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# *A fixed point theorem for Boolean networks expressed in terms of forbidden subnetworks*

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We are interested in fixed points in Boolean networks, *i.e.* functions  $f$  from  $\{0, 1\}^n$  to itself. We define the subnetworks of  $f$  as the restrictions of  $f$  to the hypercubes contained in  $\{0, 1\}^n$ , and we exhibit a class  $\mathcal{F}$  of Boolean networks, called even or odd self-dual networks, satisfying the following property: if a network  $f$  has no subnetwork in  $\mathcal{F}$ , then it has a unique fixed point. We then discuss this “forbidden subnetworks theorem”. We show that it generalizes the following fixed point theorem of Shih and Dong: if, for every  $x$  in  $\{0, 1\}^n$ , there is no directed cycle in the directed graph whose the adjacency matrix is the discrete Jacobian matrix of  $f$  evaluated at point  $x$ , then  $f$  has a unique fixed point. We also show that  $\mathcal{F}$  contains the class  $\mathcal{F}'$  of networks whose the interaction graph is a directed cycle, but that the absence of subnetwork in  $\mathcal{F}'$  does not imply the existence and the uniqueness of a fixed point.

**Keywords:** Boolean network, fixed point, self-dual Boolean function, discrete Jacobian matrix, feedback circuit.

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## 1 Introduction

A function  $f$  from  $\{0, 1\}^n$  to itself is often seen as a Boolean network with  $n$  components. On one hand, the dynamics of the network is described by the iterations of  $f$ ; for instance, with the synchronous iteration scheme, the dynamics is described by the recurrence  $x^{t+1} = f(x^t)$ . On the other hand, the “structure” of the network is described by a directed graph  $G(f)$ : the vertices are the  $n$  components, and there exists an arc from  $j$  to  $i$  when the evolution of the  $i$ th component depends on the evolution of the  $j$ th one.

Boolean networks have many applications. In particular, from the seminal works of Kauffman (1969) and Thomas (1973), they are extensively used to model gene networks. In most cases, fixed points are of special interest. For instance, in the context of gene networks, they are often seen as stable patterns of gene expression at the basis of particular biological processes.

In this paper, we are interested in sufficient conditions for the existence and the uniqueness of a fixed point for  $f$ . Such a condition was first obtained by Robert (1980), who proved that *if  $G(f)$  has no directed cycle, then  $f$  has a unique fixed point*. This result was then generalized by Shih and Dong (2005). They associated to each point  $x$  in  $\{0, 1\}^n$  a local interaction graph  $Gf(x)$ , which is a subgraph of  $G(f)$  defined as the directed graph whose the adjacency matrix is the discrete Jacobian matrix of  $f$  evaluated at point  $x$ , and they proved that *if  $Gf(x)$  has no directed cycle for all  $x$  in  $\{0, 1\}^n$ , then  $f$  has a unique fixed point*. Up

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to our knowledge, this is the weakest condition known to be sufficient for the presence and the uniqueness of a fixed point.

In this paper, we establish a sufficient condition for the existence and the uniqueness of a fixed point that is *not* expressed in terms of directed cycles. In Section 2, we defined, in a natural way, the subnetworks of  $f$  as the restrictions of  $f$  to the hypercubes contained in  $\{0, 1\}^n$ , and we introduce the class  $\mathcal{F}$  of even and odd self-dual networks. In Section 3, we prove the main result: *if  $f$  has no subnetworks in  $\mathcal{F}$ , then it has a unique fixed point*. The rest of the paper discusses this “forbidden subnetworks theorem”. In section 4, we show that it generalizes the fixed point theorem of Shih and Dong mentioned above. In section 5, we study the effect of the absence of subnetwork in  $\mathcal{F}$  on the asynchronous state graph of  $f$ , which is a directed graph on  $\{0, 1\}^n$  constructed from the asynchronous iterations of  $f$  and proposed by Thomas (1973) as a model for the dynamics of gene networks. Finally, in Section 6, we compare  $\mathcal{F}$  with the well-known class  $\mathcal{F}'$  of networks  $f$  whose the interaction graph  $G(f)$  is a directed cycle. Mainly, we show that  $\mathcal{F}' \subseteq \mathcal{F}$  and that the absence of subnetwork in  $\mathcal{F}'$  is not sufficient for the existence and the uniqueness of a fixed point.

## 2 Definitions and notations

In this section, we introduce the definitions needed to state and prove the main result. Let  $\mathbb{B} = \{0, 1\}$ , let  $n$  be a positive integer, let  $[n] = \{1, \dots, n\}$ , and let  $i \in [n]$ . The  $i$ th unit vector of  $\mathbb{B}^n$  is denoted  $e_i$  (all the components are 0, excepted the  $i$ th one which is 1). The sum modulo two is denoted  $\oplus$ . It is applied componentwise on elements of  $\mathbb{B}^n$ : for all  $x, y \in \mathbb{B}^n$ ,

$$x \oplus y = (x_1 \oplus y_1, \dots, x_n \oplus y_n) \quad \text{and} \quad x \oplus 1 = (x_1 \oplus 1, \dots, x_n \oplus 1).$$

Hence,  $x \oplus 1$  may be seen as the negation of  $x$ . The number of ones that  $x$  contains is denoted  $\|x\|$ , *i.e.*  $\|x\| = \sum_{i=1}^n x_i$ . Thus  $\|x \oplus y\|$  gives the *Hamming distance* between two points  $x$  and  $y$  of  $\mathbb{B}^n$ . We say that  $x$  is *even (odd)* if  $\|x\|$  is even (odd) (there exists  $2^{n-1}$  even (odd) points in  $\mathbb{B}^n$ ). The point of  $\mathbb{B}^n$  obtained from  $x$  by assigning the  $i$ th component to  $\alpha \in \mathbb{B}$  is denoted  $x^{i\alpha}$ , *i.e.*

$$x^{i\alpha} = (x_1, \dots, x_{i-1}, \alpha, x_{i+1}, \dots, x_n).$$

If  $n > 1$ , the point of  $\mathbb{B}^{n-1}$  obtained from  $x$  by removing the  $i$ th component is denoted  $x_{-i}$ , *i.e.*

$$x_{-i} = (x_1, \dots, x_{i-1}, x_{i+1}, \dots, x_n).$$

We call (*n-dimensional Boolean*) *networks* any function  $f$  from  $\mathbb{B}^n$  to itself.

**Definition 1 (Conjugate)** *The conjugate of  $f : \mathbb{B}^n \rightarrow \mathbb{B}^n$  is the following  $n$ -dimensional network:*

$$\tilde{f} : \mathbb{B}^n \rightarrow \mathbb{B}^n, \quad \tilde{f}(x) = x \oplus f(x) \quad \forall x \in \mathbb{B}^n.$$

Remark that  $\tilde{f}(x) = 0$  if and only if  $x$  is a fixed point of  $f$ , *i.e.*  $f(x) = x$ .

