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Fully Dynamic Maintenance of Arc-Flags in Road Networks

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Abstract. The problem of finding best routes in road networks can be solved by applying Dijkstra's shortest paths algorithm. Unfortunately, road networks deriving from real-world applications are huge yielding unsustainable times to compute shortest paths. For this reason, great research efforts have been done to accelerate Dijkstra's algorithm on road networks. These efforts have led to the development of a number of *speed-up techniques*, as for example *Arc-Flags*, whose aim is to compute additional data in a preprocessing phase in order to accelerate the shortest paths queries in an on-line phase. The main drawback of most of these techniques is that they do not work well in dynamic scenarios. In this paper we propose a new algorithm to update the Arc-Flags of a graph subject to edge weight decrease operations. To check the practical performances of the new algorithm we experimentally analyze it, along with a previously known algorithm for edge weight increase operations, on real-world road networks subject to fully dynamic sequences of operations. Our experiments show a significant speed-up in the updating phase of the Arc-Flags, at the cost of a small space and time overhead in the preprocessing phase.

1 Introduction

The problem of finding best connections in transportation networks has received a lot of attention in the last years. If travel times are assigned to the edges of the graph representing the network, this problem can be easily solved by applying Dijkstra's algorithm to find the shortest path between two points. Unfortunately, transportation networks deriving from real-world applications tend to be huge yielding unsustainable times to compute shortest paths. For this reason, great research efforts have been done over the last decade to accelerate Dijkstra's algorithm on typical instances of transportation networks, such as road or railway networks (see [5] for a recent overview). These research efforts have led to the development of a number of *speed-up techniques*, whose aim is to compute additional data in a preprocessing phase in order to accelerate the shortest paths queries during an on-line phase. However, most of the speed-up techniques developed in the literature do not work well in dynamic scenarios, when edge weights

changes occur to the network due to traffic jams or delays of trains. In other words, the correctness of these speed-up techniques relies on the fact that the network does not change between two queries. Unfortunately, such situations arise frequently in practice. In order to keep the shortest paths queries correct, the preprocessed data need to be updated. The easiest way is to recompute the preprocessed data from scratch after each change to the network. This is in general unfeasible since even the fastest methods need too much time. In general, the typical update operations that can occur on a network can be modelled as insertions and deletions of edges and edge weight changes (weight decrease and weight increase). When arbitrary sequences of the above operations are allowed, we refer to the *fully dynamic problem*, otherwise we refer to the *partially dynamic problem*; if only insertions and weight decrease (deletion and weight increase, respectively) operations are allowed, then the partially dynamic problem is known as the *incremental* (*decremental*, respectively) problem.

Related Works. We refer here only to papers on the dynamic case and refer to [5] as a survey for the static case. A number of efforts have been done in the last years to accelerate the computation of shortest paths in dynamic scenarios [1–4, 6, 13, 15]. The first of these techniques was Geometric Containers [15], whose key idea is to allow suboptimal containers after a few updates. However, this approach yields a loss in query performance. The same holds for the dynamic variant of Arc-Flags [9] proposed in [2], where, after a number of updates, the query performances get worse yielding only a low speed-up over Dijkstra’s algorithm. In [13] the authors combine ideas from highway hierarchies [12] and overlay graphs [14] yielding very good query times in dynamic road networks. The ALT algorithm, introduced in [7] works considerably well in dynamic scenarios where edge weights can increase their value. Also in this case, query performances get worse if too many edges weights change [6]. Summarizing, all above techniques work in a dynamic scenario as long as the number of updates is small. As soon as the number of updates is greater than a certain value, it is better to repeat the preprocessing from scratch. To our knowledge, the only other dynamic technique known in the literature with no loss in query performance is that in [13]. In [4], a very practically efficient algorithm has been given to compute shortest paths in continental road graphs with arbitrary metrics, whose efficiency is also due to the use of parallelism. This algorithm is fast enough to be used in dynamic scenarios for the recomputation from scratch of shortest paths. Recently, a data structure named *Road-Signs* has been introduced in [3] to compute and update the Arc-Flags of a graphs. In detail, in [3] the authors define an algorithm to preprocess Road-Signs and a decremental algorithm to update them each time that a *weight increase* operation occurs on an edge of the graph. As the updating algorithm is able to correctly update Arc-Flags, there is no loss in query performance. They also experimentally analyze this algorithm in real-world road networks showing that it yields a significant speed-up in the updating phase of Arc-Flags with respect to the recomputation from-scratch, at the price of a small space and time overhead in the preprocessing phase. However, the solution in [3] is not able to update Arc-Flags in a fully dynamic scenario.

