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On-line Ramsey Numbers for Paths and Stars

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We study on-line version of size-Ramsey numbers of graphs defined via a game played between Builder and Painter: in one round Builder joins two vertices by an edge and Painter paints it red or blue. The goal of Builder is to force Painter to create a monochromatic copy of a fixed graph H in as few rounds as possible. The minimum number of rounds (assuming both players play perfectly) is the *on-line Ramsey number* $\tilde{r}(H)$ of the graph H . We determine exact values of $\tilde{r}(H)$ for a few short paths and obtain a general upper bound $\tilde{r}(P_n) \leq 4n - 7$. We also study asymmetric version of this parameter when one of the target graphs is a star S_n with n edges. We prove that $\tilde{r}(S_n, H) \leq n \cdot e(H)$ when H is any tree, cycle or clique.

Keywords: size Ramsey number, online Ramsey games

1 Introduction

In this paper we are concerned with on-line variant of size-Ramsey numbers of graphs. For given graphs G and H we write $G \rightarrow H$ if there is a monochromatic copy of H in any red-blue coloring of the edges of G . The *size-Ramsey number* of a graph H , denoted by $\hat{r}(H)$, is the smallest possible number of edges in a graph G satisfying $G \rightarrow H$. This notion was first studied by Erdős, Faudree, Rousseau, and Schelp (8), as a variation on the usual Ramsey number $r(H)$ (which is the least n such that $K_n \rightarrow H$, where K_n is a clique on n vertices). Clearly, $\hat{r}(H)$ is at most $\binom{r(H)}{2}$ for every graph H . Somewhat surprisingly this bound is attained if $H = K_n$ (8). However, $\hat{r}(P_n) \leq cn$ for some absolute constant c , where P_n is a path on n vertices. This was proved by Beck (2) in response to a one-hundred-dollar question of Erdős. The proof uses random graphs and is nonconstructive (cf. (5)). Explicit construction based on expanders was provided by Alon and Chung (1). Since then many deep results for other natural classes of graphs were obtained (cf. (3; 7; 15; 22)).

On-line version of size-Ramsey numbers was introduced independently by Beck (4) and Kurek and Ruciński (14). It is best explained by the following game between *Builder* and *Painter*, played on a large set of vertices. In one round Builder joins two nonadjacent vertices by an edge and Painter colors this edge red or blue. The goal of Builder is to create a monochromatic copy of a fixed graph H in as few rounds as possible. Painter will try to resist doing it for as long as possible. The *on-line Ramsey number*

$\tilde{r}(H)$ of a graph H is the minimum number of rounds in which Builder achieves his goal, assuming both players play perfectly. Clearly, $\tilde{r}(H) \leq \hat{r}(H)$ and it is natural to ask how much these two parameters may differ. We study this question in several typical situations. For instance we show that $\tilde{r}(P_n) \leq 4n - 7$, while currently best bound in the off-line case is $\hat{r}(P_n) \leq 594n$ (cf. (5)). In general the problem may be hard as it is not even known if the ratio $\tilde{r}(K_n)/\hat{r}(K_n)$ tends to zero as $n \rightarrow \infty$ (cf. (14)). Other open problems and related versions of the game can be found in (10; 12; 13).

One can also consider an asymmetric version of the on-line Ramsey number. We write $F \rightarrow (G, H)$ if in any red-blue coloring of the edges of F there is a red copy of G or a blue copy of H . Similarly, $\hat{r}(G, H)$ stands for the minimum number of edges in a graph F such that $F \rightarrow (G, H)$, and $\tilde{r}(G, H)$ is defined accordingly.

Similar to the classical Ramsey numbers (see a dynamic survey of Radziszowski (21) which includes all known nontrivial values and bounds for Ramsey numbers), it is hard to compute the exact value of $\tilde{r}(G)$ unless G are trivial. In this relatively new area of small on-line Ramsey numbers, very little is known.

2 Paths

As a warm-up we start with exact determination of the on-line Ramsey numbers for a few short paths. Let $e(H)$ denote the number of edges of a graph H . Note first that in general we have $\tilde{r}(H) \geq 2e(H) - 1$, since Painter may color safely the first $e(H) - 1$ edges red, and the next $e(H) - 1$ edges blue. Since $e(P_n) = n - 1$ we have $\tilde{r}(P_n) \geq 2n - 3$. This lower bound is attained for $n = 2, 3, 4, 5$.