**Definition 2 (Self-dual networks and even/odd networks)**  *$f$  is self-dual if*

$$f(x) = f(x \oplus 1) \oplus 1 \quad \forall x \in \mathbb{B}^n.$$

*$f$  is even (odd) if the image of  $\tilde{f}$  is the set of even points of  $\mathbb{B}^n$ , *i.e.**

$$\{\tilde{f}(x) \mid x \in \mathbb{B}^n\} = \{x \mid x \in \mathbb{B}^n \text{ and } \|x\| \text{ is even (odd)}\}.$$

We say that  $f$  is *even (odd) self-dual* if it is both even (odd) and self-dual. Note that  $f(x) = f(x \oplus 1) \oplus 1$  if and only if  $\tilde{f}(x \oplus 1) = \tilde{f}(x)$ . Note also that if  $f$  is even (odd) self-dual, then for each even (odd) point  $x \in \mathbb{B}^n$ , the preimage of  $x$  by  $\tilde{f}$  is of cardinality two, *i.e.* there exists exactly two distinct points  $y, z \in \mathbb{B}^n$  such that  $\tilde{f}(y) = \tilde{f}(z) = x$ . Since  $\tilde{f}(x) = 0$  if and only if  $f(x) = x$ , we deduce that if  $f$  is even self-dual, then it has exactly two fixed points (obviously, if  $f$  is odd self-dual, then it has no fixed point).

**Definition 3 (Immediate subnetworks)** *If  $n > 1$ ,  $\alpha \in \mathbb{B}$  and  $i \in [n]$ , we call immediate subnetwork of  $f$  (obtained by fixing the  $i$ th component to  $\alpha$ ) the following  $(n - 1)$ -dimensional network:*

$$f^{i\alpha} : \mathbb{B}^{n-1} \rightarrow \mathbb{B}^{n-1}, \quad f^{i\alpha}(x_{-i}) = f(x^{i\alpha})_{-i} \quad \forall x \in \mathbb{B}^n.$$

Remark that conjugate of  $f^{i\alpha}$  is equal to the immediate subnetwork  $\tilde{f}^{i\alpha}$  of the conjugate  $\tilde{f}$  of  $f$ :

$$\widetilde{f^{i\alpha}}(x_{-i}) = x_{-i} \oplus f^{i\alpha}(x_{-i}) = x_{-i} \oplus f(x^{i\alpha})_{-i} = (x \oplus f(x^{i\alpha}))_{-i} = \tilde{f}(x^{i\alpha})_{-i} = \tilde{f}^{i\alpha}(x_{-i}).$$

**Definition 4 (Subnetworks)** *The subnetworks of  $f$  are inductively defined by: (1) if  $n = 1$ , then  $f$  has a unique subnetwork, which is  $f$  itself; and (2) if  $n > 1$ , the subnetworks of  $f$  are  $f$  and the subnetworks of the immediate subnetworks of  $f$ . A strict subnetwork of  $f$  is a subnetwork of  $f$  different than  $f$ .*

### 3 Main result

**Theorem 1 (Forbidden subnetworks theorem)** *If a network  $f : \mathbb{B}^n \rightarrow \mathbb{B}^n$  has no even or odd self-dual subnetwork, then the conjugate of  $f$  is a bijection, and in particular,  $f$  has a unique fixed point.*

The proof of Theorem 1 needs the following two lemmas.

**Lemma 1** *Let  $X$  be a non-empty subset of  $\mathbb{B}^n$  and  $V(X) = \{x \oplus e_i \mid x \in X, i \in [n]\}$ . If  $X$  and  $V(X)$  are disjoint and  $|X| \geq |V(X)|$ , then  $X$  is either the set of even points of  $\mathbb{B}^n$  or the set of odd points of  $\mathbb{B}^n$ .*

**Proof:** by induction on  $n$ . The case  $n = 1$  is obvious. So suppose that  $n > 1$  and that the lemma holds for the dimensions less than  $n$ . Let  $X$  be a non-empty subset of  $\mathbb{B}^n$  satisfying the conditions of the statement. Let  $\alpha \in \mathbb{B}$ , and consider the following subsets of  $\mathbb{B}^{n-1}$ :

$$X^\alpha = \{x_{-n} \mid x \in X, x_n = \alpha\}, \quad V(X)^\alpha = \{x_{-n} \mid x \in V(X), x_n = \alpha\}.$$

We first prove that  $V(X^\alpha) \subseteq V(X)^\alpha$  and  $X^\alpha \cap V(X^\alpha) = \emptyset$ . Let  $x \in \mathbb{B}^n$  with  $x_n = \alpha$  be such that  $x_{-n} \in V(X^\alpha)$ . To prove that  $V(X^\alpha) \subseteq V(X)^\alpha$ , it is sufficient to prove that  $x_{-n} \in V(X)^\alpha$ . Since  $x_{-n} \in V(X^\alpha)$ , there exists  $y \in \mathbb{B}^n$  with  $y_n = \alpha$  and  $i \in [n - 1]$  such that  $y_{-n} \in X^\alpha$  and  $x_{-n} = y_{-n} \oplus e_i$ . So  $x = y \oplus e_i$ , and since  $y_n = \alpha$ , we have  $y \in X$ . Hence  $x \in V(X)$  and since  $x_n = \alpha$ , we have  $x_{-n} \in V(X)^\alpha$ . We now prove that  $X^\alpha \cap V(X^\alpha) = \emptyset$ . Indeed, otherwise, there exists  $x \in \mathbb{B}^n$  with  $x_n = \alpha$  such that  $x_{-n} \in X^\alpha \cap V(X^\alpha)$ . Since  $V(X^\alpha) \subseteq V(X)^\alpha$ , we have  $x_{-n} \in X^\alpha \cap V(X)^\alpha$ , and since  $x_n = \alpha$ , we deduce that  $x \in X \cap V(X)$ , a contradiction.

Now, since  $V(X^\alpha) \subseteq V(X)^\alpha$ , we have

$$|X| = |X^0| + |X^1| \geq |V(X)| = |V(X)^0| + |V(X)^1| \geq |V(X^0)| + |V(X^1)|.$$

So  $|X^0| \geq |V(X^0)|$  or  $|X^1| \geq |V(X^1)|$ . Suppose that  $|X^0| \geq |V(X^0)|$ , the other case being similar. Since  $X^0 \cap V(X^0) = \emptyset$ , by induction hypothesis  $X^0$  is either the set of even points of  $\mathbb{B}^{n-1}$  or the

set of odd points of  $\mathbb{B}^{n-1}$ . So in both cases, we have  $|X^0| = |V(X^0)| = 2^{n-1}$ . We deduce that  $|X^1| \geq |V(X^1)|$ , and so, by induction hypothesis,  $X^1$  is either the set of even points of  $\mathbb{B}^{n-1}$  or the set of odd points of  $\mathbb{B}^{n-1}$ . But  $X^0$  and  $X^1$  are disjointed: for all  $x \in \mathbb{B}^n$ , if  $x_{-n} \in X^0 \cap X^1$ , then  $x^{n0}$  and  $x^{n1}$  are two points of  $X$ , and  $x^{n1} = x^{n0} \oplus e_n \in V(X)$ , a contradiction. So if  $X^0$  is the set of even (odd) points of  $\mathbb{B}^{n-1}$ , then  $X^1$  is the set of odd (even) points of  $\mathbb{B}^{n-1}$ , and we deduce that  $X$  is the set of even (odd) points of  $\mathbb{B}^n$ .  $\square$

**Lemma 2** *Let  $f : \mathbb{B}^n \rightarrow \mathbb{B}^n$ . Suppose that the conjugate of every immediate subnetwork of  $f$  is a bijection. If the conjugate of  $f$  is not a bijection, then  $f$  is even or odd self-dual.*

**Proof:** Suppose that  $f$  satisfies the conditions of the statement, and that the conjugate  $\tilde{f}$  of  $f$  is not a bijection. Let  $\tilde{X} \subseteq \mathbb{B}^n$  be the image of  $\tilde{f}$ , and let  $X = \mathbb{B}^n \setminus \tilde{X}$ . Since  $\tilde{f}$  is not a bijection,  $X$  is not empty. We first prove the following property:

(\*) For every  $x \in X$  and  $i \in [n]$ , the preimage of  $x \oplus e_i$  by  $\tilde{f}$  is of cardinality two.

