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# Unification and Matching in Hierarchical Combinations of Syntactic Theories

Serdar Erbatur<sup>1</sup>, Deepak Kapur<sup>2</sup>, Andrew M. Marshall<sup>3</sup>, Paliath Narendran<sup>4</sup>,  
and Christophe Ringeissen<sup>5</sup>

<sup>1</sup> Ludwig-Maximilians-Universität, München (Germany)

<sup>2</sup> University of New Mexico (USA)

<sup>3</sup> University of Mary Washington (USA)

<sup>4</sup> University at Albany, SUNY (USA)

<sup>5</sup> LORIA – INRIA Nancy-Grand Est (France)

**Abstract.** We investigate a hierarchical combination approach to the unification problem in non-disjoint unions of equational theories. In this approach, the idea is to extend a base theory with some additional axioms given by rewrite rules in such way that the unification algorithm known for the base theory can be reused without loss of completeness. Additional techniques are required to solve a combined problem by reducing it to a problem in the base theory. In this paper we show that the hierarchical combination approach applies successfully to some classes of syntactic theories, such as shallow theories since the required unification algorithms needed for the combination algorithm can always be obtained. We also consider the matching problem in syntactic extensions of a base theory. Due to the more restricted nature of the matching problem, we obtain several improvements over the unification problem.

## 1 Introduction

A critical question in matching and unification is how to obtain an algorithm for the combination of non-disjoint equational theories when there exist algorithms for the constituent theories. In recent work [7], a new “hierarchical” approach to the unification problem in the combination of non-disjoint theories is developed and classes of theories are identified for which the method can be applied. The main property of these classes is a hierarchical organization of the two equational theories  $E_1$  and  $E_2$ , where  $E_1$  is a set of axioms extending a base theory  $E_2$ . The method is successful in providing a unification method to the combination of theories for which no previous combination method was applicable. However, the main difficulty in applying the new combination method is that a new type of unification algorithm, denoted by  $A_1$ , is required to incorporate the axioms of  $E_1$ . The  $A_1$  algorithm is not actually a black box unification algorithm. Rather the algorithm constructs a specific type of solved-form which has the property of reducing an  $E_1 \cup E_2$  unification problem to one or more  $E_2$ -unification problems. A general  $A_1$  method is developed in [7], based on the general  $E$ -unification methods studied in [10, 19]. However, as with the general unification methods,

there is no general termination proof for an arbitrary theory. Termination must be proven for  $A_1$  to satisfy the restrictions of [7].

Here we are able to overcome this limitation by showing that for some classes of theories, such as shallow equational theories, the unification *algorithms* for those theories can be used for the  $A_1$  algorithm. This replaces the need to construct a dedicated unification algorithm via the method presented in [7] and allows one to use an “off the shelf” available algorithm.

In this paper, we also consider the matching problem in this new hierarchical framework and obtain a new non-disjoint combination method for the matching problem. Due to the more restricted nature of the matching problem, the combination method introduced in [7] for unification can be simplified. For the matching problem, we are not restricted to shallow theories but syntactic theories [5, 11, 12] are allowed. By assuming a resolvent presentation of a syntactic theory  $E_1$ , we are able to construct in a modular way a matching algorithm for hierarchical combination. The matching algorithm we present can be seen as an extension of the one known for disjoint unions of syntactic (finite) theories [15] and a variation of the one existing for disjoint unions of regular theories [16, 17]. A preliminary version of this algorithm has been presented in [8]. Compared to [8], we now rely on the standard definition of syntactic theory.

Let us give a brief preview of the remaining portions of the paper. Section 2 presents the preliminary background material. Section 3 presents an overview of the hierarchical combination. The overview given here is actually an improvement via simplification to the original hierarchical presentation in [7]. Section 4 applies hierarchical combination to shallow theories. Section 5 extends the hierarchical approach from the unification problem to the matching problem by considering a syntactic theory  $E_1$ .

## 2 Preliminaries

We use the standard notation of equational unification [3] and term rewriting systems [1]. Given a first-order signature  $\Sigma$  and a (countable) set of variables  $V$ , the set of  $\Sigma$ -terms over variables  $V$  is denoted by  $T(\Sigma, V)$ . The set of variables in a term  $t$  is denoted by  $Var(t)$ . A term  $t$  is *ground* if  $Var(t) = \emptyset$ . For any position  $p$  in a term  $t$ ,  $t|_p$  denotes the subterm of  $t$  at position  $p$ , and  $t[u]_p$  denotes the term  $t$  in which  $t|_p$  is replaced by  $u$ . Given a set  $E$  of  $\Sigma$ -axioms (i.e., pairs of  $\Sigma$ -terms, denoted by  $l = r$ ), the *equational theory*  $=_E$  is the congruence closure of  $E$  under the law of substitutivity. For any  $\Sigma$ -term  $t$ , the equivalence class of  $t$  with respect to  $=_E$  is denoted by  $[t]_E$ . The  $E$ -free algebra over  $V$  is denoted by  $T(\Sigma, V)/=_E$ . By a slight abuse of terminology,  $E$  will be often called an equational theory. An axiom  $l = r$  is *variable-preserving* (also called *regular*) if  $Var(l) = Var(r)$ . An axiom  $l = r$  is *linear* (resp., *collapse-free*) if  $l$  and  $r$  are linear (resp. non-variable terms). An equational theory is *variable-preserving* (resp., *regular/linear/collapse-free*) if all its axioms are variable-preserving (resp., regular/linear/collapse-free). An equational theory  $E$  is *finite* if for each term  $t$ , there are only finitely many terms  $s$  such that  $t =_E s$ .

