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Do the Fix Ingredients Already Exist? 
An Empirical Inquiry into the Redundancy Assumptions of Program Repair Approaches 

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ABSTRACT
Much initial research on automatic program repair has focused on experimental results to probe their potential to find patches and reduce development effort. Relatively less effort has been put into understanding the hows and whys of such approaches. For example, a critical assumption of the GenProg technique is that certain bugs can be fixed by copying and re-arranging existing code. In other words, GenProg assumes that the fix ingredients already exist elsewhere in the code. In this paper, we formalize these assumptions around the concept of “temporal redundancy”. A temporally redundant commit is only composed of what has already existed in previous commits. Our experiments show that a large proportion of commits that add existing code are temporally redundant. This validates the fundamental redundancy assumption of GenProg.

Categories and Subject Descriptors
D.2.5 [Software Engineering]: Testing and Debugging

General Terms
Experimentation

Keywords
Automatic program fixing; automatic software repair; mining software repositories

1. INTRODUCTION
To some extent, each program repair technique is based on an underlying assumption. GenProg’s “secret sauce” [12, 7] is the assumption that large programs contain the seeds of their own repair and thus that rearrangements of existing statements can fix most bugs. This redundancy assumption is also behind four out of PAR’s ten repair templates [5]. By contrast, SemFix [11] makes different assumptions: it is based on the idea that some bugs can be repaired by changing only one variable assignment or if conditional expression. Although extensive empirical evaluations of the bug fixing powers and costs of these techniques have been carried out, less work has focused on exploring or validating such fundamental assumptions.

In this paper we focus on the assumptions made by techniques like PAR and GenProg, and ask whether it makes sense to re-arrange existing code or code changes to fix bugs. We formalize our investigation by defining a concept of “software temporal redundancy”. A commit is said to be temporally redundant if it is only a re-arrangement of the code that has already existed in previous commits. We provide further insight into the problem by breaking down existing code along multiple dimensions, including the granularity at which it is considered (e.g., line-level vs. token-level) and the locality at which it is found (e.g., in the same file vs. anywhere in the program).

We measure temporal redundancy on a dataset of six open-source Java applications. For instance, at token level, we find that 31-52% of commits are temporally redundant. Our results show that a substantial fraction of software evolution takes place in a “closed world”, with no creation of new material. To our knowledge, this paper is the first empirical quantification linking such evidence to automated program repair.

Our experiments enable us to understand the foundational assumptions of program repair approaches based on redundancy such as GenProg or PAR. First, we validate the redundancy assumption: it makes sense to fix defects by rearranging existing statements. Second, we find that different assumptions about how patches should be formed are significant: depending on the way one picks code from elsewhere in the program, the search space is different in terms of size and potential for success. Those results are directly actionable for improving such automated program repair techniques. For example, our results indicate that only considering local redundancy at the line level would decrease the repair time without a similar reduction in success potential.

To sum up, our contributions are:
• the definition and formalization of software temporal redundancy;
• an experimental protocol to measure software temporal redundancy using versioning data;
• an empirical validation of the redundancy assumption of some program repair approaches on 16071 commits of open-source software applications;

2. SOFTWARE TEMPORAL REDUNDANCY

Research in software evolution uses time as a primary focus. For most software projects, software evolution artifacts are captured by version control systems (e.g., CVS, SVN or GIT). In this section, we present a new commit-level software evolution metric called “temporal redundancy”.

Commits.

A commit in a version control system consists of new file versions. Conceptually, a commit can be viewed in two ways: as a set of file pairs (before and after the commit), or as a set of changes (applied to the version before the commit to obtain the version after it).

In this paper, we use this change-based view of commits. We consider that a commit adds and/or deletes new source code fragments. An update is considered as the combination of a deletion and an addition. We do not consider commits done on other artifacts than source code.

Fragment Redundancy.

