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► **To cite this version:**

Saúl A. Blanco. Shortest path poset of finite Coxeter groups. Krattenthaler, Christian and Strehl, Volker and Kauers, Manuel. 21st International Conference on Formal Power Series and Algebraic Combinatorics (FPSAC 2009), 2009, Hagenberg, Austria. Discrete Mathematics and Theoretical Computer Science, DMTCS Proceedings vol. AK, 21st International Conference on Formal Power Series and Algebraic Combinatorics (FPSAC 2009), pp.189-200, 2009, DMTCS Proceedings. <hal-01185413>

HAL Id: hal-01185413

<https://hal.inria.fr/hal-01185413>

Submitted on 20 Aug 2015

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Shortest path poset of finite Coxeter groups

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Abstract. We define a poset using the shortest paths in the Bruhat graph of a finite Coxeter group W from the identity to the longest word in W , w_0 . We show that this poset is the union of Boolean posets of rank absolute length of w_0 ; that is, any shortest path labeled by reflections t_1, \dots, t_m is fully commutative. This allows us to give a combinatorial interpretation to the lowest-degree terms in the complete **cd**-index of W .

Résumé. Nous définissons un poset en utilisant le plus court chemin entre l'identité et le plus long mot de W , w_0 , dans le graph de Bruhat du groupe finie Coxeter, W . Nous prouvons que ce poset est l'union de posets Boolean du même rang que la longueur absolue de w_0 ; ça signifie que tous les plus courts chemins, étiquetés par réflexions t_1, \dots, t_m sont totalement commutatives. Ça nous permet de donner une interprétation combinatoire aux termes avec le moindre grade dans le **cd**-index complet de W .

Keywords: Coxeter group, Bruhat order, Boolean poset, complete **cd**-index.

1 Introduction

Let (W, S) be a Coxeter system, and let $T(W) = \{ws w^{-1} : s \in S, w \in W\}$ be the set of *reflections* of (W, S) . The *Bruhat graph* of (W, S) , denoted by $B(W, S)$ or simply $B(W)$, is the directed graph with vertex set W , and a directed edge $w_1 \rightarrow w_2$ between $w_1, w_2 \in W$ if $\ell(w_1) < \ell(w_2)$ and there exists $t \in T(W)$ with $tw_1 = w_2$. ℓ denotes the *length function* of (W, S) . The edges of $B(W)$ are labeled by reflections; for instance the edge $w_1 \rightarrow w_2$ is labeled with t . The Bruhat graph of an interval $[u, v]$, denoted by $B([u, v])$, is the subgraph of $B(W)$ obtained by only considering the elements of $[u, v]$. A *path* in the Bruhat graph $B([u, v])$, will always mean a *directed* path from u to v . As it is the custom, we will label these paths by listing the edges that are used.

A *reflection ordering* $<_{T(W)} = <_T$ is a total order of $T(W)$ so that $r <_T rtr <_T rtrtr <_T \dots <_T trt <_T t$ or $t <_T trt <_T trtrt <_T \dots <_T rtr <_T r$ for each subgroup $W' = \langle t, r \rangle$ where $t, r \in T(W)$. Let $\Delta = (t_1, t_2, \dots, t_k)$ be a path in $B([u, v])$, and define the *descent set* of Δ by $D(\Delta) = \{j : t_{j+1} <_T t_j\} \subset [k-1]$.

Let $w(\Delta) = x_1 x_2 \cdots x_{k-1}$, where $x_i = \mathbf{a}$ if $t_i <_{i+1}$, and $x_i = \mathbf{b}$, otherwise. In other words, set x_i to \mathbf{a} if $i \notin D(\Delta)$ and to \mathbf{b} if $i \in D(\Delta)$. In [3], Billera and Brenti showed that $\sum_{\Delta \in B([u, v])} w(\Delta)$ becomes

[†]The author was supported in part by NSF grant DMS-0555268.

a polynomial in the variables \mathbf{c} and \mathbf{d} , where $\mathbf{c} = \mathbf{a} + \mathbf{b}$ and $\mathbf{d} = \mathbf{ab} + \mathbf{ba}$. This polynomial is called the *complete \mathbf{cd} -index* of $[u, v]$, and it is denoted by $\tilde{\psi}_{u,v}(\mathbf{c}, \mathbf{d})$. Notice that the complete \mathbf{cd} -index of $[u, v]$ is an encoding of the distribution of the descent sets of each path Δ in the Bruhat graph of $[u, v]$, and thus seems to depend on \prec_T . However, it can be shown that this is not the case. For details on the complete \mathbf{cd} -index, see [3].

As an example, consider A_2 , the symmetric group on 3 elements with generators $s_1 = (1\ 2)$ and $s_2 = (2\ 3)$. Then $t_1 = s_1 < t_2 = s_1s_2s_1 < t_3 = s_2$ is a reflection ordering. The paths of length 3 are: (t_1, t_2, t_3) , (t_1, t_3, t_1) , (t_3, t_1, t_3) , and (t_3, t_2, t_1) , that encode to $\mathbf{a}^2 + \mathbf{ab} + \mathbf{ba} + \mathbf{b}^2 = \mathbf{c}^2$. There is one path of length 1, namely t_2 , which encodes simply to 1. So $\tilde{\psi}_{u,v}(\mathbf{c}, \mathbf{d}) = \mathbf{c}^2 + 1$.

Consider the paths in $B([u, v])$ of minimum length. Using these paths, we can define a ranked poset by thinking of the edges of these paths as cover relations. We call this poset $SP([u, v])$, and when the interval $[u, v]$ is the full group, we simply use the notation $SP(W)$. The rank of an element x in $SP([u, v])$ is given by its distance from u (and so if $[u, v]$ is the whole group, the rank of x is given by its distance from the identity e). Here we are interested in $SP(W)$, where W is a finite Coxeter group. To illustrate the definition consider B_2 and $SP(B_2)$ as depicted below. The rank of $SP(B_2)$ is two since that is the length of the shortest paths in $B(B_2)$.