Contribution of the paper. We propose a new *incremental* algorithm which is able to update the Arc-Flags of a graph by updating Road-Signs, during a sequence of *weight decrease* operations. Since the new incremental algorithm uses the same data structures of the decremental solution of [3], then the combination of these two solutions can be used to update Arc-Flags in a fully dynamic scenario. To check the practical usefulness of this combination we implemented the two algorithms and performed an extensive experimental study against the recomputation from scratch of Arc-Flags on fully dynamic sequences of weight increase and weight decrease operations on real world road networks. The results of our experiments can be summarized as follows: in comparison to the recomputation from-scratch of Arc-Flags, we obtained a significant speed-up in the updating phase, at the cost of a small space and time overhead in the preprocessing phase. In detail, we experimentally show that the fully dynamic algorithm is able to update the Arc-Flags between 40 and 431 times faster than the recomputation from scratch in average. However, in order to compute and store the Road-Signs, we need an overhead in the preprocessing phase and in the space occupancy. We experimentally show that such an overhead is small compared to the speed-up gained in the updating phase. In fact, the preprocessing requires, in average, 2.10 and 2.36 times the time and the space required by Arc-Flags, respectively.

2 Preliminaries

Road Graphs. A *road graph* is a weighted directed graph $G = (V, E, w)$, used to model real road networks, where nodes represent points on the network, edges represent road segments between two points and the weight function $w : E \rightarrow \mathbb{R}^+$ represents an estimate of the travel time needed to traverse road segments. Given G , we denote as $\bar{G} = (V, \bar{E})$ the *reverse graph* of G where $\bar{E} = \{(v, u) \mid (u, v) \in E\}$. A *minimal travel time route* between two crossings S and T in a road network corresponds to a *shortest path* from the node s representing S and the node t representing T in the corresponding road graph. The total weight of a shortest path between nodes s and t is called *distance* and it is denoted as $d(s, t)$. A partition of V is a family $\mathcal{R} = \{R_1, R_2, \dots, R_r\}$ of subsets of V called *regions*, such that each node $v \in V$ is contained in exactly one region. Given $v \in R_k$, v is a *boundary node* of R_k if there exists an edge $(u, v) \in E$ such that $u \notin R_k$. Minimal routes in road networks can be computed by shortest paths algorithms such as Dijkstra's one. In order to perform an s - t query, the algorithm grows a shortest path tree starting from s and stopping as soon as it visits t . A simple variation of Dijkstra's algorithm is *bidirectional Dijkstra* which grows two shortest path trees starting from both s and t . In detail, the algorithm performs a visit of G starting from s and a visit of \bar{G} starting from t . The algorithm stops as soon the two visits meet at some node in the graph.

Arc-Flags. The preprocessing phase of Arc-Flags first computes a partition $\mathcal{R} = \{R_1, R_2, \dots, R_r\}$ of V and then associates a *label* to each edge (u, v) in E . A

label contains, for each region $R_k \in \mathcal{R}$, a *flag* $A_k(u, v)$ which is true if and only if a shortest path in G towards a node in R_k starts with (u, v) . The set of flags of an edge (u, v) is called *Arc-Flags* label of (u, v) . The preprocessing phase associates also Arc-Flags to edges in the reverse graph \bar{G} . The query phase of Arc-Flags consists of a modified version of the bidirectional Dijkstra’s algorithm: the forward search only considers those edges for which the flag of the target node’s region is true, while the backward search follows only those edges having a true flag for the source node’s region. The main advantage of Arc-Flags is its easy query algorithm combined with an excellent query performance. However, preprocessing is very time-consuming since it grows a full shortest path tree from each boundary node of each region. This is unpractical in dynamic scenarios where, in order to keep correctness of queries, the preprocessing phase has to be performed from scratch any time that an edge weight changes.

Road-Signs. Given a road graph $G = (V, E, w)$, a partition $\mathcal{R} = \{R_1, R_2, \dots, R_r\}$ of V in regions, an edge $(u, v) \in E$ and a region $R_k \in \mathcal{R}$, the Road-Sign $RS_k(u, v)$ of (u, v) to R_k is the subset of boundary nodes b of R_k , such that there exists a shortest path from u to b that contains (u, v) . The Road-Signs of (u, v) are represented as a boolean vector, whose size is the overall number of boundary nodes, where the i -th element is true if and only if the i -th boundary node is contained in $RS_k(u, v)$, for some region R_k . Hence, such a data structure requires $O(|E| \cdot |B|)$ memory, where B is the set of boundary nodes of G induced by \mathcal{R} . Notice that, in [3] it has been also proposed a technique to compact road signs that requires $O((|E| - |V|) \cdot |B|)$ overall space. The Road-Signs of G can be computed in the preprocessing phase of Arc-Flags. Given an edge (u, v) and a region R_k , $A_k(u, v)$ is set to true if and only if (u, v) is an edge in at least one of the shortest path trees grown for the boundary nodes of R_k . Such a procedure can be generalized to compute also Road-Signs. In fact, it is enough to add the boundary node b to $RS_k(u, v)$ if (u, v) is in the tree grown for b .