Let  $x \in X$  and  $i \in [n]$ . By hypothesis, the conjugate of  $f^{i0}$  is a bijection, so there exists a unique point in  $\mathbb{B}^{n-1}$  whose the image by  $\tilde{f}^{i0}$  is  $x_{-i}$ . We deduce that there exists a unique point  $y \in \mathbb{B}^n$  such that  $y_i = 0$  and  $\tilde{f}^{i0}(y_{-i}) = x_{-i}$ . Then,  $\tilde{f}(y)_{-i} = \tilde{f}(y^{i0})_{-i} = \tilde{f}^{i0}(y_{-i}) = x_{-i}$ . We deduce that either  $\tilde{f}(y) = x$  or  $\tilde{f}(y) = x \oplus e_i$ . Since  $x \in X$  we have  $\tilde{f}(y) \neq x$  so  $\tilde{f}(y) = x \oplus e_i$ . Hence, we have proved that there exists a unique point  $y \in \mathbb{B}^n$  such that  $y_i = 0$  and  $\tilde{f}(y) = x \oplus e_i$ , and we prove with similar arguments that there exists a unique point  $z \in \mathbb{B}^n$  such that  $z_i = 1$  and  $\tilde{f}(z) = x \oplus e_i$ . This proves (\*).

We are now in position to prove that  $f$  is even or odd. Let  $V(X) = \{x \oplus e_i \mid x \in X, i \in [n]\}$ . We have

$$|X| + |\tilde{X}| = 2^n = |\tilde{f}^{-1}(\tilde{X})| = |\tilde{f}^{-1}(V(X))| + |\tilde{f}^{-1}(\tilde{X} \setminus V(X))| \geq |\tilde{f}^{-1}(V(X))| + |\tilde{X} \setminus V(X)|.$$

Following (\*), we have  $|\tilde{f}^{-1}(V(X))| = 2|V(X)|$  and  $V(X) \subseteq \tilde{X}$ , so

$$|X| + |\tilde{X}| \geq 2|V(X)| + |\tilde{X} \setminus V(X)| = 2|V(X)| + |\tilde{X}| - |V(X)| = |V(X)| + |\tilde{X}|.$$

Therefore,  $|X| \geq |V(X)|$ , and since  $V(X) \subseteq \tilde{X} = \mathbb{B}^n \setminus X$ , we have  $X \cap V(X) = \emptyset$ . So according to Lemma 1,  $X$  is either the set of even points of  $\mathbb{B}^n$  or the set of odd points of  $\mathbb{B}^n$ . We deduce that in the first (second) case,  $\tilde{X}$  is the set of odd (even) points of  $\mathbb{B}^n$ . Thus,  $f$  is even or odd.

It remains to prove that  $f$  is self-dual. Let  $x \in \mathbb{B}^n$ . For all  $i \in [n]$ , since  $\|\tilde{f}(x)\|$  and  $\|\tilde{f}(x) \oplus e_i\|$  have not the same parity, and since  $f$  is even or odd, we have  $\tilde{f}(x) \oplus e_i \in X$ . Thus, according to (\*), the preimage of  $(\tilde{f}(x) \oplus e_i) \oplus e_i = \tilde{f}(x)$  by  $\tilde{f}$  is of cardinality two. Consequently, there exists a point  $y \in \mathbb{B}^n$ , distinct from  $x$ , such that  $\tilde{f}(y) = \tilde{f}(x)$ . Let us prove that  $x = y \oplus 1$ . Indeed, if  $x_i = y_i = 0$  for some  $i \in [n]$ , then  $\tilde{f}^{i0}(x_{-i}) = \tilde{f}^{i0}(y_{-i}) = \tilde{f}(y)_{-i} = \tilde{f}^{i0}(y_{-i})$ . Since  $x \neq y$ , we deduce that  $\tilde{f}^{i0}$  is not a bijection, a contradiction. We show similarly that if  $x_i = y_i = 1$ , then  $\tilde{f}^{i1}$  is not a bijection. So  $x = y \oplus 1$ . Consequently,  $\tilde{f}(x \oplus 1) = \tilde{f}(x)$ , and we deduce that  $f$  is self-dual.  $\square$

**Proof of Theorem 1:** by induction on  $n$ . The case  $n = 1$  is obvious. So suppose that  $n > 1$  and that the theorem holds for the dimensions less than  $n$ . Suppose that  $f$  has no even or odd self-dual subnetwork. Under this condition,  $f$  is neither even self-dual nor odd self-dual (since  $f$  is a subnetwork of  $f$ ), and every immediate subnetwork of  $f$  has no even or odd self-dual subnetwork. So, by induction hypothesis, the dual of every strict subnetwork of  $f$  is a bijection, and we deduce from Lemma 2 that the dual of  $f$  is a

bijection. Thus, in particular, there exists a unique point  $x \in \mathbb{B}^n$  such that  $\tilde{f}(x) = 0$ , and since  $\tilde{f}(x) = 0$  if and only if  $f(x) = x$ , this point  $x$  is the unique fixed point of  $f$ .  $\square$

Clearly, if  $f$  has no even or odd self-dual subnetwork, then every subnetwork of  $f$  has no even or odd self-dual subnetwork, and according to Theorem 1, the conjugate of every subnetwork of  $f$  is a bijection. Conversely, if the conjugate of every subnetwork of  $f$  is a bijection, then  $f$  has no even or odd self-dual subnetwork, since the conjugate of an even or odd self-dual network is not a bijection. Consequently, we have the following characterization:

**Corollary 1** *The conjugate of each subnetwork of  $f$  is a bijection if and only if  $f$  has no even or odd self-dual network.*

**Example 1**  $f : \mathbb{B}^3 \rightarrow \mathbb{B}^3$  is defined by:

$$f(x_1, x_2, x_3) = (\overline{x_2} \wedge x_3, \overline{x_3} \wedge x_1, \overline{x_1} \wedge x_2).$$