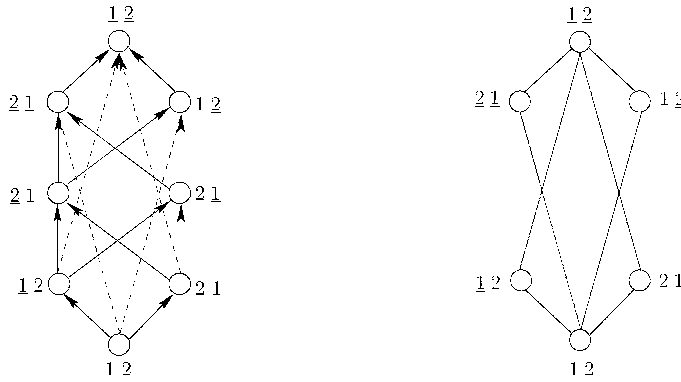


Fig. 1: $B(B_2)$ and $SP(B_2)$.

For any finite Coxeter group, there is a word w_0^W of maximal length. It is a well known fact that $\ell(w_0^W) = |T(W)|$. For any $w \in W$, one can write $t_1t_2 \cdots t_n = w$ for some $t_1, t_2, \dots, t_n \in T(W)$. If n is minimal, then we say that w is $T(W)$ -reduced, and that the *absolute length* of w is n . We write $\ell_{T(W)}(w) = n$, or simply $\ell_T(w) = n$.

Notice that for $w \in W$, if $\ell_T(w) = \ell$, then $t_1t_2 \cdots t_\ell = w$ for some reflections in $T(W)$, but this *does not* mean that $(t_1, t_2, \dots, t_\ell)$ is a (directed) path in $B([e, w])$. Nevertheless, we will show that for finite W and $w = w_0^W$, $(t_1, t_2, \dots, t_\ell)$ and any of its permutations $(t_{\tau(1)}, t_{\tau(2)}, \dots, t_{\tau(\ell)})$, $\tau \in A_{\ell-1}$, is a path in $B(W)$. To be more specific, we show the following theorem.

Theorem 1.1 *Let W be a finite Coxeter group and $\ell_0 = \ell_{T(W)}(w_0^W)$, the absolute length of the longest element of W . Then $SP(W)$ is isomorphic to the union of Boolean posets of rank ℓ_0 .*

In Section 2 we present the proof of the theorem for the infinite families (groups of type A , B , and D and Dihedral groups). In Section 3 we discussed the validity of the Theorem for the exceptional groups. Computer search was used for F_4 , H_3 , H_4 , and E_6 , and a geometric argument was used to prove the case E_7 and E_8 . We summarize the number of Boolean posets that form $SP(W)$ and the rank of $SP(W)$ for each finite Coxeter group in Table 1.

In Section 4 we discuss why Theorem 1.1 implies that the lowest-degree terms of the complete **cd**-index of W is given by $\alpha_W \tilde{\psi}(\mathcal{B}_{\ell_0})$, where $\tilde{\psi}(\mathcal{B}_{\ell_0})$ is the **cd**-index of the Boolean poset of rank $\ell_0 = \ell_{T(W)}(w_0^W)$, and α_W is the number of Boolean posets that form $SP(W)$.

The following lemma will be used in our proofs.

Lemma 1.2 (Shifting Lemma, [1], Lemma 2.5.1) *If $w = t_1 t_2 \cdots t_r$ is a $T(W)$ -reduced expression for $w \in W$ and $1 \leq i < r$, then $w = t_1 t_2 \cdots t_{i-1} (t_i t_{i+1} t_i) t_i t_{i+2} \cdots t_r$ and $w = t_1 t_2 \cdots t_{i-2} t_i (t_i t_{i-1} t_i) t_{i+1} \cdots t_r$ are $T(W)$ -reduced.*

As a consequence, there exists a $T(W)$ -reduced expression for w having t_i as last reflection (or first), for $1 \leq i \leq r$. Furthermore, for any two reflections $t_i, t_j, i < j$ there exists a $T(W)$ -reduced expression for w with t_i, t_j as the last two reflections (or the first two).

2 Groups of type A , B and D

2.1 The poset $SP(A_{n-1})$

Lemma 2.1 $\ell_{T(A_{n-1})}(w_0^{A_{n-1}}) = \lfloor \frac{n}{2} \rfloor$.

Proof: Recall that $w_0^{A_{n-1}}$ is the reverse of the identity $123 \dots n$; that is, $w_0^{A_{n-1}} = n(n-1)(n-2) \dots 21$. So $\ell_{T(A_{n-1})}(w_0^{A_{n-1}}) \geq \lfloor \frac{n}{2} \rfloor$ since a reflection in A_{n-1} is just a transposition, and thus cannot permute more than two elements of $[n]$ at a time.

For $1 \leq i \leq \lfloor \frac{n}{2} \rfloor = k$, let r_i be the transposition permuting i and $n+1-i$; that is, $r_i = (i \ n+1-i)$. Notice that $r_1 \cdots r_k = w_0^{A_{n-1}}$, and so $\ell_{T(A_{n-1})}(w_0^{A_{n-1}}) \leq \lfloor \frac{n}{2} \rfloor$. \square

Lemma 2.2 For $\sigma \in A_{n-1}$, let $k = \lfloor \frac{n}{2} \rfloor$,

$$f^A(\sigma) = \left\lfloor \frac{|\{i \in [n] \mid \sigma(i) = w_0^{A_{n-1}}(i)\}|}{2} \right\rfloor$$

and

$$g^A(\sigma) = \min\{\ell : \text{there exists } t_1, t_2, \dots, t_\ell \in T(A_{n-1}) \text{ with } t_1 t_2 \dots t_\ell \sigma = w_0^{A_{n-1}}\}.$$

Then $f^A(\sigma) = i \implies g^A(\sigma) \geq \lfloor \frac{n}{2} \rfloor - i$ for $0 \leq i \leq \lfloor \frac{n}{2} \rfloor$.

