

Operations on partially ordered sets and rational identities of type A

Adrien Boussicault

► **To cite this version:**

Adrien Boussicault. Operations on partially ordered sets and rational identities of type A. Discrete Mathematics and Theoretical Computer Science, DMTCS, 2013, Vol. 15 no. 2 (2), pp.13–32. <hal-00980747>

HAL Id: hal-00980747

<https://hal.inria.fr/hal-00980747>

Submitted on 18 Apr 2014

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Operations on partially ordered sets and rational identities of type A

Adrien Boussicault

Institut Gaspard Monge, Université Paris-Est, Marne-la-Vallée, France

received 13th February 2009, revised 1st April 2013, accepted 2nd April 2013.

We consider the family of rational functions $\psi_w = \prod (x_{w_i} - x_{w_{i+1}})^{-1}$ indexed by words with no repetition. We study the combinatorics of the sums Ψ_P of the functions ψ_w when w describes the linear extensions of a given poset P . In particular, we point out the connexions between some transformations on posets and elementary operations on the fraction Ψ_P . We prove that the denominator of Ψ_P has a closed expression in terms of the Hasse diagram of P , and we compute its numerator in some special cases. We show that the computation of Ψ_P can be reduced to the case of bipartite posets. Finally, we compute the numerators associated to some special bipartite graphs as Schubert polynomials.

Keywords: poset, rational function, Schubert polynomials

1 Introduction

For each word w without repetition, we define the rational function

$$\psi_w := \frac{1}{(x_{w_1} - x_{w_2}) \cdot (x_{w_2} - x_{w_3}) \cdots (x_{w_{n-1}} - x_{w_n})}$$

where $w_1 \dots w_n$ are the letters of w . Permutations are special words with no repetition, and summing these functions on intervals of the permutohedron (*i.e.*, for the weak order) gives remarkable properties. For example, we will show that, when a permutation σ avoids some patterns, the sum over the initial interval $[Id, \sigma]$ can be simplified to a product $\prod (x_i - x_j)^{-1}$ where the product runs over the edges (i, j) of a particular tree.

These patterns are known to characterise some families of Schubert varieties. Schubert varieties are indexed by permutations, and the varieties which are non-singular are those whose indexing permutation does not contain the pattern 2143 nor 1324. In (2), Cortez has described geometrical properties of Schubert varieties for permutations avoiding 1324 and $\overline{2143}$ ⁽ⁱ⁾. This was further clarified by Woo and Yong in (6),

⁽ⁱ⁾ A permutation σ avoids the pattern 1324 if there exist no integers $1 \leq i_1 < i_2 < i_3 < i_4 \leq n$ satisfying $\sigma_{i_1} < \sigma_{i_3} < \sigma_{i_2} < \sigma_{i_4}$. A permutation σ contains the pattern $\overline{2143}$ if for some indices $1 \leq i_1 < i_2 < i_3 < i_4 \leq n$ we have $\sigma_{i_2} \leq \sigma_{i_1} \leq \sigma_{i_4} \leq \sigma_{i_3}$ with the further restriction that there is no $i_1 \leq j \leq i_4$ such that $\sigma_{i_1} \leq \sigma_j \leq \sigma_{i_4}$.

and Butler and Bousquet-Mélou in (1). They used the fact that an initial interval $[Id, \sigma]$ is the set of linear extensions of a partially ordered set whose Hasse diagram is a tree when σ avoids 1324 and $\overline{2143}$.

Surprisingly, the same patterns occur in the study of the sum

$$\Psi_P := \sum_{w \in \mathcal{L}(P)} \psi_w,$$

where $\mathcal{L}(P)$ is the set of linear extensions of P . In fact, our work is closely connected to a study of Greene (3) on rational identities related to the Murnaghan–Nakayama formula for S_n (type A). Greene gave in (3) a closed expression for the Ψ_P of a planar poset P . He showed the equality

$$\Psi_P = \begin{cases} 0 & \text{if } P \text{ is a disconnected graph,} \\ \prod_{a,b \in P} (x_a - x_b)^{\mu_P(a,b)} & \text{if } P \text{ is a connected graph,} \end{cases}$$

where $\mu(x, y)$ is the Möbius function of the poset P . In the case of a permutation avoiding the patterns 1324 and $\overline{2143}$, the poset is planar, and the Möbius function takes only values 0 or -1 . Therefore (Corollary 4.1) the numerator of Ψ_P equals 1 if and only if the Hasse diagram of P is acyclic.

The aim of this paper consists in pointing out the connexions between some operations on posets and rational identities involving the Ψ_P . To study the rational functions Ψ_P , we introduce some operations on posets in Section 2, and describe in Section 3 the identities on the rational functions that these operations induce. Finally, we give some explicit examples in Section 4: acyclic posets, 1-cycle posets and λ -complete posets (related to Schubert polynomials).

2 Operations on partially ordered sets

We will see that the knowledge of the Hasse diagram of P gives some properties of the function Ψ_P . In this section we study four operations on posets (considered as graphs): collapse, operation of sorting a subset, contraction and suppression of extremal elements.

2.1 Basic definitions

We recall that a *partially ordered set (poset)* $P = (A, \leq)$ is a set A endowed with a reflexive, antisymmetric and transitive binary relation. For simplicity, when there is no ambiguity, we will use $a \in P$ instead of $a \in A$.

A total order T is a linear extension of a poset P if, whenever $a \leq b$ in P there holds $a \leq b$ in T . We denote by $\mathcal{L}(P)$ the set of linear extensions of P .

Classically, the *covering relation* is the minimal⁽ⁱⁱ⁾ relation whose transitive closure is \leq . This relation will be denoted by \preceq . The *Hasse diagram* of P , denoted by $H(P)$, is the oriented graph of the covering relation of P , drawn in such a way that if $a \preceq b$, then b is drawn to the right of a ⁽ⁱⁱⁱ⁾. The Hasse diagram displays the minimal set of relations generating P by transitivity.

The set of the *inner vertices* of P , that is, those being neither minimal nor maximal for \leq , will be denoted by $In(P)$. Its complementary set is called the *boundary* of P : $Bound(P) := P \setminus In(P)$

⁽ⁱⁱ⁾ as a subset of $P \times P$.

⁽ⁱⁱⁱ⁾ Usually, Hasse diagrams are drawn from bottom to top, but this representation takes more space and is less natural for our purposes.

2.2 Permutations and posets

Consider the symmetric group \mathfrak{S}_n endowed with the weak order.

The *weak order* (also called right permutohedron order) is the order on permutations obtained by defining the successors of a permutation σ as the permutations $\sigma.s_i$ if this permutation has more inversions than σ , where $s_i = (i, i + 1)$ exchanges the numbers at places i and $i + 1$ of σ .

For any pair (σ, τ) of permutations, one constructs the interval $[\sigma, \tau]$ as the set of permutations greater than or equal to σ and lower than or equal to τ .

For example, the interval $[123456, 132564]$ in \mathfrak{S}_6 contains exactly the permutations

$$[123456, 132564] = \{132564, 123564, 132546, 123546, 132456, 123456\}.$$

Proposition 2.1 *Let σ be a permutation. There exists a poset \mathbb{P}_σ such that $\mathcal{L}(\mathbb{P}_\sigma) = [Id, \sigma]$. Moreover, the Hasse diagram of \mathbb{P}_σ is obtained applying Algorithm 1.*

Proof: Let \leq_σ be the binary relation on $\{1, \dots, n\}$ defined by:

$$i \leq_\sigma j \quad \text{if and only if} \quad i \leq j \text{ and } \sigma_i \leq \sigma_j.$$

The binary relation \leq_σ is a partial order on $\{1, \dots, n\}$. Indeed, reflexivity, transitivity and antisymmetry of $>_\sigma$ are deduced from reflexivity, transitivity and antisymmetry of \leq .