*Remark that  $f$  is not self-dual, since  $f(000) = f(111) = 000$ . The immediate subnetworks of  $f$  are:*

$$\begin{aligned} f^{10}(x_2, x_3) &= (0, x_2) \\ f^{11}(x_2, x_3) &= (\overline{x_3}, 0) \\ f^{20}(x_1, x_3) &= (x_3, 0) \\ f^{21}(x_1, x_3) &= (0, \overline{x_1}) \\ f^{30}(x_1, x_2) &= (0, x_1) \\ f^{31}(x_1, x_2) &= (\overline{x_2}, 0) \end{aligned}$$

*So each immediate subnetwork  $f^{i\alpha}$  of  $f$  has one component fixed to zero, and so is not self-dual. Furthermore, each immediate subnetwork of  $f^{i\alpha}$  is the one dimensional network  $h$  defined by  $h(0) = h(1) = 0$ , which is not self-dual. So  $f$  has no self-dual subnetwork, and we deduce from Theorem 1 that the conjugate of  $\tilde{f}$  of  $f$  is a bijection, and that  $f$  has a unique fixed point. Indeed:*

$x$	$f(x)$	$\tilde{f}(x)$
000	000	000
001	100	101
010	001	011
011	001	010
100	010	110
101	100	001
110	010	100
111	000	111

## 4 Remarks on the theorem of Shih and Dong

In this section, we show that Theorem 1 implies a fixed point theorem due to Shih and Dong (2005). In order to state this theorem, we need additional definitions. Let

$$f : \mathbb{B}^n \rightarrow \mathbb{B}^n, \quad f(x) = (f_1(x), \dots, f_n(x)).$$

**Definition 5 (Discrete Jacobian matrix)** *The discrete Jacobian matrix of  $f$  evaluated at point  $x \in \mathbb{B}^n$  is the following  $n \times n$  Boolean matrix*

$$f'(x) = (f_{ij}(x)), \quad f_{ij}(x) = f_i(x^{j1}) \oplus f_i(x^{j0}) \quad (i, j \in [n]).$$

In the next definition, we represent  $f'(x)$  under the form of a directed graph, in order to use graph theoretic notions instead of matrix theoretical notions. In fact, we mainly focus on *elementary directed cycles*, that we simply call *cycles* in the following.

**Definition 6 (Local interaction graph)** *The local interaction graph of  $f$  evaluated at point  $x \in \mathbb{B}^n$  is the directed graph  $Gf(x)$  defined by: the vertex set is  $[n]$ , and for all  $i, j \in [n]$ , there exists an arc  $j \rightarrow i$  if and only if  $f_{ij}(x) = 1$ .*

The discrete Jacobian matrix of  $f$  was first defined by Robert (1983), who also introduced the notion of Boolean eigenvalue. This material allowed Shih and Ho (1999) to state a combinatorial analog of the Jacobian conjecture: *if  $f$  has the property that, for each  $x \in \mathbb{B}^n$ , all the boolean eigenvalues of  $f'(x)$  are zero, then  $f$  has a unique fixed point.* This conjecture was proved by Shih and Dong (2005). Since Robert proved that all the boolean eigenvalues of  $f'(x)$  are zero if and only if  $Gf(x)$  has no cycle, the theorem of Shih and Dong can be stated as follows.

**Theorem 2 (Shih and Dong (2005))** *If  $Gf(x)$  has no cycle  $\forall x \in \mathbb{B}^n$ , then  $f$  has a unique fixed point.*

A short prove of this theorem, independent of Theorem 1, is given in appendix. In the following of this section, we show, using Theorem 1, that the condition “*if  $Gf(x)$  has no cycle for all  $x$* ” can be weakened into a condition of the form “*if there exists “few” point  $x$  such that  $Gf(x)$  has a “short” cycle*”. The exact statement is given after the following proposition.

**Proposition 1** *If  $f$  is even or odd, then for every  $x \in \mathbb{B}^n$  the out-degree of each vertex of  $Gf(x)$  is odd. In particular,  $Gf(x)$  has a cycle.*

**Proof:** The out-degree  $d_j^+$  of any vertex  $j$  of  $Gf(x)$ , which equals the number of ones in the  $j$ th column of  $f'(x)$ , is  $d_j^+ = \|f(x^{j1}) \oplus f(x^{j0})\| = \|f(x) \oplus f(x \oplus e_j)\|$ . Since

$$\|f(x) \oplus f(x \oplus e_j)\| = \|(x \oplus \tilde{f}(x)) \oplus ((x \oplus e_j) \oplus \tilde{f}(x \oplus e_j))\| = \|\tilde{f}(x) \oplus \tilde{f}(x \oplus e_j) \oplus e_j\|,$$

the parity of  $d_j^+$  is the parity of  $\|\tilde{f}(x)\| + \|\tilde{f}(x \oplus e_j)\| + 1$ . Hence, if  $f$  is even or odd, then  $\|\tilde{f}(x)\| + \|\tilde{f}(x \oplus e_j)\|$  is even, and  $d_j^+$  is odd.  $\square$

**Corollary 2 (Extension of Shih-Dong’s fixed point theorem)** *If for  $k = 1, \dots, n$ , there exists at most  $2^k - 1$  points  $x \in \mathbb{B}^n$  such that  $Gf(x)$  has a cycle of length at most  $k$ , then the conjugate of  $f$  is a bijection. In particular,  $f$  has a unique fixed point.*

**Proof:** According to Theorem 1, it is sufficient to prove, by induction on  $n$ , that if  $f$  satisfies the conditions of the statement, then  $f$  has no even or odd self-dual subnetwork. The case  $n = 1$  is obvious. Suppose that  $n > 1$  and that  $f$  satisfies the conditions of the statement. Let  $i, j \in [n - 1]$ . For each  $x \in \mathbb{B}^n$  such that  $x_n = 0$ , we have

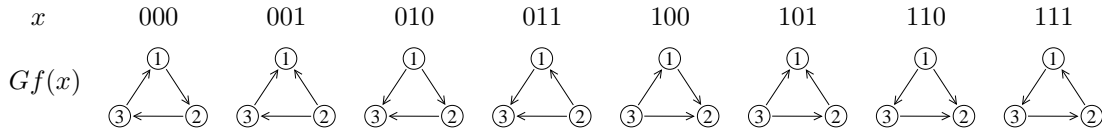
$$f_{ij}^{n0}(x_{-n}) = f_i^{n0}(x_{-n}^{j1}) \oplus f_i^{n0}(x_{-n}^{j0}) = f_i(x^{j1}) \oplus f_i(x^{j0}) = f_{ij}(x).$$

So  $Gf^{n0}(x_{-n})$  is the subgraph of  $Gf(x)$  induced by  $[n-1]$ , and we deduce that  $f^{n\alpha}$  satisfies the condition of the theorem (for every  $k \in [n-1]$ , there exists at most  $2^k - 1$  points  $x \in \mathbb{B}^{n-1}$  such that  $Gf^{n0}(x)$  has a cycle of length at most  $k$ ). Thus, by induction hypothesis,  $f^{n0}$  has no even or odd self-dual subnetwork. More generally, we prove with similar arguments, that for all  $i \in [n]$ ,  $f^{i0}$  and  $f^{i1}$  have no even or odd self-dual subnetwork. So  $f$  has no odd or even self-dual strict subnetwork. If  $f$  is itself even or odd self-dual, then by Proposition 1,  $Gf(x)$  has a cycle for every  $x \in \mathbb{B}^n$ , so  $f$  does not satisfy that conditions of the statement (for  $k = n$ ). Therefore,  $f$  has no even or odd self-dual subnetwork.  $\square$