Proof: We proceed by reverse (downward) induction. The case $i = k$ holds, since g^A is a non-negative function. Suppose that the statement holds for i . We now show that it also holds for $i-1$. Let $\sigma \in A_{n-1}$ with $f^A(\sigma) = i-1$. Consider $t_1, t_2, \dots, t_\ell \in T(A_{n-1})$ with $t_1 t_2 \cdots t_\ell \sigma = w_0^{A_{n-1}}$ and $\ell = g^A(\sigma)$. Notice that there exists a positive integer m with $f^A(t_{\ell-m+1} t_{\ell-m+2} \cdots t_\ell \sigma) = i$, since $f^A(t_1 t_2 \cdots t_\ell \sigma) = k$ and a reflection can fix at most two elements in their position in $w_0^{A_{n-1}}$, and so $f^A(t\tau) \leq f^A(\tau) + 1$ for $t \in T(A_{n-1})$ and $\tau \in A_{n-1}$. The last equality yields $g^A(\sigma) = \ell \geq k + m - i \geq k + 1 - i$. \square

We can now show the proposition below, which gives Theorem 1.1 for type A .

Proposition 2.3 Let $k = \lfloor \frac{n}{2} \rfloor$, and $R = \{r_1, r_2, \dots, r_k\}$, where $r_i = (i \ n+1-i)$ is the transposition permuting $i \in [k]$ and $n+1-i$. If $t_1 t_2 \cdots t_k = w_0^{A_{n-1}}$, then:

- (a) $\{t_1, \dots, t_k\} = R$
- (b) $t_i t_j = t_j t_i$ for all $i, j \in [k]$.
- (c) $(t_{\tau(1)}, t_{\tau(2)}, \dots, t_{\tau(k)})$ is a path in $B(A_{n-1})$ for all $\tau \in A_{k-1}$.

Proof: (a) Suppose that there exists $t_i \in T \setminus R$. Without loss of generality, using the Shifting Lemma, we can assume that $i = k$. Say $t_k = (m \ j)$ where $m < j \leq n$ and $j \neq n+1-k$. Hence $f^A(t_k) = 0$, and thus by Lemma 2.2 we have that $g^B(t_k) \geq k$. But this contradicts $t_1 t_2 \cdots t_{k-1} t_k = w_0^{A_{n-1}}$, which gives that $g^A(t_k) \leq k-1$.

(b) Notice that r_i and r_j are disjoint transpositions for $i \neq j$, and thus commute.

(c) By (b) it is enough to show that $\ell(t_1 t_2 \cdots t_m) > \ell(t_1 t_2 \cdots t_{m-1})$ for $1 < m \leq n$. To do this, we use Proposition 1.5.2 in [4]: If $w \in A_{n-1}$ then

$$\ell(w) = \text{inv}(w) = |\{(i, j) \in [n] \times [n] \mid i < j, w(i) > w(j)\}|.$$

Let $w \in A_{n-1}$, if $i < j, w(i) > w(j)$ then we say that (i, j) is an *inversion pair* of w .

Suppose that $w_m = t_1 t_2 \cdots t_m$; we now compare $\text{inv}(w_m)$ and $\text{inv}(w_{m-1})$. By (a) we have that the t_i 's are in R , so $t_m = (i \ n+1-i)$ for some $i \in [k]$. Moreover, $w_{m-1}(i) = i, w_{m-1}(n+1-i) = n+1-i$ and $w_m(l) = w_{m-1}(l)$ for all $l \in [n] \setminus \{i, n+1-i\}$. Now consider that:

1. If (l, i) is an inversion pair of w_{m-1} then $l < i$ and $w_{m-1}(l) > i$. If $w_{m-1}(l) > n+1-i$ then (l, i) and $(l, n+1-i)$ are inversion pairs of both w_{m-1} and w_m . If $w_{m-1}(l) \leq n+1-i$, then $(l, n+1-i)$ is not an inversion pair of w_{m-1} , but since $w_m(n+1-i) = i$, it is an inversion pair of w_m .
2. If $(l, n+1-i)$ is an inversion pair of w_{m-1} then $l < n+1-i$ and $w_m(l) = w_{m-1}(l) > n+1-i > i = w_m(n+1-i)$. Hence $(l, n+1-i)$ is also an inversion pair of w_m .
3. If (i, l) is an inversion pair of w_{m-1} then $i < l$ and $i > w_{m-1}(l)$. Since $w_m(i) = n+1-i > i > w_{m-1}(l) = w_m(l)$, (i, l) is also an inversion pair of w_m .
4. If $(n+1-i, l)$ is an inversion pair of w_{m-1} then $n+1-i < l$ and $n+1-i > w_{m-1}(l)$. If $i > w_{m-1}(l)$ then (i, l) and $(n+1-i, l)$ are inversion pairs of both w_{m-1} and w_m . If $i \leq w_{m-1}(l)$, then (i, l) is not an inversion pair of w_{m-1} , but since $w_m(i) = n+1-i$, it is an inversion pair of w_m .

Thus $\text{inv}(w_m) \geq \text{inv}(w_{m-1})$. To show that $\text{inv}(w_m) \geq \text{inv}(w_{m-1}) + 1$, consider the pair $(i \ n+1-i)$ which is *not* an inversion pair of w_{m-1} . But since $w_m(i) = n+1-i > i = w_m(n+1-i)$, this is an inversion pair of w_m . \square

We remark that the above proposition shows that $SP(A_{n-1})$ is isomorphic to the Boolean poset of rank k . Moreover, $SP(A_{n-1})$ is the poset of subsets of R ordered by inclusion.

2.2 The poset $SP(B_n)$

We used the combinatorial description of B_n and $T(B_n)$ in [4], Section 8.1.