A theory  $E$  is *subterm collapse-free* if and only if for all terms  $t$  it is not the case that  $t =_E u$  where  $u$  is a strict subterm of  $t$ . A theory  $E$  is *syntactic* if it has a finite *resolvent presentation*  $S$ , that is a presentation  $S$  such that each equality  $s =_E t$  has an equational proof  $s \leftrightarrow_S^* t$  with at most one step  $\leftrightarrow_S$  applied at the root position. A theory  $E$  is *shallow* if variables can only occur at a position at most 1 in axioms of  $E$ . When  $E$  is both subterm collapse-free and shallow, variables can only occur at position 1 in axioms of  $E$ . Shallow theories are known to be syntactic theories [6]. Let us recall some results connecting syntactic theories and unification. First, shallow theories admit a mutation-based unification algorithm [6]. Second, finite syntactic theories admit a mutation-based matching algorithm [16]. We will reuse both of these algorithms in this paper. In addition, any collapse-free and finitary unifying theory is syntactic [12]. For instance the Associative-Commutative ( $AC$ ) theory is syntactic, as well as the  $\mathcal{E}_{AC}$  theory introduced in [9] as a finitary unifying theory of distributive exponentiation.

A  $\Sigma$ -equation is a pair of  $\Sigma$ -terms denoted by  $s =^? t$ . An  $E$ -unification problem is a set of  $\Sigma$ -equations,  $\mathcal{S} = \{s_1 =^? t_1, \dots, s_m =^? t_m\}$ . The set of variables of  $\mathcal{S}$  is denoted by  $Var(\mathcal{S})$ . When  $t_1, \dots, t_n$  are ground,  $\mathcal{S}$  is called a matching problem, also denoted by  $\{s_1 \leq^? t_1, \dots, s_m \leq^? t_m\}$ , where  $s_i \leq^? t_i$  denotes a match-equation. A solution to an  $E$ -unification problem  $\mathcal{S}$ , called an  $E$ -unifier, is a substitution  $\sigma$  such that  $s_i \sigma =_E t_i \sigma$  for all  $1 \leq i \leq m$ . A substitution  $\sigma$  is *more general modulo  $E$*  than  $\theta$  on a set of variables  $V$ , denoted as  $\sigma \leq_E^V \theta$ , if there is a substitution  $\tau$  such that  $x \sigma \tau =_E x \theta$  for all  $x \in V$ . Two substitutions  $\theta_1$  and  $\theta_2$  are *equivalent modulo  $E$*  on a set of variables  $V$ , denoted as  $\theta_1 \equiv_E^V \theta_2$ , if and only if  $x \theta_1 =_E x \theta_2$  for all  $x \in V$ . A *Complete Set of  $E$ -Unifiers* of  $\mathcal{S}$  is a set of substitutions denoted by  $CSU_E(\mathcal{S})$  such that each  $\sigma \in CSU_E(\mathcal{S})$  is an  $E$ -unifier of  $\mathcal{S}$ , and for each  $E$ -unifier  $\theta$  of  $\mathcal{S}$ , there exists  $\sigma \in CSU_E(\mathcal{S})$  such that  $\sigma \leq_E^{Var(\mathcal{S})} \theta$ . An inference rule  $\mathcal{S} \vdash \mathcal{S}'$  is sound (resp. complete) for  $E$ -unification if  $CSU_E(\mathcal{S}) \subseteq CSU_E(\mathcal{S}')$  (resp.  $CSU_E(\mathcal{S}) \supseteq CSU_E(\mathcal{S}')$ ). An inference system is sound (resp. complete) for  $E$ -unification if all its inference rules are sound (resp. complete). A set of equations is said to be in *dag-solved form* (or *d-solved form*) if they can be arranged as a list  $x_1 =^? t_1, \dots, x_n =^? t_n$  where (a) each left-hand side  $x_i$  is a distinct variable, and (b)  $\forall 1 \leq i \leq j \leq n: x_i$  does not occur in  $t_j$ . Each  $x_i$  in this case is called a *solved variable*. We call a term (e.g., a variable) *fresh* if it is created by applying an inference rule (or a unification algorithm) and did not previously exist.

Consider the union  $\Sigma_1 \cup \Sigma_2$  of two disjoint signatures  $\Sigma_1$  and  $\Sigma_2$ . A term  $t$  is called a  $\Sigma_i$ -rooted term if its root symbol is in  $\Sigma_i$ . Variables,  $\Sigma_i$ -terms and  $\Sigma_i$ -equations are  *$i$ -pure*. We also use the notion of an *alien subterm*. An alien subterm of a  $\Sigma_i$ -rooted term  $t$  is a  $\Sigma_j$ -rooted subterm  $s$  ( $i \neq j$ ) such that all superterms of  $s$  are  $\Sigma_i$ -rooted. A *purification* procedure can be defined that uses *variable abstraction* [2] to replace any alien subterm  $u$  in a term  $t$  by a fresh variable  $x$  and adds the equation  $x =^? u$ . Then, equations  $s =^? t$  between a 1-pure term  $s$  and a 2-pure term  $t$  are split into two pure equations  $x =^? s$  and  $x =^? t$  where  $x$  is a new variable. An equation between two variables is always both 1-pure and 2-pure. Given a set of pure equations  $P$ , we denote by  $P|_{\Sigma_i}$  the

set of  $i$ -pure equations in  $P$  (for  $i = 1, 2$ ), and we say that  $P$  is in  $\Sigma_i$ -solved form (partial solved form) if  $P|_{\Sigma_i}$  is in dag-solved form.