**Example 2 (Continuation of Example 1)** Take again

$$f(x_1, x_2, x_3) = (\overline{x_2} \wedge x_3, \overline{x_3} \wedge x_1, \overline{x_1} \wedge x_2).$$

We have seen that  $f$  has no self-dual subnetwork. So it satisfies the conditions of Theorem 1, but not the conditions of Shih-Dong's theorem, since  $Gf(000)$  and  $Gf(111)$  have a cycle:

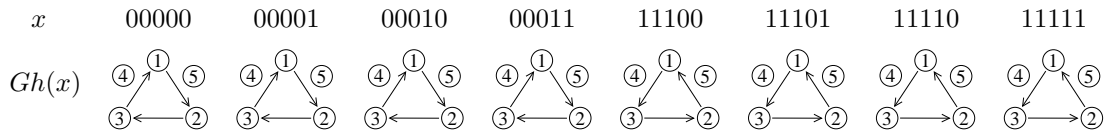


However,  $f$  satisfies the condition of Corollary 2 (there is  $0 < 2^1$  point  $x$  with a cycle of length at most 1;  $0 < 2^2$  point  $x$  such that  $Gf(x)$  has a cycle of length at most 2, and  $2 < 2^3$  points  $x$  such that  $Gf(x)$  has a cycle of length at most 3).

Now, consider the following "extension"  $h : \mathbb{B}^5 \rightarrow \mathbb{B}^5$  of  $f$ :

$$h(x_1, x_2, x_3, x_4, x_5) = (\overline{x_2} \wedge x_3, \overline{x_3} \wedge x_1, \overline{x_1} \wedge x_2, 0, 0) = (f(x_1, x_2, x_3), 0, 0)$$

Using the fact that  $f$  has no self-dual subnetwork, it's easy to see that  $h$  has no self-dual subnetwork. So  $h$  satisfies the conditions of Theorem 1. But it does not satisfy the conditions of Corollary 2. Indeed, there exists  $2^3$  points  $x$  such that  $Gh(x)$  has a cycle of length at most 3:



## 5 Remarks on asynchronous state graphs

In the following definition, we associate with  $f : \mathbb{B}^n \rightarrow \mathbb{B}^n$  a directed graph on  $\mathbb{B}^n$ , called the *asynchronous state graph* of  $f$ , which has been proposed by Thomas (1973) as a model for the dynamics of gene networks; see also Thomas and d'Ari (1990).

**Definition 7 (Asynchronous state graphs)** The asynchronous state graph of  $f$  is the directed graph  $\Gamma(f)$  defined by: the vertex set is  $\mathbb{B}^n$ , and for every  $x, y \in \mathbb{B}^n$ , there exists an arc  $x \rightarrow y$  if and only if there exists  $i \in [n]$  such that  $y = x \oplus e_i$  and  $f_i(x) \neq x_i$ .



Remark that  $\Gamma(f)$  and  $f$  share the same information. Remark also that for every  $i \in [n]$  and  $\alpha \in \mathbb{B}$ ,  $\Gamma(f^{i\alpha})$  is isomorphic to the subgraph of  $\Gamma(f)$  induced by the set of points  $x \in \mathbb{B}^n$  such that  $x_i = \alpha$ . Indeed: for every  $x, y \in \mathbb{B}^n$ ,

$$\begin{aligned} x_{-i} \rightarrow y_{-i} \text{ is an arc of } \Gamma(f^{i\alpha}) &\iff \exists j \neq i \text{ such that } y_{-i} = x_{-i} \oplus e_j \text{ and } f_j^{i\alpha}(x_{-i}) \neq x_j \\ &\iff \exists j \neq i \text{ such that } y^{i\alpha} = x^{i\alpha} \oplus e_j \text{ and } f_j(x^{i\alpha}) \neq x_j \quad (\star) \\ &\iff x^{i\alpha} \rightarrow y^{i\alpha} \text{ is an arc of } \Gamma(f). \end{aligned}$$

**Corollary 3** *If  $f$  has no even or odd self-dual subnetwork, then  $f$  has a unique fixed point  $x$ , and for all  $y \in \mathbb{B}^n$ ,  $\Gamma(f)$  contains a directed path from  $y$  to  $x$  of length  $\|x \oplus y\|$ .*

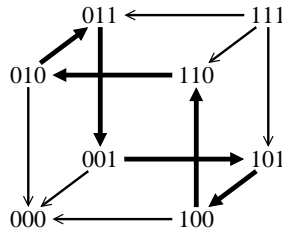
By the definition of  $\Gamma(f)$ , a path from  $x$  to  $y$  cannot be of length strictly less than  $\|x \oplus y\|$ ; a path from  $x$  to  $y$  of length  $\|x \oplus y\|$  can thus be seen has a *shortest path*.

**Proof of Corollary 3:** by induction on  $n$ . The case  $n = 1$  is obvious, so suppose that  $n > 1$  and that the corollary holds for the dimensions less than  $n$ . Let  $f : \mathbb{B}^n \rightarrow \mathbb{B}^n$ , and suppose that  $f$  has no even or odd self-dual subnetwork. By Theorem 1,  $f$  has a unique fixed point  $x$ . Let  $y \in \mathbb{B}^n$ . Suppose first that there exists  $i \in [n]$  such that  $x_i = y_i = 0$ . Then  $x_{-i}$  is the unique fixed point of  $f^{i0}$ . So, by induction hypothesis,  $\Gamma(f^{i0})$  has a path from  $y_{-i}$  to  $x_{-i}$  of length  $\|x_{-i} \oplus y_{-i}\|$ . Since  $x_i = y_i = 0$ , we deduce from  $(\star)$  that  $\Gamma(f)$  has a path from  $y$  to  $x$  of length  $\|x_{-i} \oplus y_{-i}\| = \|x \oplus y\|$ . The case  $x_i = y_i = 1$  is similar. So, finally, suppose that  $y = x \oplus 1$ . Since  $y$  is not a fixed point, there exists  $i \in [n]$  such that  $f_i(y) \neq y_i$ . Then,  $\Gamma(f)$  has an arc from  $y$  to  $z = y \oplus e_i$ . So  $z_i = x_i$ , and as previously, we deduce that  $\Gamma(f)$  has a path from  $z$  to  $x$  of length  $\|x \oplus z\|$ . This path together with the arc  $y \rightarrow z$  forms a path from  $y$  to  $x$  of length  $\|x \oplus z\| + 1 = \|x \oplus y\|$ .  $\square$

According to  $(\star)$ , the asynchronous state graph of each subnetwork of  $f$  is a subgraph of asynchronous state graph of  $f$  induced by an hypercube contained in  $\mathbb{B}^n$ . Hence, one can see the asynchronous state graphs of the subnetworks of  $f$  as “dynamical modules” of asynchronous state graph of  $f$ . The previous corollary shows that if  $f$  has no even or odd self-dual subnetwork, then the asynchronous state graph of  $f$  is “simple”: it describes a “weak convergence” toward a unique fixed point. An interpretation is then that **the asynchronous state graphs of even and odd self-dual networks are “dynamical modules” that are necessary for the “emergence” of “complex” asynchronous behaviors.**

