabets Σ and Γ . One can also consider that it is an ω -language over the finite alphabet $\Sigma \times \Gamma$. If $(u, v) \in \Sigma^\omega \times \Gamma^\omega$, one can consider this pair of infinite words as a single infinite word $(u(1), v(1)).(u(2), v(2)).(u(3), v(3)) \dots$ over the alphabet $\Sigma \times \Gamma$. We shall use this fact to investigate the topological complexity of infinitary rational relations.

3 Borel sets

We assume the reader to be familiar with basic notions of topology which may be found in [Mos80] [Kec95] [LT94] [Sta97] [PP01].

For a finite alphabet X we shall consider X^ω as a topological space with the Cantor topology. The open sets of X^ω are the sets in the form $W.X^\omega$, where $W \subseteq X^*$. A set $L \subseteq X^\omega$ is a closed set iff its complement $X^\omega - L$ is an open set. Define now the next classes of the Hierarchy of Borel sets of finite ranks:

Definition 3. The classes Σ_n^0 and Π_n^0 of the Borel Hierarchy on the topological space X^ω are defined as follows:

Σ_1^0 is the class of open sets of X^ω .

Π_1^0 is the class of closed sets of X^ω .

And for any integer $n \geq 1$:

Σ_{n+1}^0 is the class of countable unions of Π_n^0 -subsets of X^ω .

Π_{n+1}^0 is the class of countable intersections of Σ_n^0 -subsets of X^ω .

The Borel Hierarchy is also defined for transfinite levels, but we shall not need them in the present study. There are also some subsets of X^ω which are not Borel. In particular the class of Borel subsets of X^ω is strictly included into the class Σ_1^1 of analytic sets which are obtained by projection of Borel sets, see for example [Sta97] [LT94] [PP01] [Kec95] for more details.

Recall also the notion of completeness with regard to reduction by continuous functions. For an integer $n \geq 1$, a set $F \subseteq X^\omega$ is said to be a Σ_n^0 (respectively, Π_n^0, Σ_1^1)-complete set iff for any set $E \subseteq Y^\omega$ (with Y a finite alphabet): $E \in \Sigma_n^0$ (respectively, $E \in \Pi_n^0, E \in \Sigma_1^1$) iff there exists a continuous function $f : Y^\omega \rightarrow X^\omega$ such that $E = f^{-1}(F)$.

A Σ_n^0 (respectively, Π_n^0, Σ_1^1)-complete set is a Σ_n^0 (respectively, Π_n^0, Σ_1^1)-set which is in some sense a set of the highest topological complexity among the Σ_n^0 (respectively, Π_n^0, Σ_1^1)-sets. Σ_n^0 (respectively, Π_n^0)-complete sets, with n an integer ≥ 1 , are thoroughly characterized in [Sta86].

Example 4. Let $\Sigma = \{0, 1\}$ and $\mathcal{A} = (0^*1)^\omega \subseteq \Sigma^\omega$. \mathcal{A} is the set of ω -words over the alphabet Σ with infinitely many occurrences of the letter 1. It is well known that \mathcal{A} is a Π_2^0 -complete set and its complement \mathcal{A}^- is a Σ_2^0 -complete set: it is the set of ω -words over $\{0, 1\}$ having only a finite number of occurrences of letter 1.

4 Σ_3^0 -complete infinitary rational relations

We can now state the following result:

Theorem 5. *There exist some Σ_3^0 -complete infinitary rational relations.*

Proof. We shall use a well known example of Σ_3^0 -complete set which is a subset of the topological space Σ^{ω^2} .

The set Σ^{ω^2} is the set of ω^2 -words over the finite alphabet Σ . It may also be viewed as the set of (infinite) $(\omega \times \omega)$ -matrices whose coefficients are letters of Σ . If $x \in \Sigma^{\omega^2}$ we shall write $x = (x(m, n))_{m \geq 1, n \geq 1}$. The infinite word $x(m, 1)x(m, 2) \dots x(m, n) \dots$ will be called the m^{th} column of the ω^2 -word x and the infinite word $x(1, n)x(2, n) \dots x(m, n) \dots$ will be called the n^{th} row of the ω^2 -word x . Thus an element of Σ^{ω^2} is completely determined by the (infinite) set of its columns or of its rows.

