



Tag Second-preimage Attack against π -cipher

Gaëtan Leurent

► **To cite this version:**

| Gaëtan Leurent. Tag Second-preimage Attack against π -cipher. 2014. <hal-00966794v2>

HAL Id: hal-00966794

<https://hal.inria.fr/hal-00966794v2>

Submitted on 27 Mar 2014

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Tag Second-preimage Attack against π -cipher

Gaëtan Leurent

Inria, France*

Abstract. The π -cipher is one of the candidates of the CAESAR competition. One of the advertised features of the π -cipher is tag second-preimage resistance: it should be hard to generate a message with a given tag, even for the legitimate key holder (insider attack).

In this note, we show that the generalized birthday attack of Wagner gives a practical tag second-preimage attack against the π -cipher.

1 Introduction

The π -cipher [2] is an authenticated encryption algorithm submitted to the CAESAR competition. One of the extra features advertised by the designers is tag second-preimage resistance: it should be hard to produce second-preimages of a given tag, even for an adversary who knows the secret key (most authenticated encryption algorithm do not have this feature, and an insider can easily generate tag second-preimages).

As written in [2, 4.1], the tag generation of an m -block message with the π -cipher can be written as:

$$T = T'' \boxplus_8 e(1, M_1) \boxplus_8 e(2, M_2) \boxplus_8 \cdots \boxplus_8 e(m, M_m)$$

where e denotes a keyed function known to the key holder (the e-triplex), \boxplus_8 is a component-wise addition of vectors of 8 elements in \mathbb{Z}_{2^ω} , and T'' is the associated data tag (known to the insider). The word-size ω is 16, 32, or 64, depending on the security level. In a tag second-preimage attack, an insider wants to build a message M reaching a fixed tag \bar{T} . Without loss of generality, we assume $T'' = 0$ and $\bar{T} = 0$.

In the submission document of π -cipher, the tag second-preimage problem is seen as a knapsack problem, and the main attack considered is a variant of an attack by Camion and Patarin [1]. However, the generalization of this attack due to Wagner [3] can break the problem more efficiently.

2 Wagner's Generalized Birthday Attack

The generalized birthday attack of Wagner is an attack against the m -sum problem: given m lists L_1, L_2, \dots, L_m of n -bit words, one find values $l_1 \in L_1, \dots, l_m \in L_m$ such that $\bigoplus_{i=1}^m l_i = 0$. If each list contains at least $2^{n/m}$ elements there is a good probability that a solution exists, but the best known algorithm is a simple birthday attack in time and memory $\tilde{O}(2^{n/2})$. One would first build two lists L_A and L_B with all the sums of elements in $L_1, \dots, L_{m/2}$ and $L_{m/2+1}, \dots, L_m$ respectively, then sort L_A and L_B , and look for a match between the two lists (L_A and L_B contain $2^{n/2}$ elements each).

* SECRET project-team, Gaetan.Leurent@inria.fr

Wagner's algorithm has a lower complexity, but it requires more elements in the lists. For instance, with $m = 4$, it uses lists of size $2^{n/3}$ in order to find one solution using $\tilde{O}(2^{n/3})$ time and memory. The basic operation of the algorithm is the general join $\bowtie_{\tau}: L \bowtie_{\tau} L'$ consists of all the elements of $L \times L'$ that agree on their τ least significant bits. More precisely, the operation can be defined over list of values with associated data:

$$L \bowtie_{\tau} L' = \{(l \oplus l', (a, a')) \mid (l, a) \in L, (l', a') \in L', \text{low}_{\tau}(l \oplus l') = 0\}.$$

The join operation is computed efficiently by sorting the lists L and L' according to the lower τ bits, and stepping through the lists simultaneously in order to find values that agree on their low bits. Moreover, the sorting can be done in linear time using a hash table, or a radix sort.