Recall that B_n is the group of *signed permutations*; that is, the group of permutations σ of the set $[\pm n] = \{-n, -n+1, \dots, -1, 1, 2, \dots, n-1, n\}$ with the property $\sigma(-i) = -\sigma(i)$ for all $i \in [\pm n]$. We used the notation \underline{i} to denote $-i$ for $i \in [\pm n]$. We have that $w_0^{B_n} = \underline{1} \underline{2} \dots \underline{n}$. Further, $T(B_n) = \{(i \ \underline{i}) : i \in [n]\} \cup \{(i \ j)(\underline{i} \ \underline{j}) : 1 \leq i < |j| \leq n\}$. We call the set $\{(i \ \underline{i}) : i \in [n]\}$ *reflections of type I* and the set $\{(i \ j)(\underline{i} \ \underline{j}) : 1 \leq i < |j| \leq n\}$ *reflections of type II*. We now prove the analogous versions of the propositions in Section 2.1.

Proposition 2.4 $\ell_{T(B_n)}(w_0^{B_n}) = n$.

Proof: Notice that a reflection of type II changes the sign of either zero or two elements in $[n]$ and swaps them. So at least another reflection must be used to place them back in their respective order. That is, at least two reflections of type II are needed to place two elements in $[n]$ in their positions in $w_0^{B_n}$. Hence at least $2m$ reflections of type II are needed to place $2m$ elements of $[n]$ in their position in $w_0^{B_n}$, with $0 \leq m \leq \lfloor \frac{n}{2} \rfloor$, and after that $n - 2m$ reflections of type I are needed to place the remaining $n - 2m$ elements in their position in $w_0^{B_n}$. So $\ell_{T(B_n)} \geq n$.

Now, notice that $(1 \ \underline{1})(2 \ \underline{2}) \cdots (n \ \underline{n}) = w_0^{B_n}$, and so $\ell_{T(B_n)}(w_0^{B_n}) \leq n$. \square

Lemma 2.5 For $\sigma \in B_n$, let

$$f^B(\sigma) = |\{i \in [n] \mid \sigma(i) = w_0^{B_n}(i) = \underline{i}\}| + |\{(i, j) \in [n] \times [n], i < j \mid (\sigma(i), \sigma(j)) \in \{(j, i), (\underline{j}, \underline{i})\}\}|$$

and

$$g^B(\sigma) = \min\{\ell : \text{there exists } t_1, t_2, \dots, t_\ell \text{ with } t_1 t_2 \dots t_\ell \sigma = w_0^{B_n}\}.$$

Then $f^B(\sigma) = i \implies g^B(\sigma) \geq n - i$ for $0 \leq i \leq n$.

Proof: We proceed by reverse induction. The case $i = n$ holds, since g^B is a non-negative function. Suppose that the statement holds for i . We now show that it also holds for $i - 1$. Let $\sigma \in B_n$ with $f^B(\sigma) = i - 1$. Consider $t_1, t_2, \dots, t_\ell \in T(B_n)$ with $t_1 t_2 \cdots t_\ell \sigma = w_0^{B_n}$ and $\ell = g^B(\sigma)$. Notice that there exists m with $f^B(t_{\ell-m+1} \cdots t_{\ell-1} t_\ell \sigma) = i$, since $f^B(t_1 t_2 \cdots t_\ell \sigma) = n$ and a reflection can fix at most one element in its position in $w_0^{B_n}$ or create a pair (i, j) that is sent to $(\underline{j}, \underline{i})$ or (j, i) . The last equality yields $g^B(\sigma) = \ell \geq k + m - i \geq k + 1 - i$. \square

Let t_1, t_2 be two reflections of type II satisfying $\{t_1, t_2\} = \{(k \ \underline{l})(\underline{k} \ \underline{l}), (k \ \underline{l})(\underline{k} \ \underline{l})\}$ for some k, l with $1 \leq k < l \leq n$. Then we see that $t_1 t_2(k) = t_2 t_1(k) = \underline{k}$ and $t_1 t_2(l) = t_2 t_1(l) = \underline{l}$. We call the pair t_1, t_2 a *good pair*. Good pairs play a special role in the shortest paths in $B(B_n)$, as seen in the theorem below.

Proposition 2.6 If $t_1 t_2 \dots t_n = w_0^{B_n}$, then:

- For every $i \in [n]$ either t_i is of type I or there exists $j \in [n]$, $i \neq j$ so that t_i, t_j is a good pair.
- $t_i t_j = t_j t_i$ for all $i, j \in [n]$.
- $(t_{\tau(1)}, t_{\tau(2)}, \dots, t_{\tau(n)})$ is a path in $B(B_n)$ for all $\tau \in A_{n-1}$.

Proof: (a) Suppose that some reflection in $\{t_1, \dots, t_n\}$ is of type II, say $t_i = (k \ \underline{l})(\underline{k} \ \underline{l})$, and suppose that there is no t_j so that t_i, t_j is a good pair. Since $w_0^{B_n}(k) = \underline{k}$ and $w_0^{B_n}(l) = \underline{l}$, there must be another reflection t_m that is not disjoint from t_i . Without loss of generality, we can assume that $\{t_i, t_m\} = \{t_{n-1}, t_n\}$. Since t_{n-1}, t_n is not a good pair, then $f^B(t_{n-1}t_n) = 0$. Hence $g^B(t_{n-1}t_n) \geq n$, which contradicts $t_1t_2 \cdots t_n = w_0^{B_n}$.

(b) Notice that since all the reflections in $t_1 \cdots t_n = w_0^{B_n}$ of type I are distinct, they commute with each other. Furthermore, if t_i, t_j are a good pair, then they also commute. We need to verify that (i) if t_i, t_j are of type II and *not* a good pair, then they are commuting reflections, and (ii) if t_i, t_j are of mixed types, then they commute. Using the Shifting Lemma again, we can assume that the reflections in both cases are t_{n-1} and t_n . Suppose that t_{n-1} and t_n do not commute. In both (i) and (ii) we see that $f^B(t_{n-1}t_n) = 0$, and so $g^B(t_{n-1}t_n) \geq n$ by Lemma 2.5, which contradicts $t_1t_2 \cdots t_{n-1}t_n = w_0^{B_n}$.