The set Σ^{ω^2} is usually equipped with the product topology of the discrete topology on Σ (for which every subset of Σ is an open set), see [Kec95] [PP01].

This topology may be defined by the following distance d . Let x and y be two ω^2 -words in Σ^{ω^2} such that $x \neq y$, then

$$d(x, y) = \frac{1}{2^n} \quad \text{where}$$

$$n = \min\{p \geq 1 \mid \exists(i, j) \ x(i, j) \neq y(i, j) \text{ and } i + j = p\}$$

Then the topological space Σ^{ω^2} is homeomorphic to the topological space Σ^ω , equipped with the Cantor topology. Borel subsets of Σ^{ω^2} are defined from open subsets in the same manner as in the case of the topological space Σ^ω .

Σ_n^0 (respectively Π_n^0)-complete sets are also defined in a similar way.

Recall now that the set

$$S = \{x \in \{0, 1\}^{\omega^2} \mid \exists m \exists^\infty n \ x(m, n) = 1\}$$

where \exists^∞ means "there exist infinitely many", is a Σ_3^0 -complete subset of $\{0, 1\}^{\omega^2}$, [Kec95, p. 179]. It is the set of ω^2 -words over $\{0, 1\}$ having at least one column in the Π_2^0 -complete subset \mathcal{A} of $\{0, 1\}^\omega$ given in Example 4.

In order to use this example we shall firstly define a coding of ω^2 -words over Σ by ω -words over the alphabet $(\Sigma \cup \{A\}) \times (\Sigma \cup \{A\})$ where A is a new letter not in Σ . The code of $x \in \Sigma^{\omega^2}$ may be written in the form (σ_1, σ_2) , where σ_1 and σ_2 are ω -words over the alphabet $(\Sigma \cup \{A\})$. In order to describe the ω -words σ_1 and σ_2 let us call, for $x \in \Sigma^{\omega^2}$ and p an integer ≥ 2 :

$$T_{p+1}^x = \{x(p, 1), x(p-1, 2), \dots, x(2, p-1), x(1, p)\}$$

the set of elements $x(m, n)$ with $m + n = p + 1$.

The word σ_1 begins with $x(1, 1)$ followed by a letter A ; then the word σ_1 enumerates the elements of the sets T_{p+1}^x for p an **even integer**. More precisely for every even integer $p \geq 2$ the elements of T_{p+1}^x are placed before those of T_{p+3}^x and these two sets of letters are separated by an A . Moreover for each even integer $p \geq 2$ the elements of T_{p+1}^x are enumerated in the following order:

$$x(p, 1), x(p-1, 2), \dots, x(2, p-1), x(1, p)$$

The construction of σ_2 is very similar but with some modifications. It enumerates the elements of the sets T_{p+1}^x for p an **odd integer**. More precisely the word σ_2 begins with an A then for every odd integer $p \geq 3$ the elements of T_{p+1}^x are placed before those of T_{p+3}^x and these two sets of letters are separated by an A . Moreover for each odd integer $p \geq 3$ the elements of T_{p+1}^x are enumerated in the following order:

$$x(1, p), x(2, p-1), \dots, x(p-1, 2), x(p, 1)$$

Then the ω -word σ_1 and the ω -word σ_2 are in the following form:

$$\sigma_1 = x(1, 1).A.x(2, 1)x(1, 2).A.x(4, 1)x(3, 2)x(2, 3)x(1, 4).A.x(6, 1)x(5, 2) \dots$$

$$\sigma_2 = A.x(1, 3)x(2, 2)x(3, 1).A.x(1, 5)x(2, 4)x(3, 3)x(4, 2)x(5, 1).A.x(1, 7)x(2, 6) \dots$$

Let then h be the mapping from Σ^{ω^2} into $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$ such that, for every ω^2 -word x over the alphabet Σ , $h(x)$ is the code (σ_1, σ_2) of the ω^2 -word as defined above. It is easy to see, from the definition of h and of the order of the enumeration of letters $x(m, n)$ in the code of $x \in \Sigma^{\omega^2}$ (they are enumerated for increasing values of $m + n$), that h is a continuous function from Σ^{ω^2} into $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$.