Proposition 2.1 $\tilde{r}(P_n) = 2n - 3$, for $n = 2, 3, 4, 5$.

The cases $n = 2, 3$ are trivial. The case $n = 4$ is also very easy; after presenting three edges of a path P_4 , there are only two possible patterns (up to symmetry): bbr and brb . Then Builder creates a monochromatic path P_4 in the next two moves, as depicted in Figure 1. (The final edge is drawn in two colors.) To prove that $\tilde{r}(P_5) = 7$ we have to analyze more subcases. In the first four moves Builder constructs a path P_5 so that essentially one of the three possible color patterns appears, as displayed in Figure 2. Then he obtains a monochromatic P_5 in next three moves, as shown in Figure 3. (A circled number means that Painter had a choice in that move, which led to a branching into subcases.)



Fig. 1: Forcing a path P_4

The pattern breaks already for $n = 6$.

Proposition 2.2 $\tilde{r}(P_6) = 10$.

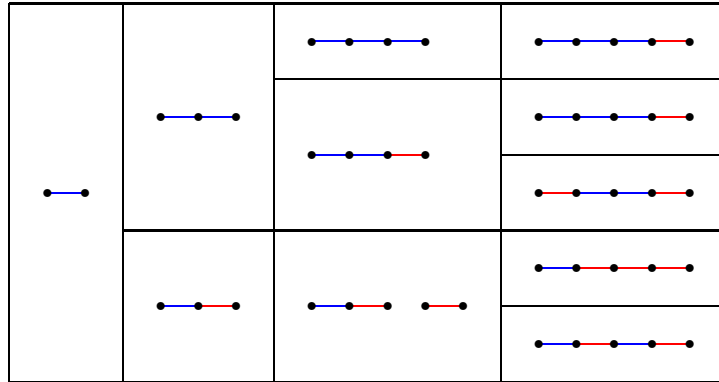


Fig. 2: Forcing a monochromatic P_5 (first four moves)

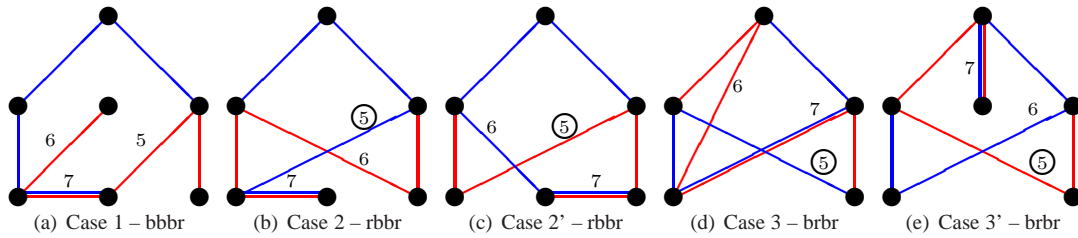


Fig. 3: Forcing a path P_5 (last three moves)

Proof: For the lower bound consider a natural strategy for Painter: color an edge blue if it does not create a blue copy of P_6 , otherwise use red. Suppose first that the graph constructed by Builder in the first four moves is not a path P_5 . Then clearly Painter will be able to use the color blue in the very next move and we are done. So, suppose that after first four rounds there is a blue path P_5 . The only way for Builder to finish the game in the total of nine rounds is to create in next four rounds a red P_5 . This is possible only by using both ends of the blue path P_5 , making the last winning move impossible.

