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Selecting polynomials for the Function Field Sieve

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Abstract

The Function Field Sieve algorithm is dedicated to computing discrete logarithms in a finite field \mathbb{F}_{q^n} , where q is a small prime power. The scope of this article is to select good polynomials for this algorithm by defining and measuring the size property and the so-called root and cancellation properties. In particular we present an algorithm for rapidly testing a large set of polynomials. Our study also explains the behaviour of inseparable polynomials, in particular we give an easy way to see that the algorithm encompasses the Coppersmith algorithm as a particular case.

1 Introduction

The Function Field Sieve (FFS) algorithm is dedicated to computing discrete logarithms in a finite field \mathbb{F}_{q^n} , where q is a small prime power. Introduced by Adleman in [Adl94] and inspired by the Number Field Sieve (NFS), the algorithm collects pairs of polynomials $(a, b) \in \mathbb{F}_q[t]$ such that the norms of $a - bx$ in two function fields are both smooth (the sieving stage), i.e having only irreducible divisors of small degree. It then solves a sparse linear system (the linear algebra stage), whose solutions, called virtual logarithms, allow to compute the discrete logarithm of any element during a final stage (individual logarithm stage).

The choice of the defining polynomials f and g for the two function fields can be seen as a preliminary stage of the algorithm. It takes a small amount of time but it can greatly influence the sieving stage by slightly changing the probabilities of smoothness. In order to solve the discrete logarithm in \mathbb{F}_{q^n} , the main required property of $f, g \in \mathbb{F}_q[t][x]$ is that their resultant $\text{Res}_x(f, g)$ has an irreducible factor $\varphi(t)$ of degree n . Various methods have been proposed to build such polynomials.

The base- m method of polynomial selection, proposed by Adleman [Adl94], consists in choosing $\varphi(t)$ an irreducible polynomial of degree n , setting $g = x - m$, where m is a power of t , and f equal to the base- m expansion of φ . He obtained a subexponential complexity of $L_{q^n}(\frac{1}{3}, c)^{1+o(1)}$ with $c = \sqrt[3]{64/9}$. Adleman and Huang [AH99] chose φ to be a sparse polynomial and obtained a constant $c = \sqrt[3]{32/9}$. They also noted that the previously known algorithm of Coppersmith [Cop84] can be seen as a particular case of the FFS. Finally, Joux and Lercier [JL02] introduced a method which, without improving the complexity, behaves better in practice. It consists in selecting a polynomial f with small degree coefficients and then of randomly testing linear polynomials $g = g_1x + g_0$ until $\text{Res}_x(f, g)$ has an irreducible factor of degree n . In [JL06] Joux and Lercier proposed two additional variants of their methods. In the first one that can be called *Two rational sides*, we add the condition that f has degree 1 in t . Its main advantage is that its description does not require the theory of function fields. The second variant, called the *Galois improvement*, applies to the case where n is composite.

In this paper we improve the method of Joux and Lercier [JL02] by showing how to select the non-linear polynomial f . For that we follow the strategy that was developed in the factorization

context. In particular Murphy [Mur99] introduced and used criteria which allow to rapidly rank any set of polynomials. See [Bai11], for recent developments in this direction.

Therefore, we introduce relevant functions for the sieving efficiency of a polynomial, taking into account a size property, a so-called root property that reflects the behaviour modulo small irreducible polynomials, and a cancellation property, which is analogous to the real roots property for the NFS. We also present efficient algorithms for quantifying these properties. A special attention is given to the particular case where f is not separable. Indeed, this is a phenomenon that has no analogue in the factorization world and that has strong repercussions on the sieving efficiency.

Recent works on composite extensions In the past few weeks, the algorithm in [JL06] was the object of further improvements in [Jou12], [GGMZ13] and [Jou13], all of them being very well adapted to the case of composite extensions. The most important of them is Joux's new algorithm which is especially suited to the fields $\mathbb{F}_{q^{2k}}$ with q close to k and whose complexity is $L(1/4 + o(1))$. Moreover, under some overhead hidden by the $o(1)$, Joux's algorithm can be adapted to the case \mathbb{F}_{q^k} when q and k are both prime.

In this context, the FFS keeps its interest for both theoretical and practical reasons. On the one hand, the FFS applies to a wider set of finite fields as the $L(1/4 + o(1))$ complexity was only obtained for constant characteristic. On the other hand, except for the composite case, the crossing point between the FFS and Joux's algorithm [Jou13] is still to be determined.

Outline The article is organized as follows. Section 2 lists the properties which improve the sieve yield. Section 3 combines the previously defined functions in order to compare arbitrary polynomials and shows how to rapidly test a large number of candidate polynomials. Section 4 focuses on the case of inseparable polynomials and, in particular, the Coppersmith algorithm. Section 5 applies the theoretic results to a series of examples. Finally, section 6 makes the synthesis.

2 Quantification functions

2.1 Size property

We start by deciding the degrees in t and x of the two polynomials f and g . The FFS has always been implemented for polynomials f of small coefficients in t and for polynomials g of degree 1 in x , like in [JL02]. It might be not obvious that this is the best choice. For instance in the case of the NFS both for factorization [Mon06, PZ11] and discrete logarithm [JL03], pairs of non-linear polynomials were used. In the following, we argue that the classical choice is indeed the best one. Let us first recall the nature of the objects we have to test for smoothness. The FFS collects coprime

pairs $(a(t), b(t)) \in \mathbb{F}_q[t]$ such that the norms of $a - bx$ in the function fields of f and g are both smooth. These norms are polynomials in t of a simple form: denoting $F(X, Y) = f(\frac{X}{Y})Y^{\deg_x f}$ the homogenization of $f(x)$, the norm of $a - bx$ is just $F(a, b)$. Similarly, we denote $G(X, Y)$ the homogenization of $g(x)$ and the second norm is $G(a, b)$. As a consequence, the polynomial selection stage can be restated as the search for polynomials f and g such that $F(a, b)$ and $G(a, b)$ are likely to be both smooth for coprime pairs (a, b) in a given range.

As a first approximation, we translate this condition into the fact that the degree of the product $F(a, b)G(a, b)$ is as small as possible. It will be refined all along this paper.

Fact 2.1. *Assume that we have to compute discrete logarithms in \mathbb{F}_{q^n} , with polynomials $f, g \in \mathbb{F}_q[t][x]$, such that $\deg_x g \leq \deg_x f$. If for bounded a and b , the polynomials f and g minimize $\max\{\deg F(a, b)G(a, b)\}$, then we have $\deg_x g = 1$.*

Argument. Let (a, b) be any pair of maximal degree e , and assume that there is no cancellation of the monomials in $F(a, b)$ and $G(a, b)$ respectively. Then one has

$$\deg(F(a, b)G(a, b)) = \deg_t g + e \deg_x g + \deg_t f + e \deg_x f. \quad (1)$$

The degree of the resultant of f and g can be bounded by $\deg \text{Res}(f, g) \leq \deg_x f \deg_t g + \deg_x g \deg_t f$. Since we need this resultant to be of degree at least n , we impose

$$\deg_x f \deg_t g + \deg_x g \deg_t f \geq n. \quad (2)$$

For a fixed value of the left hand side in Equation 2, in order to minimize the expression in Equation 1, we need to minimize $\deg_t f$. Therefore we set $\deg_t f$ as small as possible, let us call this degree ϵ . Hence the optimization problem becomes

$$\begin{aligned} \text{minimize} \quad & (\deg_x f + \epsilon + e \deg_x g + \deg_t g) \\ \text{when} \quad & \deg_x f \deg_t g + \epsilon \deg_x g \geq n. \end{aligned}$$

Since one can decrease $\deg_x g$ without changing too much the left hand side of the constraint, the choice $\deg_x g = 1$ is optimal. \square

In the rest of the article we simply write d for $\deg_x f$. The degree of g in t is then about n/d .

Remark 2.2. We decided to optimize the degree of the product of the norms. In terms of smoothness probability, it is only pertinent if both norms have similar degrees. More precisely, as the logarithm of Dickman's rho function is concave, it can be shown that it is optimal to balance the degrees of the norms. Hence sensible choices of the parameters are such that $de \approx \frac{n}{d} + e$.

We are now ready to quantify the size property for a single polynomial f . It clearly depends on the bound e on both $\deg a$ and $\deg b$. In the following definition, we also take into account the skewness improvement as implemented in [DGV13], that is, we set the skewness s to be the difference between the bounds on the degree of a and of b . We can then define sigma to match the average degree of $F(a, b)$ for (a, b) in a domain of polynomials of degrees bounded by $[e + s/2]$ and $[e - s/2]$, when no cancellation occurs among the monomials of $F(a, b)$. This translates into the following formal definition.

Definition 2.3. Let $f \in \mathbb{F}_q[t][x]$ be a polynomial, s the skewness parameter and e the sieve size parameter. We define:

$$\sigma(f, s, e) = \sum_{\substack{0 \leq d_a \leq e + s/2 \\ 0 \leq d_b \leq e - s/2}} p_{d_a, d_b} \max_{i \in [0, d]} \left(\deg(f_i) + i d_a + (d - i) d_b \right), \quad (3)$$

with $p_{d_a, d_b} = (q - 1)^2 q^{d_a + d_b} / q^{\lfloor e + s/2 \rfloor + \lfloor e - s/2 \rfloor + 2}$.

2.2 Root property

As a first approximation, for a random pair $(a, b) \in \mathbb{F}_q[t]^2$, $F(a, b)$ has a smoothness probability of the same order of magnitude as random polynomials of the same degree. Nevertheless, we shall show that for a fixed size property, some polynomials improve this probability by a factor of 2 or more.

Consider the example of $f = x(x - 1) - (t^64 - t)$ in $\mathbb{F}_2[t][x]$. For all monic irreducible polynomials ℓ of degree at most 3, ℓ divides the constant term, so f has two roots modulo ℓ . For each such ℓ and for all b non divisible by ℓ , there are 2 residues of a modulo ℓ such that $F(a, b) \equiv 0 \pmod{\ell}$. Therefore, for all ℓ of degree 3 or less, the probability that ℓ divides the norm is heuristically twice as large as the probability that it divides a random polynomial. This influences in turn the smoothness probability. This effect is quantified by the function alpha that we now introduce.

2.2.1 Definition of alpha

Introduced by Murphy [Mur99] in the case of the NFS, alpha can be extended to the case of the FFS. Let ℓ be a monic irreducible polynomial in $\mathbb{F}_q[t]$ and let us call ℓ -part of a polynomial P , the

largest power of ℓ in P . We shall prove that the quantity α_ℓ below is the degree difference between the ℓ -part of a random polynomial and the ℓ -part of $F(a, b)$ for a random coprime pair $(a, b) \in \mathbb{F}_q[t]$.

Let us first properly define the average of a function on a set of polynomials.