3 Incremental update of Arc-Flags

Given a road graph $G = (V, E, w)$ and a partition $\mathcal{R} = \{R_1, R_2, \dots, R_r\}$ of V in regions, we consider the problem of updating the Arc-Flags of G in a dynamic scenario where a sequence of only *weight decrease* operations $C = (c_1, c_2, \dots, c_h)$ occur on G . We denote as $G_i = (V, E, w_i)$ the graph obtained after i weight decrease operations, $0 \leq i \leq h$, $G_0 \equiv G$. Each operation c_i decreases the weight of one edge $e_i = (x_i, y_i)$ of an amount $\gamma_i > 0$, such that $w_i(e_i) = w_{i-1}(e_i) - \gamma_i > 0$ and $w_i(e) = w_{i-1}(e)$, for each edge $e \neq e_i$ in E . Our algorithm is based on the following proposition given in [3], which provides a straightforward method to compute the Arc-Flags of a graph given the Road-Signs of such graph.

Proposition 1. [3] *Given $G = (V, E, w)$, a partition $\mathcal{R} = \{R_1, R_2, \dots, R_r\}$ of V , an edge $(u, v) \in E$ and a region $R_k \in \mathcal{R}$, the following conditions hold:*

- (i) *if $u, v \in R_k$, then $A_k(u, v) = \text{true}$;*

- (ii) if $RS_k(u, v) \neq \emptyset$, then $A_k(u, v) = \text{true}$;
- (iii) if u or v is not in R_k and $RS_k(u, v) = \emptyset$, then $A_k(u, v) = \text{false}$.

Hence in what follows, we describe how our solution updates Road-Signs. The algorithm is denoted as INCRS and its pseudo-code is given in Figure 1. The algorithm is based on the observation that, when c_i occurs, all the shortest paths which contain (x_i, y_i) in G_{i-1} (i.e. before c_i) are shortest paths also in G_i (i.e. after c_i). Therefore, if some boundary node for some region R_k belongs to $RS_k(x_i, y_i)$ before c_i it belongs also to $RS_k(x_i, y_i)$ after c_i and no update is needed for it. However, it could happen that the shortest path from x_i to some boundary node b of R_k in G_{i-1} does not contain (x_i, y_i) but, after the weight decrease, the new shortest path in G_i contains (x_i, y_i) . In this case, b needs to be added to $RS_k(x_i, y_i)$ and removed from $RS_k(x_i, w)$, where (x_i, w) is the edge outgoing x_i whose road signs b belongs to. Same arguments can be applied to the incoming edges of x_i , (z, x_i) : if a boundary node belongs to $RS_k(z, x_i)$ and $RS_k(x_i, y_i)$ before c_i , then it belongs also to $RS_k(z, x_i)$ after c_i ; if a boundary node b is in $RS_k(x_i, y_i)$ (because it was already in $RS_k(x_i, y_i)$ or because it has been added to it as a consequence of c_i) and it does not belong to $RS_k(z, x_i)$ before c_i , then it might be added to $RS_k(z, x_i)$ in the case that the shortest path from z to b in G_i contains the sub-path (z, x_i, y_i) . We iteratively apply the same arguments to the other edges of the graph, starting from x_i and traversing its incoming edges. Note that, if at some point of the iteration we find out that the shortest path from a node z to some boundary node b does not decrease, then we do not need to add or remove b to any incoming edge of z . This allows us to reduce the search space of the algorithm.

INCRS works in two phases. In the first phase (lines 1–8) $RS_k(x_i, y_i)$ is updated by possibly adding new boundaries $b \notin RS_k(x_i, y_i)$ to it. This phase is performed for each $b \notin RS_k(x_i, y_i)$ separately (line 1). In the pseudo-code, when needed, we store distances between a node u and a boundary b in data structure $D[u, b]$ and we use an heap H to compute the minimum among the computed distances. Since each boundary node is processed separately, these data structures are overwritten at each computation, hence requiring $O(n)$ space in the worst case. In detail, at lines 2 and 3–4 we compute the distances from y_i to b and from node \bar{r} such that $b \in RS_k(x_i, \bar{r})$ to b , respectively. Then, at line 5, we check whether the weight of the path passing through y_i is smaller than that passing through \bar{r} (that is the shortest path from x_i to b in G_{i-1}). In the affirmative case, we update the road signs by adding b to $RS_k(x_i, y_i)$ and removing b from $RS_k(x_i, \bar{r})$ (lines 6–7). Moreover, we store the new distance from x_i to b in $D[x_i, b]$ (line 8) in order to use it in the next phase. In the second phase (lines 9–22) the road signs are updated for each $b \in RS_k(x_i, y_i)$ separately. The updating is done by mimicking the Dijkstra’s algorithm, that is by greedily visiting the reverse graph starting from x_i and stopping when a node does not need to update the road signs of its outgoing edges wrt b . At line 11, H is initialized by inserting x_i into it using $D[x_i, b]$ as key. If b was already in $RS_k(x_i, y_i)$, then $D[x_i, b]$ has not been computed at line 8 and hence it is computed at line 10. Until H is not empty (line 12), a node z and its distance