Notice that we have chosen $x(1, 1)$ to be the first letter of the word σ_1 . In fact, with slight modifications in the sequel, we could have chosen the letter $x(1, 1)$ to be the first letter of the word σ_2 .

Remark that the above coding of ω^2 -words resembles the use of the Cantor pairing function as it was used to construct the complete sets P_i and S_i in [SW78] (see also [Sta86] or [Sta97, section 3.4]).

We now state the following lemmas:

Lemma 6. *Let $\Sigma = \{0, 1\}$ and $S = \{x \in \{0, 1\}^{\omega^2} / \exists m \exists^\infty n x(m, n) = 1\}$. Then the set*

$$\mathcal{S} = h(S) \cup (h(\Sigma^{\omega^2}))^-$$

is a Σ_3^0 -complete subset of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$.

Proof. We just sketch the first part of the proof.

The topological space Σ^{ω^2} is compact and the function h is continuous and injective. Using these facts we can easily show that $h(S)$ is a Σ_3^0 -subset of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$.

On the other side $h(\Sigma^{\omega^2})$ is a closed subset of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$. Then its complement

$$(h(\Sigma^{\omega^2}))^- = ((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega - h(\Sigma^{\omega^2})$$

is an open (i.e. a Σ_1^0) subset of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$.

Now $\mathcal{S} = h(S) \cup (h(\Sigma^{\omega^2}))^-$ is the union of a Σ_3^0 -set and of a Σ_1^0 -set therefore it is a Σ_3^0 -set because the class of Σ_3^0 -subsets of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$ is closed under finite union.

In order to prove that \mathcal{S} is Σ_3^0 -complete it suffices to remark that $S = h^{-1}(\mathcal{S})$. This implies that \mathcal{S} is Σ_3^0 -complete because S is Σ_3^0 -complete. \square

Lemma 7. *For Σ a finite alphabet,*

$$(h(\Sigma^{\omega^2}))^- = ((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega - h(\Sigma^{\omega^2})$$

is an infinitary rational relation.

Proof. Return to the definition of coding of ω^2 -words over Σ by ω -words over the alphabet $(\Sigma \cup \{A\}) \times (\Sigma \cup \{A\})$. The code of $x \in \Sigma^{\omega^2}$ was written in the form (σ_1, σ_2) , where σ_1 and σ_2 are ω -words over the alphabet $(\Sigma \cup \{A\})$ in the form:

$$\sigma_1 = u_1.A.u_2.A.u_4.A.u_6.A.u_8.A \dots A.u_{2n}.A \dots$$

$$\sigma_2 = A.u_3.A.u_5.A.u_7.A.u_9.A \dots A.u_{2n+1}.A \dots$$

where for all integers $i \geq 1$, $u_i \in \Sigma^*$ and $|u_i| = i$.

It is now easy to see that the complement of the set $h(\Sigma^{\omega^2})$ of codes of ω^2 -words over Σ is the union of the sets \mathcal{C}_j where:

- $\mathcal{C}_1 = \{(\sigma_1, \sigma_2) / \sigma_1, \sigma_2 \in (\Sigma \cup \{A\})^\omega \text{ and } (\sigma_1 \in \mathcal{B} \text{ or } \sigma_2 \in \mathcal{B})\}$
where \mathcal{B} is the set of ω -words over $(\Sigma \cup \{A\})$ having only a finite number of letters A .
- \mathcal{C}_2 is formed by pairs (σ_1, σ_2) where
 σ_1 has not an initial segment in $\Sigma.A.\Sigma^2.A$ or
the first letter of σ_2 is not an A .
- \mathcal{C}_3 is formed by pairs (σ_1, σ_2) where
 $\sigma_1 = w_1.A.w_2.A.w_3.A.w_4 \dots A.w_n.A.u.A.z_1$
 $\sigma_2 = w'_1.A.w'_2.A.w'_3.A.w'_4 \dots A.w'_n.A.v.A.z_2$

where n is an integer ≥ 1 , for all $i \leq n$ $w_i, w'_i \in \Sigma^*$, $z_1, z_2 \in (\Sigma \cup \{A\})^\omega$ and

$$u, v \in \Sigma^* \text{ and } |v| \neq |u| + 1$$

- \mathcal{C}_4 is formed by pairs (σ_1, σ_2) where
 $\sigma_1 = w_1.A.w_2.A.w_3.A.w_4 \dots A.w_n.A.w_{n+1}.A.v.A.z_1$
 $\sigma_2 = w'_1.A.w'_2.A.w'_3.A.w'_4 \dots A.w'_n.A.u.A.z_2$

where n is an integer ≥ 1 , for all $i \leq n$ $w_i, w'_i \in \Sigma^*$, $w_{n+1} \in \Sigma^*$, $z_1, z_2 \in (\Sigma \cup \{A\})^\omega$ and

$$u, v \in \Sigma^* \text{ and } |v| \neq |u| + 1$$

Each set \mathcal{C}_j , $1 \leq j \leq 4$, is easily seen to be an infinitary rational relation $\subseteq (\Sigma \cup \{A\})^\omega \times (\Sigma \cup \{A\})^\omega$ (the detailed proof is left to the reader).

The class of infinitary rational relations is closed under finite union; this follows from the fact that they are recognized by **non deterministic** Büchi transducers.

Then

$$(h(\Sigma^{\omega^2}))^- = \bigcup_{1 \leq j \leq 4} \mathcal{C}_j$$

is an infinitary rational relation. □

We cannot show directly that $h(S) \in RAT$ so we are now looking for a rational relation $R \subseteq ((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$ (with $\Sigma = \{0, 1\}$) such that for every ω^2 -word $x \in \Sigma^{\omega^2}$, $h(x) \in R$ if and only if $x \in S$. Then we shall have $S = h^{-1}(R)$.

We shall first describe the relation R which is an ω -language over the alphabet $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))$. Every word of R may be seen as a pair $y = (y_1, y_2)$ of ω -words over the alphabet $\Sigma \cup \{A\}$ and then y is in R if and only if it is in the form

$$\begin{aligned} y_1 &= U_k.u.t(1).v_1.A.g_1.t(3).v_2.A.g_2.t(5) \dots A.g_n.t(2n+1).v_{n+1}.A \dots \\ y_2 &= V_k.u_1.t(2).z_1.A.u_2.t(4).z_2.A \dots A.u_n.t(2n).z_n.A \dots \end{aligned}$$

where k is an integer ≥ 1 , $U_k, V_k \in (\Sigma^*.A)^k$, $u = \lambda$ or $u \in \Sigma$, and for all integers $i \geq 1$, $t(i) \in \Sigma$ and $u_i, v_i, g_i, z_i \in \Sigma^*$ and

$$|v_i| = |u_i| \quad \text{and} \quad |g_i| = |z_i| + 1$$

and the ω -word $t = t(1)t(2) \dots t(n) \dots$ is in the ω -regular language \mathcal{A} given in Example 4.

Lemma 8. *The above defined relation R satisfies $S = h^{-1}(R)$, i.e.:*

$$\forall x \in \Sigma^{\omega^2} \quad h(x) \in R \iff x \in S$$

Proof. Assume first that such an $y = (y_1, y_2) \in R$ is the code $h(x)$ of an ω^2 -word $x \in \Sigma^{\omega^2}$. Then

$$u.t(1).v_1 = x(2k, 1).x(2k-1, 2) \dots x(1, 2k)$$

so if $u = \lambda$ then $x(2k, 1) = t(1)$ and $|v_1| = 2k-1$.

