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A Note on Set Systems with no Union of Cardinality 0 Modulo m

Vince Grolmusz

Department of Computer Science, Eötvös University, H-1117 Budapest, HUNGARY.
e-mail: grolmusz@cs.elte.hu

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Alon, Kleitman, Lipton, Meshulam, Rabin and Spencer (Graphs. Combin. 7 (1991), no. 2, 97-99) proved that for any hypergraph $\mathcal{F} = \{F_1, F_2, \dots, F_{d(q-1)+1}\}$, where q is a prime-power, and d denotes the maximum degree of the hypergraph, there exists an $\mathcal{F}_0 \subset \mathcal{F}$, such that $|\bigcup_{F \in \mathcal{F}_0} F| \equiv 0 \pmod{q}$. The main tool of the proof was a one-to-one correspondence between hypergraphs and polynomials. We give a direct, alternative proof to this correspondence, and also review its implications for set-systems following from the result of Barrington, Beigel and Rudich (Comput. Complexity, 4 (1994), 367-382) for certain mod 6 polynomials.

Keywords: Set systems, composite modulus, polynomials over rings

1 Introduction

Alon, Kleitman, Lipton, Meshulam, Rabin and Spencer [1] gave the following definition:

Definition 1 ([1]) For integers $d, m \geq 1$, let $f_d(m)$ denote the smallest t such that for any hypergraph $\mathcal{F} = \{F_1, F_2, \dots, F_t\}$ with maximum degree d there exists a non-empty $\mathcal{F}_0 \subset \mathcal{F}$, such that $|\bigcup_{F \in \mathcal{F}_0} F| \equiv 0 \pmod{m}$

Baker and Schmidt [2] defined the following quantity:

Definition 2 For integers $d, m \geq 1$, let $g_d(m)$ denote the smallest t such that for any polynomial $h \in \mathbb{Z}[x_1, x_2, \dots, x_t]$ of degree d , satisfying $h(\mathbf{0}) = 0$, there exists an $\mathbf{0} \neq \varepsilon \in \{0, 1\}^n$, such that $h(\varepsilon) \equiv 0 \pmod{m}$.

The following theorem was proven in [1]:

Theorem 3 ([1])

$$f_d(m) = g_d(m)$$

In the next section we give a natural one-to-one correspondence between polynomials and hypergraphs, proving Theorem 3.

For p prime, and α positive integer it is known ([1], [2], [4]) that $g_d(p^\alpha) = d(p^\alpha - 1) + 1$, so we obtain

Corollary 4 ([1]) For $\mathcal{F} = \{F_1, F_2, \dots, F_{d(q-1)+1}\}$, where q is a prime-power, and d denotes the maximum degree of the hypergraph, there exists an $\emptyset \neq \mathcal{F}_0 \subset \mathcal{F}$, such that $|\bigcup_{F \in \mathcal{F}_0} F| \equiv 0 \pmod{q}$.

This corollary is a generalization of the undergraduate exercise that from arbitrary m integers, one can choose a non-empty subset, which adds up to 0 modulo m (the $d = 1$ case).

In 1991, *Barrington, Beigel and Rudich* [3] gave an explicit construction for polynomials modulo $m = p_1^{\alpha_1} p_2^{\alpha_2} \dots p_r^{\alpha_r}$, showing that

$$g_d(m) = \Omega(d^r).$$

Since the proof of Theorem 3 (both the original and ours in the next section) gives explicit constructions for hypergraphs from polynomials, the following corollary holds:

Corollary 5 Let $m = p_1^{\alpha_1} p_2^{\alpha_2} \dots p_r^{\alpha_r}$. Then there exists an explicitly constructible hypergraph \mathcal{F} of maximum degree d , such that $|\mathcal{F}| = \Omega(d^r)$ and for each $\emptyset \neq \mathcal{F}_0 \subset \mathcal{F}$ it is satisfied that $|\bigcup_{F \in \mathcal{F}_0} F| \not\equiv 0 \pmod{m}$.

The authors of [1] gave a doubly-exponential upper bound on $f_d(m)$, which was based on a Ramsey-theoretic bound of [2]. More recently, *Tardos and Barrington* [4] showed that

$$f_d(m) = \exp(O(d^{r-1})).$$

2 Correspondence between polynomials and hypergraphs

We give here a short and direct proof for Theorem 3. Let \mathcal{Q} denote the set of rationals. It is well known that the set of functions $\{f : \{0, 1\}^t \rightarrow \mathcal{Q}\}$ forms a 2^t -dimensional vector space over the rationals. One useful basis of this vector space is the set of OR-functions $\{\bigvee_{i \in I} x_i : I \subset \{1, 2, \dots, t\}\}$, where

$$\bigvee_{i \in I} x_i = 1 - \prod_{i \in I} (1 - x_i).$$

It is easy to see that any integer-valued function on the hypercube can be written as the integer-coefficient linear combination of these OR-functions. Moreover, if the function is a degree- d polynomial, then it is enough to use OR functions with $|I| \leq d$. If we consider modulo m polynomials, then the coefficients can be restricted to the set $\{0, 1, 2, \dots, m-1\}$. It will be convenient to view modulo m polynomials as the sum of several OR functions with coefficient 1; instead of multiplying an OR function with a coefficient a we will add it up exactly a times.

Consequently, our degree- d modulo m polynomial has the following form:

$$h = S_1 + S_2 + \dots + S_\ell, \tag{1}$$

where S_i is an OR-function of degree at most d .

Now we are ready to define the one-to-one correspondence between degree- d modulo m polynomials without non-trivial zeroes on the hypercube and hypergraphs, without non-empty subhypergraphs of modulo- m union-size 0. Let h be a degree- d polynomial written in form (1), and define hypergraph $\mathcal{F} = \{F_1, F_2, \dots, F_\ell\}$, where $F_i = \{S_j : x_i \text{ appears as a variable in } S_j\}$. Clearly, the degree of this hypergraph is at most the degree of h that is, d .

On the other hand, for a hypergraph $\mathcal{F} = \{F_1, F_2, \dots, F_t\}$ on the ground-set $\{v_1, v_2, \dots, v_\ell\}$, let us define $h(x_1, x_2, \dots, x_t) = S_1 + S_2 + \dots + S_t$, where

$$S_j = \bigvee_{i: v_j \in F_i} x_i.$$

Obviously, the degree of h is at most the degree of \mathcal{F} .

Now we state that \mathcal{F} has a non-empty subhypergraph with union-size 0 modulo m if and only if there exists a $\mathbf{0} \neq \mathbf{x} : h(\mathbf{x}) \equiv 0 \pmod{m}$. The proof is as follows: For $\mathbf{x} = (x_1, x_2, \dots, x_n) \in \{0, 1\}^n$ let us denote $I(\mathbf{x}) = \{i : x_i = 1\}$. Then $S_j(\mathbf{x}) = 1$ if $S_j \in \bigcup_{i \in I(\mathbf{x})} F_i$, and $S_j(\mathbf{x}) = 0$ otherwise. Thus $h(\mathbf{x}) = |\bigcup_{i \in I(\mathbf{x})} F_i|$ holds for all $\mathbf{x} \in \{0, 1\}^n$. In particular, evaluations of h and union-sizes of subhypergraphs in \mathcal{F} become divisible by m simultaneously. \square

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