(c) By Proposition 8.1.1 in [4], if $w \in B_n$ then

$$\ell(w) = \text{inv}(w) + \text{Neg}(w)$$

where

$$\text{inv}(w) = \text{inv}(w(1), w(2), \dots, w(n)) \quad \text{and} \quad \text{Neg}(w) = - \sum_{j \in [n]: w(j) < 0} v(j).$$

For $i \in [n]$, let $w_i = t_1t_2 \cdots t_i$. Notice that from (b) it is enough to prove that $\ell(w_m) > \ell(w_{m-1})$ for $1 < m \leq n$. We have the following cases:

1. t_m is of type I, say $t_m = (j \ \underline{j})$, with $j \in [n]$. (a) and (b) give that no other reflection involves the element j , and so $w_{m-1}(j) = j$. Furthermore, we have that $w_m(k) = w_{m-1}(k)$ for $k \in [n] \setminus \{j\}$. Now,

- If (i, j) is an inversion pair of w_{m-1} , then $i < j$ and $w_{m-1}(i) > w_{m-1}(j) = j$, which gives that $w_{m-1}(i) > 0$. So $w_m(i) = w_{m-1}(i) > \underline{j} = w_m(j)$, and the pair (i, j) is also an inversion pair of w_m . Since $\text{Neg}(w_m) = \text{Neg}(w_{m-1}) + j$, we have that $\ell(w_{m-1}) < \ell(w_m)$.
- If (j, i) is an inversion pair of w_{m-1} , then $j < i$ and $w_{m-1}(j) = j > w_{m-1}(i)$. Suppose that (j, i) is not an inversion pair of w_m . There are at most $j - 1$ such inversion pairs (j, i) of w_{m-1} , since $1 < w_{m-1}(i) < j$. On the other hand, notice that $\text{Neg}(w_m) = \text{Neg}(w_{m-1}) + j$. So

$$\begin{aligned} \ell(w_m) - \ell(w_{m-1}) &= \text{inv}(w_m) + \text{Neg}(w_m) - (\text{inv}(w_{m-1}) + \text{Neg}(w_{m-1})) \\ &\geq \text{inv}(w_{m-1}) - (j - 1) + (\text{Neg}(w_{m-1}) + j) - (\text{inv}(w_{m-1}) + \text{Neg}(w_{m-1})) \\ &\geq 1 \end{aligned}$$

2. t_m is of type II but does not change any element's signs, say $t_m = (i \ \underline{j})(\underline{i} \ \underline{j})$ with $1 \leq i < j \leq n$. Then by the same argument as in the proof of Proposition 2.3 (c), we have that $\ell(w_m) > \ell(w_{m-1})$.

3. If $t_m = (i \ \underline{j})(\underline{i} \ \underline{j})$, with $1 \leq i < j \leq n$; that is, t_m swaps i and j and changes their sign. (a) and (b) give that $(w_{m-1}(i), w_{m-1}(j)) \in \{(i, j), (j, i)\}$, $(w_m(i), w_m(j)) \in \{(\underline{j}, \underline{i}), (\underline{i}, \underline{j})\}$, and $w_{m-1}(k) = w_m(k)$ for $k \in [\pm n] \setminus \{\pm i, \pm j\}$. Then

- If (k, i) is an inversion pair of w_{m-1} then $k < i$ and either $w_{m-1}(k) > i$ or $w_{m-1}(k) > j$. In either case (k, i) is also an inversion pair of w_m since $w_m(k) = w_{m-1}(k) > 0$ and $w_m(i) < 0$. Further, $\text{Neg}(w_m) = \text{Neg}(w_{m-1}) + i + j$, and so $\ell(w_{m-1}) < \ell(w_m)$.
- If (i, k) is an inversion pair of w_{m-1} then $i < k$ and either $i > w_{m-1}(k)$ or $j > w_{m-1}(k)$. If we assume that (i, k) is *not* an inversion pair of w_m , then in the former case, there are at most $i - 1$ pairs lost, and in the latter there are at most $j - 1$. However, since $\text{Neg}(w_m) = \text{Neg}(w_{m-1}) + i + j$, we still have that $\ell(w_{m-1}) < \ell(w_m)$.
- If (j, k) is an inversion pair of w_{m-1} then $j < k$ and either $j > w_{m-1}(k)$ or $i > w_{m-1}(k)$. If we assume that (j, k) is *not* an inversion pair of w_m , then in the former case, there are at most $j - 1$ pairs lost, and in the latter there are at most $i - 1$. However, since $\text{Neg}(w_m) = \text{Neg}(w_{m-1}) + i + j$, we still have that $\ell(w_{m-1}) < \ell(w_m)$.
- If (k, j) is an inversion pair of w_{m-1} then $k < j$ and either $w_{m-1}(k) > j$ or $w_{m-1}(k) > i$. In either case (k, j) is also an inversion pair of w_m since $w_m(k) = w_{m-1}(k) > 0$ and $w_m(j) < 0$. Further, $\text{Neg}(w_m) = \text{Neg}(w_{m-1}) + i + j$, and so $\ell(w_{m-1}) < \ell(w_m)$.

In all cases, we have the desired result. \square

The previous proposition says that $SP(B_n)$ is (isomorphic to) the union of Boolean posets of rank n , one for each set $\{t_1, t_2, \dots, t_n\}$ with $t_1 t_2 \cdots t_n = w_0^{B_n}$. As an example, Figure 1 illustrates that $SP(B_2)$ is the union of two Boolean posets. In general, one can compute the number of Boolean posets in $SP(B_n)$.

2.2.1 Number of Boolean posets in $SP(B_n)$

Let b_n be the number of Boolean posets in B_n . We obtain a Boolean poset for each set $\{t_1, \dots, t_n\}$ with $t_1 t_2 \cdots t_n = w_0^{B_n}$. It is easy to see that $b_1 = 1$ and $b_2 = 2$ (see Figure 1). For $n \geq 2$, notice that if $t_1 t_2 \cdots t_n(1) = \underline{1}$, then by Proposition 2.6 there are two possible cases: (i) there exists $t_j = (1 \ \underline{1})$ or there exists a good pair of reflections of the form $(1 \ \underline{k})(\underline{k} \ 1), (1 \ k)(\underline{k} \ \underline{1})$. There are b_{n-1} such reflections in case (i) and $(n-1)b_{n-2}$ in case (ii). So b_n satisfies the recurrence relation

$$b_n = b_{n-1} + (n-1)b_{n-2}$$

with initial conditions $b_1 = 1$ and $b_2 = 2$. Notice that this count is the same as the number of partitions of a set of n distinguishable elements into sets of size 1 and 2.