And if $u \in \Sigma$ then $x(2k-1, 2) = t(1)$ and $|v_1| = 2k-2$.

Next

$$u_1.t(2).z_1 = x(1, 2k+1).x(2, 2k) \dots x(2k+1, 1)$$

so if $u = \lambda$ then $|u_1| = |v_1| = 2k-1$ thus $x(2k, 2) = t(2)$ and $|z_1| = 1$.

And if $u \in \Sigma$ then $|u_1| = |v_1| = 2k-2$ thus $x(2k-1, 3) = t(2)$ and $|z_1| = 2$.

Moreover

$$g_1.t(3).v_2 = x(2k+2, 1).x(2k+1, 2) \dots x(1, 2k+2)$$

so if $u = \lambda$ then $|g_1| = |z_1| + 1 = 2$ and $t(3) = x(2k, 3)$.

And if $u \in \Sigma$ then $|g_1| = |z_1| + 1 = 3$ and $t(3) = x(2k-1, 4)$.

In a similar manner one can show by induction on integers i that if $u = \lambda$ letters $t(i)$ are successive letters of the $(2k)^{th}$ column of x and if $u \in \Sigma$ letters $t(i)$ are successive letters of the $(2k-1)^{th}$ column of x .

So assume that for some integer $i \geq 1$:
if $u = \lambda$ then $t(2i + 1) = x(2k, 2i + 1)$, and
if $u \in \Sigma$ then $t(2i + 1) = x(2k - 1, 2i + 2)$.

We know that

$$g_i.t(2i + 1).v_{i+1} = x(2k + 2i, 1).x(2k + 2i - 1, 2) \dots x(1, 2k + 2i)$$

thus $u = \lambda$ implies that $|v_{i+1}| = 2k - 1$ and $u \in \Sigma$ implies that $|v_{i+1}| = 2k - 2$.
But it holds also that

$$u_{i+1}.t(2i + 2).z_{i+1} = x(1, 2k + 2i + 1).x(2, 2k + 2i) \dots x(2k + 2i + 1, 1)$$

So if $u = \lambda$ then $|u_{i+1}| = |v_{i+1}| = 2k - 1$ and $t(2i + 2) = x(2k, 2i + 2)$ and
 $|z_{i+1}| = 2i + 1$.

And if $u \in \Sigma$ then $|u_{i+1}| = |v_{i+1}| = 2k - 2$ and $t(2i + 2) = x(2k - 1, 2i + 3)$ and
 $|z_{i+1}| = 2i + 2$.

Next

$$g_{i+1}.t(2i + 3).v_{i+2} = x(2k + 2i + 2, 1).x(2k + 2i + 1, 2) \dots x(1, 2k + 2i + 2)$$

So if $u = \lambda$ then $|g_{i+1}| = |z_{i+1}| + 1 = 2i + 2$ and $t(2i + 3) = x(2k, 2i + 3)$.

And if $u \in \Sigma$ then $|g_{i+1}| = |z_{i+1}| + 1 = 2i + 3$ and $t(2i + 3) = x(2k - 1, 2i + 4)$.

We have then proved by induction that if $u = \lambda$,

$$t = t(1).t(2) \dots t(n) \dots = x(2k, 1).x(2k, 2) \dots x(2k, n) \dots$$

and if $u \in \Sigma$,

$$t = t(1).t(2) \dots t(n) \dots = x(2k - 1, 2).x(2k - 1, 3) \dots x(2k - 1, n) \dots$$

Notice that in this second case the ω -word t begins with the second letter $x(2k - 1, 2)$ of the $(2k - 1)^{th}$ column of x and not with the first letter $x(2k - 1, 1)$ of this column. But this will not change the fact that $t \in \mathcal{A}$ or $t \notin \mathcal{A}$ because \mathcal{A} is simply the set of ω -words over the alphabet $\{0, 1\}$ with infinitely many occurrences of the letter 1.