Definition 2.4. Let v be a real function of one or two polynomial variables $v : \mathbb{F}_q[t] \rightarrow \mathbb{R}$ or $v : \mathbb{F}_q[t] \times \mathbb{F}_q[t] \rightarrow \mathbb{R}$. Let S be a subset of the domain of v . For any pair (a, b) of polynomials, we write $\deg(a, b)$ for $\max(\deg(a), \deg(b))$. If the limit below exists we call it average of v over S and denote it by $\mathbb{A}(v, S)$:

$$\mathbb{A}(v, S) = \lim_{N \rightarrow \infty} \frac{\sum_{s \in S, \deg(s) \leq N} \varphi(s)}{\#\{s \in S \mid \deg(s) \leq N\}}.$$

Definition 2.5. Let L denote the set of monic irreducible polynomials in $\mathbb{F}_q[t]$. Put $D = \{(a, b) \in \mathbb{F}_q[t]^2 \mid \gcd(a, b) = 1\}$. Take a non-constant polynomial $f \in \mathbb{F}_q[t][x]$. When the right hand members are defined, we set for all $\ell \in L$:

$$\begin{aligned} \alpha_\ell(f) &= \deg(\ell) \left(\mathbb{A}(v_\ell(P), \{P \in \mathbb{F}_q[t]\}) - \mathbb{A}(v_\ell(F(a, b)), \{(a, b) \in D\}) \right), \\ \alpha(f) &= \sum_{\ell \in L} \alpha_\ell(f), \end{aligned}$$

where v_ℓ is the valuation at ℓ . The infinite sum which defines $\alpha(f)$ must be seen as a formal notation and by its sum we denote the limit when b_0 goes to infinity of $\alpha(f, b_0) := \sum_{\ell \in L, \deg \ell \leq b_0} \alpha_\ell(f)$.

Notation 2.6. For all irreducible polynomial $\ell \in \mathbb{F}_q[t]$, $N(\ell)$ denotes the number of residues modulo ℓ , i.e. $q^{\deg \ell}$.

We call affine root of f modulo ℓ^k any $r \in \mathbb{F}_q[t]$ such that $\deg r < k \deg \ell$ and $F(r, 1) = f(r) \equiv 0 \pmod{\ell^k}$. Also we call projective root of f modulo ℓ^k the polynomials $r \in \mathbb{F}_q[t]$ such that $\ell \mid r$, $\deg r < k \deg \ell$ and $F(1, r) \equiv 0 \pmod{\ell^k}$. Note that, when $k = 1$ any affine and projective roots can be seen as an element of $\mathbb{P}^1(\mathbb{F}_q(t))$, hence we denote them $(r : 1)$ and $(1 : r)$ respectively.

Notation 2.7. We denote by $S(f, \ell)$ the set of affine and projective roots modulo ℓ^k for any k :

$$S(f, \ell) = \left\{ (r, k) \mid k \geq 1, r \text{ affine or projective root of } f \text{ modulo } \ell^k \right\}. \quad (4)$$

Proposition 2.8. *Let $f \in \mathbb{F}_q[t][x]$ and $\ell \in \mathbb{F}_q[t]$ a monic irreducible polynomial. Then α_ℓ exists and we have*

$$\alpha_\ell(f) = \deg \ell \left(\frac{1}{N(\ell) - 1} - \frac{N(\ell)}{N(\ell) + 1} \sum_{(r, k) \in S(f, \ell)} \frac{1}{N(\ell)^k} \right). \quad (5)$$

Proof. In order to prove the convergence of Equation 5, note that some elements of $S(f, \ell)$ group into infinite sequences $\{(r^{(k)}, k)\}_k$ with $r^{(k)} \equiv r^{(k-1)} \pmod{\ell^{k-1}}$. Since each infinite sequence defines a root of f in the ℓ -adic completion of $\mathbb{F}_q(t)$, there are at most d such sequences, whose contributions converge geometrically. There are only finitely many remaining elements of $S(f, \ell)$ because otherwise one could extract an additional ℓ -adic root. This proves the convergence.

In order to show the equality, note that we have $\mathbb{A}(v_\ell(P), \mathbb{F}_q[t]) = \sum_{k=1}^{\infty} \frac{1}{N(\ell)^k} = \frac{1}{N(\ell)-1}$. Indeed, for all $k \in \mathbb{N}$, the density of the set of polynomials divisible by ℓ^k is the inverse of the number of residues modulo ℓ^k , which is $q^{\deg(\ell)k} = N(\ell)^k$.

Let us compute $a_{\text{hom}}^{(\ell)} := \mathbb{A}(v_\ell(F(a, b)), \{(a, b) \in \mathbb{F}_q[t]^2 \mid \gcd(a, b) \not\equiv 0 \pmod{\ell}\})$. This corresponds to the contribution of $F(a, b)$ in the definition of α_ℓ . The condition $\gcd(a, b) = 1$ has been replaced by a local condition. Since we are only interested in ℓ -valuations, this does not change the result. The number of ℓ -coprime pairs (a, b) of degree less than $k \deg \ell$ is $N(\ell)^{2k} - N(\ell)^{2k-2}$. Such a pair (a, b) satisfies $\ell^k \mid F(a, b)$ if and only if $\ell \nmid b$ and ab^{-1} is an affine root modulo ℓ^k or $\ell \mid b$ and ba^{-1} is a projective root modulo ℓ^k . For each affine root r , the number of ℓ -coprime pairs (a, b) such that $a \equiv br \pmod{\ell^k}$ equals the number of choices for b , which is $N(\ell)^k - N(\ell)^{k-1}$. Similarly, for

each projective root, the number of coprime pairs (a, b) such that $b \equiv ar \pmod{\ell^k}$ is the number of choices for a , which is $N(\ell)^k - N(\ell)^{k-1}$. Hence, it follows that for each $(r, k) \in S(f, \ell)$, the probability that the residues of a coprime pair (a, b) modulo ℓ^k match the root (r, k) is given by the formula below, where $(a : b) \equiv r \pmod{\ell^k}$ is short for $r \equiv ab^{-1} \pmod{\ell^k}$ when $\ell \nmid b$ or $r \equiv ba^{-1} \pmod{\ell^k}$ when $\ell \mid b$.

$$\mathbb{P}((a : b) \equiv r \pmod{\ell^k}) = \frac{1}{N(\ell)^k} \frac{N(\ell)}{N(\ell) + 1}, \quad (6)$$

The following equation is intuitive so that we postpone its proof until Lemma A.1 in the Appendix. The technicalities arise due to the fact that calling the quantities “probability” is formally incorrect, and must be replaced by natural densities.

$$\begin{aligned} a_{\text{hom}}^{(\ell)}(f) &= \sum_{k=1}^{\infty} k \mathbb{P}(v_{\ell}(F(a, b)) = k) \\ &= \sum_{k=1}^{\infty} \mathbb{P}(v_{\ell}(F(a, b)) \geq k) = \sum_{(r, k) \in S(f, \ell)} \mathbb{P}((a : b) \equiv r \pmod{\ell^k}). \end{aligned} \quad (7)$$

Replacing $a_{\text{hom}}^{(\ell)}$ in formula $\alpha_{\ell}(f) = \frac{\deg \ell}{N(\ell) - 1} - \deg(\ell) a_{\text{hom}}^{(\ell)}(f)$ and using Equation 6 yields the desired result. \square

If f has only simple roots modulo ℓ , then by Hensel’s Lemma f has the same number of roots modulo every power of ℓ . We obtain the following formula.

Corollary 2.9. *Let $f \in \mathbb{F}_q[x][t]$ and ℓ a monic irreducible polynomial in $\mathbb{F}_q[t]$ such that the affine and projective roots of f modulo ℓ are simple and call n_{ℓ} their number. Then*

$$\alpha_{\ell}(f) = \frac{\deg \ell}{N(\ell) - 1} \left(1 - \frac{N(\ell)}{N(\ell) + 1} n_{\ell} \right). \quad (8)$$

2.2.2 The case of linear polynomials

Showing that $\alpha(f)$ converges is not trivial and, to our knowledge, it is not proven in the NFS case. Let us first show that $\alpha(g)$ converges for linear polynomials g . This requires the following classical identity.

Lemma 2.10. *(Chapter I, [Apo90]) Let μ denote Möbius’ function and let x be such that $|x| < 1$. Then we have*

$$\sum_{h \geq 1} \mu(h) \frac{x^h}{1 - x^h} = x.$$

Notation 2.11. We denote by I_k the number of monic degree- k irreducible polynomials in $\mathbb{F}_q[t]$.

Theorem 2.12. *Let $g \in \mathbb{F}_q[t][x]$ be such that $\deg_x g = 1$. Then*

$$\alpha(g) = \frac{1}{q - 1}.$$

Proof. Let L be the set of monic irreducible polynomials in $\mathbb{F}_q[t]$. We shall prove that, when b_0 goes to infinity, $\sum_{\ell \in L, \deg \ell \leq b_0} \alpha_{\ell}(g)$ tends to $\frac{1}{q-1}$. In Equation 8 one has $n_{\ell} = 1$ and therefore

$$\sum_{\ell \in L, \deg \ell \leq b_0} \alpha_{\ell}(g) = \sum_{\ell \in L, \deg \ell \leq b_0} \frac{\deg(\ell)}{N(\ell)^2 - 1} = \sum_{k \leq b_0} \frac{k I_k}{q^{2k} - 1}. \quad (9)$$

Since $k I_k = \sum_{h|k} \mu(h) q^{\frac{k}{h}}$, the series transforms into a double series for which we shall prove the absolute convergence and shall compute the sum:

$$\sum_{k \geq 1} \sum_{h|k} \frac{1}{q^{2k} - 1} \mu(h) q^{\frac{k}{h}}.$$

The absolute value of the term $\left| \mu(h) \frac{q^{\frac{k}{h}}}{q^{2k}-1} \right|$ is bounded by $\frac{q^{\frac{k}{h}}}{q^{2k}-1}$. It follows easily that the sum is bounded by $\frac{4}{q-1}$. Therefore, we can change the summation order:

$$\sum_{k \geq 1} \sum_{h|k} \mu(h) \frac{q^{\frac{k}{h}}}{q^{2k}-1} = \sum_{i \geq 1} \left(\sum_{h \geq 1} \mu(h) \frac{q^i}{q^{2hi}-1} \right). \quad (10)$$

Applying Lemma 2.10 to $x = \frac{1}{q^{2i}}$ leads to $\sum_{h \geq 1} \mu(h) \frac{1}{q^{2hi}-1} = \frac{1}{q^{2i}}$. This shows that, when b_0 goes to infinity, $\sum_{\deg \ell \leq b_0} \alpha_\ell(g)$ tends to $\sum_{i \geq 1} \frac{1}{q^i} = \frac{1}{q-1}$. \square

We conclude the subsection by showing the convergence of alpha. Let now \mathcal{C} be a (singular or non-singular) projective curve. Call $\mathcal{P}_k(\mathcal{C})$ the set of points of \mathcal{C} with coefficients in \mathbb{F}_{q^k} and $P_k(\mathcal{C})$ its cardinality. Next, call $\mathcal{P}'_k(\mathcal{C})$ the set of points in $\mathcal{P}_k(\mathcal{C})$ whose t -coordinate does not belong to a strict subfield of \mathbb{F}_{q^k} . Finally, we denote by $P'_k(\mathcal{C})$ its cardinality.