Let \mathbb{P}_σ be the poset defined by:

$$\mathbb{P}_\sigma = (\{1, \dots, n\}, \leq_\sigma).$$

The interval $[Id, \sigma]$ is the set of permutations obtained from σ by permuting recursively all consecutive decreasing pairs of letters of σ . It is well known that $[id, \sigma]$ is exactly the set

$$[Id, \sigma] = \{\tau \in \mathfrak{S}_n \mid \text{if } i \leq j \text{ and } \sigma_i^{-1} \leq \sigma_j^{-1} \text{ then } \tau_i^{-1} \leq \tau_j^{-1}\}$$

which is, by definition, the set of linear extensions of \mathbb{P}_σ .

The Hasse diagram of such a poset is built by using Algorithm 1. Indeed, the first ten steps draw each vertices of the graph satisfying that if $i < j$ and $\sigma_i < \sigma_j$ then σ_i is draw to the right of σ_j . Step 11 draws the edges of the Hasse diagram. \square

Example 2.2 *Figure 1 shows an execution of Algorithm 1.*

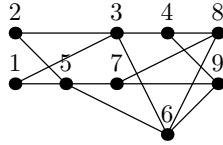


Fig. 1: Hasse diagram of $\mathbb{P}_{215736498}$

Algorithm 1 construction of \mathbb{P}_σ

Let σ be a permutation and i the index of the current letter.

- 1: $i \leftarrow |\sigma|$
 - 2: Write σ_i .
 - 3: **while** $i > 0$ **do**
 - 4: **if** $\sigma_{i-1} > \sigma_i$ **then**
 - 5: write σ_{i-1} in the same column as σ_i
 - 6: **else**
 - 7: write σ_{i-1} to the left of σ_i
 - 8: **end if**
 - 9: $i \leftarrow i - 1$.
 - 10: **end while**
 - 11: Draw an edge between each pair (σ_i, σ_j) satisfying
 - a) $\sigma_i < \sigma_j$,
 - b) σ_j is to the right of σ_i ,
 - c) there is no vertex σ_k such that $\sigma_i < \sigma_k < \sigma_j$ and k is to the left of j and to the right of i .
 - 12: The graph obtained is the Hasse diagram of the poset \mathbb{P}_σ .
-

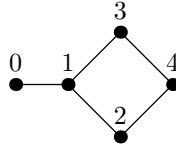


Fig. 2: 0 is a free vertex and $(0, 1)$ is a non-inner edge

2.3 Collapses

Hasse diagrams can be viewed as non-oriented graphs. We will use notations and definitions given in (9). In his book (9), Giblyn defined a *free vertex* as a vertex belonging to exactly one edge. An edge with no free vertex is an *inner edge*.

Example 2.3 The poset P (see Figure 2) has only one free edge: 0. The only non-inner edge is $(0, 1)$.

The action which consists in removing a non-inner edge together with one of its free vertices is an *elementary collapse*.

A *collapse* is a sequence of elementary collapses. We will denote by $Coll(\mathcal{G})$ the collapsed graph \mathcal{G} , that is the unique (up to relabeling) maximal subgraph of \mathcal{G} without free vertices.

Example 2.4 Figure 3 shows the collapsed graph of Figure 2.

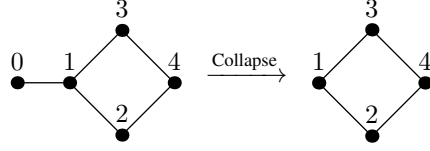


Fig. 3: A collapse

2.4 Subposets and linear extensions

Let $P = (A, \leq)$ be a poset. Each subset $A' \subset A$ can be naturally endowed with the partial order induced by P which is the projection on A' of the order of P . Let us denote such a poset by P' . Classically, we will say that P' is a *subposet* of P whose *support* is A' . If w is a linear extension of P' , the transitive relation generated by \leq and \leq_w is a partial order on A and the corresponding poset will be denoted by P_w .

Lemma 2.5 *Let $P = (A, \leq)$ be a poset, $A' \subset A$ and $P' = (A', \leq)$, we have*

$$\mathcal{L}(P) = \bigsqcup_{w \in \mathcal{L}(P')} \mathcal{L}(P_w)$$

where \bigsqcup denotes the disjoint union.

Proof: Let us prove first that $\mathcal{L}(P) = \bigcup_{w \in \mathcal{L}(P')} \mathcal{L}(P_w)$. By definition of P_w one has the inclusion $\bigcup_{w \in \mathcal{L}(P')} \mathcal{L}(P_w) \subset \mathcal{L}(P)$. Conversely, each linear extension of P induces a total order on P' . Hence, $\mathcal{L}(P) = \bigcup_{w \in \mathcal{L}(P')} \mathcal{L}(P_w)$.

Furthermore, if w and w' are two distinct extensions of P' then, obviously, $\mathcal{L}(P_w) \cap \mathcal{L}(P_{w'}) = \emptyset$. \square

Example 2.6 *The linear extensions of the poset P (see Figure 4) can be partitioned in disjoint subsets indexed by the linear extensions of the subposet P' .*

$$\begin{aligned} \mathcal{L}(P) &= \mathcal{L}(P_{456}) \sqcup \mathcal{L}(P_{465}) \\ \mathcal{L}(P_{456}) &= \{123456, 213456, 124356, 214356\} \\ \mathcal{L}(P_{465}) &= \{123465, 213465, 124365, 214365\} \end{aligned}$$

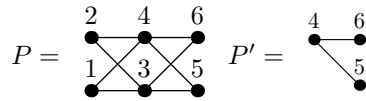


Fig. 4: A poset P with a subposet P' .

2.5 Contractions

We describe in Proposition 2.10 an operation on posets called contraction. This operation is fundamental since it allows to simplify posets by decreasing the number of edges. Before defining contraction, we need to consider the Lemmas 2.7, 2.8 and 2.9.

We denote the negation of a binary relation \leq by $\not\leq$.

Lemma 2.7 *Let \leq_1 and \leq_2 be two antisymmetric relations on A such that $a \leq_2 b$ implies $b \not\leq_1 a$ or $b = a$. The binary relation $a \leq_3 b$ defined by $a \leq_1 b$ or $a \leq_2 b$ is antisymmetric.*

Proof: Let a, b be two elements of A such that $a \leq_3 b$ and $b \leq_3 a$. One needs to examine four cases.

- i) if $a \leq_1 b$ and $b \leq_1 a$ then $a = b$ because \leq_1 is antisymmetric.
- ii) if $a \leq_1 b$ and $b \leq_2 a$ then, by hypothesis, $a \leq_1 b$ and $(a \not\leq_1 b$ or $b = a)$ which implies $b = a$.
- iii) The case $a \leq_2 b$ and $b \leq_1 a$ is similar with the previous statement.
- iv) if $a \leq_2 b$ and $b \leq_2 a$ then $a = b$ because \leq_2 is antisymmetric. □

Let $P = (A, \leq)$ be a poset and c, d two elements of A satisfying $c \preceq d$. We will denote by $a \leq_2 b$ the relation defined by $((a \leq c) \text{ or } (a \leq d))$ and $((c \leq b) \text{ or } (d \leq b))$.