Thus if a code $h(x)$ of an ω^2 -word $x \in \Sigma^{\omega^2}$ is in R then x has a column in \mathcal{A} , i.e. $x \in S$. Conversely it is easy to see that every code $h(x)$ of $x \in S$ may be written in the above form (y_1, y_2) of a word in R . Then we have proved that $S = h^{-1}(R)$. \square

Remark that the *non determinism* of a transducer recognizing R (or of a 2-tape finite automaton accepting R) will be used to guess the integer k and whether $u = \lambda$ or $u \in \Sigma$.

Intuitively, if \mathcal{T} is such a transducer recognizing R then, during a successful computation accepting the code $h(x)$ of an ω^2 -word x in Σ^{ω^2} , the *non determinism* of \mathcal{T} is used to guess a column of the ω^2 -word x in order to simulate on this

column the behaviour of a finite Büchi automaton accepting the Π_2^0 -complete ω -regular language \mathcal{A} .

Lemma 9. *The above defined relation R is an infinitary rational relation.*

Proof. This is easy to see from the definitions of R and of an infinitary rational relation. The infinitary rational relation R is recognized by the following Büchi transducer $\mathcal{T} = (K, (\{0, 1, A\}), (\{0, 1, A\}), \Delta, q_0, F)$, where

$$K = \{q_0, q_1, q_2, q_3, q_1^0, q_1^1, q_2^0, q_2^1\}$$

is a finite set of states, $\{0, 1, A\} = \Sigma \cup \{A\}$ is the input *and* the output alphabet (with $\Sigma = \{0, 1\}$), q_0 is the initial state, and $F = \{q_1^1, q_2^1\}$ is the set of accepting states. Moreover $\Delta \subseteq K \times (\Sigma \cup \{A\})^* \times (\Sigma \cup \{A\})^* \times K$ is the finite set of transitions, containing the following transitions:

(q_0, λ, a, q_0) and (q_0, a, λ, q_0) , for all $a \in \Sigma$,
 (q_0, A, A, q_0) ,
 (q_0, A, A, q_1) ,
 (q_1, u, λ, q_2) , for all $u \in \Sigma \cup \Sigma^2$,
 (q_2, a, b, q_2) , for all $a, b \in \Sigma$,
 $(q_2, A, 0, q_1^0)$ and $(q_2, A, 1, q_1^1)$,
 (q, a, λ, q_3) , for all $a \in \Sigma$ and $q \in \{q_1^0, q_1^1\}$,
 (q_3, a, b, q_3) , for all $a, b \in \Sigma$,
 $(q_3, 0, A, q_2^0)$ and $(q_3, 1, A, q_2^1)$,
 $(q, \lambda, \lambda, q_2)$, for all $q \in \{q_2^0, q_2^1\}$. □

Return now to the proof of Theorem 5 and consider the set $\mathcal{R} = R \cup (h(\Sigma^{\omega^2}))^-$. It turns out that

$$\mathcal{R} = \mathcal{S} = h(S) \cup (h(\Sigma^{\omega^2}))^-$$

because $S = h^{-1}(R)$. But we have proved that $(h(\Sigma^{\omega^2}))^-$ and R are infinitary rational relations thus $\mathcal{R} = R \cup (h(\Sigma^{\omega^2}))^-$ is the union of two infinitary rational relations hence $\mathcal{R} \in RAT$. Lemma 6 asserts that $\mathcal{R} = \mathcal{S}$ is a Σ_3^0 -complete subset of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$ and this ends the proof. □

Remark 10. *With a slight modification we could have replaced the set S by the set of ω^2 -words over Σ having at least one column in a given Π_2^0 -complete ω -regular language.*

5 Π_3^0 -complete infinitary rational relations

We can now state our next result:

Theorem 11. *There exist some Π_3^0 -complete infinitary rational relations.*

Proof. We are going to sketch the proof but we cannot give here all details because of limited space for this paper.