We shall need the following intermediate result.

Lemma 2.13. *Let \mathcal{C} be projective plane curve of degree d_0 defined over \mathbb{F}_q and let $g_0 = (d_0 - 1)(d_0 - 2)/2$ be its arithmetic genus. Then, for all $k \geq 1$,*

$$\left| P'_k(\mathcal{C}) - (q^k + 1) \right| < (4g_0 + 6)q^{\frac{k}{2}}. \quad (11)$$

Proof. Let $\tilde{\mathcal{C}}$ be a non-singular model for \mathcal{C} and let g be its geometric genus. We apply the Hasse-Weil Theorem and obtain:

$$\left| P_k(\tilde{\mathcal{C}}) - (q^k + 1) \right| \leq 2g \cdot q^{\frac{k}{2}}. \quad (12)$$

Next, according to Chapter VI in [Ful69],

$$|P_k(\tilde{\mathcal{C}}) - P_k(\mathcal{C})| \leq g_0 - g. \quad (13)$$

Every point (t_0, x_0) of $\mathcal{P}_k(\mathcal{C}) \setminus \mathcal{P}'_k(\mathcal{C})$ is determined by the choice of: a) a strict divisor d of k , b) an element $t_0 \in \mathbb{F}_{q^{\frac{k}{d}}}$ and c) a root x_0 of $f(t_0, x)$ in \mathbb{F}_{q^k} . Therefore:

$$\left| P'_k(\mathcal{C}) - P_k(\mathcal{C}) \right| \leq \sum_{d|k, 1 < d \leq k} q^{\frac{k}{d}} \deg_x(f). \quad (14)$$

On the one hand $\deg_x f < g_0 + 3$; on the other hand, if one calls d_1 the smallest proper divisor of k , $\sum_{d|k, 1 < d \leq k} q^{\frac{k}{d}} = q^{\frac{k}{d_1}} \sum_{d|k, 1 < d \leq k} q^{\frac{k}{d} - \frac{k}{d_1}}$. Since $d | k$ and $d \geq d_1$, $\frac{k}{d} - \frac{k}{d_1}$ is a negative or null integer. Therefore this sum is bounded by $q^{\frac{k}{d_1}} \sum_{i \geq 0} q^{-i} \leq q^{\frac{k}{d_1}} \sum_{i \geq 0} 2^{-i} = 2q^{\frac{k}{d_1}}$. But $d_1 \geq 2$, so

$$\left| P'_k(\mathcal{C}) - P_k(\mathcal{C}) \right| \leq 2q^{\frac{k}{2}} \deg_x f < 2(g_0 + 3)q^{\frac{k}{2}}. \quad (15)$$

The result follows from Equations 12, 13 and 15. \square

The following theorem shows that $\alpha(f)$ is well defined. More specifically, call L the set of irreducible monic polynomials in $\mathbb{F}_q[t]$. We show the existence of alpha by showing the convergence of $v_{b_0} := \sum_{\ell \in L, \deg \ell \leq b_0} (\alpha_\ell(f) - \alpha_\ell(x))$.

Theorem 2.14. *Let $f \in \mathbb{F}_q[t][x]$ an absolutely irreducible separable polynomial. Then the sequence v_{b_0} defined above converges. If one defines alpha by $\alpha(f) = \lim_{b_0 \rightarrow \infty} \sum_{\deg \ell \leq b_0, \ell \in L} \alpha_\ell(f)$, then there exist explicit bounds on α that depend only on $\deg_x f$, $\deg_t f$ and q .*

Proof. Call $d_0 = \deg f$ the degree of f as a polynomial in two variables and $g_0 = (d_0 - 1)(d_0 - 2)/2$. Let L_0 be the set of irreducible divisors of $\text{Disc}(f) \cdot f_d$. Call b_0 the largest degree of elements in L_0 . Let $k > b_0$. We are in the case of Corollary 2.9, hence Equation 8 gives

$$\begin{aligned} & \sum_{\ell \in L, \deg \ell = k} \alpha_\ell(f) - \alpha_\ell(x) = \\ &= \frac{kq^k}{q^{2k} - 1} (\#\{(\ell, r) \mid \deg(\ell) = k, f(r) \equiv 0 \pmod{\ell}\} - I_k). \end{aligned}$$

Each pair (ℓ, r) as in the equation above corresponds to exactly k points on the curve \mathcal{C} associated to f . Indeed, each ℓ has exactly k distinct roots in \mathbb{F}_{q^k} . Hence we have $I_k = \frac{1}{k}P'_k(\mathbb{P}^1(\mathbb{F}_q))$ and $\#\{(\ell, r) \mid f(r) \equiv 0 \pmod{\ell}\} = \frac{1}{k}P'_k(\mathcal{C})$, and further:

$$|\#\{(\ell, r) \mid f(r) \equiv 0 \pmod{\ell}\} - I_k| = \frac{1}{k} \left| P'_k(\mathcal{C}) - P'_k(\mathbb{P}^1(\mathbb{F}_q)) \right|. \quad (16)$$

Finally, Lemma 2.13 applied to \mathcal{C} and $\mathbb{P}^1(\mathbb{F}_q)$ respectively gives

$$\left| P'_k(\mathcal{C}) - (q^k + 1) \right| \leq (4g_0 + 6)\sqrt{q^k} \quad (17)$$

$$\left| P'_k(\mathbb{P}^1(\mathbb{F}_q)) - (q^k + 1) \right| \leq 6\sqrt{q^k}, \quad (18)$$

where g_0 is the arithmetic genus of \mathcal{C} .

Hence $|\#\{(\ell, r) \mid f(r) \equiv 0 \pmod{\ell}\} - I_k| \frac{kq^k}{q^{2k} - 1} \leq (4g_0 + 12) \frac{q^k}{q^{2k} - 1} \sqrt{q^k}$. The series $\sum_{k \geq 1} \frac{q^k}{q^{2k} - 1} \sqrt{q^k}$ is equivalent to the series $\sum_k \sqrt{q}^{-k}$ which converges. Therefore the sequence v_{b_0} converges when b_0 tends to infinity.

For a given pair d_0 and q , one can clearly bound the set L_0 . For all $\ell \in L_0$, by Proposition 2.8, $S(f, \ell)$ is formed by at most $\deg_x f \leq d_0$ infinite sequences and a finite number of additional elements. We are thus left with finding a bound for the roots which do not extend into ℓ -adic roots. By Hensel's Lemma, if a root (r, k) does not lift to an ℓ -adic one, we have $f'(r) \equiv 0 \pmod{\ell^k}$. This implies $\text{Disc}(f) \equiv 0 \pmod{\ell^k}$ which gives a bound on k . Therefore, alpha admits an effective bound depending exclusively on q and d_0 . \square

Example 2.15. Following the proof of the previous theorem, we can not only find a bound on α , but also evaluate the speed of convergence. Take $q = 2$, and let $f \in \mathbb{F}_q[t][x]$ such that $\deg_x f = 6$ and $\tilde{g} = 19$ and suppose that L_0 contains only polynomials of degree less than 15. Using Equations 16 and 17 in the proof above and the exact formula for I_k we can prove that $\alpha(f)$ is computed up to an error of 0.567 if we sum polynomials ℓ up to degree 15 and we reduce the error to 0.097 if we go to degree 20.

2.3 Cancellation property-Laurent roots

Consider the polynomial $f = x^3 + t^2x + 1 \in \mathbb{F}_2[t][x]$. For all $(a, b) \in \mathbb{F}_2[t]^2$, if no cancellation occur, the degree of $F(a, b) = a^3 + t^2ab^2 + b^3$ is $\max(\deg a^3, \deg(t^2ab^2), \deg b^3)$. One can easily check that the degree of $F(a, b)$ is lower than this value if and only if $\deg a - \deg b$ equals 1 or -2 . Moreover, in the first case the decrease is at least 2 while in the second case it is at least 1. These conditions can be better explained thanks to the Laurent series.

We call Laurent series (in $1/t$) over \mathbb{F}_q any series $\sum_{n \geq n_0} a_n \frac{1}{t^n}$ with $n_0 \in \mathbb{Z}$ and coefficients a_n in \mathbb{F}_q . We make the common convention to call degree of a rational fraction f_1/f_2 with $f_1, f_2 \in \mathbb{F}_q[t]$ the difference $\deg f_1 - \deg f_2$. The degree of a Laurent series is then defined as the degree of any of its nonzero truncations e.g $t + 1 + 1/t^2 + \dots$ has degree 1. Equivalently, it is the opposite of the valuation of the Laurent series in $1/t$. We call Laurent polynomial a pair (r, m) such that $r \in \mathbb{F}_q(t)$, m is an integer and r is the sum of a Laurent series whose terms are null starting from index $m + 1$. We may also use $r + O(\frac{1}{t^{m+1}})$ for writing the Laurent polynomial (r, m) .

Formally, a pair (a, b) has a “decrease in degree” if $\max_i \deg(f_i a^i b^{d-i})$ is strictly larger than $\deg F(a, b)$. Note that $F(a, b)$ has a decrease in degree if and only if $b \neq 0$ and the first terms of the Laurent series $\frac{a}{b}$ match those of a Laurent polynomial r with the property given in the following definition.

Definition 2.16. Let $f \in \mathbb{F}_q[t][x]$ be a polynomial and call d its degree in x . Let (r, m) be a Laurent polynomial. We say that (r, m) is a Laurent root of f if

$$\max_{i \in [0, d]} \deg(f_i r^i) - \deg f(r) > 0. \quad (19)$$

We call gap of (r, m) the least value in the left hand side of the inequality above when we replace r by any Laurent series extending r . A Laurent series such that all its truncations are Laurent roots is called an infinite Laurent root.

In the example above, $\frac{1}{t^2} + \frac{1}{t^8} + O(\frac{1}{t^9})$ is a Laurent roots of gap 7 and it extends into an infinite Laurent root. Also $t + O(1)$ is a Laurent root of gap 2 that is not the truncation of any Laurent root with a larger m . It also shows that the gap is not directly connected to the number of terms in the Laurent polynomial.