A *term rewriting system* (TRS) is a pair  $(\Sigma, R)$ , where  $\Sigma$  is the signature and  $R$  is a finite set of rewrite rules of the form  $l \rightarrow r$  such that  $l, r$  are  $\Sigma$ -terms,  $l$  is not a variable and  $\text{Var}(r) \subseteq \text{Var}(l)$ . Given a TRS  $R = (\Sigma, R)$ , the signature  $\Sigma$  is often partitioned into two disjoint sets  $\Sigma := C \uplus D$ , where  $D := \{f | f(t_1, \dots, t_n) \rightarrow r \in R\}$  and  $C := \Sigma \setminus D$ . Symbols in  $C$  are called *constructors*, and symbols in  $D$  are called *defined functions*. A term  $s$  *rewrites* to a term  $t$ , denoted by  $s \rightarrow_R t$  (or simply  $s \rightarrow t$ ), if there exist a position  $p$  of  $s$ ,  $l \rightarrow r \in R$ , and substitution  $\sigma$  such that  $s|_p = l\sigma$  and  $t = s[r\sigma]_p$ . A term  $s$  is a *normal form with respect to the relation  $\rightarrow_R$*  (or simply a normal form), if there is no term  $t$  such that  $s \rightarrow_R t$ . This notion is lifted to substitutions as follows: a substitution  $\sigma$  is *normalized* if, for every variable  $x$  in the domain of  $\sigma$ ,  $x\sigma$  is a normal form. A TRS  $R$  is *terminating* if there are no infinite reduction sequences with respect to  $\rightarrow_R$ . A TRS  $R$  is *confluent* if, whenever  $t \rightarrow_R s_1$  and  $t \rightarrow_R s_2$ , there exists a term  $w$  such that  $s_1 \rightarrow_R^* w$  and  $s_2 \rightarrow_R^* w$ . A confluent and terminating TRS is called *convergent*. In a convergent TRS, we have the existence and the unicity of  $R$ -normal forms, denoted by  $t \downarrow_R$  for any term  $t$ . We define an *inner constructor* to be a constructor  $f$  that satisfies the following additional restrictions: (1)  $f$  does not appear on the left-hand side on any rule in  $R$ ; (2)  $f$  does not appear as the root symbol on the right-hand side of any rule in  $R$ ; (3) there are no non-constant function symbols below  $f$  on the right-hand side of any rule in  $R$ . We consider a  $\Sigma_1 \cup \Sigma_2$ -theory  $E_1$  and a  $\Sigma_2$ -theory  $E_2$ . The  $\Sigma_1 \cup \Sigma_2$ -theory  $E_1$  is given by a  $\Sigma_1$ -rooted TRS  $R_1$  such that for each  $l \rightarrow r \in R_1$ ,  $l$  and  $r$  are  $\Sigma_1$ -rooted terms. Moreover,  $\Sigma_2$ -symbols do not occur in left-hand sides of  $R_1$ . We use the notion of *convergence modulo an equational theory*. When  $R_1$  is convergent modulo  $E_2$ , we have that for any terms  $s$  and  $t$ ,  $s =_{R_1 \cup E_2} t$  if and only if  $s \downarrow_{R_1} =_{E_2} t \downarrow_{R_1}$ .

### 3 Hierarchical Combination

Our hierarchical combination has been introduced in [7]. We present a simplified version in Definition 1. The motivation is to simplify the relationships between the theories and allows for an easier notation. However, the new definition follows the one presented in [7]. For completeness, we also repeat several results proven in [7] on hierarchical combination. These results will prove useful in showing the applicability of hierarchical combination to shallow theories, and then more generally to syntactic theories.

**Definition 1.** A hierarchical combination is a pair  $(E_1, E_2)$  such that

- $\Sigma_1 \cap \Sigma_2 = \emptyset$ ;
- $E_1$  is a subterm collapse-free equational  $\Sigma_1 \cup \Sigma_2$ -theory given by a TRS  $R_1$  which is  $\Sigma_1$ -rooted and convergent modulo  $E_2$ ;
- $E_2$  is a finite equational  $\Sigma_2$ -theory;
- $\Sigma_2$ -symbols are inner constructors in  $R_1$ .

A hierarchical combination  $(E_1, E_2)$  is finite (resp., shallow/syntactic) if  $E_1$  is finite (resp., shallow/syntactic). The equational theory of  $(E_1, E_2)$  is  $E_1 \cup E_2$ .

Since a finite theory is subterm collapse-free, we have that both  $E_1$  and  $E_2$  are subterm collapse-free.

**Proposition 1.** [7] *If  $(E_1, E_2)$  is a hierarchical combination, then  $E_1 \cup E_2$  is subterm collapse-free.*

Let us now introduce a key notion of great interest to relate combined equational proofs with pure ones. The notion of variable abstraction is widely used in the context of combined equational theories [2, 4, 18]. In the following, we use a slight adaptation to abstract ground terms by constants ( $\mathcal{C}$ ) and non-ground terms by variables ( $\mathcal{Y}$ ). According to our assumptions, theories are regular and so a ground term cannot be equal to a non-ground one.