As in the last section, we shall use a well known example of Π_3^0 -complete set which is a subset of the topological space Σ^{ω^2} with $\Sigma = \{0, 1\}$.

Recall that the set

$$P = \{x \in \{0, 1\}^{\omega^2} / \forall m \exists^{<\infty} n \ x(m, n) = 1\}$$

where $\exists^{<\infty}$ means "there exist only finitely many", is a Π_3^0 -complete subset of $\{0, 1\}^{\omega^2}$, [Kec95, p. 179]. $P = \{0, 1\}^{\omega^2} - S$ so " P is Π_3^0 -complete" follows directly from " S is Σ_3^0 -complete".

P is the set of ω^2 -words having all their columns in the Σ_2^0 -complete subset \mathcal{A}^- of $\{0, 1\}^\omega$ where \mathcal{A} is the Π_2^0 -complete ω -regular language given in Example 4.

We shall use the same coding $h : x \rightarrow h(x)$ for ω^2 -words over the alphabet $\Sigma = \{0, 1\}$ as in preceding section.

Lemma 12.

$$\mathcal{P} = h(P) \cup (h(\Sigma^{\omega^2}))^-$$

is a Π_3^0 -complete subset of $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))^\omega$.

Proof. It is similar to proof of Lemma 6. □

We are going to find an infinitary rational relation R_1 such that $P = h^{-1}(R_1)$. We define now the relation R_1 . It is an ω -language over the alphabet $((\Sigma \cup \{A\}) \times (\Sigma \cup \{A\}))$. Every word of R_1 may be seen as a pair $y = (y_1, y_2)$ of ω -words over the alphabet $\Sigma \cup \{A\}$ and then y is in R_1 if and only if it is in the form

$$\begin{aligned} y_1 &= U_k.u.t(1).v_1.A.g_1.t(3).v_2.A.g_2.t(5) \dots A.g_n.t(2n+1).v_{n+1}.A \dots \\ y_2 &= V_k.u_1.t(2).z_1.A.u_2.t(4).z_2.A \dots A.u_n.t(2n).z_n.A \dots \end{aligned}$$

where k is an integer ≥ 1 , $U_k, V_k \in (\Sigma^*.A)^k$, $u \in \Sigma^*$, and for all integers $i \geq 1$, $t(i) \in \Sigma$ and

$$u_i, v_i \in 0^* \text{ and } g_i, z_i \in \Sigma^* \text{ and}$$

$$|v_i| = |u_i| \quad \text{and} \quad [|g_i| = |z_i| + 1 \text{ or } |g_i| = |z_i|]$$

and there exist infinitely many integers i such that $|g_i| = |z_i|$.

Lemma 13. *The above defined relation R_1 satisfies $P = h^{-1}(R_1)$, i.e.:*

$$\forall x \in \Sigma^{\omega^2} \quad h(x) \in R_1 \iff x \in P$$

Lemma 14. *The above defined relation R_1 is an infinitary rational relation.*

Return to the proof of theorem 11. By Lemma 13 the infinitary relation R_1 satisfies $P = h^{-1}(R_1)$ thus we shall have

$$\mathcal{P} = h(P) \cup (h(\Sigma^{\omega^2}))^- = R_1 \cup (h(\Sigma^{\omega^2}))^-$$

But by Lemma 14 R_1 is rational hence \mathcal{P} is the union of two infinitary rational relations thus $\mathcal{P} \in RAT$ and is Π_3^0 -complete by Lemma 12. \square

From Theorems 5 and 11 we can now easily infer the following result:

Corollary 15. *There exists some Δ_4^0 (i.e. $\Sigma_4^0 \cap \Pi_4^0$) infinitary rational relations which are not in $(\Sigma_3^0 \cup \Pi_3^0)$.*

The question naturally arises whether there exist some infinitary rational relations $R \subseteq \Sigma^\omega \times \Gamma^\omega$ which are Σ_4^0 -complete or Π_4^0 -complete or even higher in the Borel hierarchy.

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