2.3.1 Computation

One can compute every Laurent root in two steps. First, one computes the Laurent roots of type λt^δ with λ is in \mathbb{F}_q and δ is an integer. For this, call Newton polygon of f , with respect to valuation $-\deg$, the convex hull of $\{(d-i, \deg(f_i)) \mid i \in [0, d]\}$. Chapter II in [Neu99] shows that δ must be an integer slope of the Newton polygon of f . Next, to extend a Laurent root $r = a_{n_0} t^{n_0} + \dots + a_m \frac{1}{t^m}$ with $a_m \neq 0$ to a root with precision larger than m , one computes the Laurent roots λt^δ of $f(x+r)$ for which δ is an integer such that $\delta < -m$. Note that this corresponds to make a Hensel lift with respect to the valuation $-\deg$.

In order to compute the gap of a Laurent root (r, m) , note that in Equation 19 the term $\max_{i \in [0, d]} \deg(f_i r^i)$ depends only on the leading term of r . Hence the problem is reduced to that of computing the maximal degree of $f(R)$ for the Laurent series R which extend r . For this, one sets an upper bound and then tests Laurent polynomials with increasingly more terms and reduces the upper bound until they produce a certificate.

2.3.2 Definition of α_∞ .

For each Laurent root (r, m) , we can compute the proportion of pairs (a, b) on a sieve domain such that the first terms of a/b match (r, m) . Recall that a sieve domain of sieve parameter e and skewness s corresponds to all the pairs (a, b) with $\deg a \leq \lfloor e + s/2 \rfloor$ and $\deg b \leq \lfloor e - s/2 \rfloor$.

Lemma 2.17. Let $r + O(\frac{1}{t^{m+1}})$ be a Laurent root of $f \in \mathbb{F}_q[t][x]$. Call $N_r = \deg(r) + m$, the number of terms of r other than the leading one. Then the proportion of pairs (a, b) on a domain of sieve parameter $e \geq N_r + |\deg r|$ and skewness parameter s such that the Laurent series a/b matches $r + O(\frac{1}{t^{m+1}})$ is:

$$\mathbb{P} \left(\frac{a}{b} = r + O \left(\frac{1}{t^{m+1}} \right) \right) = \frac{q^{-N_r - |s - \deg r|}}{q + 1} (1 + O_{e \rightarrow \infty}(1/q^{2e})).$$

Proof. The proportion of pairs (a, b) such that $\deg a - \deg b = \deg r$ is approximated by $(\frac{q-1}{q})^2 \cdot \sum_{i \geq 0} \frac{1}{q^{2i}} q^{-|s - \deg r|} = \frac{q-1}{q+1} q^{-|s - \deg r|}$. See Figure 1 for an illustration. The relative error made is $O(1/q^{2e})$ and corresponds to the fact that the series $\sum \frac{1}{q^{2i}}$ must be truncated at $i = e - |\deg r|$ and to the fact that for $\deg a, \deg b < |s|$ there is no pair (a, b) such that $\deg a - \deg b = s$. Next, only a fraction $\frac{1}{q-1}$ of these pairs have leading coefficients such that $a/b = r(1 + O(\frac{1}{t}))$. Finally, when $a/b = r(1 + O(\frac{1}{t}))$, the condition $a/b = r + O(\frac{1}{t^{m+1}}) = r(1 + O(\frac{1}{t^{N_r+1}}))$ can be expressed as a system

of N_r linear equations, that is triangular on the variables of a . Therefore, a pair with $\deg a - \deg b = \deg r$ and $a/b = r(1 + O(1/t))$ has probability q^{-N_r} to satisfy $a/b = r(1 + O(\frac{1}{t^{N_r+1}}))$. This leads to the proportion announced in the statement. Note that the condition $e > N_r + |\deg r|$ guarantees that the polynomials a and b have sufficiently many coefficients so that the linear conditions make sense. \square

We can now define a function which, for large sieve domains, measures the average degree gained due to the cancellations.

Definition 2.18. For any Laurent root $r + O(\frac{1}{t^{m+1}})$ we call $\text{trunc}(r)$ the Laurent polynomial obtained by deleting the term in $\frac{1}{t^m}$. If $\text{trunc}(r) \neq 0$ we write $\gamma(r, m)$ for the gap of $r + O(\frac{1}{t^{m+1}})$ minus the gap of $\text{trunc}(r) + O(\frac{1}{t^m})$. Otherwise $\gamma(r, m)$ is the gap of $r + O(\frac{1}{t^{m+1}})$. We call alpha infinity the following quantity

$$\alpha_\infty(f, s) := \sum_{(r,m) \text{ Laurent root}} -\gamma(r, m) \frac{1}{q+1} q^{-N_r - |s - \deg(r)|}. \quad (20)$$

Consider the case when all the Laurent roots of f extend infinitely into the Laurent series r_1, r_2, \dots, r_h . If, for all i , each new term of a/b which matches r_i increases the gap by one, then we obtain the simpler formula below.

$$\alpha_\infty(f, s) = \frac{-q}{q^2 - 1} \sum_{i=1}^h q^{-|s - \deg r_i|}. \quad (21)$$

As a particular case, it is clear that the polynomials of degree 1 in x have exactly one infinite Laurent root. If one sets the skewness $s = \deg_t f$, then $\alpha_\infty(f, s) = -\frac{q}{q^2 - 1}$. As a second example, a degree-6 polynomial f over \mathbb{F}_2 which has infinite Laurent roots of degrees 3, 2, 1, 0, -1 and -2 has $\alpha_\infty(f, 0) = -1.75$.

Example 2.19. A special class of polynomials are those corresponding to $\mathcal{C}_{a,b}$ curves, which were proposed for the FFS in [Mat99]. If f is a $\mathcal{C}_{a,b}$ polynomial and if we denote $a = \deg_x f$ and $b = \deg_t f_0 - \deg f_a$, then for all $i \in [0, a - 1]$ we have $\deg f_i < \deg f_a + (a - i)\frac{b}{a}$.

Suppose that a $\mathcal{C}_{a,b}$ polynomial f had a Laurent root r . If $\deg r < \frac{b}{a}$, then $\max\{\deg f_i r^i \mid i \in [1, a]\} < \deg f_0$. If $\deg r \geq \frac{a}{b}$ then $\deg f_a r^a$ dominates all the other terms of $f(r)$, so f has no Laurent roots. Hence, for any s , $\alpha_\infty(f, s) = 0$.

2.3.3 Constructing polynomials with many Laurent roots

One can easily check that, if a polynomial $f = \sum_i f_i x^i$ satisfies $\deg(f_d) = 0$, $\deg(f_{d-1}) = s$ for some $s > 0$ and $\deg(f_i) < (d - i)s$ for all $i \in [0, d - 2]$, then f has a Laurent root of degree s . This can be generalized to up to d roots.

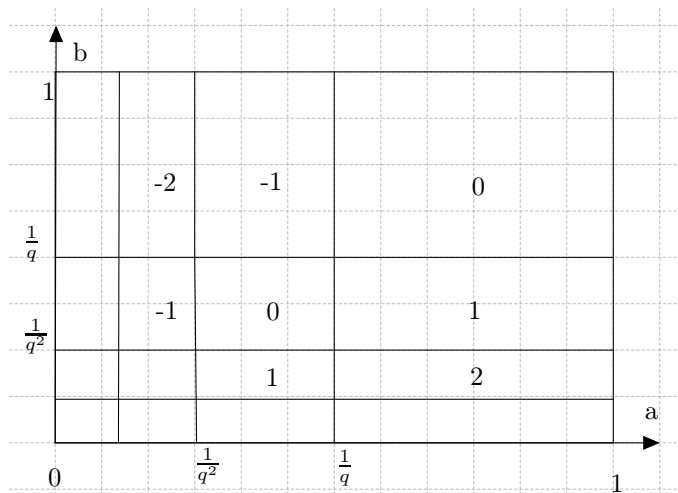
For any edge $(v_i, i) \leftrightarrow (v_j, j)$ of the Newton polygon we call length the quantity $|i - j|$.

Proposition 2.20. (Section II.6, [Neu99]) *If $f \in \mathbb{F}_q[t][x]$ is a polynomial, each edge of length 1 in the Newton polygon corresponds to an infinite Laurent root for f .*

For example, the polynomial $f = x^7 + t^3 x^6 + t^5 x^5 + (t^6 + 1)x^4 + t^6 x^3 + (t^5 + t + 1)x^2 + t^3 x + 1$ has a Newton polygon with 7 edges of length 1 and therefore 7 infinite Laurent roots.

Note however that the converse of the Proposition 2.20 is false in general: a polynomial might have infinite Laurent roots which cannot be counted using the Newton polygon, e.g. for the polynomial f above, $f^1 := f(x + t^4)$ also has 7 infinite Laurent roots, although its Newton polygon has no edge of length 1.

Figure 1: A domain of pairs (a, b) in lexicographical order having skewness S . We write the quantity $\deg a - \deg b - S$ for each region.



3 The combined effect of the properties which influence the sieving efficiency

The previous sections identified three elements which affect the sieve and defined associated measures: σ for the size property, α for the root property and α_∞ for the cancellation property. The list might be extended with other properties and the quantifying functions can be combined in different fashions in order to compare arbitrary polynomials. In this section we define two general purpose functions based on the three properties above and show their relevance through experimentation.

3.1 Adapting Murphy's \mathbb{E} to the FFS

As a first function which compares arbitrary polynomials, we adapt Murphy's \mathbb{E} , already used for the NFS algorithm (Equation 5.7, [Mur99]). Heuristically, \mathbb{E} uses Dickman's ρ to approximate the number of relations found by f and g on a sieving domain.

Definition 3.1. Let $f, g \in \mathbb{F}_q[t][x]$ be two irreducible polynomials, s an integer called skewness parameter, e a half integer called sieve parameter and β an integer called smoothness bound. Let $D(s, e)$ be the set of coprime pairs $(a, b) \in \mathbb{F}_q[t]^2$ such that $0 \leq \deg(a) \leq \lfloor e + \frac{s}{2} \rfloor$ and $0 \leq \deg(b) \leq \lfloor e - \frac{s}{2} \rfloor$. We define:

$$\mathbb{E}(f, g, s, e, \beta) = \sum_{(a,b) \in D(s,e)} \rho\left(\frac{\deg F(a, b) + \alpha(f)}{\beta}\right) \cdot \rho\left(\frac{\deg G(a, b) + \alpha(g)}{\beta}\right).$$

Unlike the situation in the NFS case, where \mathbb{E} must be approximated by numerical methods, in the case of the FFS one can compute \mathbb{E} in polynomial time with respect to $\deg(f)$, $\deg(g)$ and $e + |s|$. Note that ρ can be evaluated in polynomial time to any precision on the interval which is relevant in this formula.

