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# On the Minimum Number of Completely 3-Scrambling Permutations

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A family  $\mathcal{P} = \{\pi_1, \dots, \pi_q\}$  of permutations of  $[n] = \{1, \dots, n\}$  is *completely  $k$ -scrambling* [Spencer, 1972; Füredi, 1996] if for any distinct  $k$  points  $x_1, \dots, x_k \in [n]$ , permutations  $\pi_i$ 's in  $\mathcal{P}$  produce all  $k!$  possible orders on  $\pi_i(x_1), \dots, \pi_i(x_k)$ . Let  $N^*(n, k)$  be the minimum size of such a family. This paper focuses on the case  $k = 3$ . By a simple explicit construction, we show the following upper bound, which we express together with the lower bound due to Füredi for comparison.

$$\frac{2}{\log_2 e} \log_2 n \leq N^*(n, 3) \leq 2 \log_2 n + (1 + o(1)) \log_2 \log_2 n.$$

We also prove the existence of  $\lim_{n \rightarrow \infty} N^*(n, 3)/\log_2 n = c_3$ . Determining the value  $c_3$  and proving the existence of  $\lim_{n \rightarrow \infty} N^*(n, k)/\log_2 n = c_k$  for  $k \geq 4$  remain open.

## 1 Introduction and Summary

Following Spencer [Sp72] and Füredi [Fü96], call a family  $\mathcal{P} = \{\pi_1, \dots, \pi_q\}$  of permutations of  $[n]$  *completely  $k$ -scrambling* if for any distinct  $x_1, x_2, \dots, x_k \in [n]$ , there exists a permutation  $\pi_i \in \mathcal{P}$  such that  $\pi_i(x_1) < \pi_i(x_2) < \dots < \pi_i(x_k)$ ; or equivalently,  $\pi_i$ 's applied to  $x_1, x_2, \dots, x_k$  produce all  $k!$  orders. This paper focuses on the case  $k = 3$ . Following Füredi [Fü96], say that a family  $\mathcal{P}$  is *3-mixing* if for any distinct  $x, y, z \in [n]$ , there is a permutation  $\pi_i \in \mathcal{P}$  that places  $x$  between  $y$  and  $z$ , i.e., there is a permutation  $\pi_i$  such that either  $\pi_i(y) < \pi_i(x) < \pi_i(z)$  or  $\pi_i(z) < \pi_i(x) < \pi_i(y)$ .

Let  $N^*(n, k)$  be the minimum  $q$  such that completely  $k$ -scrambling  $q$  permutations exist for  $[n]$ . The best known bounds for  $N^*(n, k)$  can be expressed as follows. For arbitrary fixed  $k \geq 3$ , as  $n \rightarrow \infty$ ,

$$\left( \frac{1}{\log_2 e} (k-1)! + o(1) \right) \log_2 n \leq N^*(n, k) \leq \frac{k}{\log_2(k!/(k!-1))} \log_2 n. \quad (1)$$

The coefficient of the upper bound in (1) is  $\Theta(k \cdot k!)$ ; thus the gap between the coefficients of the lower and upper bounds in (1) is  $\Theta(k^2)$ . The upper bound in (1) was shown by Spencer [Sp72] by a probabilistic argument, where one considers the probability that some order among some  $x_1, \dots, x_k$  is never produced by  $q$  independent random permutations. The lower bound in (1) was first proved by Füredi [Fü96] for  $k = 3$ , and was proved for  $k \geq 3$  by Radhakrishnan [Ra03]; entropy arguments are used in both work;

the factor  $\log_2 e$  in the lower bound comes from the fact that  $\int_0^1 H(x)dx = (\log_2 e)/2$ , where  $H(x)$  is the binary entropy function.