Lemma 2.8 *The relation $a \leq_2 b$ is antisymmetric on $A \setminus \{d\}$.*

Proof: By definition, $a \leq_2 b$ and $b \leq_2 a$ is equivalent to

$$(a \leq c \text{ or } a \leq d) \text{ and } (c \leq a \text{ or } d \leq a) \text{ and } (b \leq c \text{ or } b \leq d) \text{ and } (c \leq b \text{ or } d \leq b).$$

Since $c \preceq d$, we deduce that $(a = c \text{ or } a = d)$ and $(b = c \text{ or } b = d)$. Remarking that $a \neq d$ and $b \neq d$, we obtain $a = b = c$. □

Lemma 2.9 *Let a, b be two elements of $A \setminus \{d\}$. Then $a \leq_2 b$ implies $(b \not\leq a \text{ or } a = b)$.*

Proof: Let a, b be two elements of $A \setminus \{d\}$ satisfying $a \leq_2 b$. By definition of \leq_2 we have $((a \leq c) \text{ or } (a \leq d))$ and $((c \leq b) \text{ or } (d \leq b))$. Since $c \preceq d$, the previous formula implies $a \leq b$ or $(a \leq d \text{ and } c \leq b)$. If $b \leq a$, we deduce that $c \leq b \leq a \leq d$. Since $c \preceq d$, $a \neq d$ and $b \neq d$, we have $a = b = c$ and we obtain

$$a \leq_2 b \Rightarrow (a \leq b \text{ or } a, b \text{ are not comparable}) \Rightarrow (b \not\leq a \text{ or } a = b).$$

□

We can now define the contraction.

Proposition 2.10 *Let $P = (A, \leq)$ be a poset and c, d be two elements of A . If $c \preceq d$, then the relation $\leq_{d=c}$ defined on $A \setminus \{d\}$ by*

$$a \leq_{d=c} b \Leftrightarrow (a \leq b) \text{ or } ((a \leq c) \text{ or } (a \leq d)) \text{ and } ((c \leq b) \text{ or } (d \leq b))$$

is a partial order.

Proof: The reflexivity and transitivity of $\leq_{d=c}$ follow straightforwardly from the reflexivity and transitivity of \leq .

By Lemmas 2.7 and 2.8, the relation \leq_2 is antisymmetric, and $a \leq_2 b$ implies $b \not\leq_2 a$ or $a = b$. We deduce (cf. Lemma 2.9) that $\leq_{d=c}$ is antisymmetric. \square

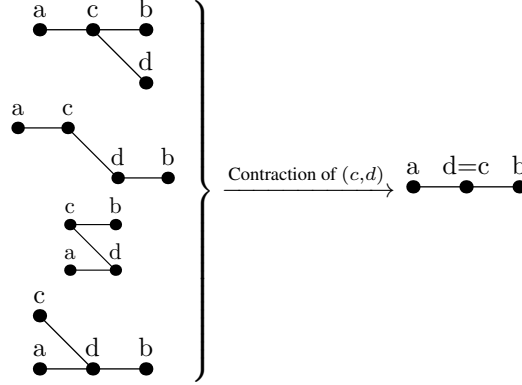


Fig. 5: Contraction of the covering relation (c, d) .

When c covers d , we denote by $P_{d=c} = (A \setminus \{d\}, \leq_{d=c})$ the *contraction of the edge* (c, d) .

Proposition 2.11 *Let P be a poset and (c, d) be an edge in $H(P)$. Then $w'cdw''$ is a linear extension of P if and only if $w'cw''$ is a linear extension of $P_{d=c}$.*

Proof: We will denote by i_z the position of z in the word $w'cdw''$ and by j_z the position of z in the word $w'cw''$.

Suppose that $w'cdw''$ is a linear extension of $P = (A, \leq)$. Let a, b be two elements of $A \setminus \{d\}$ such that $a \leq_{P_{d=c}} b$ and $a \neq b$. By definition of $\leq_{P_{d=c}}$, we have $(a \leq b)$ or $((a \leq c$ or $a \leq d)$ and $(c \leq b$ or $d \leq b))$. Since $w'cdw''$ is a linear extension, $(a \leq b)$ implies $i_a \leq i_b$. But $i_d = i_c + 1$, $a \neq d$ and $c \neq d$, hence $((a \leq c$ or $a \leq d)$ and $(c \leq b$ or $d \leq b))$ implies $i_a \leq i_c \leq i_b$. Finally, in all the cases, if $a \leq_{P_{d=c}} b$ then $i_a \leq i_b$ and equivalently $w'cw'' \in \mathcal{L}(P_{d=c})$.

Conversely, suppose that $w'cw''$ is a linear extension of $P_{d=c}$. Let a, b be two elements of A such that $a \leq b$ and $a \neq b$. If $a \neq d$ and $b \neq d$ then $a \leq b$ implies $a \leq_{P_{d=c}} b$ and $j_a \leq j_b$. If $a = d$ or $b = d$ then $a \leq d$ and $d \leq b$ which implies $a \leq_{P_{d=c}} b$ and $j_a \leq j_b$. Then in all the cases, $a \leq b$ implies $j_a \leq j_b$. Equivalently, $w'cdw''$ is a linear extension of P . \square

Example 2.12 *The edge $(4, 5)$ of the poset P (see Figure 6) can be contracted. The linear extensions of $P_{5=4}$ are*

$$\mathcal{L}(P_{5=4}) = \{1234, 2134\},$$

and, the set obtained by removing all words of $\mathcal{L}(P)$ with no factor 45 has the same size than $\mathcal{L}(P_{5=4})$.

$$\{w'45w'' \in \mathcal{L}(P)\} = \{12345, 21345\}$$

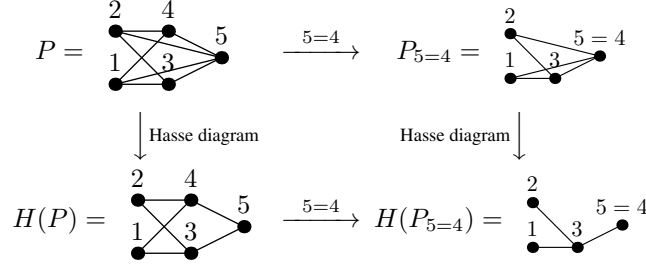


Fig. 6: The poset P and $P_{5=4}$, their Hasse diagrams and their linear extensions.

2.6 Decontraction

In the previous section, we have described the operation of contraction. We will now show that all posets can be obtained from bipartite posets. A poset is *bipartite* if it has only extremal elements.

Theorem 2.13 *Each poset can be obtained from a bipartite poset by applying a succession of contractions.*

Proof: Consider a poset P . We construct a new poset \bar{P} by duplicating each vertex of $In(P)$

$$\bar{P} = P \cup \{\bar{a} \mid a \in In(P)\}.$$

The poset \bar{P} is endowed with the relation \leq' defined by $b \leq' c$ if and only if one of the following statements is true:

1. $b = c$,
2. $c = \bar{a}$ with $a \in In(P)$, $b \in P$ and $b \leq a$,
3. $b \in P$, $c \in Bound(P)$ and $b \leq c$,
4. $c = \bar{b}$ with $b \in In(P)$.

The relation \leq' is reflexive because of the first rule of construction. By construction, each element of \bar{P} is either minimal or maximal. We obtain that \leq' is transitive and antisymmetric. Finally, when contracting the edges $\{(a, \bar{a}) \mid a \in In(P)\}$, the poset P is recovered. \square

Let P be a poset. We will call the poset \bar{P} obtained in the proof of Theorem 2.13 the *decontraction* of P .