The upper bound is achieved by case-by-case analysis shown in Figure 4 (first five moves) and Figure 5 (last five moves). □

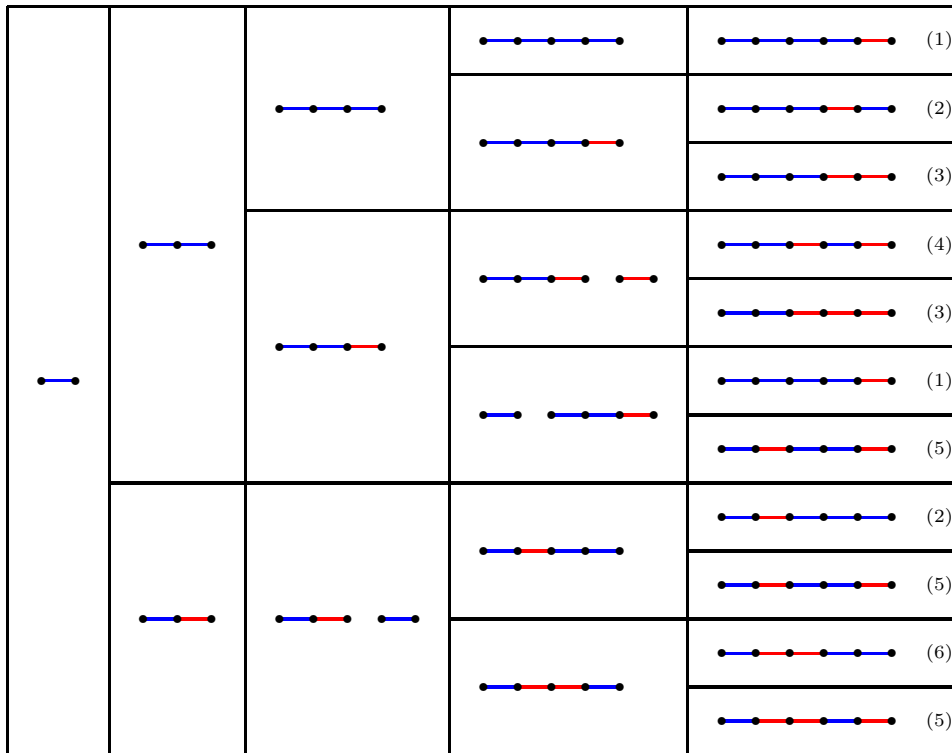


Fig. 4: Forcing a monochromatic P_6 (first five moves)

Further computer assisted results for small paths can be found in (18; 19). (It has been shown that $\tilde{r}(P_7) = 12$, $\tilde{r}(P_8) = 15$, and $\tilde{r}(P_9) = 17$. The values for an asymmetric version of the on-line Ramsey number for paths of lengths at most 8 are also determined.) In general we have the following upper bound.

Theorem 2.3 *Let $k \geq 1$. In $2k - 1$ rounds Builder can force Painter to create two vertex disjoint monochromatic paths (the red one and the blue one) the sum of whose lengths is equal to k .*

In particular, $\tilde{r}(P_n) \leq 4n - 7$ for all $n \geq 2$.