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Procedure: INCRS( $G_{i-1}, c_i, R_k$ )
Input : Graph  $G_{i-1}$ , weight decrease  $c_i$  on edge  $(x_i, y_i)$  and region  $R_k \in \mathcal{R}$ 
Output: Updated  $RS_k(u, v) \forall (u, v) \in E$ 

1 foreach  $b \notin RS_k(x_i, y_i)$  do
2   Compute  $D[y_i, b]$ ;
3   Let  $\bar{r} : (x_i, \bar{r}) \in E$  and  $b \in RS_k(x_i, \bar{r})$ ;
4   Compute  $D[\bar{r}, b]$ ;
5   if  $w(x_i, y_i) + D[y_i, b] < w(x_i, \bar{r}) + D[\bar{r}, b]$  then
6      $RS_k(x_i, \bar{r}) := RS_k(x_i, \bar{r}) \setminus \{b\}$ ;
7      $RS_k(x_i, y_i) := RS_k(x_i, y_i) \cup \{b\}$ ;
8      $D[x_i, b] := w(x_i, y_i) + D[y_i, b]$ ;
9 foreach  $b \in RS_k(x_i, y_i)$  do
10  Compute  $D[x_i, b]$  if it has not been already computed at line 8;
11   $H.Insert(x_i, D[x_i, b])$ ;
12  while  $H \neq \emptyset$  do
13     $(z, D[z, b]) := H.Delete\_Min()$ ;
14    foreach  $z_i : (z_i, z) \in E$  do
15      Let  $\bar{r} : (z_i, \bar{r}) \in E$  and  $b \in RS_k(z_i, \bar{r})$ ;
16      Compute  $D[\bar{r}, b]$ ;
17      if  $w(z_i, z) + D[z, b] < w(z_i, \bar{r}) + D[\bar{r}, b]$  or  $\bar{r} = z$  then
18         $RS_k(z_i, \bar{r}) := RS_k(z_i, \bar{r}) \setminus \{b\}$ ;
19         $RS_k(z_i, z) := RS_k(z_i, z) \cup \{b\}$ ;
20         $D[z_i, b] := w(z_i, z) + D[z, b]$ ;
21        if  $z_i \in H$  then  $H.Decrease\_Key(z_i, D[z_i, b])$ ;
22        else  $H.Insert(z_i, D[z_i, b])$ ;

```

Fig. 1. Pseudo-code of algorithm INCRS.

$D[z, b]$ is extracted from H. Then, for each node z_i such that $(z_i, z) \in E$, at lines 14–22 we perform the same steps done for x_i and (x_i, y_i) : we compute the distance from \bar{r} such that $b \in RS_k(z_i, \bar{r})$ to b ; we check whether the weight of the path from z_i to b passing through z is smaller than that passing through \bar{r} and, in the affirmative case, we update the road signs by adding b to $RS_k(z_i, z)$ and removing it from $RS_k(z_i, \bar{r})$; finally, we store the new distance from z_i to b in $D[z_i, b]$ and insert z_i into H or decrease its key if it already belongs to H.

The following theorem states the correctness of INCRS.

Theorem 1. *Given $G = (V, E, w)$ and a partition $\mathcal{R} = \{R_1, R_2, \dots, R_r\}$ of V , for each $(u, v) \in E$ and $R_k \in \mathcal{R}$, INCRS correctly updates $RS_k(u, v)$ and $A_k(u, v)$ after a weight decrease operation on an edge of G .*

Proof. Let us consider a region $R_k \in \mathcal{R}$ and a weight decrease operation c_i on edge (x_i, y_i) . From Proposition 1, it is enough to show that INCRS correctly updates $RS_k(u, v)$ after c_i , for each $(u, v) \in E$. Given an edge (u, v) , we denote as $RS_k(u, v)$ and $RS'_k(u, v)$ the road-signs of (u, v) before and after c_i , respectively.

As c_i decreases the weight of (x_i, y_i) , then $RS_k(x_i, y_i) \subseteq RS'_k(x_i, y_i)$, moreover for each $(u, v) \in E$, $RS_k(u, v)$ can be modified by adding or removing

boundary nodes in $RS'_k(x_i, y_i)$, that is $RS'_k(u, v) \setminus RS_k(u, v) \subseteq RS'_k(x_i, y_i)$ and $RS_k(u, v) \setminus RS'_k(u, v) \subseteq RS'_k(x_i, y_i)$. It follows that phase one (lines 1–8) correctly updates the road-signs of edges outgoing from x_i . The road-signs of the remaining edges are updated in phase two, whose correctness is shown separately for each boundary node $b \in RS'_k(x_i, y_i)$ and derives from the following facts.

F1 *The nodes are extracted from H at line 13 in non-decreasing order of keys.*