We recall that we focus in this work on the case where the polynomial g is linear, as in [JL02]. More precisely, we assume that $g = g_1x + g_0$ is chosen with g_0 of much higher degree than g_1 . In this case, the algorithm goes as follows:

1. compute the Laurent roots of f up to $\lfloor d(e + \frac{s}{2}) + \deg_t f \rfloor$ terms;

Table 1: Choosing the best skewness using \mathbb{E} . The parameters are set to $e = 24.5$ and $\beta = 22$.

s	-1	1	3	5	7
$10^{-5}E(f, g, s, e, \beta)$	2.54	3.31	3.46	2.88	2.12

2. for each $d_a \leq \lfloor e + \frac{s}{2} \rfloor$, $d_b \leq \lfloor e - \frac{s}{2} \rfloor$ and $i \in \mathbb{N}$, use Lemma 2.17 to compute the number $n(d_a, d_b, i)$ of pairs (a, b) such that $\deg(a) = d_a$, $\deg(b) = d_b$ and $\deg(F(a, b)) = i$;

3. compute

$$\sum_{d_a, d_b, i} n(d_a, d_b, i) \rho\left(\frac{i + \alpha(f)}{\beta}\right) \cdot \rho\left(\frac{d_b + \deg g_0 + \alpha(g)}{\beta}\right). \quad (22)$$

One can use Murphy's \mathbb{E} to choose the optimal skewness corresponding to a pair (f, g) of polynomials.

Example 3.2. Consider for instance the two polynomials used for the computation of the discrete logarithm in $GF(2^{619})$: $f = x^6 + (t^2 + t + 1)x^5 + (t^2 + t)x + 0x152a$ and $g = x - t^{104} - 0x6dbb$ written in hexadecimal¹ notation [BBD⁺12]. They used the smoothness bound 22 and most of the computations were done using special-Q's (q, r) with $\deg q = 25$. The pairs (a, b) considered for each special-Q were $(ia_0 + ja_1, ib_0 + jb_1)$ with (a_0, b_0) and (a_1, b_1) two pairs on the special-Q lattice and i, j were polynomials of degree at most 12. Hence, the pairs (a, b) considered were such that $\deg a + \deg b = \deg q + 24$, so $e = (\deg a + \deg b)/2 = 24.5$. Note that, if a and b have maximal degree on our set, the difference $\deg a - \deg b$ cannot be even. Table 1 shows that the best skewness value is 3.

Note though that in [BBD⁺12] one started by experimentally choosing the best skewness for polynomials of a given bound on the degrees. Then they selected polynomials which, for a given value of s , minimize the value of epsilon, the function that we define below.

An alternative to \mathbb{E} : Epsilon Recall that σ is the degree of the norm when no cancellation occurs, α is the degree gained due to the modular roots and α_∞ is the degree gained thanks to cancellations. It seems natural that their sum is the degree of a polynomial which has the same skewness probability as $F(a, b)$ for an "average" pair (a, b) on the sieving domain.

Definition 3.3. For a polynomial $f \in \mathbb{F}_q[t][x]$, a skewness parameter s and a sieve parameter e , we call epsilon the following average degree

$$\epsilon(f, s, e) = \alpha(f) + \alpha_\infty(f, s) + \sigma(f, s, e).$$

Epsilon can be used to estimate the speedup of a polynomial with good properties. For example if the smoothness bound is 28 and two polynomials have the value of epsilon equal to 107 and 109 respectively, then we expect a speedup of $\rho(107/28)/\rho(109/28) \approx 1.19$.

Comparing ϵ and \mathbb{E} Since the subroutines necessary in the computation of ϵ are equally used when evaluating \mathbb{E} , in practice epsilon is faster to compute than \mathbb{E} . The advantage of \mathbb{E} is that it is more precise, but the experiments of the next section will show that epsilon is reliable enough.

¹Each polynomial ℓ of $\mathbb{F}_2[t]$, $\ell = \sum_i \ell_i t^i$ with $\ell_i \in \{0, 1\}$, is represented by base-16 notation of the integer $\sum_i \ell_i 2^i$.

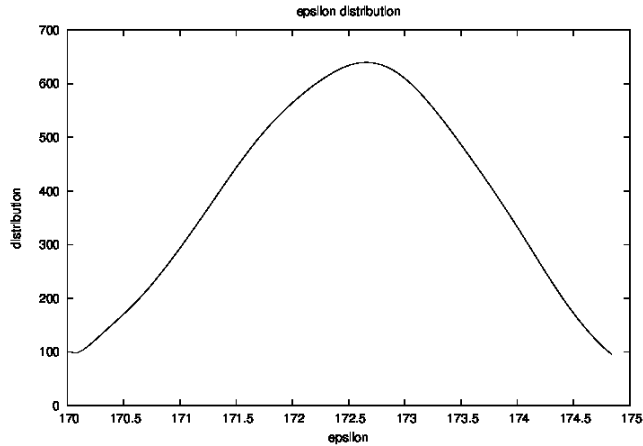


Figure 2: Distribution of epsilon on a sample of 20000 polynomials of $\mathbb{F}_2[t][x]$ of degree 6 in x and 12 in t .

3.2 Experimental validation

The implementation used in the experiments is the one described in [DGV13], which is freely available at [BFG⁺]. To our knowledge, no other press-a-button implementation of the FFS is publicly available. In addition, this implementation does relatively few modifications which could loose relations, making a theoretical study inexact.

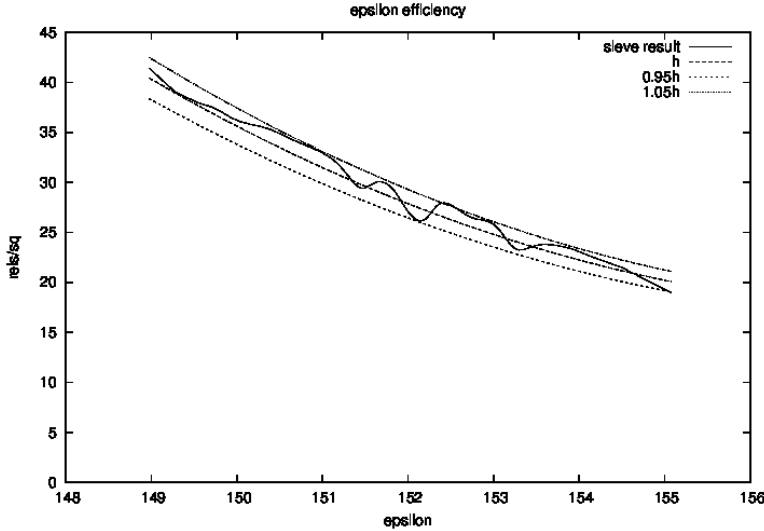
The real-life efficiency of a polynomial is measured either by the number of relations per second, by the total number of relations, or as the average number of relations per special-Q. We kept the last one as a measure of efficiency since the software offers an option to reliably measure it and because it considers only the polynomial properties rather than the implementation quality.

Experiment 1. We selected a sample of polynomials f after evaluating epsilon for a range of polynomials considered one after another in lexicographical order starting from $x^6 + (t^2 + t + 1)x^5 + (t^3 + t^2 + t + 1)x^3 + tx + t^{11}$. Note that the choice of the starting point guarantees that the polynomials considered have at least one infinite Laurent root. Since the distribution of epsilon was that of Figure 2, most of the polynomials tested had values of epsilon in a narrow interval. This lead us to select only one polynomial in each interval of length 0.01, to a total of 119 polynomials. Next we extended the sample with 60 polynomials starting from $x^6 + t^3x^5 + (t^5 + 1)x^4 + t^6x^3 + t^6 + 1$. For each polynomial f we associated a random monic linear polynomial g suited to the FFS, having degree in t equal to 104. Indeed, as shown in Theorem 2.12 and in section 2.3.2 respectively, linear polynomial have the same values of α and α_∞ respectively.

We set the parameters as follows: $I = J = 12$, $\text{fbb0} = \text{fbb1} = 22$, $\text{lpb0} = \text{lpb1} = 28$, $\text{thresh0} = \text{thresh1} = 100$, $\text{sqside} = 1$. The polynomials g used in the special-Q technique were the first irreducible ones starting from t^{25} . We called the option “reliablenrels” which tests as many values of q as needed in order to obtain a measurement error of $\pm 3\%$ with a confidence level of 95%. The skewness parameter was set to $S = 3$ because, for the finite fields where the degree-6 polynomials are optimal, this is a sensible choice. Finally, the parameter sqt was set to 1 so that, for most special-Q’s, the sieving domain was such that $\deg a \leq 26$ and $\deg b \leq 23$.

The results plotted in Figure 3 indicate that the sieve efficiency is not far from a strictly decreasing function of epsilon. To illustrates this, we plotted a decreasing function that fits our results, such that the relative error of our measurements is always less than 5%.

Figure 3: Epsilon and sieve efficiency for the polynomials f in Experiment 1. The function h is a function of type $a + bx + c \log x$, with no special significance.



Finally, one can see that a sensible choice of the polynomial can save a factor 2 in the sieve time when compared to a bad choice.

3.3 Correlation between f and g

A standard heuristic states that the probabilities of $F(a, b)$ and $G(a, b)$ to be smooth are independent, e.g. Murphy's \mathbb{E} multiplies the two probabilities. The inexactness of this approximation could be called correlation property. According to Experiment 1, the correlation property has a small effect on the sieve, so that we bound ourselves to illustrate it by an example and a practical experiment.

Example 3.4. Let $f = x^2 - t^2(t + 1)$, $g_1 = x - (t^2 + t)^7$ and $g_2 = x - (t^2 + t + 1)^7$. Let F , G_1 and G_2 be the homogenizations of f , g_1 and g_2 respectively. Note that, for coprime pairs $(a, b) \in \mathbb{F}_q[t]$, $F(a, b)$ is divisible by t if and only if $a \equiv 0 \pmod{t}$. For these pairs we have $G_1(a, b) \equiv 0 \pmod{t}$ whereas $G_2(a, b) \not\equiv 0 \pmod{t}$. In short, g_1 increases the number of doubly smooth pairs whereas g_2 that of pairs which are smooth on the rational or the algebraic side, but not on both.

If $\deg_x g = 1$, then every prime power ℓ^k of $\text{Res}(f, g)$ implies a correlation between the events $\ell^k \mid F(a, b)$ and $\ell^k \mid G(a, b)$ on a domain of pairs (a, b) . Table 2 summarizes an experiment in which we compared different pairs (f, g) with f having similar values of ϵ . We selected three polynomials f of the form $f = \tilde{f} + f_0$ with $\tilde{f} = x^6 + tx^5 + (t+1)x^4 + (t^2+t+1)x^3$ and we associated to each one a linear polynomial g of the form $g = \tilde{g} + g_{00}$ with $\tilde{g} = x - t^{104} - t^{14} + t^{13} + t^{11} + t^{10} + t^8$. Instead of imposing that $\text{Res}(f, g)$ has an irreducible factor of degree 619 as in the previous experiment, we aimed to find polynomials g such that $\text{Res}(f, g)$ has first no, then many, small factors. The experiment indicated that the correlation property explains part of the error observed in Experiment 1, bringing it close to 3%, which is equal to our measurement error.