Example 2.14 *The posets P' (see Figure 7) are obtained by decontraction of the poset P .*

2.7 Suppression of extremal elements

Another way to simplify posets consists in removing extremal vertices. Proposition 2.15 shows that linear extensions beginning with an extremal element c are equivalent to linear extensions of a new poset obtained by removing c from P . We will see in Subsection 3.2 that this property admits an interpretation in terms of rational functions.

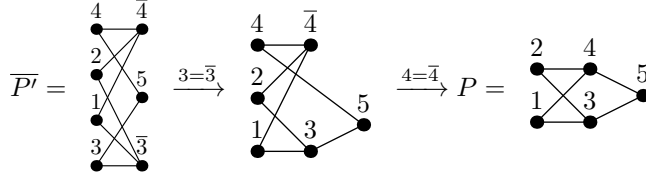


Fig. 7: The poset P is the contraction of the bipartite poset \overline{P} .

Proposition 2.15 *Assume that c is a minimal (resp. maximal) vertex. Then, cw (resp. wc) is a linear extension of P if and only if w is a linear extension of $P \setminus \{c\}$.*

Proof: Suppose that cw is a linear extension of P and let a, b be two elements of $P \setminus \{c\}$. Hence, if $a \leq_{P \setminus \{c\}} b$ then $a \leq_P b$ which implies $a \leq_w b$ and $w \in \mathcal{L}(P \setminus (\{c\}))$.

Conversely, suppose that w is a linear extension of $P \setminus \{c\}$. Let $a, b \in P$ be such that $a \leq_P b$. If $a \neq c$ then, since c is minimal, we have $b \neq c$. This implies $a \leq_{P \setminus \{c\}} b$ and $a \leq_w b$. So we have $a \leq_{cw} b$. If $a = c$, then trivially $a \leq_{cw} b$. We deduce that $cw \in \mathcal{L}(P)$. \square

Example 2.16 *The vertex 4 of the poset P (see Figure 8) can be deleted. The linear extensions of $P \setminus 4$ are*

$$\mathcal{L}(P \setminus 4) = \{123, 213\},$$

and, the set obtained by setting all word of $\mathcal{L}(P)$ having 4 in the last position has the same size than $\mathcal{L}(P \setminus 4)$.

$$\{w'4 \in \mathcal{L}(P)\} = \{1234, 2134\}$$

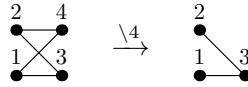


Fig. 8: Suppression of the extremal element 4.

3 Operations on rational functions

In the previous section we have described some operations on posets. We will now see the connexions with the rational functions Ψ_P .

3.1 Residues and contractions

In this subsection, we show that the fraction $\Psi_{P_{c=d}}$ can be obtained from Ψ_P by applying a residue.

Theorem 3.1 *Let c and d be two elements of a poset P . We have*

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) \cdot \Psi_P = \begin{cases} \Psi_{P_{d=c}} & \text{if } (c, d) \text{ is an edge of the Hasse diagram of } P, \\ 0 & \text{otherwise.} \end{cases}$$

Proof: Taking the residue at $x_c = x_d$ in Ψ_P , we get

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_P = \lim_{x_d \rightarrow x_c} \sum_{w \in \mathcal{L}(P)} (x_c - x_d) \psi_w. \quad (1)$$

If neither cd nor dc are factors of w , then $(x_c - x_d)$ is a factor neither of $Den(\psi_w)$ nor $Num(\psi_w)$. It follows that $\lim_{x_d \rightarrow x_c} (x_c - x_d) \psi_w = 0$.

If either cd or dc are factors of w , then $(x_c - x_d)$ is a factor of $Den(\psi_w)$ with multiplicity 1. It follows that the $\lim_{x_d \rightarrow x_c} (x_c - x_d) \psi_w$ converges.

We conclude that sum and limit can be permuted in Equation (1) to obtain

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_P = \sum \lim_{x_d \rightarrow x_c} (x_c - x_d) \psi_w, \quad (2)$$

where the sum runs over the linear extensions w of P having cd and dc as factor (when considered as a word).

We need to consider three cases:

- 1) If c and d are not comparable. Obviously, the word $w'cdw''$ is a linear extension of P if and only if $w'dcw''$ is also a linear extension of P . Hence, by considering the pairs $\psi_{w'cdw''}$ and $\psi_{w'dcw''}$ in (2), we obtain

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_P = \sum_{w'cdw'' \in \mathcal{L}(P)} \lim_{x_d \rightarrow x_c} (x_c - x_d) [\psi_{w'cdw''} + \psi_{w'dcw''}]. \quad (3)$$

Setting $f(x_c, x_d) = (x_c - x_d) \psi_{w'cdw''}$, one has $f(x_d, x_c) = -(x_c - x_d) \psi_{w'dcw''}$. Hence,

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) [\psi_{w'cdw''} + \psi_{w'dcw''}] = \lim_{x_d \rightarrow x_c} (f(x_c, x_d) - f(x_d, x_c)) = 0. \quad (4)$$

Finally, by using Equations (3) and (4), we deduce that $\lim_{x_c \rightarrow x_d} (x_c - x_d) \Psi_P = 0$.

- 2) If c and d are comparable but $c \not\leq d$ and $d \not\leq c$. Assuming that $c \leq d$ (the other case being similar), there is at least one element a such that $c \leq a \leq d$. Then $\mathcal{L}(P)$ contains no word having neither cd nor dc as a factor and the residue $\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_P$ equals 0.
- 3) If $c \leq d$ (the case $d \leq c$ is similar), one has

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_P = \sum \lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_w,$$

where the sum runs over the linear extensions w of P satisfying $w = w'cdw''$. From Proposition 2.11, one has

$$\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_P = \sum_{w'cw'' \in \mathcal{L}(P_{d=c})} \lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_{w'cdw''}.$$

A straightforward computation gives $\lim_{x_d \rightarrow x_c} (x_c - x_d) \Psi_{w'cdw''} = \Psi_{w'cw''}$. The result follows. \square

Theorem 3.1 and Proposition 2.13 show that the knowledge of the fraction Ψ_P for each bipartite poset is sufficient to compute any Ψ_P by applying a sequence of residues. In fact, we can construct other algorithms like the decontraction to obtain bipartite posets with such properties. However, we will show that the decontraction implies a stronger result (Proposition 3.10, Section 3.4).

3.2 Limits and suppression of extremal elements

The deletion of a vertex admits also an interpretation in terms of rational functions.