As for the case  $k = 3$ , Füredi [Fü96] has shown that

$$\frac{2}{\log_2 e} \log_2 n \leq N^*(n, 3) \leq \left( \frac{10}{\log_2 7} \right) \log_2 n + O(1), \quad (2)$$

where the coefficients of  $\log_2 n$  are  $1.38\dots$  and  $3.56\dots$  in (2). The lower bound in (2) is in fact a lower bound for the case where we only require a family to be 3-mixing. No better lower bound for completely 3-scrambling families is known. If a family  $\mathcal{P} = \{\pi_1, \dots, \pi_q\}$  is 3-mixing, by adding to  $\mathcal{P}$  the  $q$  reverse permutations of  $\pi_i$ 's mapping  $x \mapsto n+1-\pi_i(x)$ , we can obtain completely 3-scrambling  $2q$  permutations. Ishigami [Is95] has given an efficient recursive construction of 3-mixing families starting with a 3-mixing family of five permutations of  $\{1, \dots, 7\}$ . Füredi [Fü96] gave the upper bound in (2) by making these observations and doubling the size of Ishigami's 3-mixing family.

In this paper, we first give an improved upper bound for  $N^*(n, 3)$  by a simple construction. Let  $f(q)$  be the maximum  $n$  such that completely 3-scrambling  $q$  permutations exist for  $[n]$ .

### Theorem 1

$$f(q) \geq \binom{\lfloor q/2 \rfloor}{\lfloor q/4 \rfloor}.$$

The following upper bound on  $N^*(n, 3)$  readily follows.

### Corollary 1

$$N^*(n, 3) \leq 2 \log_2 n + (1 + o(1)) \log_2 \log_2 n.$$

It seems natural to conjecture that for every fixed  $k \geq 3$ , as  $n \rightarrow \infty$ ,  $N^*(n, k) = (c_k + o(1)) \log_2 n$  for some  $c_k$ . We show the existence of limit for the case  $k = 3$ :

### Theorem 2

$$\lim_{q \rightarrow \infty} \frac{\log_2 f(q)}{q} = C \text{ exists.}$$

The following immediately follows.

### Corollary 2

$$\lim_{n \rightarrow \infty} \frac{N^*(n, 3)}{\log_2 n} = 1/C = c_3 \text{ exists.}$$

## 2 Proofs

We can identify in a natural way a total order  $\phi$  on  $[n]$  and the permutation of  $[n]$  induced by  $\phi$ ; thus we speak interchangeably in terms of permutations and total orders. In fact for an arbitrary finite set  $U$  with  $n$  elements, we can assume for our purposes that  $U$  is identified with  $[n]$  in an arbitrary fixed way, and speak about permutations of  $U$  in terms of total orders on  $U$ .

**Proof of Theorem 1.** Put  $r = \lfloor q/2 \rfloor$  and let  $\mathcal{F} = \{A_1, A_2, \dots, A_m\}$  be a family of subsets of  $\{1, \dots, r\}$  such that  $A_i \not\subseteq A_j$  for all  $i \neq j$ ; i.e.,  $\mathcal{F}$  is an antichain.

For each point  $x \in \{1, \dots, r\}$ , define two orders  $\phi_x$  and  $\psi_x$  on  $\mathcal{F}$ . In both orders  $\phi_x$  and  $\psi_x$ , the sets  $A_i$  containing the point  $x$  are smaller than all the sets  $A_k$  not containing  $x$ . Among the sets containing  $x$  and among the sets not containing  $x$ : in the order  $\phi_x$ ,  $A_i < A_j$  precisely when  $i < j$ ; in the order  $\psi_x$ , this is reversed, and  $A_i < A_j$  precisely when  $i > j$ .

We claim that for arbitrary distinct  $i, j, k \in [m]$ , there exists an order  $\theta \in \{\phi_1, \psi_1, \phi_2, \psi_2, \dots, \phi_r, \psi_r\}$  such that  $A_i < A_j < A_k$  in the order  $\theta$ . To see the claim fix a point  $x \in (A_i - A_k) \neq \emptyset$ , i.e.,  $x \in A_i$  and  $x \notin A_k$ . Depending on whether  $x \in A_j$  or  $x \notin A_j$ , we specify an order  $\theta$  that produces the ordering  $A_i < A_j < A_k$ .

Case  $x \in A_j$ : Let  $\theta = \phi_x$  if  $i < j$  and let  $\theta = \psi_x$  if  $i > j$ .