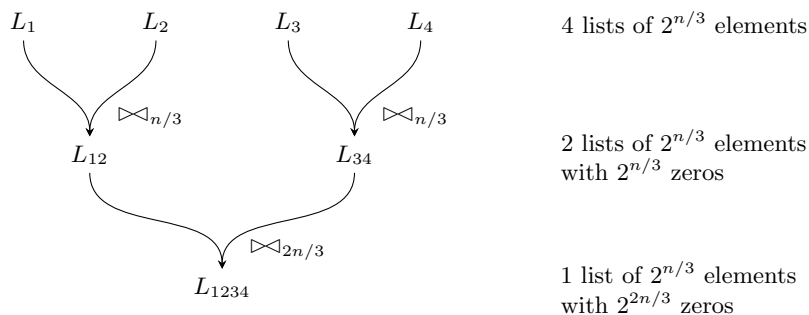


Fig. 1. Wagner's algorithm for $m = 4$

The generalized birthday algorithm for $m = 4$ is described by Figure 1. We first build the lists $L_{12} = L_1 \bowtie_{n/3} L_2$ and $L_{34} = L_3 \bowtie_{n/3} L_4$, containing about $2^{n/3}$ elements. Next, we build $L_{1234} = L_{12} \bowtie_{2n/3} L_{34}$. Since the elements of L_{12} and L_{34} already agree on their $n/3$ lower bits, we are only matching bits $n/3$ to $2n/3$, so we still expect to find $2^{n/3}$ elements. Finally, we expect one of the elements of L_{1234} to be zero. This can be generalized to any m that is a power of two, using a binary tree: if $m = 2^a$, we need m lists of $2^{n/(a+1)}$ elements and the time and memory used by the algorithm is $2^a \cdot r 2^{n/(a+1)}$. The algorithm for $m = 8$ is shown by Figure 2.

3 Application to the π -cipher

In order to apply this attack to the π -cipher, we need to solve the m -sum problem for the word-wise modular addition \boxplus_8 , instead of the exclusive-or \oplus . Wagner showed how to solve the generalized birthday problem with a modular addition, and his trick also works for the word-wise modular addition. More precisely, we have to modify the join operator to:

$$L \blacktriangleright_{\tau} L' = \{(l \boxplus_8 l', (a, a')) \mid (l, a) \in L, (l', a') \in L', \text{low}_{\tau}(l \boxplus_8 l') = 0\}.$$

Since the word-wise modular addition \boxplus_8 only has carries from the low order bits to the high order bits, when x and y have their τ low-order bits set to zero, $x \boxplus_8 y$

also has τ low-order bits set to zero. Moreover, the join \bowtie can still be computed efficiently. We first negate the list L and define $-L = \{(-l, a) \mid (l, a) \in L\}$, where $-l$ is the additive inverse with regard to the word-wise addition, *i.e.* $l \boxplus_8 (-l) = 0$. Then we sort $-L$ and L' according to their lower τ bits, and step through the lists in parallel. When an element of $-L$ and an element of L' agree on their low bit, the corresponding sum will have its low bits equal to zero. Therefore, this variant of Wagner's algorithm is suitable for a tag second-preimage attack on the π -cipher.

We give a full description of an attack with $\omega = 16$ in Algorithm 1; this attack uses 8 lists of size 2^{32} (illustrated by Figure 2), *i.e.* we consider an 8-block message, with 2^{32} possibilities for each block. This gives a complexity of 2^{35} . More generally, we can apply Wagner's attack to different versions of π -cipher (*i.e.* with different values of ω), and several trade-offs between the message length and the attack complexity are possible. We give some parameters in Table 1.

Table 1. Attack parameters

ω	Optimal parameters			Short messages		
	m	$ L $	Complexity	m	$ L $	Complexity
16	2^{11}	2^{11}	2^{22}	2^3	2^{32}	2^{35}
32	2^{16}	2^{15}	2^{31}	2^7	2^{32}	2^{39}
64	2^{22}	2^{23}	2^{45}	2^{15}	2^{32}	2^{47}