Let us consider two nodes u and v extracted from H at times t_u and t_v with keys $D[u, b]$ and $D[v, b]$, respectively. By contradiction, suppose that $t_u < t_v$ and $D[u, b] > D[v, b]$. Since at line 13 the node with the minimum key is extracted, at time t_u , $D[u, b]$ was minimum and hence either u has been extracted into H after t_u or its key has been decreased after t_u . In both cases, the algorithm passed the test at line 17 which implies that there exists a node v_1 and a time $t_{v_1} < t_v$ such that $D[v_1, b] < D[v, b] < D[u, b]$, where $D[v_1, b]$ is the key of v_1 at time t_{v_1} when v_1 is extracted from H. If we iteratively repeat this arguments for node v_1 , we obtain a sequence of nodes $v \equiv v_0, v_1, v_2, \dots, v_k$, where $v_k \equiv x_i$, such that, if we denote as $D[v_j, b]$ the key of v_j when it is extracted from H at time t_{v_j} , then $D[v_{j+1}, b] < D[v_j, b]$ and $t_{v_{j+1}} < t_{v_j}$ for each $j = 0, 1, \dots, k-1$. For the condition at line 12, one of the nodes in the sequence v_j has to belong to H with key $D[v_j, b] < D[u, b]$ at time t_u which contradicts the fact that, at time t_u , node u is the node with the minimum key.

F2 *A node is extracted from H at line 13 at most once.*

Suppose that a node u is extracted from H at two different times $t_1 < t_2$. Then, node u has been inserted into H at two different times, denoted as \bar{t}_1 and \bar{t}_2 , when it does not belong to H. It follows that $[\bar{t}_1, t_1] \cap [\bar{t}_2, t_2] = \emptyset$. Further, let us denote as $D_{t_1}[u, b]$ ($D_{t_2}[u, b]$, resp.) the key of u at time \bar{t}_1 (\bar{t}_2 , resp.) which is equal to that at time t_1 (t_2 , resp.). Let us consider the two (possibly different) nodes v_1 and v_2 which are extracted from H immediately before times \bar{t}_1 and \bar{t}_2 , respectively. Let us analyze the keys extracted from H. At time \bar{t}_1 , v_1 is extracted with key $D[v_1, b]$, $D[u, b]$ is set to $D_{t_1}[u, b] = w(u, v_1) + D[v_1, b]$ and b is added to $RS_k(u, v_1)$. At time t_1 , u is extracted with key $D_{t_1}[u, b]$. At time \bar{t}_2 , v_2 is extracted with key $D[v_2, b]$, and $D[u, b]$ is set to $D_{t_2}[u, b] = w(u, v_2) + D[v_2, b]$, it follows that the test at line 17 returned true and, as $b \in RS_k(u, v_1)$, $w(u, v_2) + D[v_2, b] < w(u, v_1) + D[v_1, b]$. Hence $D_{t_2}[u, b] < D_{t_1}[u, b]$. At time t_2 , u is extracted with key $D_{t_2}[u, b]$ and, since $t_1 < t_2$, this contradicts Fact F1.

F3 *For each edge (u, v) such that $RS'_k(u, v) \neq RS_k(u, v)$, u is inserted into H.*

We show a stronger statement that is: if a node changes its distance to $b \in RS'_k(u, v)$ it is inserted into H. By contradiction, let us consider the node u such that: it changes its distance to b , it is not inserted into H, and its distance to b after c_i is minimal among the nodes with the same properties. By this last property, the node v on the shortest path from u to b after c_i is inserted into H. When v is extracted from H, either b belongs to $RS_k(u, v)$ or $w(u, v) + D[v, b] < w(u, \bar{r}) + D[\bar{r}, b]$, where \bar{r} is the node such that $b \in RS_k(u, \bar{r})$. In both cases, u is inserted into H at line 22.

F4 *When a node u is extracted from H at line 13, for each $(u, v) \in E$, $RS_k(u, v)$ is correctly updated.*

| graph | n. of nodes | n. of edges | %MOT | %NAT | %REG | %URB |
|-------|-------------|-------------|-------|------|-------|-------|
| LUX | 30 647 | 75 576 | 0.55 | 1.95 | 14.77 | 82.71 |
| DNK | 469 110 | 1 090 148 | 24.02 | 3.06 | 0.48 | 72.45 |
| BEL | 458 403 | 1 164 046 | 22.90 | 2.92 | 0.52 | 73.62 |
| AUT | 722 512 | 1 697 902 | 27.60 | 5.33 | 1.71 | 65.21 |
| ESP | 695 399 | 1 868 838 | 33.22 | 6.34 | 1.51 | 58.87 |
| NED | 892 027 | 2 278 824 | 0.40 | 0.56 | 5.16 | 93.86 |
| SWE | 1 546 705 | 3 484 378 | 19.54 | 2.86 | 0.45 | 77.10 |

Table 1. Tested road graphs. 1st col.: the graph; 2nd and 3rd col.s: number of nodes and edges in the graph; 4th–7th col.s: percentage of edges into categories: motorways (MOT), national roads (NAT), regional roads (REG), and urban streets (URB).