Since it is easy to associate many linear polynomials g to a unique f and since linear polynomials have the same value of epsilon, it is interesting to select g such that $\text{Res}(f, g)$ has many small factors. Nevertheless, f and g are chosen such that $\text{Res}(f, g)$ has degree $d(\lfloor \frac{n}{7} \rfloor + 1)$ and we require

Table 2: Influence of the linear polynomials g on the sieve efficiency. $f = \tilde{f} + f_0$ and $g = \tilde{g} + g_{00}$ with \tilde{f} and \tilde{g} given in section 3.3.

f_0	g_{00}	small factors of $\text{Res}(f, g)$	efficiency (rels/sq)
0x19	0xb2	–	3.3
0x12	0xbf	–	3.4
0x12	0xb8	–	3.4
0x12	0xae	$t \cdot (t + 1)$	3.5
0x12	0xa0	$t \cdot (t + 1)$	3.5
0x12	0xbb	$t \cdot (t + 1) \cdot (t^3 + t^2 + 1)$	3.5
0x1e	0xbb	$t \cdot (t + 1) \cdot (t^3 + t + 1) \cdot (t^6 + \dots)$	3.6

an irreducible factor of degree n , leaving little room for additional factors. Moreover, when f imposes an extra factor to $\text{Res}(f, g)$ (for example by having 1 projective and q affine roots modulo t), depending on the congruence of n modulo d , it can be impossible to choose g of optimal degree in t . See 5.2 for an example.

3.4 A sieve algorithm for alpha

After we showed the relevance of epsilon, the polynomial selection comes to evaluating epsilon on a large set of polynomials. One can try various ranges of polynomials $f = \sum f_i x^i$, given by some degree bounds on their coefficients f_i , which optimize sigma and/or impose a number of Laurent roots, as shown in 2.3.3. The most time-consuming part of the computations, the evaluation of alpha, can be done on each range by a sieving procedure.

The idea is that, for each irreducible polynomial $\ell \in \mathbb{F}_q[t]$, we compute α_ℓ for all the residue polynomials f of $\mathbb{F}_q[t][x]$ modulo ℓ and then we update the values of α_ℓ for all the polynomials f in the range.

Let d, e_0, \dots, e_{d-1} and e_d be integers. We consider the range of the polynomials $f = \sum_{i=0}^d f_i x^i \in \mathbb{F}_q[t][x]$ such that for $i \in [0, d]$, $\deg_t f_i \leq e_i$. Call H the set of values taken by the tuple (f_d, \dots, f_2) and T those taken by (f_1, f_0) . Let L be the set of irreducible polynomials up to a given bound. Let k_{\max} be a parameter and let us suppose that, for all $\ell \in L$, all the roots $r \pmod{\ell^k}$ with $\deg(\ell^k) \geq k_{\max}$ extend indefinitely.

For an irreducible polynomial ℓ , Algorithm 1 below computes $\alpha_\ell(f)$ for all f in the range and can be a subroutine to computing $\alpha(f)$ for the same range. We denote by $\text{residues}(\ell^k)$ the set of polynomials in $\mathbb{F}_q[t]$ of degree at most $k \deg \ell - 1$.

The correctness of Algorithm 1 follows from Proposition 2.8. For a fixed value of k_{\max} , the complexity per polynomial is $O(1)$, as the most time-consuming steps are those in lines 8 and 10. For comparison, in the naive algorithm, for each polynomial, one needs to find the roots modulo ℓ , which takes a non-constant polynomial time in $d + \deg(\ell)$. In practice, Algorithm 1 showed to be much faster, as Paul Zimmermann used it to compute $\alpha_\ell(f)$, for all the irreducible polynomials ℓ with $\deg \ell \leq 6$, on the range of the 2^{48} monic polynomials $f \in \mathbb{F}_2[t][x]$ of degree 6 such that for $i \in [0, 6]$, $\deg_i f_i \leq 12 - 2i$.

4 Sieving with inseparable polynomials

4.1 Particularities of the inseparable polynomials

Despite the possibility of adding new technicalities, the inseparable polynomials have been preferred in two record computations [HSW⁺10], [HSST12]. Moreover, the Coppersmith algorithm, implemented in [Tho03], can be seen as a particular case of the FFS, using inseparable polynomials.

Algorithm 1 The alpha sieve

```
1: Initialize  $\alpha_\ell$  to a vector of value  $\deg(\ell)/(\mathbf{N}(\ell) - 1)$ 
2:  $k_0 \leftarrow \lceil k_{\max}/\deg \ell \rceil$ 
3: for  $(f_d, f_{d-1}, \dots, f_2)$  in  $H$  do
4:   for  $k$  in  $[1..k_0]$  and  $r$  in residues( $\ell^k$ ) do
5:     for  $(f_1, f_0)$  in  $T$  such that  $f_1 r + f_0 \equiv -\sum_{i=2}^d f_i r^i \pmod{\ell^k}$  do
6:        $f \leftarrow \sum_{i=0}^d f_i x^i$ 
7:       if  $k < k_0$  then
8:          $\alpha_\ell(f) \leftarrow \alpha_\ell(f) - \deg(\ell) \mathbf{N}(\ell)^{1-k}/(\mathbf{N}(\ell) + 1)$ 
9:       else
10:         $\alpha_\ell(f) \leftarrow \alpha_\ell(f) - \deg(\ell) \mathbf{N}(\ell)^{2-k}/(\mathbf{N}(\ell)^2 - 1)$ 
11:       end if
12:     end for
13:   end for
14: end for
```

In order to present the Coppersmith algorithm from this point of view and in order to compare inseparable polynomials to separable ones, we start with their definition, followed by their main properties.

Definition 4.1. An irreducible non-constant polynomial $f \in \mathbb{F}_q[t][x]$ is said inseparable if $f' = 0$, where f' denotes the derivative with respect to x .

For every inseparable polynomial f , there exists a power of the characteristic of \mathbb{F}_q , d , and a polynomial $\hat{f} \in \mathbb{F}_q[t][x]$ such that $f = \hat{f}(x^d)$ and $\hat{f}' \neq 0$. This simple property allows us to factor any irreducible polynomial ℓ in the function field of f in two steps. First we factor ℓ in the function field of \hat{f} , then we further factor every prime ideal \mathfrak{l} of \hat{f} . The main advantage is that some prime ideal factorization algorithms work only for separable polynomials (for example Magma implements the function fields only in the case of separable polynomials [BCP97]). The factorization of the ideals \mathfrak{l} of \hat{f} in the function field of f is easy using the following result.

Proposition 4.2. (Corollary X.1.8, [Lor96]) Let $p > 0$ be a prime and q and d two powers of p . Let $\hat{K}/\mathbb{F}_q(t)$ be a function field. Let K/\hat{K} be an extension of polynomial $x^d - \theta_1$ with $\theta_1 \in \hat{K}$. Then every prime ideal \mathfrak{l} of \hat{K} decomposes as

$$\mathfrak{l}\mathcal{O}_K = \mathfrak{L}^d \tag{23}$$

for some prime ideal \mathfrak{L} such that $\mathfrak{L} \cap \mathcal{O}_{\hat{K}} = \mathfrak{l}$.

In the FFS algorithm, it is required to compute for each smooth element $a - b\theta$ of the function field of f , the valuation of every prime ideal \mathfrak{L} in the factor base. For this, we start by factoring $(a - b\theta)^d$ in the integer ring of the function field \hat{K} of \hat{f} :

$$(a - b\theta)^d \mathcal{O}_{\hat{K}} = (a^d - b^d \theta_1) \mathcal{O}_{\hat{K}} = \prod_i \mathfrak{l}_i^{e_i}$$

and then we obtain $(a - b\theta)\mathcal{O}_K = \prod_i \mathfrak{L}_i^{e_i}$ where the \mathfrak{L}_i are such that $\mathfrak{l}_i \mathcal{O}_K = \mathfrak{L}_i^d$.

4.2 Speed-up in the FFS due to the inseparability

Definition 4.3. Let f and g be two polynomials of $\mathbb{F}_q[t][x]$ such that $\text{Res}(f, g)$ has an irreducible factor of degree n . Assume that $\deg_x g = 1$ and write $f = \hat{f}(x^d)$ for some separable polynomial \hat{f} and some integer d which is either 1 or a power of $\text{char}(\mathbb{F}_q)$. We call free relation any irreducible polynomial $\ell \in \mathbb{F}_q[t]$ such that $\ell \nmid \text{Disc}(\hat{f})f_d$ and $(f \pmod{\ell\mathbb{F}_q[t][x]})$ splits into degree-1 factors.

Table 3: Number of free relations of a pair f, g with $f = \hat{f}(x^d)$, \hat{f} separable and $\deg(g) = 1$. N is the number of irreducible monic polynomials of degree below the smoothness bound. The computations assume that $\#\text{Gal}(\hat{f}) = (\deg \hat{f})!$.

$\text{char}(\mathbb{F}_q)$	d	$\deg(\hat{f})$	$\#\{\text{factor base}\}$	$\#\{\text{free relations}\}$
any	1	6	$2N$	$N/720$
2	2	3	$2N$	$N/6$
3	3	2	$2N$	$N/2$
any	1	8	$2N$	$N/40320$
2	2	4	$2N$	$N/24$
2	4	2	$2N$	$N/2$
2	8	1	$2N$	N

Clearly each free relation of norm less than the smoothness bound creates an additive equation between the virtual logarithms of the ideals in the factor base.

The number of free relations is given by Chebotarev's Theorem as follows. First note that, due to Proposition 4.2, a polynomial ℓ is a free relation for f if and only if it is a free relation for \hat{f} . Then, the proportion of free relations among the irreducible polynomials is, according to Chebotarev's Theorem, asymptotically equal to the inverse of the cardinality of the Galois group of the splitting field of \hat{f} . Call N the number of monic irreducible polynomials in $\mathbb{F}_q[t]$ of degree less than the smoothness bound. Then, the number of free relation is:

$$\#\{\text{free relations}\} = \frac{N}{\#\text{Gal}(\hat{f})}. \quad (24)$$

We compare this to the cardinality of the factor base. Since the cardinality of the rational side is N and f has as many ideals as \hat{f} , it is enough to evaluate the cardinality of the algebraic side. According to Chebotarev's Theorem, the number of pairs (ℓ, r) such that $f(r) \equiv 0 \pmod{\ell}$, $\deg r < \deg \ell$ and $\deg \ell$ is less than the smoothness bound is χN where χ is the average number of roots of \hat{f} fixed by the automorphisms of the splitting field of \hat{f} . It can be checked that each root of \hat{f} is fixed by a fraction $1/\deg(\hat{f})$ of the automorphisms, so $\chi = 1$. Hence, asymptotically the factor base has $2N + o(N)$ elements.