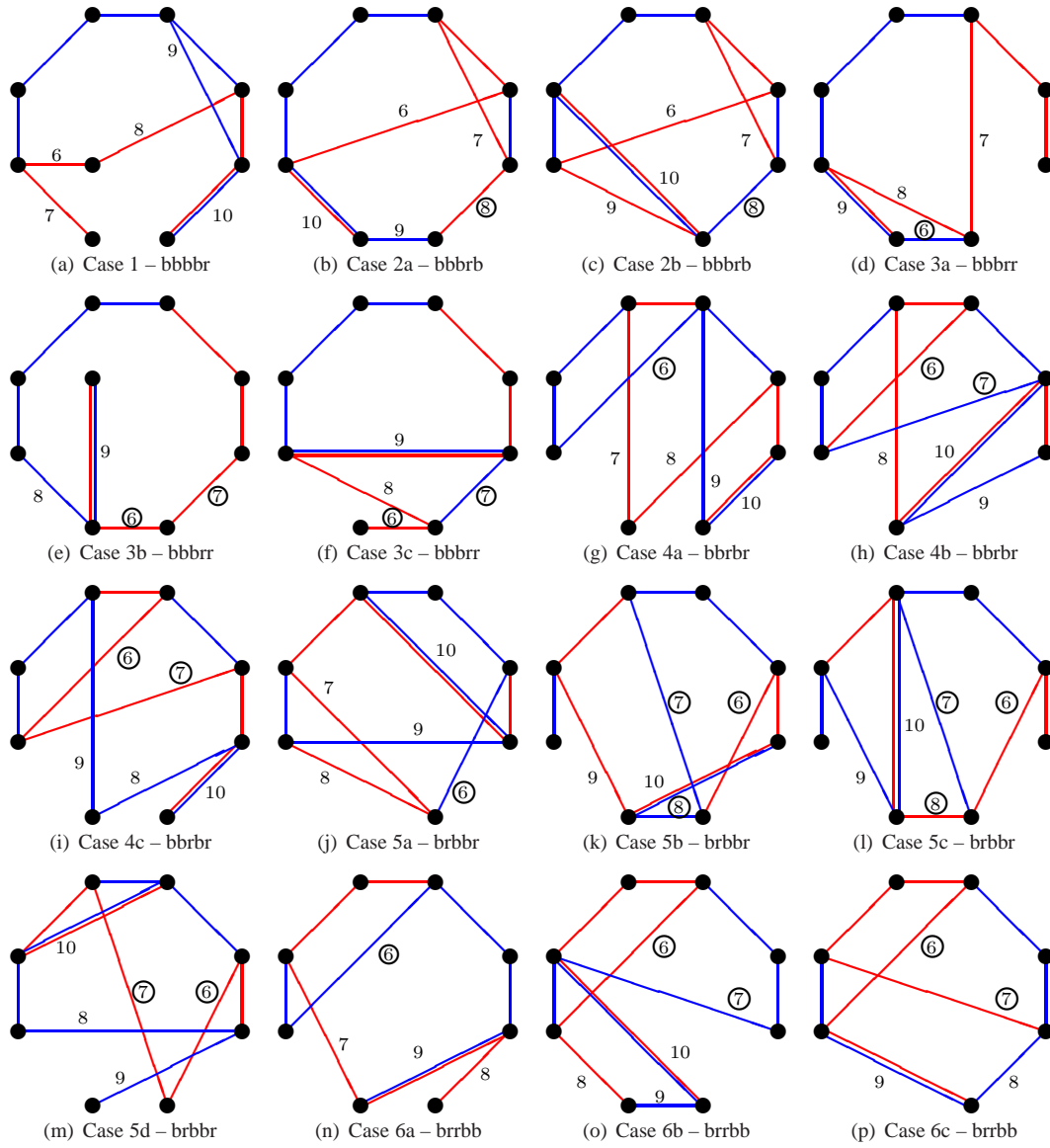


Fig. 5: Forcing a path P_6 (last five moves)

Proof: The proof goes by induction on k . The basis step $k = 1$ is trivial. For the inductive step assume that the statement holds for some $k \geq 1$, that is after $2k - 1$ rounds there is a red path R and a blue path B such that $V(R) \cap V(B) = \emptyset$ and $|E(R) \cup E(B)| = k$. Let v_R and v_B denote any end vertices of R and B , respectively (take any isolated vertex as v_R or v_B if the corresponding path has no edge). If there is no edge $v_R v_B$, then Builder draws the edge $v_R v_B$ and Painter does her job (note that it is only to Builder's advantage if it already exists).

Suppose Painter colors this edge red. Then Builder draws an edge from vertex v_B to a new vertex v . If Painter uses red color, then there are two vertex disjoint monochromatic paths $R' = R + v_B + v$ and $B' = B - v_B$, with $k + 1$ edges in total. If Painter colors $v_B v$ blue, the paths $R' = R$ and $B' = B + v$ satisfy the assertion. If initially Painter colors the edge $v_B v_R$ blue, then Builder draws an edge from v_R to a new vertex v instead, and the situation is symmetric.

In order to finish the proof notice that after $4n - 7 = 2(2n - 3) - 1$ rounds there must be a monochromatic path with at least $n - 1$ edges, that is, a monochromatic P_n . \square

3 Trees

Let $f(n)$ be the maximum value of $\tilde{r}(T)$ taken over all trees with n edges. In (9) Erdős and Graham proved that $r(T) \leq 4e(T) + 1$ for every tree T . From this it follows that $f(n) = O(n^2)$. We show that this bound cannot be improved asymptotically, that is, $f(n) = \Theta(n^2)$. First we derive a general lower bound for $\tilde{r}(H)$ for arbitrary graphs and then present an example that proves the statement. Let $\tau(H)$ denote the smallest number of vertices covering all edges of H . The size-Ramsey number of H can be bounded from below in terms of $\tau(H)$ and the maximum degree $\Delta(H)$ of a graph H . A similar bound holds in the on-line case.