We use fragment to denote a substring of source code. For instance, the source code line “for (int i=0;i<n;i++)” is a fragment. Fragments are always defined according to a level of granularity. In this paper we consider two different levels: lines (as separated by line breaks) and tokens (as separated by lexing rules).

A fragment $F$ is snapshot redundant at time $T$ if another instance of that same fragment $F$ exists elsewhere in the program at time $T$. This is the redundancy studied by Gabel and Su [1] and used by GenProg [12]. In this paper we consider a richer notion of redundancy that includes historical context.

A fragment $F$ is temporally redundant at time $T$ if that same fragment $F$ has already been seen during the history of the software under analysis (i.e., at time $T' < T$). For instance, literal “42” might be added in version $\#1$, be removed in version $\#2$ and reused again in version $\#3$. In the context of version $\#3$, the “42” fragment is temporally redundant. Thus, in this paper, once a fragment has appeared it is always subsequently viewed as a potential source of redundancy. We consider the first version of a program to be created by insertions from an empty initial program.

Commit Redundancy.

A source code commit is composed of added and/or deleted fragments. Using a cooking metaphor, the added fragments are the “ingredients” of the commit.

We define a temporally redundant commit as a commit for which all added fragments are (individually) temporally redundant. More formally, we define a commit $C_i$ performed at the time $T_i$ as a set $S_i$ of added fragments and a set $R_i$ of removed fragments. Let $C$ be the set of all commits for a program. Then a temporally redundant commit $C_i$ satisfies

$$\forall f \in S_i \exists C_i \in C \mid T_i < T_j \land f \in S_j$$

For such commits, no new fragments are invented and no fresh material is introduced; the commit is only a rearrangement of insertions that have already been seen in previous commits.

Scope of Temporal Redundancy.

A fragment is temporally redundant if that same fragment appeared in a previous commit. This kind of redundancy has a global scope: the location of the previous fragment instance does not matter.

We now define a more restricted local scope notion of temporal redundancy. A fragment is locally temporally redundant if that same fragment has been used in a previous commit to the same file.

For example, consider two files, $F_1$ and $F_2$, each containing 3 fragments: $F_1 = \{a, b, c\}$, $F_2 = \{d, e, f\}$. Suppose commit $C_1$ adds fragment $c$ to file $F_2$. For that commit, fragment $c$ is global temporally redundant (already available in $F_1$), but not locally temporally redundant (never previously available in $F_2$). Suppose commit $C_2$ introduces another version of $F_2$ replacing fragment $e$ with $d$. In that commit, fragment $d$ is locally redundant (since $d$ was previously available in $F_2$).

3. MEASURING REDUNDANCY

We now present an experimental design to measure the fraction of relevant commits to a program that are temporally redundant.

Experimental Protocol.

Given a level of granularity and a scoping level, our experimental protocol to measure the temporal redundancy of the evolution of a program consists of the following phases:

a) Retrieving commits. All commits of the program under analysis are collected from the repository.

b) Filtering commit files. Only commits to executable code are retained; commits to test cases are discarded.

c) Fragmenting commits. We split each relevant file into fragments at a given level of granularity (e.g., lines or tokens). This results in a before-commit and an after-commit sequence of fragments. We use the Myers differencing algorithm [9] to compare both fragment sequences and obtain the added fragments of the commit.

d) Filtering fragments. We filter out whitespace and comments. In this paper we are only interested in the evolution of executable code and not in indentation or documentation.

e) Selecting acceptable commits. We select those commits that introduce at least one fragment after filtering. We call such commits acceptable.

f) Indexing fragments. We consider each added fragment in each acceptable commit in ascending temporal order. If a fragment has not been encountered previously at the given scoping level (i.e., global or local), we index it with the date of its first introduction.

g) Measuring temporal redundancy. The temporal redundancy of the entire program’s evolution is the fraction of acceptable commits that are temporally redundant (see Section 2). For example, a temporal redundancy of 0.1 means that 10% of acceptable commits are temporally redundant commits.

Dataset.