Theorem 3.2 *Let c be an element of a poset P ,*

$$\lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_P) = \begin{cases} -\Psi_{P \setminus c} & \text{if } c \text{ is maximal,} \\ \Psi_{P \setminus c} & \text{if } c \text{ is minimal,} \\ 0 & \text{otherwise.} \end{cases}$$

Proof: We have

$$\lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_P) = \lim_{x_c \rightarrow +\infty} \sum_{w \in \mathcal{L}(P)} (x_c \cdot \Psi_w).$$

We split the sum into two parts, the first one being the sum over all the words beginning or finishing by c

$$\lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_P) = \lim_{x_c \rightarrow +\infty} \left(\sum_{\substack{w \in \mathcal{L}(P) \\ w = cw' \text{ or } w = w'c}} (x_c \cdot \Psi_w) + \sum_{\substack{w \in \mathcal{L}(P), w = w'cw'' \\ \text{and } w', w'' \text{ are non-trivial}}} (x_c \cdot \Psi_w) \right). \quad (5)$$

If c is the first or the last letter in w , then the degree of x_c in ψ_w is -1 and $\lim_{x_c \rightarrow +\infty} (x_c \psi_w) = \pm \psi_{w'}$. If c is not in the first and last position in w , then the degree of x_c in ψ_w is -2 and $\lim_{x_c \rightarrow +\infty} (x_c \psi_w) = 0$. We deduce that limit and sums in Equation (5) can be interchanged to obtain

$$\lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_P) = \sum_{\substack{w \in \mathcal{L}(P) \\ w = cw' \text{ or } w = w'c}} \lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_w).$$

We get

$$\lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_P) = \sum_{\substack{w \in \mathcal{L}(P) \\ w = cw' \text{ or } w = w'c}} \Psi_{w'}.$$

By Proposition 2.15, we have

$$\lim_{x_c \rightarrow +\infty} (x_c \cdot \Psi_P) = \pm \Psi_{P \setminus \{c\}}.$$

□

3.3 Connectivity and vanishing conditions

In this subsection, we investigate the interpretation of connectivity in terms of rational functions. In particular, Corollary 3.6 shows that $\Psi_P = 0$ if and only if P is disconnected.

Definition 3.3 (Greene (3)) *A poset P is planar if its Hasse diagram may be ordered-embedded in $\mathbb{R} \times \mathbb{R}$ without edge crossings, even when extra maximal and minimal elements are added.*

Example 3.4 *Following the definition of Greene (3), Figures 9 and 10 are examples of non-planar and planar posets, respectively.*

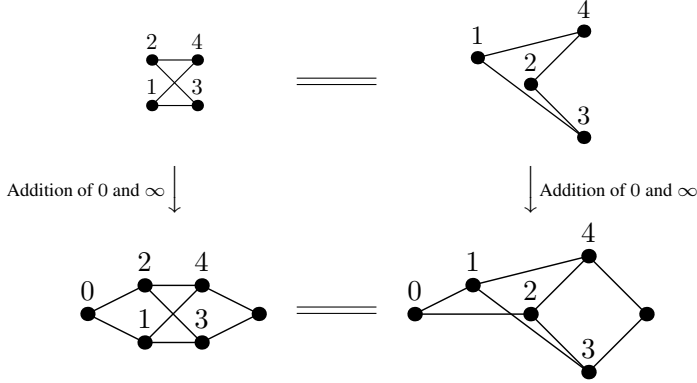


Fig. 9: The poset \mathbb{P}_{2143} is not planar.

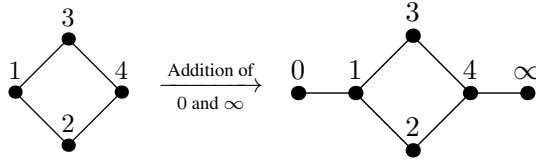


Fig. 10: The poset \mathbb{P}_{1324} is planar.

Theorem 3.5 (Greene (3)) *Let P be a planar poset, then*

$$\Psi_P = \begin{cases} 0 & \text{if } P \text{ is not connected,} \\ \prod_{\substack{a,b \in P \\ a < b}} (x_a - x_b)^{\mu_P(a,b)} & \text{if } P \text{ is connected.} \end{cases},$$

where μ_P denotes the Möbius function of P .

As a consequence, one has:

Corollary 3.6 *Let P be a poset, the Hasse diagram of P is connected if and only if $\Psi_P \neq 0$.*

Proof: Suppose first that P is connected and $\Psi_P = 0$. Since contractions preserve the connectivity of a poset, we contract successively the edges of P until we obtain a new poset with only two elements $c \leq d$. Using Theorem 3.1, we get Ψ_{c-d} from Ψ_P by applying a succession of residues. It follows that $\Psi_{c-d} = 0$. This is in contradiction with the direct computation $\Psi_{c-d} = \frac{1}{x_c - x_d}$. Hence, $\Psi_P \neq 0$.

Conversely, we consider the case of a disconnected poset P . Let C_1 and C_2 be two disconnect components of P . Applying Lemma 2.5, one gets:

$$\mathcal{L}(P) = \bigsqcup_{w_1 \in \mathcal{L}(C_1), w_2 \in \mathcal{L}(C_2)} \mathcal{L}((P_{w_1})_{w_2}).$$

The disjoint union interpreted as a sum, when stated in terms of rational function, gives

$$\Psi_P = \sum_{w_1 \in \mathcal{L}(C_1), w_2 \in \mathcal{L}(C_2)} \Psi_{(P_{w_1})_{w_2}}.$$

Since each poset $(P_{w_1})_{w_2}$ is planar and disconnected, Greene's theorem (Theorem 3.5) gives $\Psi_P = 0$. \square

3.4 Reduced fractions and Hasse diagrams

In this subsection, we present two main results on Ψ_P . We characterise the denominator of the reduced fraction Ψ_P with the help of the Hasse diagram of P , and we give an interpretation of the degree of the numerator of Ψ_P in terms of cycles.

We call *cycles* (resp. *cyclomatic number*) of a poset P , the cycles (resp. cyclomatic number) of the Hasse diagram of P . We recall that the cyclomatic number of a graph G is the minimal number of edges we have to remove from G to obtain a graph without cycle (see e.g. (8) or (9)).

We denote by $Den(\Psi_P)$ the denominator of the reduced fraction Ψ_P and by $Num(\Psi_P)$ its numerator.

Corollary 3.7 *Let P be a connected poset, then:*

$$Den(\Psi_P) = \prod_{c \prec d} (x_c - x_d).$$

Proof: Theorem 3.1 implies that $\prod_{c \prec d} (x_c - x_d)$ is a factor of $Den(\Psi_P)$. Since contractions preserve the connectivity (Corollary 3.6), we deduce that $Den(\Psi_P)$ is exactly $\prod_{c \prec d} (x_c - x_d)$. \square

This result shows that the Hasse diagram is a relevant notion in our context. Moreover, the following corollary confirms the special status of the Hasse diagram and in particular of its cycles.

Corollary 3.8 *Let P be a connected poset, the degree of $Num(\Psi_P)$ is equal to the cyclomatic number of the Hasse diagram of P .*

Proof: Let P be a connected poset with n elements. Let m be the number of edges in $H(P)$ and o its cyclomatic number. By Corollary 3.7, we deduce that the degree of the numerator of the reduced fraction is at most equal to $m - n + 1$ that is, from the Euler formula, (see e.g. (8) or (9)), the cyclomatic number of $H(P)$. The polynomial $Num(P)$ being homogeneous, it equals either zero or its degree is o . Since P is connected, Corollary 3.6 finishes the proof. \square

Example 3.9 *The cyclomatic number of the Hasse diagram of P_{132546} (see Figure 11) is 3. So the degree of its numerator equals 3.*

$$Num(\Psi_{\mathbb{P}_{132546}}) = x_1 \cdot x_2 \cdot x_3 - x_1 \cdot x_2 \cdot x_6 - x_2 \cdot x_3 \cdot x_4 - x_1 \cdot x_3 \cdot x_6 - x_1 \cdot x_4 \cdot x_5 - x_2 \cdot x_3 \cdot x_5 + x_1 \cdot x_4 \cdot x_6 + x_2 \cdot x_3 \cdot x_6 + x_2 \cdot x_4 \cdot x_5 + x_1 \cdot x_5 \cdot x_6 + x_3 \cdot x_4 \cdot x_5 - x_4 \cdot x_5 \cdot x_6$$

$$Den(\Psi_{\mathbb{P}_{132546}}) = (x_1 - x_3) \cdot (x_1 - x_2) \cdot (x_5 - x_6) \cdot (x_4 - x_6) \cdot (x_3 - x_5) \cdot (x_2 - x_5) \cdot (x_3 - x_4) \cdot (x_2 - x_4)$$

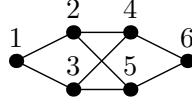


Fig. 11: Numerator and denominator of the reduced fraction $\Psi_{\mathbb{P}_{132546}}$.