Proposition 3.1 $\tilde{r}(H) \geq \frac{1}{2}\tau(H)(\Delta(H) - 1) + e(H)$ for every graph H .

Proof: Painter's strategy is to color an edge red as long as the red subgraph has maximum degree strictly less than $\Delta(H)$; otherwise she colors an edge blue. So, a red copy of H will never appear. Suppose a blue copy of H has appeared in a graph constructed by Builder. Then at least one end vertex of each blue edge must be incident with exactly $\Delta(H) - 1$ red edges. Hence the red subgraph must have at least $\tau(H)$ vertices of degree $\Delta(H) - 1$ and thus the number of red edges is at least $\frac{1}{2}\tau(H)(\Delta(H) - 1)$. \square

Now, we are ready to show that $f(n) = \Omega(n^2)$. Let B_n be a graph with n edges, obtained from a star $S_{\lfloor n/2 \rfloor} = K_{1, \lfloor n/2 \rfloor}$ by subdividing each edge, and with an additional edge attached to the center of the star if n is odd.

Corollary 3.2 $f(n) \geq \tilde{r}(B_n) \geq \frac{n^2}{8} + \frac{3n}{4}$, for every $n \geq 1$.

Proof: We have $\Delta(B_n) = \tau(B_n) = \lceil n/2 \rceil$ as each cover must contain a non-center vertex from each subdivided edge and one extra vertex for an additional edge if n is odd. \square

If Builder knows that Painter will follow the strategy from Proposition 3.1, then he can force B_n in $\binom{\lceil n/2 \rceil}{2} + n = \frac{n^2}{8} + O(n)$ moves. He starts by building a $\lceil n/2 \rceil$ -clique, which Painter colors entirely red. From now on, Painter will color any vertex incident to a clique vertex blue. So Builder can force a blue B_n in n moves by using the clique vertices as subdivision vertices of B_n . Painter can delay her agony a

little by coloring edges that are incident to the clique and do not create forbidden graph with red, but her real problem was allowing the red clique.

On the other hand, at most $\frac{n^2}{4} + O(n)$ moves are enough to force Painter to create monochromatic B_n .

Proposition 3.3 $\tilde{r}(B_n) \leq \frac{n^2}{4} + \frac{9}{2}n - 7$ for every even $n = 2k \geq 4$.

Proof: Builder starts by playing $6k - 5$ independent edges. Painter is forced to create a monochromatic matching $M = \{x_i y_i : i \in [3k - 2]\}$; suppose the edges of M are blue. Add a new vertex v and edges $v x_i$ for $i \in [3k - 2]$. To avoid creating B_n , Painter must color $2k - 1$ of these edges with red. Say she colors $x_i y_i$ red for $i \in [2k - 1]$. In the next k^2 steps, Builder introduces k new vertices z_l ($l \in [k]$) and edges $z_l x_m$ (for some values of m), forcing Painter to either create a blue star with center at some z_l (thus finishing a blue B_n), or a red matching $M' = \{z_l x_{m_l} : l \in [k]\}$ (thus finishing a red B_n). To do this, he simply plays k edges between z_l and k vertices of the form x_m , where x_m is not incident with an edge in M' .

This takes $(6k - 5) + (3k - 2) + k^2$ edges. Builder can save some edges by being more careful at the beginning. \square

If we restrict to trees of bounded degree then even the size-Ramsey numbers are linearly bounded (cf. (11)).

4 Stars versus trees, cycles, and cliques

In this section we focus on asymmetric version of the on-line Ramsey number where one of the graphs is a star S_n .

It is clear that the Proposition 3.1 can be easily generalized to the asymmetric version we consider. Since the proof is almost identical, we omit it here.