Heuristically, $\text{Gal} \hat{f}$ is the full symmetric group for all but a negligible set of polynomials \hat{f} , so most often we have $\#\text{Gal}(\hat{f}) = \deg(\hat{f})!$. We list the results for $\deg f$ equal to 6 and 8 in Table 3. The case in which $d = 3$ and $\deg \hat{f} = 2$ brought a $\frac{4}{3}$ -fold speedup in [HSW⁺10].

Coppersmith algorithm The case $d = 8$ and $\deg \hat{f} = 1$ corresponds to the Coppersmith algorithm. Indeed, since half of the relations are free relations, the sieve is accelerated by a factor of 2. Moreover, since $\deg(\hat{f}) = 1$, the free relations are particularly simple, linking exactly one element in the rational side to one element in the algebraic side of the factor base (Proposition 4.2). Therefore one can rewrite the relations using only the elements in the rational side, hence speeding up the linear algebra step by a factor of 4.

4.3 Root property of inseparable polynomials

Despite the fact that the inseparable polynomials are relatively few, being possible to exhaustively test them, it has its own interest to understand why inseparable polynomials have a bad root property and in particular, why the alpha value of many polynomials used in the Coppersmith algorithm is 2. Note that our proof that alpha converges covers only the case of separable polynomials. In this section we give some results on their root property.

First, the number of pairs (ℓ, r) with ℓ irreducible and r a polynomial of degree less than $\deg(\ell)$ such that $f(r) \equiv 0 \pmod{\ell}$ has a narrower range of values than it does for the separable polynomials. Indeed, as shown by the following result, this number corresponds to the number of roots of \hat{f} . The bounds in Theorem 2.14, when written explicitly, are narrower for polynomials of degree $\deg(\hat{f})$ than for those of degree $\deg f$. For example, if \hat{f} is linear, this number is a constant.

Lemma 4.4. *Let $\hat{f} \in \mathbb{F}_q[t][x]$ be a polynomial, d a power of the characteristic of \mathbb{F}_q and $f = \hat{f}(x^d)$. Let ℓ be an irreducible polynomial in $\mathbb{F}_q[t]$. Then there is a bijection between the sets $\{\hat{r} \in \mathbb{F}_q[t] \mid \deg \hat{r} < \deg \ell, \hat{f}(\hat{r}) \equiv 0 \pmod{\ell}\}$ and $\{r \in \mathbb{F}_q[t] \mid \deg r < \deg \ell, f(r) \equiv 0 \pmod{\ell}\}$.*

Proof. The non-null residues of ℓ form a group of cardinality $N(\ell) - 1$, which is coprime to q and hence to d . Therefore any root \hat{r} accepts one and only one d^{th} root modulo ℓ . \square

The second reason for having bad values of alpha is that most of the roots modulo irreducible polynomials ℓ do not lift to roots modulo ℓ^2 . Recall the following classical result.

Lemma 4.5. *Let $\ell \in \mathbb{F}_q[t]$ be an irreducible polynomial. Write $(\mathbb{F}_q[t]/\langle \ell^2 \rangle)^*$ for the group of residues modulo ℓ^2 which are not divisible by ℓ . Put $U = \{e^{N(\ell)} \mid e \in (\mathbb{F}_q[t]/\langle \ell^2 \rangle)^*\}$ and $V = \{1 + \ell w \mid \deg w < \deg \ell\}$. Then we have*

$$(\mathbb{F}_q[t]/\langle \ell^2 \rangle)^* \simeq U \times V.$$

The group U has order $N(\ell) - 1$ which is coprime to d , so d^{th} roots always exist and are unique in U . On the other hand, in V , only the neutral element is a d -th power. As a consequence only a fraction $1/\#V = 1/N(\ell)$ of the residues \hat{r} modulo ℓ^2 can have d^{th} roots modulo ℓ^2 . Let us make the heuristic that the roots of \hat{f} modulo ℓ^2 are random elements of $\mathbb{F}_q[t]/\langle \ell^2 \rangle$. Then only a small fraction of the roots of f lift modulo squares of irreducible polynomials and, for a non negligible fraction of polynomials f no root r modulo some irreducible polynomial ℓ lifts modulo ℓ^2 .

Among the Coppersmith polynomials f , i.e. such that \hat{f} is linear, many f are such that no modular root of f lifts modulo squares. Let us compute the value of alpha in this situation.

Lemma 4.6. *Let \hat{f} be a linear polynomial of $\mathbb{F}_q[t][x]$, d a power of the characteristic of \mathbb{F}_q and put $f = \hat{f}(x^d)$. Assume that there is no pair of polynomials ℓ and r with ℓ irreducible and r of degree less than that of ℓ^2 such that $f(r) \equiv 0 \pmod{\ell^2}$. Then we have*

$$\alpha(f) = \frac{2}{q-1}.$$

Proof. By Lemma 4.4, for all ℓ , f has exactly $n_\ell = 1$ affine or projective roots modulo ℓ . By Corollary 2.9, for all irreducible polynomial ℓ we have

$$\alpha_\ell(f) = \deg \ell \left(\frac{1}{N(\ell) - 1} - \frac{1}{N(\ell)} \frac{N(\ell)}{N(\ell) + 1} \right) = 2 \frac{\deg \ell}{q^{2 \deg \ell} - 1}.$$

Hence we obtain $\alpha(f) = 2 \left(\sum_{k \geq 1} \frac{k I_k}{q^{2k} - 1} \right)$ with I_k the number of irreducible monic polynomials of degree k in $\mathbb{F}_q[t]$. The sum in the parenthesis was computed in Proposition 2.12 and equals $1/(q-1)$. This completes the calculations. \square

5 Applications to some examples in the literature

5.1 Thomé's record using the Coppersmith algorithm

Thomé [Tho03] solved the discrete logarithm problem in $\mathbb{F}_{2^{607}}$ using the Coppersmith algorithm. Following this algorithm, one sets $g = x - t^{152}$ and $f = x^4 + t\lambda$ for some polynomial $\lambda \in \mathbb{F}_q[t]$ such

	$\alpha(f)$	$\alpha_\infty(f, s)$	$\sigma(f, s, e)$	$\epsilon(f, s, e)$	$E(f, g, s, e, \beta)$	efficiency
f_0	1.27	0	108.12	109.39	$1.82 \cdot 10^8$	15.2
f_1	-1.05	0	108.42	107.36	$2.10 \cdot 10^8$	18.8

Table 4: Coppermith polynomials for $\mathbb{F}_{2^{607}}$. The parameters in the table are $s = 7$, $e = 24.5$ and $\beta = 28$. The efficiency, measured in rels/sq, uses the parameters in Experiment 1.

f, g	$\alpha(f)$	$\alpha_\infty(f, s)$	$\sigma(f, s, e)$	$\epsilon(f, s, e)$	$E(f, g, s, e, \beta)$	efficiency
f_2, g_2	2.15	0	122.33	124.46	$8.54 \cdot 10^8$	66.0
f_3, g_3	-0.24	0	123.66	123.36	$8.64 \cdot 10^8$	73.8
f_4, g_4	-0.10	0	123.66	123.42	$9.49 \cdot 10^8$	76.0

Table 5: Classical FFS polynomials for $\mathbb{F}_{2^{607}}$. The parameters are $s = 1$, $e = 24.5$, $\beta = 28$. The efficiency, measured in rels/sq, uses the parameters in Experiment 1.

that $t^{607} + \lambda$ is irreducible. The polynomial $\lambda_0 = t^9 + t^7 + t^6 + t^3 + t + 1$ used by Thomé minimizes the degree of λ . If one searches for an alternative, it is necessary to increase $\deg_t f$, but this is possible without affecting much the size property. Indeed, the sensible choice is to set the skewness s to 7, so sigma does not vary much if one increases $\deg f_0$. By testing the polynomials λ with $\deg \lambda \leq 18$, we determined that the best alpha corresponds to $f_1 = x^4 + t(t^{16} + t^{12} + t^{11} + t^7 + t^4 + 1)$. We compare the two polynomials in Table 4 using the functions defined in this article as well as the sieve efficiency measured with the implementation of [DGV13] and the parameters in Experiment 1.

5.2 Joux-Lercier’s implementation of the classical variant of the FFS

Joux and Lercier [JL02, JL07] considered the fields \mathbb{F}_{2^n} with $n = 521, 607$ and 613 . For $n = 607$ they set $f_2 = x^5 + x + t^2 + 1$ and $g_2 = (t^{121} + t^8 + t^7 + t^5 + t^4 + 1)x + 1$. If one searches for an alternative, the sensible choice is to improve the root property without changing the size property. Since $\deg_t f_2 = 2$ we tested all the polynomials whose degree in t is 1 or 2.

Experiment 2. There are 2^{18} polynomials f such that $\deg_t f \leq 2$, out of which 2^{12} have $\deg_t f \leq 1$. There were 1776 irreducible polynomials with $\deg_t f \leq 1$ whose alpha is below 3. There were 650 irreducible polynomials f , with $\deg_t f \leq 2$, whose alpha is negative and such that the partial sum of alpha up to degree 6 is less than 0.5.

The best 10 values for ϵ with skewness $s = 0$ and sieve parameter $e = 24.5$ were all obtained for polynomials f with $\deg_t f = 2$. The best value was that of $f_3 = (t^2 + t)x^5 + (t^2 + t + 1)x^4 + (t + 1)x^3 + t^2x^2 + t^2x + t^2$. We could not associate a linear polynomial g with $\deg_t g = 121$ because $\text{Res}(f, g)$ is always divisible by t (see 3.3 for more details), hence we took $g_3 = x + t^{122} + t^{13} + t^{11} + t^6 + t^5 + t^3 + t^2$. The best f for which we could select a linear polynomial g of degree 121 was $f_4 = (t^2 + t + 1)x^5 + (t^2 + t + 1)x^4 + x^3 + (t^2 + t + 1)x^2 + (t^2 + t + 1)x + t^2 + t$, for which we took $g_4 = x + t^{121} + t^{12} + t^{11} + t^8 + t^6 + t^2 + 1$. In Table 5 we compare (f_3, g_3) and (f_4, g_4) to (f_2, g_2) . Note also that all the polynomials f tested have a small genus, which could explain the small variance of alpha when $\deg_t f \leq 2$.