From Theorem 3.1 and Theorem 2.13, we know that any rational function of a poset P can be calculated from the bipartite poset \bar{P} by applying a sequence of residues. In fact, due to the construction of \bar{P} , the following proposition gives a stronger result.

Proposition 3.10 *Let P be a poset and \bar{P} its decontraction, then any sequence of limits setting x_a to $x_{\bar{a}}$ for all elements $a \in In(P)$ send $Num(\Psi_{\bar{P}})$ to $Num(\Psi_P)$, i.e.,*

$$Num(\Psi_P) = \lim_{\substack{x_{\bar{a}} \rightarrow x_a \\ a \in In(P)}} Num(\Psi_{\bar{P}})$$

Proof: By construction of \bar{P} (see Theorem 2.13), we can apply to \bar{P} any sequence of contractions of edges (a, \bar{a}) where a is an element of $In(P)$. The poset obtained is exactly P and by Theorem 3.1, we have:

$$\Psi_P = \lim_{\substack{x_{\bar{a}} \rightarrow x_a \\ a \in In(P)}} \prod_{a \in In(P)} (x_a - x_{\bar{a}}) \cdot \Psi_{\bar{P}}.$$

From Corollary 3.7, we have

$$\Psi_P = \lim_{\substack{x_{\bar{a}} \rightarrow x_a \\ a \in In(P)}} \frac{Num(\Psi_{\bar{P}}) \cdot \prod_{a \in In(P)} (x_a - x_{\bar{a}})}{\prod_{(a,b) \in H(\bar{P})} (x_a - x_b)}. \quad (6)$$

By construction of \bar{P} ,

$$(H(\bar{P}) \setminus \{(a, \bar{a}) | a \in In(P)\})|_{\bar{a}=a} = H(\bar{P})|_{\bar{a}=a} = H(P). \quad (7)$$

Consequently,

$$\frac{\prod_{a \in In(P)} (x_a - x_{\bar{a}}) \cdot Num(\Psi_{\bar{P}})}{\prod_{(a,b) \in H(\bar{P})} (x_a - x_b)} = \frac{Num(\Psi_{\bar{P}})}{\prod_{(a,b) \in H(\bar{P}) \setminus \{(a, \bar{b}) | a \in In(P)\}} (x_a - x_b)}. \quad (8)$$

Substituting Equations (7) and (8) in Equation (6), we obtain

$$\Psi_P = \frac{\lim_{\substack{x_{\bar{a}} \rightarrow x_a \\ a \in In(P)}} Num(\Psi_{\bar{P}})}{\prod_{(a,b) \in H(P)} (x_a - x_b)}.$$

Moreover, using Corollary 3.7, we conclude

$$\text{Num}(\Psi_P) = \lim_{\substack{x_a \rightarrow x_a \\ a \in \text{In}(P)}} \text{Num}(\Psi_{\overline{P}}).$$

□

Example 3.11 The rational functions $\Psi_{\overline{\mathbb{P}_{21435}}}$ and $\Psi_{\mathbb{P}_{21435}}$ (see Figure 12) have the following numerators and denominators:

$$\begin{aligned} \text{Num}(\Psi_{\overline{\mathbb{P}_{21435}}}) &= x_5 \cdot x_3 - x_1 \cdot x_4 - x_2 \cdot x_3 - x_2 \cdot x_4 - x_3 \cdot x_3 - x_3 \cdot x_4 - x_4 \cdot x_3 - x_4 \cdot x_4 - x_1 \cdot x_3 + \\ & x_5 \cdot x_4 + x_3^2 + x_4^2 + x_3 \cdot x_4 + x_1 \cdot x_2 + x_1 \cdot x_3 + x_1 \cdot x_4 + x_2 \cdot x_3 - x_1 \cdot x_5 + x_2 \cdot x_4 - x_2 \cdot x_5 \end{aligned}$$

$$\text{Num}(\Psi_{\mathbb{P}_{21435}}) = x_1 \cdot x_2 + x_3 \cdot x_5 + x_4 \cdot x_5 - x_1 \cdot x_5 - x_2 \cdot x_5 - x_3 \cdot x_4$$

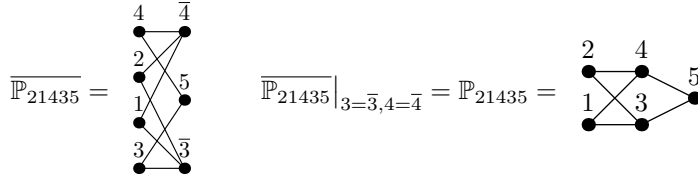


Fig. 12: We obtain the fraction $\Psi_{\mathbb{P}_{21435}}$ from $\Psi_{\overline{\mathbb{P}_{21435}}}$.

3.5 Collapses and factorisations

In general, when a is an extremal element of P , the knowledge of $\Psi_{P \setminus a}$ is not sufficient to compute Ψ_P . However, collapse is a special suppression which allows us to calculate Ψ_P from $\Psi_{P \setminus a}$.

Proposition 3.12 Let a be an element of a connected poset P such that a is a free vertex in the Hasse diagram of P . Let b be the unique vertex such that either $b \leq a$ or $a \leq b$. Then we have

$$\Psi_P = \begin{cases} \Psi_{P \setminus \{a\}} \cdot \frac{1}{x_a - x_b} & \text{if } a \text{ is minimal,} \\ \Psi_{P \setminus \{a\}} \cdot \frac{1}{x_b - x_a} & \text{if } a \text{ is maximal.} \end{cases}$$

Proof: Let a be a free vertex in the Hasse diagram of P . We write the numerator of Ψ_P as a polynomial in x_a with coefficients in $\mathbb{C}[\mathbb{X} \setminus x_a]$:

$$\text{Num}(\Psi_P) = \sum_i C_i x_a^i.$$

Since a is a free vertex, a is either maximal or minimal in P . Theorem 3.2 shows that

$$\lim_{x_a \rightarrow +\infty} (x_a \cdot \Psi_P) = \begin{cases} -\Psi_{P \setminus a} & \text{if } a \text{ is maximal,} \\ \Psi_{P \setminus a} & \text{if } a \text{ is minimal,} \end{cases}$$

which implies that $C_0 \neq 0$ and for all $i \geq 1$, $C_i = 0$. Hence, $\text{Num}(\Psi_P) = C_0 = \text{Num}(\Psi_{P \setminus \{a\}})$. □

As a straightforward consequence, we obtain the following result.

Corollary 3.13 $Num(\Psi_P) = Num(\Psi_{Coll(P)})$, where, as before, $Coll(P)$ denotes the collapse of P .

4 Examples

In general we do not know how to compute Ψ_P . But, for some special cases, we can give a closed formula.