We use six open-source Java projects to measure the temporal redundancy. They are: Apache Log4j, JUnit, Picocontainer, Apache Commons Collections, Apache Commons...
Math, and Apache Commons Lang. The inclusion criterion is as follows: the Apache projects were used in previous research on automatic program repair [5], while the remaining two are Java projects mentioned in previous research on software evolution [3]. After applying the filters presented above on the 16071 commits of the dataset, we obtain 7076 acceptable commits.

4. EMPIRICAL RESULTS

We now present our empirical results on the temporal redundancy of software (as defined in Section 2) following the experimental design presented in Section 3.

4.1 Line-Level Temporal Redundancy

Research question: What is the amount of line-based temporal redundancy?

For each application of our dataset, we measured the total number of acceptable commits within the analysis timespan (the complete data is available at http://goo.gl/k0rZWc) and the global-scope line-level temporal redundancy. Columns #2 and #3 of Table 1 report the results. For instance, for log4j (the first row), there are 1681 commits which add at least one executable line. Only 9% of those 1681 are temporally redundant commits.

Overall, at the level of lines, 3–17% of the accepted commits are temporally redundant commits. Their basic ingredients are only previously-inserted code.

This has additional implications for automatic repair and code synthesis: for synthesizing those commits, the search-space has a finite number of atomic building blocks (previously-observed line-level fragments). Theoretically, a redundancy-based approach should be able to synthesize all the temporally-redundant commits.

The interval 3–17% is large and the reasons behind this variation are not obvious. One reason could be that the commit conventions used by the developers of a project are different. For instance, some projects prefer to have small and atomic commits (one bug fix or feature per commit). Other projects are less restrictive on this point. This has a direct impact on the redundancy: small and atomic commits are more likely to be redundant. A deeper analysis of this question is future work.

4.2 Token-Level Temporal Redundancy

Research question: Is there a difference between line-based and token-based temporal redundancy?

Before answering this question, we note that, analytically, all temporally redundant commits at the level of lines are necessarily temporally redundant commits at the level of tokens. Furthermore, a unique new line might be exclusively composed of existing tokens. Consequently, the token-based temporal redundancy must be equal to or greater than line-based temporal redundancy. We now measure the temporal redundancy at the line and token level.

Table 1 reports global scope line-level (column #3) and token-level temporal redundancy (column #7). For instance, in log4j, there is a line-level temporal redundancy of 9% but a token-level temporal redundancy of 39%. Overall, at the token level, 29–52% of commits are temporally redundant.

For all projects, token-based temporal redundancy exceeds line-based. This follows from the analytical argument above and gives confidence in the experiment’s construct validity.

Overall, token-level temporal redundancy (between 29% and 52%) is much higher than the line-level temporal redundancy.

This is very good news. For automated repair and code synthesis, a high temporal redundancy implies a smaller search space. This holds for both the line and token level of granularity. Our token-level temporal redundancy measurements imply that for between 29% and 52% of accepted commits, synthesis and repair need never invent a new token. For instance, repair or synthesis of arithmetic code need only consider recombining existing literals and operations for one-third to one-half of commits.

We thus propose that the repair search space is composed of two components: the search space of fragments (the atomic building blocks) and the search space of their combinations. Our experiment enables us to precisely measure the former: we can count the total number of fragments seen up to a given point in time. We call this a fragment pool. There is one fragment pool per level of granularity. This is conceptually similar to the pool of four DNA bases or the pool of amino acids in biology.

Table 1 gives the size of the global scope fragment pool for line-level (column #4) and token-level (column #8) analyses. For instance, in the considered slice of history of log4j, there are 43313 different lines and 14294 different tokens that are involved in the software evolution. For all applications under study, the token pool at the point in time of the last commit is much smaller than the line pool.