Case  $x \notin A_j$ : Let  $\theta = \phi_x$  if  $j < k$  and let  $\theta = \psi_x$  if  $j > k$ .

Clearly under the order  $\theta$ ,  $A_i < A_j < A_k$ . Hence the  $2r$  orders thus defined on  $[m]$  are completely 3-scrambling. We obtain the theorem by taking  $\mathcal{F}$  to be the family of all subsets of  $\{1, \dots, r\}$  with cardinality  $\lfloor r/2 \rfloor = \lfloor q/4 \rfloor$ .  $\square$

**Proof of Theorem 2.** Our proof of Theorem 2 will be basically similar to Füredi's proof [Fü96] of the existence of  $\lim_{q \rightarrow \infty} (\log_2 g(q)) / q$ , where  $g(q)$  is the maximum  $n$  such that 3-mixing  $q$  permutations exist for  $[n]$ . To make a recursive construction go through for scrambling permutations, we introduce and use red-blue colored doubly reversing permutations: Call a family  $\mathcal{P} = \{\pi_1, \dots, \pi_q\}$  of permutations of  $[n]$  2-reversing if there is a coloring  $\chi : \{\pi_1, \dots, \pi_q\} \rightarrow \{\text{red}, \text{blue}\}$  such that for every distinct  $i, j \in [n]$ , there are red  $\pi_\kappa$ , red  $\pi_\lambda$ , blue  $\pi_\mu$ , and blue  $\pi_\nu$  satisfying

$$\pi_\kappa(i) < \pi_\kappa(j), \pi_\lambda(i) > \pi_\lambda(j); \pi_\mu(i) < \pi_\mu(j), \pi_\nu(i) > \pi_\nu(j).$$

For a permutation  $\pi$  of  $[n]$ , let  $\text{reverse}(\pi)$  be the permutation of  $[n]$  mapping  $x \mapsto n + 1 - \pi(x)$ . Let  $\mathcal{P}$  be a family of permutations of  $[n]$  with  $|\mathcal{P}| \geq 3$ . We can easily transform  $\mathcal{P}$  to a 2-reversing family by adding at most two permutations as follows. Arbitrarily fix two distinct permutations  $\sigma, \tau \in \mathcal{P}$  such that  $\tau \neq \text{reverse}(\sigma)$ ; such  $\sigma$  and  $\tau$  exist since  $|\mathcal{P}| \geq 3$ ; add  $\text{reverse}(\sigma)$  and  $\text{reverse}(\tau)$  to  $\mathcal{P}$ ; color  $\sigma$  and  $\text{reverse}(\sigma)$  red; color  $\tau$  and  $\text{reverse}(\tau)$  blue; color the remaining permutations arbitrarily.

Let  $f^*(q)$  be the maximum  $n$  such that completely 3-scrambling and 2-reversing  $q$  permutations exist for  $[n]$ . By definition and from the discussion above we have

$$f^*(q) \leq f(q) \leq f^*(q + 2). \tag{3}$$

### Claim 1

$$f^*(q + r) \geq f^*(q)f^*(r).$$

For the moment we assume that Claim 1 holds and go on to derive Theorem 2.

The sequence  $(1/q) \log_2 f^*(q)$  is bounded above. From this and Claim 1 it follows by classical calculus (Fekete's theorem) that

$$\lim_{q \rightarrow \infty} \frac{1}{q} \log_2 f^*(q) = \limsup_{q \rightarrow \infty} \frac{1}{q} \log_2 f^*(q).$$

From (3) it now follows that

$$\lim_{q \rightarrow \infty} \frac{1}{q} \log_2 f(q) = \lim_{q \rightarrow \infty} \frac{1}{q} \log_2 f^*(q).$$

Thus we are left to prove Claim 1.

Let  $\mathcal{S} = \{\sigma_1, \dots, \sigma_q\}$  and  $\mathcal{T} = \{\tau_1, \dots, \tau_r\}$  be completely 3-scrambling and 2-reversing families of permutations of  $[l]$  and  $[m]$  respectively. Assume that both families are validly red-blue colored. Let  $U = \{(i, j) : 1 \leq i \leq l, 1 \leq j \leq m\}$ ; think of  $U$  as a matrix with  $l$  rows and  $m$  columns. We will show that we can define  $q + r$  orders on  $U$  that are completely 3-scrambling and 2-reversing. Note that from this Claim 1 follows.