4.1 Acyclic posets

Proposition 4.1 *The Hasse diagram $H(P)$ has no cycle if and only if $Num(P) = 1$.*

Proof: This result is a direct consequence of Greene's theorem (Theorem 3.5) and Corollary 3.8. \square

A permutation σ avoids the patterns 1324 if there exist no integers $1 \leq i_1 < i_2 < i_3 < i_4 \leq n$ such that $\sigma_{i_1} < \sigma_{i_3} < \sigma_{i_2} < \sigma_{i_4}$. A permutation σ contains the pattern $\overline{2143}$ if for some indices $1 \leq i_1 < i_2 < i_3 < i_4 \leq n$ we have $\sigma_{i_2} \leq \sigma_{i_1} \leq \sigma_{i_4} \leq \sigma_{i_3}$ with the further restriction that there is no $i_1 \leq j \leq i_4$ such that $\sigma_{i_1} \leq \sigma_j \leq \sigma_{i_4}$.

Butler and Bousquet-Mélou have shown in (1) that the Hasse diagram of a poset associated to a permutation avoiding 1324 and $\overline{2143}$ has no cycle. As a consequence we have the following corollary.

Corollary 4.2 *We have $Num(\Psi_{\mathbb{P}_\sigma}) = 1$ if and only if σ avoids the patterns 1324 and $\overline{2143}$.*

4.2 1-cycle posets

Proposition 4.3 *Let P be a connected 1-cycle poset, then*

$$Num(P) = \sum_{a \in \min(Coll(P))} x_a - \sum_{a \in \max(Coll(P))} x_a.$$

Proof: Consider the poset $P' = \overline{Coll(P)}$ obtained by the construction given in Theorem 2.13 applied to $Coll(P)$.

Since P' is bipartite with only 1 cycle, by Corollary 3.8 we obtain $Num(\Psi_{P'}) = \sum_a \beta_a x_a$, where $\beta_a \in \mathbb{Z}$. Let a be a minimal element in P' . Since $P' \setminus \{a\}$ is acyclic and connected, $Num(\Psi_{P' \setminus \{a\}}) = 1$. Theorem 3.2 implies that $\beta_a = 1$ if a is maximal and $\beta_y = -1$ if a is minimal. So, we have:

$$Num(\Psi_{P'}) = \sum_{a \in \min(P')} x_a - \sum_{a \in \max(P')} x_a.$$

By Proposition 3.10, we have:

$$Num(\Psi_P) = \lim_{\substack{x_{\bar{a}} \rightarrow x_a \\ a \in In(P')}} Num(\Psi_{P'}).$$

Hence,

$$\begin{aligned} Num(\Psi_P) &= \sum_{a \in \min(Call(P))} x_a - \sum_{a \in \max(Call(P))} x_a + \sum_{a \in In(Call(P))} \lim_{\substack{x_{\bar{a}} \rightarrow x_a \\ a \in In(Call(P))}} (x_a - x_{\bar{a}}) \\ &= \sum_{a \in \max(Coll(P))} x_a - \sum_{a \in \max(Coll(P))} x_a. \end{aligned}$$

\square

Example 4.4 The numerator of the 1-cycle poset P in Figure 13 is

$$\text{Num}(\Psi_P) = x_1 + x_2 - x_4 - x_7.$$

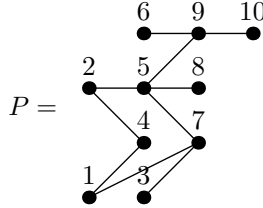


Fig. 13: 1-cycle poset.

4.3 λ -complete posets and Schubert polynomials

We have seen in Subsection 3.1 that the bipartite posets are fundamental for the description of the functions Ψ_P . In this section, we compute Ψ_P for special cases of bipartite posets.

Let λ be a partition (a weakly decreasing sequence of non-negative integers). We call a bipartite poset with elements $\{a_1, \dots, a_{l(\lambda)}, b_1, \dots, b_{\lambda_1}\}$ and relation \leq_λ a λ -complete poset, if the relation \leq_λ is defined by $a_i \leq_\lambda b_j$ if and only if $j \leq \lambda_i$. We will write P^λ for this poset.

For example, Figure 14 shows a 54331-complete poset.

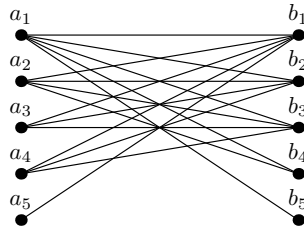


Fig. 14: The 54331-complete poset.

Let $f(a_1, \dots, a_n)$ be a polynomial in the alphabet $A = \{a_1, \dots, a_n\}$. The *divided differences* ∂_i (see (4)) is the operator defined by

$$f \cdot \partial_i = \frac{f - f^{s_i}}{x_{a_i} - x_{a_{i+1}}},$$

where f^{s_i} is the function $f(a_1, \dots, a_{i+1}, a_i, \dots, a_n)$.

We recall that the definition of Schubert polynomials as it can be found in (4).

Schubert polynomials $\mathbb{Y}_v(A, B)$ are functions in the alphabets $A = \{a_1, a_2, \dots\}$ and $B = \{b_1, b_2, \dots\}$ indexed by vectors $v \in \mathbb{N}^\infty$ having a finite number of non-zero components and defined recursively by:

1. if v is weakly decreasing then,

$$\mathbb{Y}_v(\mathbb{A}, \mathbb{B}) := \prod_{i=1}^{l(v)} \prod_{j=1}^{v_i} (x_{a_i} - x_{b_j});$$

2. if there exists $i \in \mathbb{N}$ such that $v_i < v_{i+1}$ then,

$$\mathbb{Y}_v \cdot \partial_i := \mathbb{Y}_{v \cdot \partial_i},$$

where $v \cdot \partial_i = (v_1, \dots, v_{i-1}, v_{i+1}, v_i - 1, v_{i+2}, \dots)$.

Proposition 4.5 *Let $\lambda = \lambda_1 \dots \lambda_n$ be a partition with $\lambda_n > 0$. For any λ -complete poset \mathcal{P}^λ , the numerator of $\Psi_{\mathcal{P}^\lambda}$ is the Schubert polynomial*

$$\text{Num}(\Psi_{\mathcal{P}^\lambda}) = \mathbb{Y}_{0, \lambda_2 - 1, \dots, \lambda_n - 1}.$$

Proof: Let A and B be two alphabets, and let X be a subalphabet of A . Let λ be a partition and k an integer with $\lambda_k > 1$ and $\lambda_{k+1} \leq 1$. Let $(A^i)_{i \in [1, k-1]}$ be the family of alphabets defined by $A^i = \{a_i, \dots, a_k\}$. Let $(\beta^i)_{i \in [1, k-1]}$ be the family of partitions defined by $\beta^i = [\lambda_{i+1}, \lambda_{i+1}, \dots, \lambda_k]$. We will

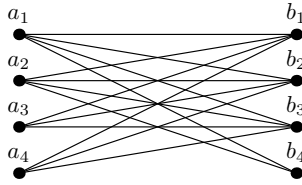


Fig. 15: Poset \mathcal{P}^{β^1} for $\lambda = 54331$ (which implies $\beta^1 = 4433$).

denote by $\mathcal{P}^\lambda(X)$ the λ -complete poset on X and B :

$$\mathcal{P}^\lambda(X) := \mathcal{P}^\lambda(X, B).$$

We will prove the following recursive formula:

$$\text{Num}(\Psi_{\mathcal{P}^{\beta^i}(A^i)}) = (\mathbb{Y}_{\lambda_{i+1}}(a_i) \cdot \text{Num}(\Psi_{\mathcal{P}^{\beta^{i+1}}}(A^{i+1}))) \partial_i \quad (9)$$

We denote by P' the expression $\mathcal{P}^{\beta^i}(A^i)$. By Lemma 2.5, we sort a_i and a_{i+1} to obtain

$$\Psi_{P'} = \Psi_{P'_{a_i a_{i+1}}} + \Psi_{P'_{a_{i+1} a_i}}.$$

The Hasse diagram of $P'_{a_i a_{i+1}}$ is obtained from the Hasse diagram of P' by removing all the edges having a_i as vertex and adding the edge (a_i, a_{i+1}) . In the same way, the Hasse diagram of $P'_{a_{i+1} a_i}$ can be obtained by applying the transposition s_i on $P'_{a_i, a_{i+1}}$.