By contradiction, let us consider the first node u whose outgoing edges have wrong road-signs when u is extracted from H . Let us consider the node v such that $b \in RS_k(u, v)$ when u is extracted, that is v is the node that was extracted from H immediately before the last time that either u is inserted into H or its key is decreased. As u is the first node whose outgoing edges have wrong road-signs when it is extracted from H , then the road-signs of the edges outgoing from v are correctly updated. Moreover, also the road-signs of edges outgoing from w , for each $(u, w) \in E$ are correctly updated. In fact two cases may arise: if w changes the road-signs of its outgoing edges then, by Fact F3, it is inserted into H , by facts F1 and F2, it is extracted before u and hence, by hypothesis, it correctly updates the road-signs of its outgoing edges; otherwise the road-signs of its outgoing edges are already correct. In any case, when u is inserted into H the distances used in the test of line 17 are correctly computed and hence the road signs are correctly updated. \square

From a theoretical point of view, INCRS requires a computational complexity which is, in the worst case, comparable to that of the recomputation from scratch of Arc-Flags. However, INCRS focuses the computation only on the nodes that change shortest paths to a subset of the boundary nodes (and possibly on the neighbors of such nodes). In contrast, the recomputation from scratch computes all the shortest paths from any boundary node to each other node of the network. This difference is difficult to be captured by a worst case analysis and this motivates the experimental study of the next section.

4 Experimental study

In this section, we compare the performances of the incremental algorithm proposed in this paper and the decremental algorithm of [3], whose combination is named here DYNRS, on fully dynamic sequences of weight change operations against the recomputation from scratch of Arc-Flags. We used the implementation of Arc-Flags of [2]. Furthermore, it has been shown in [10] that the best query performances for Arc-Flags are achieved when partitions are computed by using arc-separator algorithms. For this reason, we used arc-separators obtained

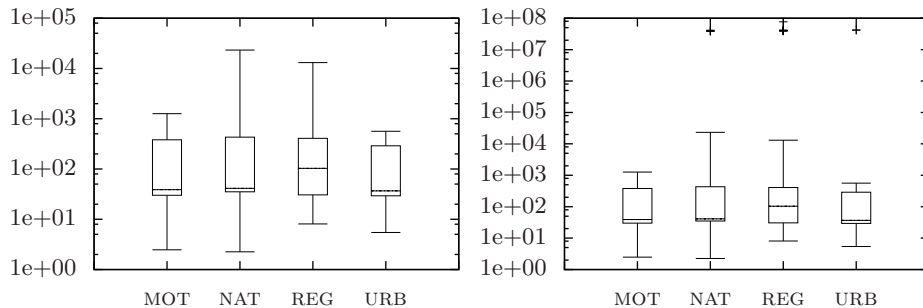


Fig. 2. Speed-up factors of the fully dynamic algorithm for the road network of Sweden, without (left) and with (right) outliers.

by the METIS library [8]. Our experiments are performed on a workstation equipped with a Quad-core 3.60 GHz Intel Xeon X5687 processor, with 12MB of internal cache and 24 GB of main memory. The program has been compiled with GNU g++ compiler 4.4.3 under Linux (kernel 2.6.32).

We consider seven road graphs available from PTV [11] representing the road networks of Luxembourg, Denmark, Belgium, Austria, Spain, Netherlands and Sweden, denoted as LUX, DNK, BEL, AUT, ESP, NED and SWE, respectively. In each graph, edges are classified into four categories according to their speed limits: motorways (MOT), national roads (NAT), regional roads (REG) and urban streets (URB). The main characteristics of these graphs are reported in Table 1. We consider fully dynamic sequences of updates simulating disruptions on road networks built as follows. The most significant operation that can occur on a road segment is the increase of a weight, which simulates a delay in the travel time on that segment due, for instance, to a traffic jam. This operation is usually followed by a weight decrease on the same road segment which simulates the restore from the delay. Hence, for each operation in the sequence that increases the weight of an edge (x_i, y_i) of a quantity γ_i , there is a corresponding