Proposition 4.1 $\tilde{r}(G, H) \geq \frac{1}{2} \max\{\tau(H)(\Delta(G) - 1) + e(H), \tau(G)(\Delta(H) - 1) + e(G)\}$ for arbitrary graphs G and H .

In particular, $\tilde{r}(S_n, H) \geq \frac{1}{2}\tau(H)(n - 1) + e(H)$.

Before we state a general upper bound we need a few more definitions. Let π be a permutation of the vertices of G . For a vertex $v \in V(G)$, let $N_\pi^+(v)$ be the number of neighbors of v that follow it in the permutation. Similarly, we define $N_\pi^-(v)$. Let $d_\pi(G)$ be the maximum ‘‘back degree’’, that is,

$$d_\pi(G) = \max\{N_\pi^-(v) : v \in V(G)\},$$

and $l_\pi(G)$ be the number of vertices with no ‘‘forward neighbors’’, that is,

$$l_\pi(G) = \sum_{v \in V(G)} \delta(N_\pi^+(v) = 0)$$

where $\delta(A)$ is the Kronecker delta function

$$\delta(A) = \begin{cases} 1, & \text{if } A \text{ is true} \\ 0, & \text{otherwise.} \end{cases}$$

Finally, let $n(G)$ denote the number of vertices in a graph G .

Proposition 4.2 $\tilde{r}(S_n, H) \leq \min_{\pi} d_{\pi}(H)(n(H) - l_{\pi}(H))(n - 1) + e(H)$ for any graph H .

Proof: Fix a permutation π of vertices of H . We present Builder's strategy to construct H in an order yielded by permutation π . When we consider a new vertex v of H , we choose a new vertex w in the graph we are constructing and add all edges from w to previously assigned vertices that correspond to the neighbors of v in H . If Painter uses blue color only, we assign w to v and move to the next vertex in H . Otherwise we try again starting from a new vertex w . In total we can have at most $(n(H) - l_{\pi}(H))(n - 1)$ retries since that is exactly an upper bound for the number of red edges when Painter is trying to avoid red S_n and each retry can be associated with a unique red edge generated in the previous try. In each retry we use at most $d_{\pi}(H)$ edges. \square

Now, we are ready to prove bounds for $\tilde{r}(S_n, H)$ in a few special cases where graph H is any tree, cycle or clique.

Proposition 4.3 Let T be any tree with $e(T) \geq 2$ edges and l vertices of degree 1 (leaves). Then for every $n \in \mathbb{N}$

$$\tilde{r}(S_n, T) \leq e(T)n - (l - 1)(n - 1).$$

The proof follows immediately from Proposition 4.2 since there is a permutation π with $d_{\pi}(G) = 1$ and $l_{\pi}(G) = l$. We present an alternative proof below.

Proof: Let $r \in V(T)$ be any vertex of degree at least 2, and select this vertex as a root of the tree T making this tree a rooted one. Let $d(v)$ denote the degree of a vertex v in the tree T . In the first $d(r) + n - 1$ moves Builder draws a star $S_{d(r)+n-1}$ forcing Painter to create a blue star $S_{d(r)}$. This is the first level of T . Then for every leaf v of the blue star $S_{d(r)}$, which is not a leaf of T , Builder draws a star $S_{(d(v)-1)+(n-1)}$ with center in v . This creates the second level of T (see Figure 6). And so on, until a complete blue copy of T (or a red copy of S_n) appears. The total number of edges drawn by Builder is

$$(n - 1)(|V(T)| - l) + e(T) = e(T)n - (l - 1)(n - 1)$$

as asserted. \square

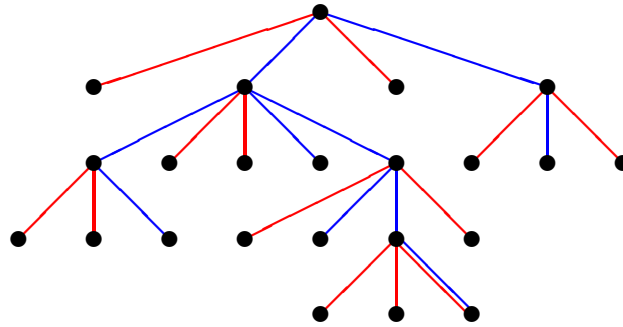


Fig. 6: Forcing a red $K_{1,n}$ or blue T

In next proposition we show that the on-line Ramsey number for a star S_n versus a fixed cycle grows linearly in n . This is different than in the off-line case, where for odd cycles we have a quadratic growth (cf. (16; 17)).