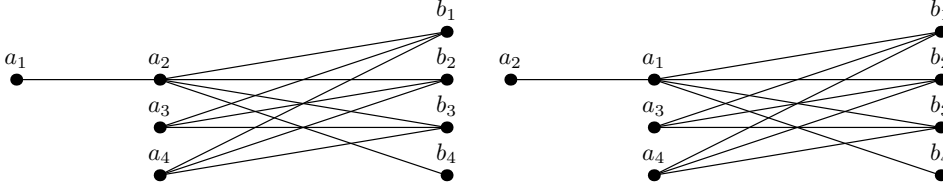


Fig. 16: a_1 and a_2 are sorted in β^1 -complete poset, where $\lambda = 54331$ (which implies $\beta^1 = 4433$).

The vertex a_i (respectively a_{i+1}) is free in $P'_{a_i a_{i+1}}$ (respectively $P'_{a_{i+1} a_i}$) and can be collapsed. By Proposition 3.12 and Corollary 3.7 we obtain

$$\Psi_{P'} = \frac{1}{x_{a_i} - x_{a_{i+1}}} \cdot (\Psi_{\mathcal{P}^{[\beta_2^i \dots \beta_k^i]}(A_i \setminus \{a_i\})} - \Psi_{\mathcal{P}^{[\beta_2^i \dots \beta_k^i]}(A_i \setminus \{a_{i+1}\})}). \quad (10)$$

Observing

$$\mathbb{Y}_{\lambda_{i+1}}(x) = \prod_{j=1}^{\lambda_{i+1}} x - x_{b_j},$$

and using Corollary 3.7 in Equation 10, we obtain:

$$\Psi_{P'} = \frac{1}{x_{a_i} - x_{a_{i+1}}} \cdot \frac{\mathbb{Y}_{\lambda_{i+1}}(a_i) \cdot \text{Num}(\Psi_{\mathcal{P}^{[\beta_2^i \dots \beta_k^i]}(A \setminus \{a_i\})}) - \mathbb{Y}_{\lambda_{i+1}}(a_{i+1}) \cdot \text{Num}(\Psi_{\mathcal{P}^{[\beta_2^i \dots \beta_k^i]}(A \setminus \{a_{i+1}\})})}{\prod_{c \leq d} (x_c - x_d)}.$$

Hence,

$$\text{Num}(\Psi_{\mathcal{P}^{\beta^i}(A^i)}) = \left(\mathbb{Y}_{\lambda_{i+1}}(a_i) \cdot \text{Num}(\Psi_{\mathcal{P}^{[\beta_2^i \dots \beta_k^i]}(A^{i+1})}) \right) \partial_i. \quad (11)$$

If we collapse the poset $\mathcal{P}^{[\beta_2^i \dots \beta_k^i]}(A^{i+1})$, we obtain the poset $\mathcal{P}^{\beta^{i+1}}(A^{i+1})$ and we can conclude that

$$\text{Num}(\Psi_{\mathcal{P}^{\beta^i}(A^i)}) = \left(\mathbb{Y}_{\lambda_{i+1}}(a_i) \cdot \text{Num}(\Psi_{\mathcal{P}^{\beta^{i+1}}(A^{i+1})}) \right) \partial_i. \quad (12)$$

From Equation (12) we obtain

$$\text{Num}(\Psi_{\mathcal{P}^{\beta^1}}) = \left(\mathbb{Y}_{\lambda_2}(a_1) \cdots \mathbb{Y}_{\lambda_{k-1}}(a_{k-2}) \text{Num}(\Psi_{\mathcal{P}^{\lambda_k \lambda_k}(\{a_{k-1}, a_k\})}) \right) \partial_{k-2} \cdots \partial_1.$$

Using Equation (11), we deduce that

$$\text{Num}(\Psi_{\mathcal{P}^{\lambda_k \lambda_k}(\{a_{k-1}, a_k\})}) = \mathbb{Y}_{\lambda_k}(a_{k-1}) \text{Num}(\Psi_{\mathcal{P}^{\lambda_k}(\{a_k\})}).$$

Since $\mathcal{P}^{\lambda_k}(\{a_k\})$ has no cycle, $\text{Num}(\Psi_{\mathcal{P}^{\lambda_k}(\{a_k\})}) = 1$ and

$$\text{Num}(\Psi_{\mathcal{P}^{\beta^1}}) = (\mathbb{Y}_{\lambda_2}(a_1) \cdots \mathbb{Y}_{\lambda_k}(a_{k-1})) \partial_{k-1} \cdots \partial_1.$$

Finally, we obtain the result

$$\text{Num}(\Psi_{\mathcal{P}^{\beta^1}}) = \mathbb{Y}_{0, \lambda_2 - 1, \dots, \lambda_k - 1}.$$

It is easy to see that $\mathcal{P}^\lambda = \text{Coll}(\mathcal{P}^{\beta^1})$. We deduce that,

$$\text{Num}(\Psi_{\mathcal{P}^\lambda}) = \mathbb{Y}_{0, \lambda_2-1, \dots, \lambda_k-1} = \mathbb{Y}_{0, \lambda_2-1, \dots, \lambda_n-1}.$$

□

Acknowledgements

The author is grateful to A. Lascoux for his suggestion to work on rational functions and his introduction to Schur polynomials. The author thanks J.G. Luque for useful discussions and revisions of this paper. The author thanks T. Gomez-Diaz for critical reading of this article.

References

- [1] S. Butler and M. Bousquet-Mélou, *Forest-like permutations*, Ann. Comb. **11** (2007), 335–354.
- [2] A. Cortez, *Singularités génériques et quasi-résolutions des variétés de Schubert pour le groupe linéaire*, Adv. Math. **178** (2003), 396–445.
- [3] C. Greene, *A rational function identity related to the Murnaghan–Nakayama formula for the characters of \mathcal{S}_n* , J. Alg. Comb. **1** 3 (1992), 235–255.
- [4] A. Lascoux, *Symmetric Functions and Combinatorial Operators on Polynomials*, Conference Board of the Mathematical Sciences Series, Amer. Math. Soc. **99** (2003).
- [5] A.M. Fu and A. Lascoux, *q -Identities from Lagrange and Newton interpolation*, Adv. App. Math. **31** (2003), 527–531.
- [6] A. Woo and A. Yong, *When is a Schubert variety Gorenstein?*, Adv. Math. **207** (2006), 205–220.
- [7] I. G. Macdonald, *Symmetric Functions and Hall Polynomials*, Oxford University Press, 2nd edition (1998).
- [8] R. Diestel, *Graph Theory*, Graduate Texts in Mathematics **173** (2005), Springer–Verlag Heidelberg.
- [9] P. J. Giblin, *Graphs, Surface and Homology — An Introduction to Algebraic Topology* (1981), London and New York Chapman and Hall.