subsequent operation which decreases the weight of edge (x_i, y_i) of the same amount γ_i . We execute, for each graph considered and for each road category, random sequences of 100 weight-change operations as described above. The weight-change amount for each operation is chosen uniformly at random in [25%, 75%] of the weight of the edge involved in that operation. As a performance indicator, we choose the time used by the algorithms to complete a single update during the execution of a sequence. We measure, as speed-up factor, the ratio between the time required by the recomputation from scratch of Arc-Flags and that required by DYNRS. The results are reported in Fig. 2–3, and in Table 2.

Fig.s 2 and 3 show two box-plot diagrams representing the values of the speed-up factors obtained for SWE and NED. For each category, we represent minimum value, first quartile, median value, third quartile, and maximum value. In both figures, the diagram on the left does not show outlier values while the diagram

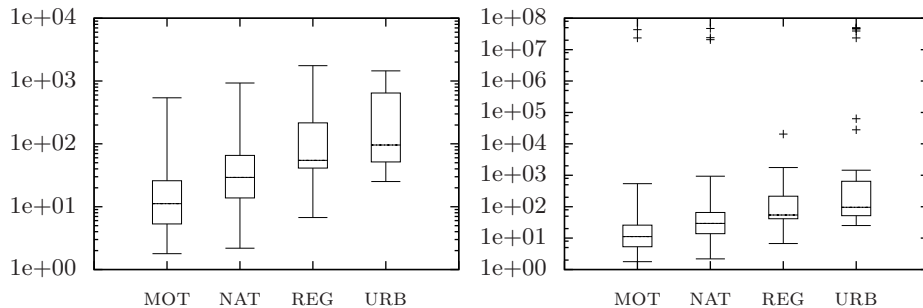


Fig. 3. Speed-up factors of the fully dynamic algorithm for the road network of the Netherlands, without (left) and with (right) outliers.

on the right does. Outlier values occur when DYNRS performs much better than Arc-Flags because the number of Road-Signs changed is very small. We consider a test as outlier if the speed-up factor is 1000 times the speed-up factor median value. Even without outliers, the speed-up gained by DYNRS is high.

Notice that, in the case of SWE (Fig. 2), the speed-up factors are quite similar on the different categories, thus highlighting an independency from categories. This is the typical behavior of road networks, as shown also for BEL, DNK, ESP and AUT in Table 2, where we report the *average* time of the recomputation from-scratch of Arc-Flags and the *average* time of DYNRS, the *average ratios* between these quantities and the speed-up factors. The only exceptions are LUX and NED, where the percentage of motorways is very low. This is the reason why we highlight the behavior of DYNRS on NED in Fig. 3, where the speed-up factor reaches the highest values when update operations occur on urban edges, while it is smaller when they occur on motorway edges. In fact, when an update operation occurs on urban edges, the number of shortest paths that change is small compared to the case that an update operation occurs on motorways edges. This implies that DYNRS, which selects the nodes that change such shortest paths and focus the computation only on such nodes, performs better than the recomputation from-scratch of the shortest paths from any boundary node.

We note that, the speed-up factor increases with the size of the network. This can be explained by the fact that, when an edge update operation occurs, it affects only a part of the graph, hence only a subset of the edges in the graph need to update their Arc-Flags or Road-Signs. In most of the cases, this part is small compared to the size of the network and, with high probability, it corresponds to the subnetwork close to the edge increased or closely linked to it. In other words, it is unlikely that a traffic jam in a certain part of the network affects the shortest paths of another part which is far or not linked to the first one. Clearly, this fact is more evident when the road network is big. In conclusion, it is evident from Table 2, that DYNRS outperforms the recomputation from-scratch by far and that it requires reasonable computational time.

| graph | cat. | AF (s) | | DYNRS (s) | | ratio | | avg. speed-up | |
|-------|------|----------|----------|-----------|-------|-------|-------|---------------|--------|
| LUX | MOT | 5.50 | 5.52 | 2.28 | 1.57 | 2.42 | 6.16 | 11.09 | 40.32 |