**Example 3 (Continuation of Example 1)** *Take again the 3-dimensional network  $f$  defined in Example 1, which has no self-dual subnetwork. The asynchronous state graph  $\Gamma(f)$  of  $f$  is the following:*

$x$	$f(x)$	$\tilde{f}(x)$
000	000	000
001	100	101
010	001	011
011	001	010
100	010	110
101	100	001
110	010	100
111	000	111



In agreement with Corollary 3, there exists, from any initial point, a shortest path leading to the unique fixed point of  $f$  (the point 000): the asynchronous state graph describes a “weak asynchronous convergence” (by shortest paths) toward a unique fixed point. However,  $\Gamma(f)$  has a cycle (of length 6), so every path does not lead to the unique fixed point: the condition “has no even or odd self-dual subnetworks” does not ensure a “strong asynchronous convergence” toward a unique fixed point.

## 6 Remarks on positive and negative cycles

In this section, we show that positive (negative) circular networks, *i.e.* Boolean networks whose the global interaction graph reduces to a positive (negative) cycle, are simple instances of even (odd) circular networks. From this fact and existing results about positive and negative cycles, we will see that natural ideas of generalizations of Theorem 1 arise, but that none of these generalizations is true.

Let us begin with additional definitions. A *signed directed graph* is a directed graph in which each arc is either *positive*, *negative* or *unsigned*. In such a graph, a cycle is *positive (negative)* if it contains an unsigned arc or an even (odd) number of negative arcs (a directed cycle may be both positive and negative).

**Definition 8 (Global interaction graph)** *The global interaction graph of  $f : \mathbb{B}^n \rightarrow \mathbb{B}^n$  is the signed directed graph  $G(f)$  defined by: the vertex set is  $[n]$ , and for all  $i, j \in [n]$ , there exists an arc  $i \rightarrow j$  if and only if  $f_i(x^{j1}) \neq f_i(x^{j0})$  for at least one  $x \in \mathbb{B}^n$ ; and an arc  $j \rightarrow i$  of  $G(f)$  is: positive if  $f_i(x^{j1}) \geq f_i(x^{j0})$  for all  $x \in \mathbb{B}^n$ ; negative if  $f_i(x^{j1}) \leq f_i(x^{j0})$  for all  $x \in \mathbb{B}^n$ ; and unsigned in the other cases.*

Remark that  $G(f)$  has an arc  $j \rightarrow i$  if and only if  $f_i$  depends on the  $j$ th variable  $x_j$  (and that  $f_i(x^{j1}) \neq f_i(x^{j0})$  if and only if  $f_{ij}(x) = 1$ ).

**Definition 9 (Positive and negative circular networks)**  *$f$  is a positive (negative) circular network if  $G(f)$  is a positive (negative) cycle.*

The dynamics of positive and negative circular networks has been widely studied; see Remy et al. (2003) and Demongeot et al. (2010). Here, we prove that they are simple instances of even and odd self-dual networks.

**Proposition 2** *Every positive (negative) circular network is even (odd) and self-dual.*

**Proof:** Let  $f$  be a circular network. Without loss of generality, suppose that the  $n$  arcs of  $G(f)$  are  $i + 1 \rightarrow i$  for all  $i \in [n]$ ;  $n + 1$  being identified to 1 (here and in the rest of the proof). Then  $f_i$  depends only on  $x_{i+1}$ , so either  $f_i(x) = x_{i+1}$  (and  $i + 1 \rightarrow i$  is positive), or  $f_i(x) = x_{i+1} \oplus 1$  (and  $i + 1 \rightarrow i$  is negative); in the first case, we set  $s_i = 0$ , and in the second case, we set  $s_i = 1$  (so that  $f_i(x) = x_{i+1} \oplus s_i$  in both cases). Let  $s = (s_1, \dots, s_n) \in \mathbb{B}^n$ . By construction,  $f$  is positive if  $\|s\|$  is even, and negative if  $\|s\|$  is odd. Furthermore,

$$f(x) = (x_2, x_3, \dots, x_n, x_1) \oplus s \quad \forall x \in \mathbb{B}^n.$$

Hence

$$f(x \oplus 1) = (x_2 \oplus 1, \dots, x_n \oplus 1, x_1 \oplus 1) \oplus s = (x_2, \dots, x_n, x_1) \oplus 1 \oplus s = f(x) \oplus 1.$$

So  $f$  is self-dual. Also, we have  $\tilde{f}(x) = x \oplus (x_2, \dots, x_n, x_1) \oplus s$  so the parity of  $\tilde{f}(x)$  is the parity of  $\|x\| + \|(x_2, \dots, x_n, x_1)\| + \|s\|$ . Since  $\|x\| = \|(x_2, \dots, x_n, x_1)\|$ , we deduce that the parity of  $\tilde{f}(x)$  is the parity of  $\|s\|$ . So if  $f$  is positive (negative) then the image of  $\tilde{f}$  only contains even (odd) points.

It remains to prove that if  $f$  is positive (negative) then each even (odd) point is in the image of  $\tilde{f}$ . Suppose that  $f$  is positive (negative), and let  $z$  be an even (odd) point of  $\mathbb{B}^n$ . Let  $x \in \mathbb{B}^n$  be recursively defined by

$$x_1 = z_n, \quad x_{i+1} = z_i \oplus s_i \oplus x_i \quad \text{for all } i \in [n-1].$$

Then, for every  $i \in [n-1]$ , we have

$$\tilde{f}_i(x) = x_i \oplus f_i(x) = x_i \oplus x_{i+1} \oplus s_i = x_i \oplus (z_i \oplus s_i \oplus x_i) \oplus s_i = z_i.$$

It remains to prove that  $\tilde{f}_n(x) = z_n$ . By the definition of  $x$ , we have

$$\begin{aligned} x_n &= (z_{n-1} \oplus s_{n-1}) \oplus x_{n-1} \\ &= (z_{n-1} \oplus s_{n-1}) \oplus (z_{n-2} \oplus s_{n-2}) \oplus x_{n-2} \\ &\quad \vdots \\ &= (z_{n-1} \oplus s_{n-1}) \oplus (z_{n-2} \oplus s_{n-2}) \oplus \dots \oplus (z_1 \oplus s_1) \oplus z_n \\ &= (z_1 \oplus z_2 \oplus \dots \oplus z_n) \oplus (s_1 \oplus s_2 \oplus \dots \oplus s_{n-1}). \end{aligned}$$

So  $z$  and  $(s_1, s_2, \dots, s_{n-1}, x_n)$  have the same parity, and since  $z$  and  $s$  have the same parity, we deduce that  $x_n = s_n$ . Thus  $\tilde{f}_n(x) = x_n \oplus f_n(x) = s_n \oplus x_1 \oplus s_n = x_1 = z_n$ , and we deduce that  $\tilde{f}(x) = z$ . So  $f$  is even (odd) self-dual.  $\square$