Proposition 4.4 For every $n, k \in \mathbb{N}, k \geq 3$

$$\frac{k}{4}(n + 3) \leq \tilde{r}(S_n, C_k) \leq nk.$$

Proof: The lower bound follows immediately from Proposition 4.1 since $\tau(C_k) \geq k/2$.

For the upper bound, Builder first forces a blue path with $k - 2$ edges. By Proposition 4.3 this takes at most $n(k - 3) + 1$ rounds. Let v_1 and v_{k-1} be the two end vertices of this path. Builder draws a star S_{2n-1} with center in v_{k-1} and new vertices u_1, \dots, u_{2n-1} as its leaves. Clearly at least n edges of this star must be blue. Then he joins these vertices to the vertex v_1 . If there is no red star yet, then there must be a vertex u_i such that both edges $u_i v_1$ and $u_i v_{k-1}$ are blue. Hence a blue cycle is created (see Figure 7). Total number of edges is equal to $((k - 3)n + 1) + (2n - 1) + n = kn$. \square

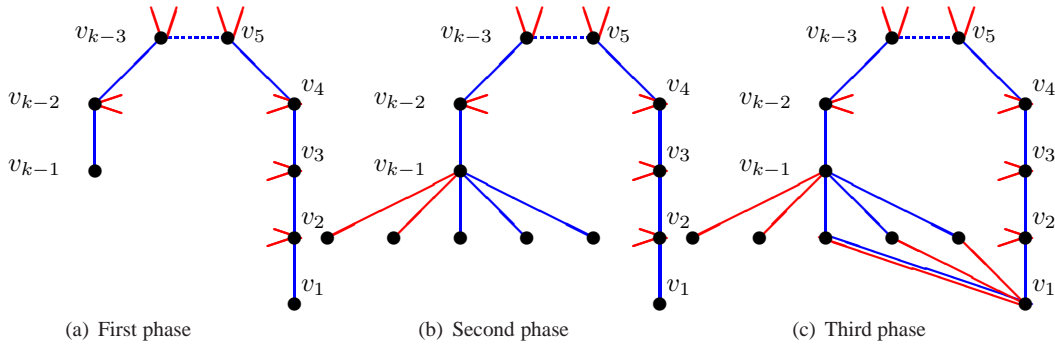


Fig. 7: Forcing a red $K_{1,n}$ or blue C_k

It is clear that better bounds can be achieved with more effort. For example let us mention the following small improvement. If Painter creates red S_{n-1} in the first phase, then Builder can start growing a path from the other direction so that vertex v_1 is attached to the red star. Now in the second phase we can generate one blue edge only (instead of n as before). This saves at least $\min\{n - 1, k - 3\}$ edges. Since there is no hope for a general exact formula, we are happy with the bound which is optimal up to a multiplicative constant.

A similar bounds hold for cliques K_k . The bound is optimal up to a factor of ck for some constant $c > 0$.

Proposition 4.5 For every $n, k \in \mathbb{N}, k \geq 2$

$$\frac{1}{2}(k - 1)(n - 1) + \binom{k}{2} \leq \tilde{r}(S_n, K_k) \leq \binom{k}{2}n.$$

Proof: Again, the lower bound is a simple consequence of Proposition 4.1 since $\tau(K_k) = k - 1$.