For automated repair or code synthesis, there is a tension between working with the line pool or the token pool. To some extent, the temporally redundant commits correspond to the number of commits that can be synthesized. With the line pool, the combination of lines is much smaller (the combination space is smaller) but fewer commits can be synthesized (~10%). With the token pool, more commits can

<table>
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<tr>
<th>Program</th>
<th>Acceptable Commits</th>
<th>Line granularity</th>
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<th>Token granularity</th>
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Table 1: The temporal redundancy of six open-source applications.
be synthesized (~40%), but at the price of exploring a much bigger combination space.

4.3 Redundancy Scope Experiment

Research question: Do file scope restrictions impact the amount of temporal redundancy?

We now measure the temporal redundancy available at the local scope in the same file (as defined in Section 2).

Table 1 reports global and local scope temporal redundancy in our dataset. For each granularity, there is one column “Global” and one column “Local” corresponding to the different scope. For instance, at the line level, column #3 is the temporal redundancy at the global scope and column #5 gives it when considering a local scope.

As discussed in Section 4.1, at the line granularity, there are between 3% and 17% of temporally redundant commits at the global scope. At the local scope, there are between 2% and 16%.

The temporal redundancy of both scopes is of the same order of magnitude. In all projects, more than half of the temporally redundant commits actually have local temporal redundancy. Consequently, at line granularity, most of the temporal redundancy is localized in the same file.

At token-level granularity, the results are similar: we find a large amount of token-level local-scope temporal redundancy. Tokens are likely to be reused in commits impacting the same file. This further indicates that the fragment locality matters during software evolution. We note that the difference of redundancy between global and local scopes is slightly higher at the token level (col. #7 vs #9) than at the line level (col. #3 vs #5).

This is again a promising result with respect to the search space of automatic repair or synthesis. First, at the line level, the local scope pool is able to seed the same order of magnitude of commits as the global one. In other words, it is almost as fertile as the global pool. Second, when one considers the local scope pool, the search space is much smaller. For instance, for log4j, the median local pool size at the line level is 57 lines, compared to 43313 at the global scope level. Restricting attention to the local scope reduces the search space greatly while still enabling the synthesis of a large number of commits. Those results are directly actionable for improving GenProg and other redundancy-based approaches: our results indicate that only considering local redundancy would decrease the repair time while keeping a high repair success potential. This is in line with our body of previous experiments.

5. RELATED WORK

Some work on code clone detection also focuses on software redundancy, studying the differences between line- and token-level granularity [4]. Kim et al. [6] considered code clones via a temporal perspective, linking clones together across versions. By contrast, we study whether changes only contain existing code. Most importantly, a change composed of many redundant fragments may not be a clone in itself: each redundant fragment could come from different, unrelated, locations or appear in a rearranged order.

In previous work [8], we discussed the topology of the repair space and efficient manners to navigate it. In this paper, we focus on the relation between commits and the repair space, which has not been previously addressed.

Gabel and Su [1] studied the uniqueness of source code. Their “syntactic redundancy” calculates the degree to which portions of software applications are redundant. Hindle and colleagues [2] studied the repetitiveness and predictability of code. Both consider software redundancy from a spatial viewpoint. By contrast, we study temporal redundancy. Our measures are not present in, and cannot be inferred from, their experiments.

Nguyen et al. [10] presented a study of repetitiveness of code changes in software evolution. Where they measure the repetitiveness of code changes by abstracting over commits and literals, we measure the redundancy of code changes at the commit level considering real, non-abstracted code (e.g., we care about the actual values of literals).

6. CONCLUSION

In this paper, we define a notion of software temporal redundancy and empirically measure its presence in source code. For example, as many as 52% of commits are composed entirely of previously-existing tokens. Our results have direct application to automated program repair. For example, we find that searches for line-level repairs can focus on a single file, dramatically reducing the search space without substantially reducing repair success potential.

Our future work will focus on measuring temporal redundancy on bug fix commits and on studying different granularities beyond lines and tokens (e.g., at the AST level, focusing on some kind of source code elements).

7. ACKNOWLEDGMENTS

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8. REFERENCES