It is easy to see that

$$b_n = 1 + \sum_{j=1}^{\lfloor \frac{n}{2} \rfloor} \frac{1}{j!} \prod_{i=0}^{j-1} \binom{n-2i}{2}.$$

2.3 The poset $SP(D_n)$

As in the previous section, we used the combinatorial description of D_n and $T(D_n)$ in [4] Section 8.2.

D_n ($n > 1$) is the group of *signed permutations* with an *even* number of negative elements (e.g. $\underline{1} \ \underline{2} \ 3$ is an element in D_3 whereas $\underline{1} \ \underline{2} \ \underline{3}$ is not). Like B_n , if $\sigma \in D_n$ then $\sigma(-i) = -\sigma(i)$ for $i \in [\pm n]$. Moreover, $w_0^{D_n} = \underline{1} \ \underline{2} \ \dots \ \underline{n}$ if n is *even* and $w_0^{D_n} = 1 \ \underline{2} \ \dots \ \underline{n}$ if n is *odd*. Further, $T(D_n) = \{(i \ j)(\underline{i} \ \underline{j}) : 1 \leq i < |j| \leq n\}$; that is, the reflections of D_n are the reflections of B_n of type II.

Proposition 2.7 $\ell_{T(D_n)}(w_0^{D_n}) = n$ if n is even, and $\ell_{T(D_n)} = n - 1$ if n is odd.

Proof: Same as for Proposition 2.4, but only using reflections of type II. Notice that that for n even, $r_1 r'_1 r_2 r'_2 \cdots r_k r'_k = w_0^{D_n}$, where $k = n/2$ and $r_i = (2i-1 \ 2i)(2i \ 2i-1)$, $r'_i = (2i-1 \ 2i)(2i \ 2i-1)$ $1 \leq i \leq n/2$. Similarly, for n odd, we have that $t_1 t'_1 t_2 t'_2 \cdots t_k t'_k = w_0^{D_n}$, where $k = (n-1)/2$ and $r_i = (2i \ 2i+1)(2i+1 \ 2i)$, $r'_i = (2i \ 2i+1)(2i+1 \ 2i)$, $1 \leq i \leq (n-1)/2$. \square

Lemma 2.8 For $\sigma \in D_n$, for n even, define

$$f^D(\sigma) = |\{i \in [n] \mid \sigma(i) = w_0^{D_n}(i) = \underline{i}\}| + |\{(i, j) \in [n] \times [n], i < j \mid (\sigma(i), \sigma(j)) \in \{(j, i), (\underline{j}, \underline{i})\}\}|$$

and for n odd, define

$$f^D(\sigma) = |\{i \in [n] \setminus \{1\} \mid \sigma(i) = w_0^{D_n}(i) = \underline{i}\}| + |\{(i, j) \in [n] \times [n], i < j \mid (\sigma(i), \sigma(j)) \in \{(j, i), (\underline{j}, \underline{i})\}\}|.$$

Moreover, let

$$g^D(\sigma) = \min\{\ell : \text{there exists } t_1, t_2, \dots, t_\ell \text{ with } t_1 t_2 \dots t_\ell \sigma = w_0^{D_n}\}.$$

Then $f^D(\sigma) = i \implies g^D(\sigma) \geq m - i$ for $0 \leq i \leq m$ and $m = n$ if n is even, $m = n - 1$ if n is odd.

Proof: Same as in Lemma 2.5, using only reflections of type II. \square

Proposition 2.9 Suppose that $t_1 t_2 \dots t_m = w_0^{D_n}$, where $m = n$ if n is even and $m = n - 1$ if n is odd. Then:

- (a) For every $i \in [m]$ there exists $j \in [m]$, $i \neq j$ so that t_i, t_j is a good pair.
- (b) $t_i t_j = t_j t_i$ for all $i, j \in [m]$.
- (c) $(t_{\tau(1)}, t_{\tau(2)}, \dots, t_{\tau(m)})$ is a path in $B(D_n)$ for all $\tau \in A_{m-1}$.

Proof: The proof for (a) and (b) is the same as in Proposition 2.6, but only using reflections of type II.

For (c), even though the length function is not the same as described in the Section 2.2, we recall that $B(D_n)$ is the induced graph of $B(B_n)$ on the elements of D_n by Proposition 8.2.6 in [4]. \square

2.3.1 Number of Boolean posets in $SP(D_n)$

Let d_n be the number of Boolean posets in $SP(D_n)$ for each set $\{t_1, t_2, \dots, t_n\} \subset T(D_n)$ with $t_1 t_2 \cdots t_n = w_0^{D_n}$. Counting these subsets is equivalent to counting the partitions of $[n]$, if n is even, or $[n-1]$, if n is odd, into subsets of two elements (these represents the good pairs). That is,

$$d_m = \frac{1}{[\frac{m}{2}]!} \prod_{i=0}^{[\frac{m}{2}]-1} \binom{m-2i}{2}$$

where $m = n$ if n is even, and $m = n - 1$ if n is odd. Since m is even, notice that this is the same as counting the number of partitions of $[m]$ into sets of size 2.

2.4 Finite Dihedral groups

Let $I_2(m)$, $m \geq 1$ be the dihedral group of order $2m$ with generating set $\{s_1, s_2\}$, and let $T = T(I_2(m))$ its reflection set. If n is odd, then

$$w_0^{I_2(m)} = \underbrace{s_1 s_2 s_1 \cdots s_1}_m = \underbrace{s_2 s_1 s_2 \cdots s_2}_m$$

is a reflection, and so $\ell_T(w_0) = 1$. Hence $SP(I_2(m))$ is isomorphic to the Boolean poset of rank 1, if m is odd.