Let  $x = (i, j)$  and  $y = (i', j')$  be distinct elements of  $U$ . For  $k = 1, \dots, q$ , define the order  $\tilde{\sigma}_k$  using  $\sigma_k$  in a row-major form as follows: if  $i \neq i'$ , order  $x$  and  $y$  according to the order of  $\sigma_k(i)$  and  $\sigma_k(i')$ . When  $i = i'$ : if  $\sigma_k$  is red,  $(i, j) < (i, j') \iff j < j'$ ; if  $\sigma_k$  is blue,  $(i, j) < (i, j') \iff j > j'$ . Similarly for  $k = 1, \dots, r$ , define the order  $\tilde{\tau}_k$  on  $U$  in a column-major form: when  $j \neq j'$ ,  $x < y \iff \tau_k(j) < \tau_k(j')$ ; when  $j = j'$ : if  $\tau_k$  is red,  $(i, j) < (i', j) \iff i < i'$ ; if  $\tau_k$  is blue,  $(i, j) < (i', j) \iff i > i'$ . As for colors, let  $\tilde{\sigma}_k$  and  $\tilde{\tau}_k$  inherit the colors of  $\sigma_k$  and  $\tau_k$ .

**Claim 2** The family  $\mathcal{F} = \{\tilde{\sigma}_1, \dots, \tilde{\sigma}_q, \tilde{\tau}_1, \dots, \tilde{\tau}_r\}$  is completely 3-scrambling and 2-reversing.

To see Claim 2, let  $x_1 = (i_1, j_1), x_2 = (i_2, j_2), x_3 = (i_3, j_3)$  be distinct elements of  $U$ . If  $i_1, i_2, i_3$  are all distinct,  $\sigma_k$ 's produce all six orderings of  $i_1, i_2, i_3$ , and hence  $\tilde{\sigma}_k$ 's produce all six orderings of  $x_1, x_2, x_3$ . Similar arguments with  $\tau_k$ 's and  $\tilde{\tau}_k$ 's apply for the case when  $j_1, j_2, j_3$  are all distinct.

The remaining case is when  $|\{i_1, i_2, i_3\}| = |\{j_1, j_2, j_3\}| = 2$ . We write, e.g., 231 to express the ordering  $x_2 < x_3 < x_1$ . Assume that

$$x_1 = (i, j), x_2 = (i, j'), x_3 = (i', j), \quad i \neq i', j \neq j'.$$

We will see that all six orderings of  $x_1, x_2, x_3$  are produced by checking that (1) all the four orders in which  $x_3$  is smallest or largest, i.e., 312, 321, 123, 213 are produced and that (2) all the four orders in which  $x_2$  is smallest or largest are produced.

A red  $\tilde{\sigma}_\kappa$  and a blue  $\tilde{\sigma}_\mu$  satisfying  $\sigma_\kappa(i) < \sigma_\kappa(i')$  and  $\sigma_\mu(i) < \sigma_\mu(i')$  produce 123 and 213 respectively. Similarly, a red  $\tilde{\sigma}_\lambda$  and a blue  $\tilde{\sigma}_\nu$  satisfying  $\sigma_\lambda(i) > \sigma_\lambda(i')$  and  $\sigma_\nu(i) > \sigma_\nu(i')$  produce 312 and 321 respectively. Thus all the four orders in which  $x_3$  is smallest or largest are produced. Similarly, two red  $\tilde{\tau}$ 's and two blue  $\tilde{\tau}$ 's ordering  $j$  and  $j'$  in both directions produce the four orders in which  $x_2$  is smallest or largest.

Finally, if  $x = (i, j)$  and  $y = (i', j')$  are distinct points in  $U$ , either (i)  $i \neq i'$  or (ii)  $j \neq j'$ . The 2-reversing condition is satisfied by  $\tilde{\sigma}_k$ 's in case (i) and by  $\tilde{\tau}_k$ 's in case (ii).  $\square$

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