For the upper bound, let V_0 be a set of $|V_0| = (k - 1)n + 1$ vertices. First Builder proposes $(k - 1)n$ edges of a star $S_{(k-1)n}$ on the set V_0 . Since Painter is trying to avoid a red star with n edges, at least $(k - 2)n + 1$ edges must be colored blue. Let $V_1 \subset V_0, |V_1| = (k - 2)n + 1$ be a set of vertices of degree

1 connected with a blue edge. In the second phase Builder is drawing a star $S_{(k-2)n}$ on V_1 forcing Painter to use blue color at least $(k-3)n+1$ times. Thus, there is a set $V_2 \subset V_1, |V_2| = (k-3)n+1$ of vertices of degree 2 incident with a blue edges only. Repeating this strategy in each phase, Builder eventually forces Painter to draw either a red star S_n or a blue clique K_k . An example of the game in case of $K_{1,2}$ and K_4 is shown on Figure 8. Note that the game consists of $k-1$ phases and during the i th round, the *Builder* is drawing a star with $(k-i)n$ edges. Thus

$$\tilde{r}(S_n, K_k) \leq \sum_{i=1}^{k-1} (k-i)n = \frac{k(k-1)}{2}n$$

and the assertion follows. \square

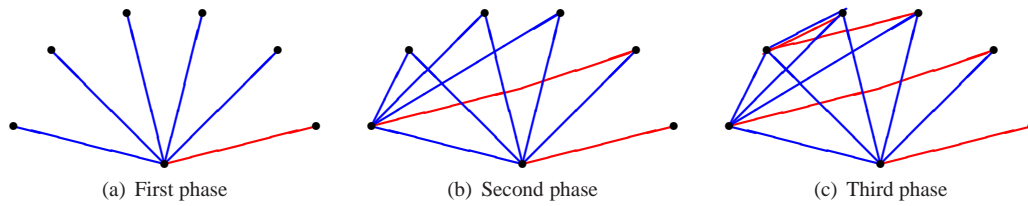


Fig. 8: Forcing a red $K_{1,n}$ or blue K_k

The above results suggest that perhaps $\tilde{r}(S_n, H) \leq n \cdot e(H)$ for every graph H .

5 Discussion

We conclude the paper with some questions for future consideration. Let \mathcal{F} be a fixed family of graphs and let $\tilde{r}(n) = \max\{\tilde{r}(H) : H \in \mathcal{F}, |V(H)| = n\}$. Also let $\hat{r}(n)$ and $r(n)$ be defined analogously for size-Ramsey and Ramsey numbers of graphs from \mathcal{F} , respectively. Clearly for any class \mathcal{F} we have

$$\tilde{r}(n) \leq \hat{r}(n) \leq \binom{r(n)}{2}.$$

A general problem is to determine the asymptotic growth of the function $\tilde{r}(n)$ for some basic graph classes. In particular, to decide whether

$$\lim_{n \rightarrow \infty} \frac{\hat{r}(n)}{\tilde{r}(n)} = \infty. \quad (1)$$

It is natural to expect that (1) should hold for most reasonable families of graphs. Clearly (1) cannot hold if $\hat{r}(n)$ is linear, but, as we demonstrated, it also does not hold for trees, where $\hat{r}(n)$ grows quadratically. We do not know what happens for graphs of bounded degree. A well known theorem of Chvátal, Rödl, Szemerédi, and Trotter (6) asserts that $r(n)$ grows linearly for every fixed Δ (hence $\hat{r}(n)$ is at most quadratic). On the other hand, as proved by Rödl and Szemerédi (22), there are cubic graphs with quadratic size-Ramsey number.

It is obvious that the real challenge is in computing the on-line Ramsey numbers for cliques. Therefore, the following intriguing question was posed by Kurek and Ruciński (14): does (1) hold for the family of all cliques? A natural strategy for Builder is to mimic a recursive construction giving the upper bound for classical Ramsey number $r(K_n)$. This gives for instance $\tilde{r}(K_3) \leq 8$ (in fact, $\tilde{r}(K_3) = 8$; the proof of the lower bound is elementary but nontrivial). Moreover, it has been shown that $\tilde{r}(K_k) \leq 2k \binom{2k-2}{k-1} \sim \frac{1}{2\sqrt{\pi}} \sqrt{k} 4^k$. The third author of this note, with computer support, showed that an asymmetric version of the on-line Ramsey number $\tilde{r}(K_3, K_4)$ is equal to 17, provided a general upper bound for $\tilde{r}(K_k, K_l)$, which gives a slightly better asymptotic upper bound of $\frac{3}{8\sqrt{\pi}} \frac{4^k}{\sqrt{k}}$ for a symmetric version of the on-line Ramsey numbers (see (20) for more details).

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