**Definition 2 (Variable Abstraction).** *Let  $\mathcal{Y}$  be a countably infinite set of variables,  $\mathcal{C}$  be a countably infinite set of constants such that  $V$ ,  $\mathcal{Y}$  and  $\mathcal{C}$  are pairwise disjoint. We consider a bijection*

$$\pi : (T(\Sigma_1 \cup \Sigma_2 \cup V) \downarrow_{R_1} \setminus V) / =_{E_2} \longrightarrow \mathcal{Y} \cup \mathcal{C}$$

such that  $\pi([t \downarrow_{R_1}]_{E_2}) \in \mathcal{C}$  if and only if  $t$  is ground. For  $i = 1, 2$ , the  $\pi_i$ -abstraction of  $t$  is denoted by  $t^{\pi_i}$  and defined as follows:

- If  $t \in V$ , then  $t^{\pi_i} = t$ .
- If  $t = f(t_1, \dots, t_n)$  and  $f \in \Sigma_i$ , then  $t^{\pi_i} = f(t_1^{\pi_i}, \dots, t_n^{\pi_i})$ .
- Otherwise,  $t^{\pi_i} = \pi([t \downarrow_{R_1}]_{E_2})$ .

An inverse mapping of  $\pi$  is a mapping  $\pi^{-1} : \mathcal{Y} \cup \mathcal{C} \longrightarrow (T(\Sigma_1 \cup \Sigma_2 \cup V) \setminus V)$  such that  $\pi([\pi^{-1}(z) \downarrow_{R_1}]_{E_2}) = z$  for any  $z \in \mathcal{Y} \cup \mathcal{C}$ .

In a hierarchical combination, a key feature is the ability to reuse an  $E_2$ -unification algorithm, say  $A_2$ , to solve 2-pure unification problems without loss of completeness.

**Lemma 1.** [7] *Let  $(E_1, E_2)$  be any hierarchical combination.*

- For any terms  $s$  and  $t$ , if  $s \longleftrightarrow_{E_1 \cup E_2} t$ , then  $s^{\pi_2} =_{E_2} t^{\pi_2}$ .
- $E_2$ -unification is a sound and complete method to solve 2-pure  $E_1 \cup E_2$ -unification problems.

Let us now present the combination procedure.

### 3.1 Combination Procedure for Hierarchical Combination

We describe the combination procedure in an abstract way. By applying variable abstraction, we split the input set of equations  $P$  into two sets of 1-pure and 2-pure equations, denoted respectively by  $P_1$  and  $P_2$ . Two algorithms  $A_1$  and  $A_2$  aim at solving  $P_1$  and  $P_2$ :

- $A_1$  returns a  $\Sigma_1$ -solved form of 1-pure  $E_1 \cup E_2$ -unification problems. Thus,  $A_1$  is a special type of algorithm that returns a “partial” solved form. We address the problem of building such a procedure in Section 4.
- $A_2$  is an  $E_2$ -unification algorithm.

Lemma 1 shows that an  $E_2$ -unification algorithm is the right tool for solving 2-pure problems. For unification problems that are not 2-pure, the idea is to use a rule-based unification procedure  $A_1$ . The role of  $A_1$  is to reduce the problem to a form for which  $A_2$  can be applied. The approach taken is as follows:

1. Run  $A_1$  on  $P_1$  to obtain  $P'_1$ .
2. Run  $A_2$  on  $P_2$  to obtain  $P'_2$ .
3. Combine  $P'_1$  and  $P'_2$  into a set of equations  $P'$ .
4. If  $P'$  is not solved, then iterate the procedure with  $P'$  as input.

In comparison with the combination algorithm presented in [7], the above description represents the core procedure of hierarchical combination. The algorithm of [7] can be seen as a single iteration of the loop in the above procedure. To ensure a single iteration is sufficient there are some additional technical restrictions present in [7], such as restricting the occurrence of “ping-pong variables”. Additionally, some machinery, such as variable identification, is also required.

## 4 Unification in Shallow Hierarchical Combination

In this section we show that shallow theories are another class of theories for which a hierarchical combination approach can be applied. Let  $E_1$  be a shallow equational theory [6] and  $E_2$  an equational theory, both satisfying Definition 1. Thus,  $R_1$  is a rewrite system where variables occur in rules at depth 1 (a variable cannot occur at depth 0 since  $R_1$  is collapse-free). In this particular case, an inner constructor can only occur in a right hand side of a rule and as the root symbol of a ground flat term. By applying constant abstraction on right-hand sides of rules, we can build an “equivalent” shallow rewrite system which is disjoint from  $E_2$ . As shown below, we can rely on a unification algorithm for shallow equational theories [6, 13, 14] to build an  $A_1$  procedure.

Let  $R_1^\pi = \{l \rightarrow r^{\pi_1} \mid l \rightarrow r \in R_1\}$ . Obviously,  $R_1^\pi$  is also a shallow rewrite system.

**Lemma 2.** *Let  $s$  be a  $\Sigma_1$ -rooted term such that its alien subterms are  $R_1$ -normalized. If  $s \rightarrow_{R_1} t$ , then*

- $t$  is a  $\Sigma_1$ -rooted term such that its alien subterms are  $R_1$ -normalized,
- $s^{\pi_1} \rightarrow_{R_1^\pi} t^{\pi_1}$ .

*Proof.* Let  $s \rightarrow_{R_1} t$ .