5.3 Joux-Lercier’s two rational side variant

We recall briefly the “two rational sides” variant of [JL06], and study its properties according to our criteria. This variant selects two polynomials $f = \gamma_1(x) - t$ and $g = x - \gamma_2(t)$ for some γ_1 and γ_2 in $\mathbb{F}_q[t]$. Then one collects coprime pairs $(a(t), b(t)) \in \mathbb{F}_q[t]$ such that both

$$a(\gamma_1(x)) - xb(\gamma_1(x)) \text{ and } a(t) - \gamma_2(t)b(t) \quad (25)$$

f, g	$\alpha(f)$	$\alpha_\infty(f, s)$	$\sigma(f, s, e)$	$\epsilon(f, s, e)$	$E(f, g, s, e, \beta)$	efficiency
f_i	1.33	0	94.00	95.33	$1.03 \cdot 10^8$	14.6
f_s	0.29	0	94.75	95.04	$1.23 \cdot 10^8$	17.0
$f_{s'}$	-3.67	0	96.75	93.03	$1.61 \cdot 10^8$	21.3

Table 6: Polynomials f for fields of characteristic 3. The last column was obtained using same software as in Experiment 1 and with parameters $\text{fbb0}=\text{fbb1}=14$, $\text{lpb0}=\text{lpb1}=17$, $S=1$ and $q_0 = t^{15}$.

are β -smooth for some smoothness bound β . It can be easily checked that the expressions in Equation 25 have precisely the same degrees as in the classical FFS where we consider the norms of $a - bx$ with respect to the function fields of f and g when $\deg_t f = 1$ and $\deg_x g = 1$. Therefore, this variant does not overpass the classical one in terms of size property.

As for the root properties, note that the expression in Equation 25 which is a polynomial in $\mathbb{F}_q[t]$ is the norm of $a - bx$ with respect to the linear polynomial g , so its alpha value is constant and equal to one of the linear side of the classical variant. On the other hand, the root property of the polynomial in $\mathbb{F}_q[x]$ in Equation 25 can not be directly measured with our definition of alpha. Still, the number of polynomials f among which are selected to have a good root property is small, so we can not expect a large deviation; e.g. there are 2^6 values of f such that $\deg_x f = 6$.

5.4 Records on pairing-friendly curves

The fields $\mathbb{F}_{3^{6n}}$ are of particular interest in cryptography as one can break the cryptosystems which use pairing-friendly curves over \mathbb{F}_{3^n} by solving the discrete logarithm problem in $\mathbb{F}_{3^{6n}}$. The recently proposed algorithm of Joux [Jou13] proved to be very fast for the fields of composite degree. It rendered FFS obsolete in this case and drastically reduced the security of these curves. This section is also interesting for illustrating the behaviour separable polynomials.

These fields allow us to run the FFS with a base field \mathbb{F}_{3^d} with $d = 2, 3$ or 6 . Hence we collect coprime pairs (a, b) of polynomials in $\mathbb{F}_{3^d}[t]$ such that both $F(a, b)$ and $G(a, b)$ factor into small degree polynomials of $\mathbb{F}_{3^d}[t]$. The Galois variant [JL06] consists in choosing the polynomials f and g to have their coefficients in $\mathbb{F}_3[t]$ rather than $\mathbb{F}_{3^d}[t]$. Its main advantage is that, due to Galois properties, the factor base is reduced by a factor of d .

Hayashi *et al.* [HSW⁺10] used the base field \mathbb{F}_{3^6} and the polynomial $f_i = x^6 + t$ to break curves over $\mathbb{F}_{3^{71}}$. Two years later, Hayashi *et al.* [HSST12] used the base field \mathbb{F}_{3^3} and again $f_i = x^6 + t$ to break cryptosystems over $\mathbb{F}_{3^{97}}$.

Since the polynomial f_i is inseparable, as explained in 4.2, one quarter of the relations collected by the FFS are free. This roughly translates into a 4/3-fold speedup with respect to the separable polynomials having the same sieve efficiency. Note that f_i has the best epsilon among the 486 inseparable polynomials f in $\mathbb{F}_3[t][x]$ with $\deg_t f \leq 1$ and $\deg_x f = 6$.

A better choice can be only a separable polynomial with a better efficiency. Since the efficiency of a polynomial depends on the base field of the factor base, we distinguish the case of \mathbb{F}_3 from the case of \mathbb{F}_{3^d} with $d = 2, 3$ or 6 .

Experiment 3. Since $\deg_t(f_i) = 1$ we can use any of the $8 \cdot 3^{12}$ polynomials f in $\mathbb{F}_3[t][x]$ such that $\deg_t f \leq 1$, without changing the size property. The best alpha with respect to \mathbb{F}_3 corresponded to $f_s = tx^6 - tx^4 + (-t + 1)x^3 + (t - 1)x + t$.

Since alpha has a small variance on the polynomials tested in Experiment 3, we also consider the separable polynomial $f_{s'} = x^6 - x^2 + (t^8 + t^6 - t^4 + t^2 + 1)$, which is well suited when the skewness parameter is set to 1. Table 6 compares f_s and $f_{s'}$ to f_i for some randomly chosen linear polynomials.

In the case when the base field is \mathbb{F}_{3^6} and \mathbb{F}_{3^3} the evaluation of alpha is slower, with a factor of 200 compared to the case of \mathbb{F}_{3^6} . Note first that the polynomial f_s , whose root property over \mathbb{F}_3 is

f	$\bar{\alpha}(f, \mathbb{F}_3)$	$\bar{\alpha}(f, \mathbb{F}_{3^2})$	$\bar{\alpha}(f, \mathbb{F}_{3^3})$	$\bar{\alpha}(f, \mathbb{F}_{3^6})$
f_i	2.11	0.35	0.53	0.03
f_s	0.46	0.16	0.21	0.08

Table 7: The values of alpha with respect to different base field. The notation $\bar{\alpha}(f, \mathbb{F}_q)$ denotes $\log(q)/\log(2) \cdot \alpha(f)$ with alpha corresponding to the field \mathbb{F}_q .

better than that of f_i , has a poorer value of alpha when the base field is \mathbb{F}_{3^6} . In Table 7 we use $\frac{\log q}{\log 2} \times \alpha$ as an alternative of alpha which allows us to compare polynomials f with coefficients in different rings $\mathbb{F}_q[t]$. The values of alpha are approximated by considering only the contribution of at most 1000 irreducible polynomials in \mathbb{F}_{3^d} with $d = 1, 2, 3$ or 6 . Note that the values of $\frac{\log q}{\log 2} \times \alpha$ are close to each other when $q = 3^6$. This opens the question of how does the distribution of alpha evolve when we compute it with respect to a factor base in $\mathbb{F}_{3^6}[t]$ but for which the polynomials f are in $\mathbb{F}_3[t][x]$.

6 Conclusions and open questions

Improving on Joux and Lercier’s method of polynomial selection [JL02], we noted that a unique polynomial f can be used to solve the discrete logarithm problem on a range of inputs. Since the selection of f can be seen as a precomputation, we developed a series of functions which compare arbitrary polynomials and which are much faster than directly testing the sieve efficiency. In particular we obtained a sieving procedure for computing alpha, the function which measures the root property and we defined a function for measuring the cancellation property.

The case of inseparable polynomials was of particular interest as it has no equivalent notion in the NFS world. We showed that inseparable polynomials have the advantage of a large number of free relations, but most of the inseparable polynomials have a bad root property. The last section applied the new functions to some records in the literature.

The paper also opened some questions. First, thanks to the polynomial selection proposed in [JL03] this discussion could be adapted to the prime field computations by NFS. Secondly, the proof of the convergence of alpha seems to indicate that the distribution of alpha is influenced by the genus of the function fields. Finally, in the case of the Galois variant, it is interesting to know how does the variance of alpha evolve when restricted to Galois polynomials. If the variance is small enough, the sensible choice for the Galois variant seems to be the inseparable polynomials.

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A Appendix

Lemma A.1. *Under the notations of Proposition 2.8, Equation 7 holds.*

Proof. First, the equality below holds as we can change the summation order of an absolutely convergent series.

$$\sum_{k=1}^{\infty} \mathbb{P} \left(v_{\ell}(F(a, b)) \geq k \right) = \sum_{(r, k) \in S(f, \ell)} \mathbb{P} \left((a : b) \equiv r \pmod{\ell^k} \right).$$

Secondly, since for all k , $\mathbb{P}(v_{\ell}(F(a, b)) = k) = \mathbb{P}(\ell^k \mid F(a, b)) - \mathbb{P}(\ell^{k+1} \mid F(a, b))$ and $k \mathbb{P}(\ell^k \mid F(a, b)) \rightarrow 0$, we have

$$\sum_{k=1}^{\infty} k \mathbb{P} \left(v_{\ell}(F(a, b)) = k \right) = \sum_{k=1}^{\infty} \mathbb{P} \left(v_{\ell}(F(a, b)) \geq k \right).$$

Finally, let us prove that $a_{\text{hom}}(f)$ exists and equals $\sum_{k=1}^{\infty} k \mathbb{P}(v_{\ell}(F(a, b)) = k)$. For each N and k put

$$a_{\text{hom}}(f; N, k) = \frac{\sum \{ \min(v_{\ell}(F(a, b)), k) \mid \gcd(a, b) \not\equiv 0 \pmod{\ell}, \deg a, \deg b \leq N \}}{\#\{ (a, b) \mid \gcd(a, b) \not\equiv 0 \pmod{\ell}, \deg a, \deg b \leq N \}}.$$

Call $a_{\text{hom}}(f; N)$ the expression above when $\min(v_{\ell}(F(a, b)), k)$ is replaced with $v_{\ell}(F(a, b))$. On the one hand, for any $k_0 \in \mathbb{N}$, we have $a_{\text{hom}}(f; N, k_0) \leq a_{\text{hom}}(f; N)$, so

$$\sum_{k=1}^{k_0} k \mathbb{P}(v_{\ell}(F(a, b)) = k) \leq \liminf_{N \rightarrow \infty} a_{\text{hom}}(f; N).$$

On the other hand, let k_1 be large enough such that all the affine and projective roots modulo ℓ^{k_1} are simple and put $N = k_1 \deg \ell$. Since $N = k_1 \deg \ell$, the proportion of pairs (a, b) of degree at most N such that $(a : b) \equiv r \pmod{\ell^k}$ for some root (r, k) equals the probability $\mathbb{P}((a : b) \equiv r \pmod{\ell^k})$. Hence $a_{\text{hom}}(f; N, N/\deg \ell) = \sum_{k=1}^{k_1} k \mathbb{P}(v_{\ell}(F(a, b)) = k)$. Now, since the roots modulo ℓ^{k_1} are simple there are at most $\deg f$ of them. Also, for a and b of degree bounded by N , the norm

$F(a, b)$ has degree bounded by $N \deg f + |f|$, with $|f|$ the maximal degree of the coefficients of f . Hence

$$|a_{\text{hom}}(f; N) - a_{\text{hom}}(f; N, N/\deg \ell)| \leq \frac{\deg f(N \deg f + |f|)}{q^{N/\deg \ell}}.$$

This further implies

$$\limsup_{N \rightarrow \infty} a_{\text{hom}}(f; N) \leq \sum_{k \in \mathbb{N}} k \mathbb{P}(v_\ell(F(a, b)) = k).$$

We conclude that $a_{\text{hom}}(f; N)$ converges to $\sum_{k \in \mathbb{N}} k \mathbb{P}(v_\ell(F(a, b)) = k)$. □