The case where m is even is more interesting, as $w_0^{I_2(m)} \notin T$. We readily see that $\ell_T(w_0) = 2$, since for instance $w_0^{I_2(m)} = s_1 \underbrace{s_2 s_1 \cdots s_2}_{m-1}$. Thus $SP(I_2(m))$ is the union of Boolean posets of rank 2, if m is even.

Fix w_0 to start with s_1 . We now count number of Boolean posets in $SP(I_2(m))$ for m even. This number is the same as the number of sets $\{t_1, t_2\}$ with $t_1 t_2 = w_0^{I_2(m)}$. There is one such set for each element of odd rank that starts with s_1 , since for each such element t_1 there exists a unique element t_2 with $t_1 t_2 = w_0^{I_2(m)}$. Since there are $\frac{m}{2}$ such elements, there are $\frac{m}{2}$ Boolean posets in $SP(I_2(m))$.

3 Exceptional Coxeter groups

3.1 F_4, H_3, H_4 , and E_6

We were able to verify through computer search that the the results in the previous sections also worked for the following exceptional groups: F_4, H_3, H_4, E_6 . That is, the shortest path poset for these groups form a union of Boolean posets of rank the absolute length of the longest word w_0^W . We summarize the results in Table 1. The computer search was done using Stembridge's `coxeter` Maple package [7], and it basically consisted of finding all shortest paths and verifying the analogues of Propositions 2.3, 2.6, 2.9 for those groups; that is, that the paths are given by reflections that are fully commutative.

An interesting observation is that the 3 Boolean posets that form $SP(E_6)$ are almost disjoint, sharing only e and $w_0^{E_6}$ (the bottom and top elements of each poset).

3.1.1 E_7 and E_8

For E_7 and E_8 we were not able to verify by computer that the shortest paths form a union of Boolean posets, since it involved more computer power (or a better code) than was available to us. However, we can argue that this is indeed the case using geometric arguments. Let (W, S) be Coxeter system, and consider the *geometric representation* of W , $\sigma : W \hookrightarrow GL(V)$, where V is a vector space with basis $\Pi = \{\alpha_s \mid s \in S\}$ (Π is called the set of *simple roots*). It is shown in [6] Section 5.4 that σ is a faithful representation.

The *root system* of the Coxeter system (W, S) is the set $\Phi = \{\sigma(w)(\alpha_s) : s \in S, w \in W\}$. Let $\beta \in \Phi$, then $\beta = \sum_{s \in S} c_s \alpha_s$. It is a well-known result that either $c_s \geq 0$ or $c_s \leq 0$ for all $s \in S$. In the former case we say that β is a *positive root*, and in the latter case we say that β is *negative root*. The set of positive roots is denoted by Φ^+ and the set of negative roots is denoted by Φ^- . It is also a well known fact (Proposition 4.4.5 in [4]) that there is a bijection between the set of reflections of W , $T(W)$ and Φ^+ given by $t = w s w^{-1} \mapsto \sigma(w)(\alpha_s)$.

Finally, we shall use the fact that $\sigma(w_0^{E_n}) = -\mathbf{id}$, where \mathbf{id} is the identity matrix of dimension n , and $n = 7, 8$. We point out that $\sigma(w_0^{E_n}) \neq -\mathbf{id}$, and thus $\text{rank}(SP(E_6)) < 6$. For details, see [2] Chapter VI, §4.10 and §.11. With this in mind we can show

Proposition 3.1 *For E_n , where $n = 7, 8$ we have that:*

- (a) $\ell_T(w_0^{E_n}) = n$.
- (b) If $w_0^{E_n} = t_1 t_2 \cdots t_n$ then $t_i t_j = t_j t_i$ for all $i, j \in [n]$.
- (c) $(t_{\tau(1)}, t_{\tau(2)}, \dots, t_{\tau(n)})$ is a path in $B(E_n)$ for all $\tau \in A_{n-1}$.

Proof: (a) Since a reflection fixes a hyperplane, the product of k reflections fixes the intersection of the k hyperplanes that are fixed by each reflection. This intersection has codimension at most k , and so it's not empty unless $k \geq n$. In particular, $\sigma(w_0^{E_n}) = -\mathbf{id}$ leaves no points fixed (except for $\mathbf{0}$) and so cannot be written as a product of fewer than n reflections; that is $\ell_T(w_0^{E_n}) \geq n$. Moreover by Carter's Lemma (Lemma 2.4.5 in [1]), we have that $\ell_T(w_0^{E_n}) \leq n$. Thus $\ell_T(w_0^{E_n}) = n$.

(b) Now consider $-\mathbf{id} = s_{t_1} s_{t_2} \cdots s_{t_n}$, where $\sigma(t_i) = s_{t_i}$ for $1 \leq i \leq n$ are the reflections (in V) with respect to the hyperplanes $\mathcal{H}_1, \mathcal{H}_2, \dots, \mathcal{H}_n$ that are perpendicular to the unit vectors $\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n$. The space fixed by the product $s_{t_1} s_{t_2} \cdots s_{t_{n-1}}$ is $\mathbb{R}\mathbf{v}_n$ (since the product of everything is $-\mathbf{id}$) which has co-dimension $n - 1$ and then by the previous argument,

$$\bigcap_{i < n} \mathcal{H}_i = \mathbb{R}\mathbf{v}_n,$$

that is, $v_n \in \mathcal{H}_i$ for all $i < n$. Hence, $\mathbf{v}_i \perp \mathbf{v}_n$, which means that t_n commutes with t_i for $i < n$. By the Shifting Lemma, we have that any two reflections t_i, t_j commute.