1.  $E_1$  is subterm collapse free and  $\Sigma_2$  symbols in  $R_1$  appear as inner constructors. Thus, if  $s$  is  $\Sigma_1$ -rooted then  $t$  is  $\Sigma_1$ -rooted.

2. If  $s = s[l\sigma]_p \rightarrow_R s[r\sigma]_p = t$  for some  $R$ -normalized  $\sigma$  and position  $p$ , then  $p$  must not occur in an alien subterm of  $s$  since we assume  $s$  has  $R_1$ -normalized alien subterms. This implies,  $s^{\pi_1} = s^{\pi_1}[(l\sigma)^{\pi_1}]_p$  and  $t^{\pi_1} = s^{\pi_1}[(r\sigma)^{\pi_1}]_p$ . Let  $(\sigma)^{\pi_1} := \{x \mapsto (x\sigma)^{\pi_1} \mid x \in \text{Var}(l)\}$ . Note,  $l$  contains only  $\Sigma_1$  symbols and any  $\Sigma_2$  symbols in  $r$  must be roots of ground terms. Thus,  $(l\sigma)^{\pi_1} = l(\sigma)^{\pi_1}$  and  $(r\sigma)^{\pi_1} = r^{\pi_1}(\sigma)^{\pi_1}$ . This implies

$$s^{\pi_1} = s^{\pi_1}[(l\sigma)^{\pi_1}]_p = s^{\pi_1}[l(\sigma)^{\pi_1}]_p \rightarrow_{R_1^\pi} s^{\pi_1}[r^{\pi_1}(\sigma)^{\pi_1}]_p = s^{\pi_1}[(r\sigma)^{\pi_1}]_p = t^{\pi_1} \quad \square$$

**Lemma 3.** *Let  $s$  and  $t$  be  $\Sigma_1$ -rooted terms. If  $s =_{E_2} t$  then  $s^{\pi_1} = t^{\pi_1}$ .*

*Proof.* Let  $s =_{E_2} t$ . For each pair of alien subterms  $s' \in s$  and  $t' \in t$  such that  $s' \leftrightarrow_{E_2}^* t'$ ,  $\pi(s') = \pi(t')$ . In addition,  $s$  and  $t$  have the same root symbol. Therefore,  $s^{\pi_1} = t^{\pi_1}$ .  $\square$

**Proposition 2.** *Let  $s =^? t$  be a  $\Sigma_1$ -equation, and  $\sigma$  be a  $R_1$ -normalized substitution. If  $s\sigma =_{E_1 \cup E_2} t\sigma$ , then  $s\sigma^{\pi_1} =_{R_1^\pi} t\sigma^{\pi_1}$ .*

*Proof.* Since  $s\sigma =_{E_1 \cup E_2} t\sigma$ , we have  $s\sigma \rightarrow_{R_1}^* (s\sigma)\downarrow_{R_1} =_{E_2} (t\sigma)\downarrow_{R_1} \leftarrow_{R_1}^* t\sigma$ . By Lemma 2, we have  $s\sigma^{\pi_1} \rightarrow_{R_1^\pi}^* ((s\sigma)\downarrow_{R_1})^{\pi_1}$  and  $t\sigma^{\pi_1} \rightarrow_{R_1^\pi}^* ((t\sigma)\downarrow_{R_1})^{\pi_1}$ . By Lemma 3,  $((s\sigma)\downarrow_{R_1})^{\pi_1} = ((t\sigma)\downarrow_{R_1})^{\pi_1}$ , and so  $s\sigma^{\pi_1} =_{R_1^\pi} t\sigma^{\pi_1}$ .  $\square$

Proposition 2 shows that the problem of solving 1-pure equations can be reduced to a  $R_1^\pi$ -unification problem. Let us now show that the hierarchical combination approach initiated in [7] can be applied to the theories under consideration. It is shown in [7] that if three restrictions are satisfied by the combined theory, there exists a combined unification algorithm. We present those restrictions below, rephrased as to conform to the current presentation.

**Restriction 1** (*Algorithm  $A_1$* ) *Let  $P_1$  be a set of 1-pure equations. Algorithm  $A_1$  applied to  $P_1$  computes a set of problems  $\{\mathcal{Q}_k\}_{k \in K}$  such that*

$$\bigcup_{k \in K} CSU_{E_1 \cup E_2}(\mathcal{Q}_k) \text{ is a } CSU_{E_1 \cup E_2}(P_1) \text{ and for each } k \in K:$$

- (i)  $\mathcal{Q}_k$  is in  $\Sigma_1$ -solved form.
- (ii) No fresh variable occur under  $\Sigma_2$ -rooted terms.

**Restriction 2** (*Algorithm  $A_2$* )

*Algorithm  $A_2$  computes a finite complete set of  $E_1 \cup E_2$ -unifiers of 2-pure unification problems.*

**Restriction 3** (*Errors*)

- (i) A  $\Sigma_2$ -rooted term cannot be  $E_1 \cup E_2$ -equal to a  $\Sigma_1$ -rooted term.
- (ii)  $E_1 \cup E_2$  is subterm collapse-free. Therefore, an  $E_1 \cup E_2$ -unification problem including a cycle has no solution.

**Lemma 4.** *Shallow hierarchical combination  $(E_1, E_2)$  satisfies Restrictions (1) through (3).*



*Proof.* Let us consider the different restrictions.