**Remark 1** *There are  $2^{n-1}!$   $n$ -dimensional even (odd) self-dual networks, but “only”  $(n-1)!2^{n-1}$   $n$ -dimensional positive (negative) circular networks. Since  $2^{n-1}! = (n-1)!2^{n-1}$  for  $n = 1, 2$ , we deduce that every one or two-dimensional even (odd) self-dual network is a positive (negative) circular network.*

Since the class of positive and negative circular networks is contained in the class of even and odd self-dual networks, it is natural to think about the following generalization of Theorem 1: *if  $f$  has no positive or negative circular networks, then  $f$  has a unique fixed point.* However, this is false, as showed by the following example. Hence, **Theorem 1 becomes false if “has no even or odd self-dual subnetwork” is replaced by “has no positive or negative circular subnetwork”.**

**Example 4**  $f : \mathbb{B}^4 \rightarrow \mathbb{B}^4$  is defined by

$$\begin{aligned} f_1(x) &= (\overline{x_2} \wedge x_3 \wedge \overline{x_4}) \vee ((\overline{x_2} \vee x_3) \wedge x_4) \\ f_2(x) &= (\overline{x_3} \wedge x_1 \wedge \overline{x_4}) \vee ((\overline{x_3} \vee x_1) \wedge x_4) \\ f_3(x) &= (\overline{x_1} \wedge x_2 \wedge \overline{x_4}) \vee ((\overline{x_1} \vee x_2) \wedge x_4) \\ f_4(x) &= (x_2 \wedge x_3 \wedge \overline{x_1}) \vee ((x_2 \vee x_3) \wedge x_1) \end{aligned}$$

The table of  $f$  and  $\tilde{f}$ , and the asynchronous state graph of  $f$  are as follow:

$x$	$f(x)$	$\tilde{f}(x)$
0000	0000	0000
0010	1000	1010
0100	0010	0110
0110	0011	0101
1000	0100	1100
1010	1001	0011
1100	0101	1001
1110	0001	1111
0001	1110	1111
0011	1010	1001
0101	0110	0011
0111	1011	1100
1001	1100	0101
1011	1101	0110
1101	0111	1010
1111	1111	0000

One can see that  $f$  is even self-dual. The immediate subnetworks of  $f$  are the following:

$$\begin{aligned}
 f^{10}(x_2, x_3, x_4) &= (\overline{x_3} \wedge x_4, x_2 \vee x_4, x_2 \wedge x_3) \\
 f^{11}(x_2, x_3, x_4) &= (\overline{x_3} \vee x_4, x_2 \wedge x_4, x_2 \vee x_3) \\
 f^{20}(x_1, x_3, x_4) &= (x_3 \vee x_4, \overline{x_1} \wedge x_4, x_3 \wedge x_1) \\
 f^{21}(x_1, x_3, x_4) &= (x_3 \wedge x_4, \overline{x_1} \vee x_4, x_3 \vee x_1) \\
 f^{30}(x_1, x_2, x_4) &= (\overline{x_2} \wedge x_4, x_1 \vee x_4, x_2 \wedge x_1) \\
 f^{31}(x_1, x_2, x_4) &= (\overline{x_2} \vee x_4, x_1 \wedge x_4, x_2 \vee x_1) \\
 f^{40}(x_1, x_2, x_3) &= (\overline{x_2} \wedge x_3, \overline{x_3} \wedge x_1, \overline{x_1} \wedge x_2) \quad (\text{as in Examples 1-3}) \\
 f^{41}(x_1, x_2, x_3) &= (\overline{x_2} \vee x_3, \overline{x_3} \vee x_1, \overline{x_1} \vee x_2)
 \end{aligned}$$

Proceeding as in Example 1, one can check that none immediate subnetwork of  $f$  has a self-dual subnetwork (actually, it is sufficient to check this for each  $f^{i0}$  since  $f^{i1}(x) = f^{i0}(x \oplus 1) \oplus 1$ ,  $1 \leq i \leq 4$ ). So  $f$  has no circular strict subnetwork, and since  $f$  is not circular,  $f$  has no circular subnetwork, but it has not a unique fixed points. Note that for  $1 \leq i \leq 4$ , the 4-dimensional network  $h$  defined by  $h(x) = f(x) \oplus e_i$  is odd self-dual, has no circular subnetwork, and no fixed point.

Now, consider the following three fundamental theorems about cycles and fixed points (the last two theorems result from two conjectures of Thomas; see Remy et al. (2008); Richard (2010) and the references therein).

**Theorem 3 (Robert (1980))** *If  $G(f)$  has no cycle, then  $f$  has a unique fixed point.*

**Remark 2** *Clearly, each local interaction graph  $Gf(x)$  is a subgraph of the (unsigned version of the) global interaction graph  $G(f)$ . Hence, the condition “ $G(f)$  has no cycle” of Robert’s theorem is (much*

more) stronger than the condition “ $Gf(x)$  has no cycle for every  $x$ ” of Shih-Dong’s Theorem. Consequently, Shih-Dong’s theorem is a generalization of Robert’s theorem. Thus, Theorem 1 is also a generalization of Robert’s theorem.

**Remark 3** Actually, Robert proved, in Robert (1980) and Robert (1995), that if  $G(f)$  has no cycle, then  $f$  has a unique fixed point  $x$  and: (1) the synchronous iteration  $x^{t+1} = f(x^t)$  converges toward  $x$  in at most  $n$  steps for every initial point  $x^0 \in \mathbb{B}^n$ ; (2) every path of  $\Gamma(f)$  leads to  $x$  in at most  $n$  steps (“strong asynchronous convergence by shortest paths toward a unique fixed points”). These results shows the necessity of cycles for obtaining “complex” synchronous or asynchronous behaviors (e.g. multiple fixed points, cyclic attractors, long transient phases...).

**Theorem 4 (Remy et al. (2008))** If  $G(f)$  has no positive cycle, then  $f$  has at most one fixed point.

**Remark 4** Actually, by saying that an arc  $j \rightarrow i$  of  $Gf(x)$  is positive if  $f_i(x^{j1}) > f_i(x^{j0})$  and negative if  $f_i(x^{j1}) < f_i(x^{j0})$ , Remy et al. (2008) proved the following more general statement: if  $Gf(x)$  has no positive cycle for all  $x \in \mathbb{B}^n$ , then  $f$  has at most one fixed point.