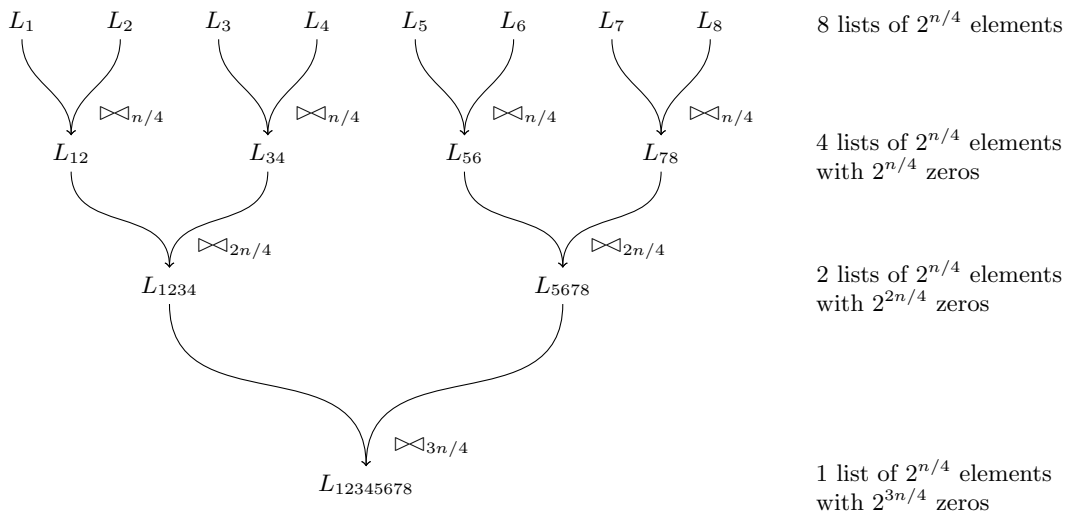


Fig. 2. Wagner's algorithm for $m = 8$

References

1. Camion, P., Patarin, J.: The knapsack hash function proposed at crypto'89 can be broken. In: Davies, D.W. (ed.) EUROCRYPT. Lecture Notes in Computer Science, vol. 547, pp. 39–53. Springer (1991)
2. Gligoroski, D., Mihajloska, H., Samardjiska, S., Jacobsen, H., El-Hadedy, M., Jensen, R.E.: π -Cipher. Submission to CAESAR. Available from: <http://competitions.cr.yu.to/round1/picipherv1.pdf> (v1) (March 2014)
3. Wagner, D.: A generalized birthday problem. In: Yung, M. (ed.) CRYPTO. Lecture Notes in Computer Science, vol. 2442, pp. 288–303. Springer (2002)

Algorithm 1 Short message attack with $\omega = 16$ and $m = 8$.

```
for  $0 \leq i < 8$  do
  for  $0 \leq j < 2^{32}$  do
     $L[i][j] \leftarrow (e(i, [j]), j)$ 
  end for
end for
 $L[8] \leftarrow \text{MERGE}(L[0], L[1], 32)$ 
 $L[9] \leftarrow \text{MERGE}(L[2], L[3], 32)$ 
 $L[10] \leftarrow \text{MERGE}(L[4], L[5], 32)$ 
 $L[11] \leftarrow \text{MERGE}(L[6], L[7], 32)$ 
 $L[12] \leftarrow \text{MERGE}(L[8], L[9], 64)$ 
 $L[13] \leftarrow \text{MERGE}(L[10], L[11], 64)$ 
 $L[14] \leftarrow \text{MERGE}(L[12], L[13], 96)$ 
for all  $(l, ((a_1, a_2), (a_3, a_4), ((a_5, a_6), (a_7, a_8)))) \in L[14]$  do
  if  $l = 0$  then
    return  $[a_1] \parallel [a_2] \parallel [a_3] \parallel [a_4] \parallel [a_5] \parallel [a_6] \parallel [a_7] \parallel [a_8]$ 
  end if
end for

function MERGE( $L, L', \tau$ )
  SORT( $L, -\text{low}_\tau$ )
  SORT( $L', \text{low}_\tau$ )
   $i \leftarrow 0$ 
   $j \leftarrow 0$ 
   $M \leftarrow \emptyset$ 
  while  $i < |L|$  and  $j < |L'|$  do
     $(l, a) \leftarrow L[i]$ 
     $(l', a') \leftarrow L'[j]$ 
    if  $\text{low}_\tau(-l) = \text{low}_\tau(l')$  then
       $M \leftarrow M \cup \{(l \boxplus_8 l', (a, a'))\}$ 
    else if  $\text{low}_\tau(-l) < \text{low}_\tau(l')$  then
       $i \leftarrow i + 1$ 
    else
       $j \leftarrow j + 1$ 
    end if
  end while
  return  $M$ 
end function
```