- Restriction 1: Since  $R_1^\pi$  is shallow,  $R_1^\pi$ -unification is finitary and a unification algorithm for shallow theories is known. Given a set of  $\Sigma_1$ -equations  $P_1$ , it computes a finite complete set of  $R_1^\pi$ -unifiers of  $P_1$ , say  $\{\sigma_k\}_{k \in K}$ . By Proposition 2,  $\{\sigma_k \pi^{-1}\}_{k \in K}$  is a  $CSU_{E_1 \cup E_2}(P_1)$ . By purification, each  $\sigma_k$  can be equivalently written as a unification problem  $\mathcal{Q}_k$  such that  $\mathcal{Q}_k$  is in  $\Sigma_1$ -solved form, and  $\mathcal{Q}_k$  contains only ground  $\Sigma_2$ -rooted terms.
- Restriction 2: This restriction follows directly from Lemma 1.
- Restriction 3:  $E_1 \cup E_2$  is subterm collapse-free due to Proposition 1. Since  $E_1$  and  $E_2$  are both subterm collapse-free,  $\Sigma_2$ -rooted terms cannot be  $E_1 \cup E_2$ -equal to  $\Sigma_1$ -rooted terms without contradicting Lemma 1.  $\square$

Directly from Lemma 4 and [7] we have the following.

**Theorem 1.** *For any shallow hierarchical combination  $(E_1, E_2)$ , there exists a combined  $E_1 \cup E_2$ -unification algorithm, provided that an  $E_2$ -unification algorithm is known.*

*Proof.* In [6], unification is shown decidable and finitary in shallow theories. In addition, they provide a method for constructing a unification algorithm for an arbitrary shallow theory. Thus, we can construct a finitary unification algorithm,  $A_1$ , from the shallow equational theory of  $R_1^\pi$ . By Lemma 4, the result follows.  $\square$

*Example 1.* Consider the (ground) hierarchical combination  $(E_1, E_2)$  where the TRS  $R_1$  associated with  $E_1$  is  $\{f(a, b) \rightarrow g(a + b), f(b, a) \rightarrow g(b + a)\}$  and  $E_2 = \{x + y = y + x\}$ . The TRS  $R_1$  is convergent modulo  $E_2$ . The TRS  $R_1^\pi$  is  $\{f(a, b) \rightarrow g(c), f(b, a) \rightarrow g(c)\}$  where  $c = \pi([a + b]_{E_2})$ . A  $R_1^\pi$ -unification algorithm can be obtained by adding to a syntactic unification algorithm some mutation rules [6]. Then, most general  $E_1 \cup E_2$ -unifiers of a 1-pure unification problem can be derived from most general  $R_1^\pi$ -unifiers by replacing the constant  $c$  with  $a + b$ .

Shallow theories are particular syntactic theories admitting a terminating mutation-based unification procedure. In [13], Lynch and Morawska have investigated a larger class of syntactic theories admitting a terminating mutation-based unification procedure. In Section 5, we consider a mutation-based approach for the matching problem. This particular unification problem is particularly interesting with respect to our combination method since there will be no termination issue due to the fact that solutions are necessarily ground.

## 5 Matching in Finite Syntactic Hierarchical Combination

In this section we consider the matching problem in any finite syntactic hierarchical combination. Due to the more restricted nature of the matching problem

we obtain several improvements over the unification problem. One of the improvements is that we are able to relax several restrictions we assumed for the unification problem. In the unification setting it was necessary to restrict variables which could cause reapplication of the first unification algorithm, denoted as “ping pong” variables in [7]. This restriction can be easily fulfilled if most general solutions can be expressed without any new variable. Considering matching problems in regular theories, there are only ground solutions, and so no ping pong variables. Since subterm collapse-free theories are regular, the theories we are interested in seem well-suited for a hierarchical approach to the matching problem.

The combination algorithm we propose is similar to the one existing for matching in disjoint unions of regular theories [16, 17]. In that context, matching algorithms  $A_1$  and  $A_2$  are applied repeatedly until reaching normal forms that correspond to most general solutions. The key principle of the combination algorithm for matching is to purify only the left-hand sides of matching problems. Thus, this purification introduces a “pending” equation  $s =^? X$  that will lead to a match-equation: since  $X$  occurs in a match-equation solved by  $A_1$  or  $A_2$ ,  $X$  will be instantiated by a ground term, say  $u$ , transforming eventually  $s =^? X$  into a match-equation  $s \leq^? u$ .

**Definition 3.** *Let  $P$  be a conjunction of match-equations and equations.*

*The set of ground-solved variables in  $P$  is the smallest set of variables in  $P$  such that*

1. *variables occurring in left-hand sides of match-equations are ground-solved.*
2. *if  $t =^? t'$  is an equation in  $P$  such that variables in  $t'$  are ground-solved, then variables in  $t$  are also ground-solved.*

*When  $P$  contains only ground-solved variables, it is called an extended matching problem. A matching problem is an extended matching problem containing no equations. For any match-equation in an extended matching problem, we assume that the right-hand side is  $R_1$ -normalized. Hence, there will be an implicit normalization of right-hand sides when new match-equations are generated by some inference rules.*

In the following, we present a rule-based combination method to solve  $E_1 \cup E_2$ -matching problems. Let us briefly introduce the main steps of that combination method. On the one hand, solving match-equations with 2-pure left-hand sides generate solved match-equations. On the other hand, solving match-equations whose left-hand sides are  $\Sigma_1$ -rooted generate new equalities while preserving the syntactic form of an extended matching problem. Then, there are some combination rules to remove successively the equations introduced during the purification and the solving phases. Eventually, we obtain a matching problem in solved form. As shown below, this method is sound and complete when  $(E_1, E_2)$  is a finite syntactic hierarchical combination, which means that  $E_1, E_2$  are finite and  $E_1$  admits a resolvent presentation denoted by  $S_1$ . If  $E_1$  and  $E_2$  are finite, then  $E_1 \cup E_2$  is finite according to the result below.