(c) Let $t_1 \cdots t_n = w_0^{E_n}$. We are going to show that $\ell(t_1 t_2 \cdots t_k) > \ell(t_1 t_2 \cdots t_{k-1})$ for $1 < k \leq n$. As before, let $s_{t_i} = \sigma(t_i)$ be the reflection on V corresponding to t_i about the hyperplane \mathcal{H}_i , and let \mathbf{v}_i be the normal vector to \mathcal{H}_i . Since $\mathbf{v}_i \perp \mathbf{v}_j$ for all $i \neq j$, we have that $s_{t_1} s_{t_2} \cdots s_{t_{i-1}}(\alpha_i) = \alpha_i$, where $\alpha_i \in \Phi^+$ is the positive root corresponding to t_i . Thus by Proposition 4.4.6 in [4], we have that $\ell(t_1 t_2 \cdots t_i) > \ell(t_1 t_2 \cdots t_{i-1})$ for $1 \leq i \leq n$. \square

As a consequence of the above theorem, $SP(E_7)$ and $SP(E_8)$ are both formed by the union of Boolean posets that share at least the bottom and top elements. We are now done with the proof of Theorem 1.1.

3.1.2 Number of Boolean posets in $SP(E_7)$ and $SP(E_8)$

To count the number of paths (chains) in $SP(E_n)$ where $n = 7, 8$ we simply count the number n -tuples of perpendicular roots, since $\sigma(w_0^{E_n}) = -\mathbf{id}$. Each one of these n -tuples up to signs and permutations represents a Boolean poset. Direct computation yields 135 Boolean posets in $SP(E_7)$ and 2025 Boolean posets in $SP(E_8)$. These results are included in Table 1.

Remark 3.2 *The above geometric argument can be used to obtain the results that were proven in Section 2. As was the case in our proofs, each group type requires its own argument, since $\sigma(w_0^W)$ is different for each case. However we believe that the combinatorial proofs are more appropriate for the FPSAC audience.*

4 Lowest-degree terms of the complete **cd**-index of finite Coxeter groups

For any Eulerian poset P , one can define the **cd**-index of P . This polynomial encodes the flag h -vector. The interested reader is referred to [5] for more information on the **cd**-index of Eulerian posets. Since Bruhat intervals are Eulerian and the reflection ordering has the property of having a unique chain (path in the Bruhat graph) with no descents for every interval $[u, v]$, the highest-degree terms of the complete **cd**-index coincide with the **cd**-index.

Let $\tilde{\psi}(\mathcal{B}_n)$ be **cd**-index of the Boolean poset \mathcal{B}_n (so \mathcal{B}_n is the poset of subsets of $[n]$ ordered by inclusion). We can use Theorem 5.2 in [5] to compute $\tilde{\psi}(\mathcal{B}_n)$. First $\tilde{\psi}(\mathcal{B}_1) = 1$ and for $n > 1$,

$$\tilde{\psi}(\mathcal{B}_n) = \tilde{\psi}(\mathcal{B}_{n-1}) \cdot \mathbf{c} + G(\tilde{\psi}(\mathcal{B}_{n-1})) \quad (1)$$

where G is the derivation $G(\mathbf{c}) = \mathbf{d}$ and $G(\mathbf{d}) = \mathbf{cd}$. In particular, we have that $\tilde{\psi}(\mathcal{B}_2) = \mathbf{c}$, $\tilde{\psi}(\mathcal{B}_3) = \mathbf{c}^2 + \mathbf{d}$, $\tilde{\psi}(\mathcal{B}_4) = \mathbf{c}^3 + 2\mathbf{cd} + 2\mathbf{dc}$, and so on.

Propositions 2.3, 2.6 and 2.9, and the results and computer search of Section 3 give that for a finite Coxeter group W , the corresponding $SP(W)$ is the union of Boolean posets (that share at least the bottom and top elements). So any interval in $SP(W)$ belongs to a Boolean poset corresponding to a set $R = \{t_1, t_2, \dots, t_\ell\} \subset T(W)$ with $\ell_{T(W)}(w_0^W) = \ell$ and $t_1 t_2 \cdots t_\ell = w_0^W$. Thus any interval of $SP(W)$ (thought of as paths in $B(W)$ labeled with $T(W)$, where $T(W)$ is ordered by a reflection ordering) has a unique chain (path) with empty descent set. Hence counting descent sets in the chains given by R is the same as counting the flag h -vector of the Boolean poset of rank ℓ .

As a consequence, the lowest-degree terms in the complete **cd**-index of W add up to a multiple N of the **cd**-index of the Boolean poset of ranks $\ell_{T(W)}(w_0^W)$. N is the number of Boolean posets in $SP(W)$; that is, the number of sets $\{t_1, \dots, t_{\ell_{T(W)}(w_0^W)}\}$ with $t_1 t_2 \cdots t_{\ell_{T(W)}(w_0^W)} = w_0^W$. These terms can be computed using (1) and Table 1. So we have

Theorem 4.1 *Let W be a finite Coxeter group, α_W is the number of Boolean posets that form $SP(W)$ and $\ell_0 = \ell_{T(W)}(w_0^W)$. Then lowest degree terms of $\tilde{\psi}_{e, w_0^W}$ are given by $\alpha_W \tilde{\psi}(\mathcal{B}_{\ell_0})$.*

In particular, the lowest-degree terms of $\tilde{\psi}_{e, w_0^W}$ are minimized by $\tilde{\psi}(\mathcal{B}_{\ell_0})$.

Acknowledgements

I thank Professor Louis J. Billera for insightful conversations and suggestions, including simplifying and correcting proofs. Martin Kassabov helped me (very patiently and diligently) understand the geometric argument used in Section 3. Finally, Mihai Bailesteanu helped with the translation of the abstract to French.

Tab. 1: Finite coxeter groups W , $\text{rank}(SP(W))$, and the number of Boolean posets in $SP(W)$

W	$\text{rank}(SP(W))$	# of Boolean posets in $SP(W)$
A_{n-1}	$\lfloor \frac{n-1}{2} \rfloor$	1
B_n	n	b_n
D_n	n if n is even; $n - 1$ if n is odd	d_n
$I_2(m)$	2 if m is even; 1 if m is odd	$\frac{m}{2}$ if m is even; 1 if m is odd
F_4	4	24
H_3	3	5
H_4	4	75
E_6	4	3
E_7	7	135
E_8	8	2025

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