**Theorem 5 (Richard (2010))** If  $G(f)$  has no negative cycle, then  $f$  has at least one fixed point.

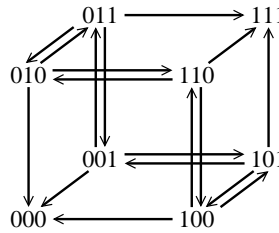
Hence, Theorems 4 and 5 give a nice proof “by dichotomy” of Robert’s theorem: the absence of positive cycle gives the uniqueness, and absence of negative cycle gives the existence. Seeing the relationship between positive (negative) circular networks and even (odd) self-dual networks, one may ask if a “proof by dichotomy” occurs for Theorem 1, i.e., if the absence of even self-dual subnetwork gives the uniqueness, and if the absence of odd self-dual network gives the existence. The following example shows that both cases are false. Hence: **if  $f$  has no even (odd) self-dual subnetworks, then it has not necessarily at most (at least) one fixed point.**

**Example 5**  $f : \mathbb{B}^3 \rightarrow \mathbb{B}^3$  is defined by

$$\begin{aligned} f_1(x) &= (\overline{x_1} \wedge (x_2 \vee x_3)) \vee (x_2 \wedge x_3) \\ f_2(x) &= (\overline{x_2} \wedge (x_3 \vee x_1)) \vee (x_3 \wedge x_1) \\ f_3(x) &= (\overline{x_3} \wedge (x_1 \vee x_2)) \vee (x_1 \wedge x_2) \end{aligned}$$

The table of  $f$  and  $\tilde{f}$ , and the asynchronous state graph of  $f$  are as follow:

$x$	$f(x)$	$\tilde{f}(x)$
000	000	000
001	110	111
010	101	111
011	100	111
100	011	111
101	010	111
110	001	111
111	111	000



$f$  is self-dual, but not even since  $\|f(001)\|$  is odd. The immediate subnetworks of  $f$  are:

$$\begin{aligned} f^{10}(x_2, x_3) &= (\overline{x_2} \wedge x_3, \overline{x_3} \wedge x_2) \\ f^{11}(x_2, x_3) &= (\overline{x_2} \vee x_3, \overline{x_3} \vee x_2) \\ f^{20}(x_1, x_3) &= (\overline{x_1} \wedge x_3, \overline{x_3} \wedge x_1) \\ f^{21}(x_1, x_3) &= (\overline{x_1} \vee x_3, \overline{x_3} \vee x_1) \\ f^{30}(x_1, x_2) &= (\overline{x_1} \wedge x_2, \overline{x_2} \wedge x_1) \\ f^{31}(x_1, x_2) &= (\overline{x_1} \vee x_2, \overline{x_2} \vee x_1) \end{aligned}$$

So each  $f^{i\alpha}$  is not circular, and according to Remark 1, it is not even and self-dual. Furthermore, each strict subnetwork  $h$  of  $f^{i\alpha}$  is either constant or defined by  $h(0) = 1$  and  $h(1) = 0$  (in the second case,  $h$  is odd and self-dual). So  $f^{i\alpha}$  has no strict even self-dual subnetwork. We deduce that  $f$  has no even self-dual subnetwork. But it has two fixed points.

Now consider the network  $f : \mathbb{B}^3 \rightarrow \mathbb{B}^3$  is defined by

$$\begin{aligned} f_1(x) &= \overline{x_2} \\ f_2(x) &= \overline{x_3} \\ f_3(x) &= (x_3 \wedge (\overline{x_1} \vee x_2)) \vee (\overline{x_1} \wedge x_2) \end{aligned}$$

The table of  $f$  and  $\tilde{f}$ , and the asynchronous state graph of  $f$  are as follow:

$x$	$f(x)$	$\tilde{f}(x)$
000	110	110
001	101	100
010	011	001
011	001	010
100	110	010
101	100	001
110	010	100
111	001	110

$f$  is self-dual, but not odd since  $\|f(000)\|$  is even. The immediate subnetworks of  $f$  are:

$$\begin{aligned} f^{10}(x_2, x_3) &= (\overline{x_3}, x_3 \vee x_2) \\ f^{11}(x_2, x_3) &= (\overline{x_3}, x_3 \wedge x_2) \\ f^{20}(x_1, x_3) &= (1, x_3 \wedge \overline{x_1}) \\ f^{21}(x_1, x_3) &= (0, x_3 \vee \overline{x_1}) \\ f^{30}(x_1, x_2) &= (\overline{x_2}, 1) \\ f^{31}(x_1, x_2) &= (\overline{x_2}, 0) \end{aligned}$$

So each  $f^{i\alpha}$  is not circular, and according to Remark 1, it is not odd and self-dual. Furthermore, each strict subnetwork  $h$  of  $f^{i\alpha}$  is either constant or defined by  $h(0) = 0$  and  $h(1) = 1$  (in the second case,  $h$  is even and self-dual). So  $f^{i\alpha}$  has no strict odd self-dual subnetwork. We deduce that  $f$  has no odd self-dual subnetwork. But it has no fixed point.

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## A A short proof of the theorem of Shih and Dong

The “trick” consists in proving, by induction on  $n$ , the following more general statement:

(\*) If  $Gf(x)$  has no cycle for all  $x \in \mathbb{B}^n$ , then the conjugate of  $f$  is a bijection (so that  $f$  has a unique fixed point).

The case  $n = 1$  is obvious, so suppose that  $n > 1$  and that (\*) holds for the dimensions less than  $n$ . Suppose that  $Gf(x)$  has no cycle for all  $x \in \mathbb{B}^n$ . Let  $i, j \in [n - 1]$ , and  $x \in \mathbb{B}^n$  such that  $x_n = 0$ . We have

$$f_{ij}^{n0}(x_{-n}) = f_i^{n0}(x_{-n}^{j1}) \oplus f_i^{n0}(x_{-n}^{j0}) = f_i(x^{j1}) \oplus f_i(x^{j0}) = f_{ij}(x).$$

So  $Gf^{n0}(x_{-n})$  is the subgraph of  $Gf(x)$  induced by  $[n - 1]$ , and thus, it has no cycle. We deduce that  $f^{n0}$  satisfies the conditions of (\*). Thus, by induction hypothesis, the conjugate of  $f^{n0}$  is a bijection. We prove with similar arguments that  $\tilde{f}^{i0}$  and  $\tilde{f}^{i1}$  are bijections for all  $i \in [n]$ .

Now, suppose, by contradiction, that  $\tilde{f}$  is not a bijection. Then, there exists two distinct points  $x, y \in \mathbb{B}^n$  such that  $\tilde{f}(x) = \tilde{f}(y)$ . Let us prove that  $x = y \oplus 1$ . Indeed, if  $x_i = y_i = \alpha$  for some  $i \in [n]$ , then  $\tilde{f}^{i\alpha}(x_{-i}) = \tilde{f}(x)_{-i} = \tilde{f}(y)_{-i} = \tilde{f}^{i\alpha}(y_{-i})$ . Thus  $\tilde{f}^{i\alpha}$  is not a bijection, a contradiction. So  $x = y \oplus 1$ . Since  $Gf(x)$  has no cycle, it contains at least one vertex of out-degree 0. In other words, there exists  $i \in [n]$  such that  $f(x^{i1}) = f(x^{i0})$ . Thus  $\tilde{f}(x^{i1})_{-i} = \tilde{f}(x^{i0})_{-i} = \tilde{f}(x)_{-i}$ . Hence, setting  $\alpha = y_i$ , we obtain

$$\tilde{f}^{i\alpha}(x_{-i}) = \tilde{f}(x^{i\alpha})_{-i} = \tilde{f}(x)_{-i} = \tilde{f}(y)_{-i} = \tilde{f}(y^{i1})_{-i} = \tilde{f}^{i1}(y_{-i}).$$

So  $\tilde{f}^{i\alpha}$  is not a bijection, a contradiction. Thus  $\tilde{f}$  is a bijection and (\*) is proved.

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