**Proposition 3.**  $(E_1, E_2)$  is a finite hierarchical combination if and only if  $E_1 \cup E_2$  is finite.

*Proof.* (If direction) Straightforward.

(Only if direction) Assume that this is not the case, that is, there exists a term  $t \in T(\Sigma_1 \cup \Sigma_2, \mathcal{X})$  with an infinite set of terms  $S$  such that  $t =_{E_1 \cup E_2} s$  for each  $s \in S$ . Since there exists a convergent rewrite system  $R_1$  which is equivalent to  $E_1$ , we can consider the unique  $R_1$ -normal form of each term  $s \in S$ . Let  $S' = \{s \downarrow_{R_1} \mid s \in S\}$ . The set of terms  $S'$  is also infinite, otherwise it would contradict that  $E_1$  is finite. Since  $R_1$  is convergent modulo  $E_2$ ,  $t \downarrow_{R_1} =_{E_2} s'$  for each  $s' \in S'$ . This implies the existence of an infinite  $E_2$  equivalence class  $[t \downarrow_{R_1}]_{E_2}$ , which contradicts our assumption that  $E_2$  is finite.  $\square$

Since matching is finitary in a finite theory, we could take a brute force approach to construct an  $E_1 \cup E_2$ -matching algorithm [15]. Actually, a match-equation  $s \leq_{E_1 \cup E_2}^? t$  has the same set of (ground)  $E_1 \cup E_2$ -solutions as the syntactic matching problem

$$\bigvee_{\{t' \mid t' =_{E_1 \cup E_2} t\}} s \leq_{\emptyset}^? t'$$

where  $\emptyset$  denotes the empty equational theory. However, similarly to [15], we can also proceed in a modular way, by using some “pure” matching algorithms, say  $A_1$  and  $A_2$  dedicated respectively to  $E_1$  and  $E_2$ . In the context of matching,  $A_1$  aims at handling match-equations whose left-hand sides are  $\Sigma_1$ -rooted, whilst  $A_2$  denotes an  $E_2$ -matching algorithm to solve match-equations whose left-hand sides are 2-pure.

**Restriction 4**  $(E_1, E_2)$  is a finite syntactic hierarchical combination with the following constituent algorithms:

1.  $A_1$  is given by the mutation-based algorithm  $\mathcal{M}$  depicted in Fig. 1, where  $S_1$  denotes a resolvent presentation of the finite syntactic theory  $E_1$ .
2.  $A_2$  is an  $E_2$ -matching algorithm.

**Lemma 5.**  $E_2$ -matching is a sound and complete method to solve  $E_1 \cup E_2$ -matching problems whose left-hand sides are 2-pure, provided that the right-hand sides are replaced by their  $\pi_2$ -abstractions.

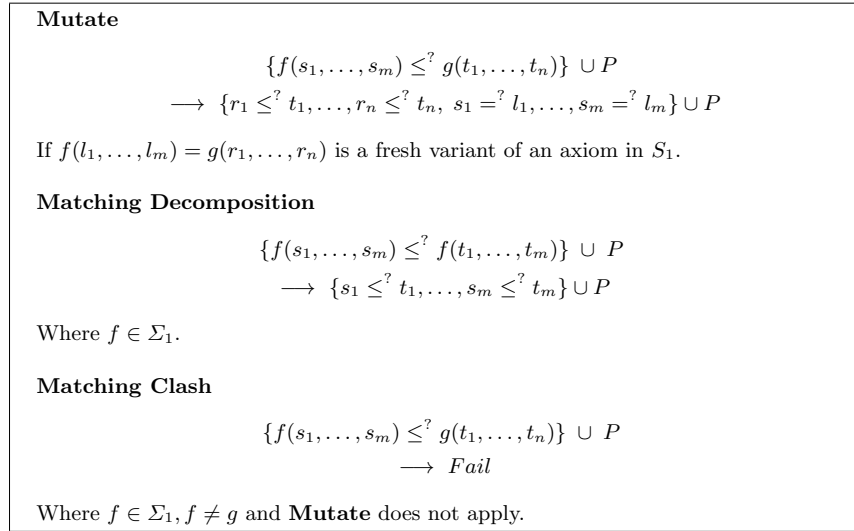
*Proof.* Let  $s \leq^? t$  be a match-equation such that  $s$  is 2-pure term, and consider the corresponding 2-pure match-equation  $s \leq^? t^{\pi_2}$ , where  $t^{\pi_2}$  can be effectively computed (up to a renaming). The soundness is obvious: any  $E_2$ -solution  $\sigma$  of  $s \leq^? t^{\pi_2}$  leads to an  $E_1 \cup E_2$ -solution  $\sigma\pi^{-1}$  of  $s \leq^? t$ . For the completeness, consider an  $E_1 \cup E_2$ -solution  $\sigma$  of  $s \leq^? t$  where  $s$  is 2-pure. By Lemma 1 (and induction on the length of derivation),  $s\sigma =_{E_1 \cup E_2} t$  implies  $s\sigma^{\pi_2} =_{E_2} t^{\pi_2}$  where  $\sigma =_{E_1 \cup E_2} \sigma^{\pi_2}\pi^{-1}$ .  $\square$

## 5.1 Mutation-based Procedure

We now consider the question of applying a mutation-based inference system to simplify match-equations whose left-hand sides are  $\Sigma_1$ -rooted. The mutation inference is known to be complete for syntactic theories [5, 11, 12], and so for  $E_1$ . The next lemma shows that this property is preserved when considering  $E_1 \cup E_2$ .