| | NAT | 5.51 | | 2.64 | | 2.09 | | 32.81 | |
| | REG | 5.56 | | 0.99 | | 5.61 | | 24.79 | |
| | URB | 5.52 | | 0.38 | | 14.50 | | 92.60 | |
| DNK | MOT | 542.55 | 542.46 | 17.05 | 14.90 | 31.84 | 37.40 | 449.10 | 431.41 |
| | NAT | 542.83 | | 11.24 | | 48.31 | | 430.16 | |
| | REG | 542.22 | | 16.16 | | 33.61 | | 330.64 | |
| | URB | 542.25 | | 15.13 | | 35.84 | | 515.72 | |
| BEL | MOT | 644.48 | 644.34 | 23.16 | 23.73 | 27.83 | 29.18 | 195.53 | 238.23 |
| | NAT | 644.30 | | 28.93 | | 22.26 | | 195.81 | |
| | REG | 644.31 | | 28.00 | | 23.03 | | 421.83 | |
| | URB | 644.28 | | 14.85 | | 43.58 | | 139.73 | |
| AUT | MOT | 935.16 | 940.83 | 39.31 | 26.25 | 23.76 | 39.04 | 108.12 | 166.10 |
| | NAT | 934.16 | | 19.08 | | 48.99 | | 185.48 | |
| | REG | 956.62 | | 27.70 | | 33.74 | | 183.09 | |
| | URB | 937.38 | | 18.89 | | 49.65 | | 187.71 | |
| ESP | MOT | 736.81 | 737.31 | 22.21 | 21.41 | 33.15 | 34.91 | 392.19 | 373.98 |
| | NAT | 737.30 | | 21.44 | | 34.39 | | 562.27 | |
| | REG | 736.50 | | 24.46 | | 30.12 | | 264.78 | |
| | URB | 736.32 | | 17.54 | | 41.98 | | 276.66 | |
| NED | MOT | 1 607.36 | 1 606.29 | 206.36 | 90.55 | 7.78 | 41.31 | 31.45 | 169.79 |
| | NAT | 1 606.22 | | 107.27 | | 14.97 | | 107.20 | |
| | REG | 1 609.60 | | 30.85 | | 52.03 | | 181.84 | |
| | URB | 1 601.96 | | 17.71 | | 90.45 | | 358.67 | |
| SWE | MOT | 2 681.54 | 2 681.16 | 113.90 | 76.12 | 23.65 | 38.14 | 180.36 | 316.64 |
| | NAT | 2 681.94 | | 68.98 | | 38.87 | | 519.96 | |
| | REG | 2 678.81 | | 52.79 | | 50.75 | | 394.99 | |
| | URB | 2 682.34 | | 68.82 | | 39.29 | | 171.24 | |

Table 2. Avg update times and speed-up factors of DYNRS. 1st col.: graph; 2nd col.: category where the weight changes occur; 3rd and 4th col.s: avg computational time for Arc-Flags and DYNRS, resp.; 5th col.: ratio between the values of the 3rd and the 4th col.s; 6th col.: avg speed-up factors of DYNRS against Arc-Flags.

Regarding the preprocessing phase, in Table 3 we report the computational time and the space occupancy required by Arc-Flags and DYNRS. Table 3 shows that, for computing Road-Signs along with Arc-Flags, we need about twice the computational time required for computing only Arc-Flags, which is a small overhead compared to the speed-up gained in the updating phase. The same observation can be done regarding the space occupancy. In fact, Table 3 also shows that the space required for storing both Road-Signs and Arc-Flags is between 1.44 and 3.48 times that required to store only Arc-Flags.

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| graph | reg. | AF (s) | AF+RS (s) | t. ratio | AF (B) | AF+RS (B) | s. ratio |
|-------|------|----------|-----------|----------|------------|-------------|----------|
| LUX | 64 | 5.52 | 11.79 | 2.14 | 1 209 216 | 1 744 531 | 1.44 |
| DNK | 128 | 542.46 | 1 104.16 | 2.04 | 17 442 368 | 43 932 508 | 2.52 |
| BEL | 128 | 644.34 | 1 373.11 | 2.13 | 18 624 736 | 64 759 010 | 3.48 |
| AUT | 128 | 940.83 | 1 995.93 | 2.12 | 27 166 432 | 78 720 055 | 2.90 |
| ESP | 128 | 737.31 | 1 446.65 | 1.96 | 29 901 408 | 67 664 666 | 2.26 |
| NED | 128 | 1 606.29 | 3 214.34 | 2.00 | 36 461 184 | 64 612 836 | 1.77 |
| SWE | 128 | 2 681.16 | 5 986.62 | 2.23 | 55 750 048 | 163 151 588 | 2.93 |

Table 3. Preprocessing time and space. 1st col.: graph; 2nd col.: number of regions; 3rd col.: preprocessing time, in seconds, of Arc-Flags; 4th col.: preprocessing time, in seconds, of Arc-Flags and Road-Signs; 5th col.: ratio between the values reported in the 4th and the 3rd column; 6th col.: space required, in Bytes, to store Arc-Flags; 7th col.: space required, in Bytes, to store Arc-Flags and Road-Signs; 8th col.: ratio between the values of the 7th and the 6th column.

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