**Lemma 6.** *Let  $S_1$  be a resolvent presentation of  $E_1$ , and let  $s$  and  $t$  be  $\Sigma_1$ -rooted terms such that  $t$  is  $R_1$ -normalized. If  $s =_{E_1 \cup E_2} t$ , then there exists an equational proof of  $s =_{E_1 \cup E_2} t$  with at most one  $S_1$ -equational step applied at the root position.*

Lemma 6 justifies the eager  $R_1$ -normalization of the right-hand sides of match-equations, and it implies the completeness of the inference system  $\mathcal{M}$  (Fig. 1).



**Fig. 1.** Mutation-based inference system  $\mathcal{M}$  for Matching

**Lemma 7.**  *$\mathcal{M}$  (Fig. 1) is sound, complete and terminating.*

*Proof.* The soundness is straightforward. The completeness is a consequence of Lemma 6. Let us now prove the termination. The multiset consisting of the sizes of the right-hand sides of match-equations can be used as a complexity measure. The three rules of  $\mathcal{M}$  decrease this complexity measure, and so  $\mathcal{M}$  terminates.  $\square$

**Corollary 1.** *Let  $P$  be a set of match-equations whose left-hand sides are  $\Sigma_1$ -rooted. There exists a finite derivation  $P \vdash \dots \vdash P'$  in  $\mathcal{M}$  such that  $P$  and  $P'$  have the same set of solutions, and  $P'$  is in normal form w.r.t.  $\mathcal{M}$ .*

**Lemma 8.** *If  $P$  is an extended matching problem in normal form w.r.t.  $\mathcal{M}$ , then  $P$  does not contain match-equations whose left-hand sides are  $\Sigma_1$ -rooted.*

*Proof.* By contradiction. Let  $P_1$  be the set of match-equations in  $P$  whose left-hand sides are  $\Sigma_1$ -rooted. If  $P_1$  is non-empty, then some rule in  $\mathcal{M}$  must apply, which means that  $P$  is not a normal form w.r.t.  $\mathcal{M}$ .  $\square$

To conclude this section we introduce a complexity measure, which will be useful to prove the termination of the combined matching procedure (Fig. 2). Since  $E_1 \cup E_2$  is finite, any extended matching problem  $P$  has only finitely many ground solutions, and each non-ground right-hand side of an equation in  $P$ , say  $t$ , can only be instantiated among a (possibly empty) finite set of ground terms  $Gnd(t)$ . Then, the set of match-equations *encoded by*  $P$  is inductively defined as follows:

$$\begin{aligned} ms(\{s =^? t\} \cup P) &= \bigcup_{t' \in Gnd(t)} \{s \leq^? t'\} \cup ms(P) \text{ if } t \text{ is non-ground} \\ ms(\{s =^? t\} \cup P) &= \{s \leq^? t\} \cup ms(P) \text{ if } t \text{ ground} \\ ms(\{s \leq^? t\} \cup P) &= \{s \leq^? t\} \cup ms(P) \end{aligned}$$

To compare match-equations we use the ordering:  $(s \leq^? t) \prec (s' \leq^? t')$  if

- $s$  is a strict subterm of  $s'$ ,
- or  $Var(s) \cap Var(s') = \emptyset$  and  $t$  is a strict subterm of  $t'$ ,
- or  $t = t'$  and  $s$  strictly subsumes  $s'$ .

The ordering  $\prec$  is well-founded and so its multiset extension  $\prec_m$  is well-founded.

**Lemma 9.** *Any rule in  $\mathcal{M}$  strictly decreases  $ms$  with respect to  $\prec_m$ .*

The above lemma is another way to prove the termination of  $\mathcal{M}$ .

## 5.2 Combination Procedure

Based on the restrictions related to  $A_1$  and  $A_2$ , we can give a new matching procedure for the hierarchical combination.

Consider  $\mathfrak{C}_{\mathcal{M}}$  the inference system depicted in Fig. 2, with the following inferences rules **{Solve<sub>1</sub>, Solve<sub>2</sub>, VA, RemEq, Rep, Merge, Clash}**. We can easily verify that each rule in  $\mathfrak{C}_{\mathcal{M}}$  preserves the set of  $E_1 \cup E_2$ -solutions. This is clear for the rules in **{VA, RemEq, Rep, Merge, Clash}**. Moreover, this is true by Lemma 7 for **Solve<sub>1</sub>**, and by Lemma 5 for **Solve<sub>2</sub>**. The inference system  $\mathfrak{C}_{\mathcal{M}}$  aims at computing matching problems in solved form.

**Lemma 10.** *Normal forms with respect to  $\mathfrak{C}_{\mathcal{M}}$  are matching problems in solved form.*

*Proof.* Assume  $P$  is an extended matching problem which is not a matching problem in solved form. Then we can always show that a rule in  $\mathfrak{C}_{\mathcal{M}}$  can be applied:

<b>Solve<sub>1</sub></b> :	$\frac{P_1 \cup P}{A_1(P_1) \cup P}$	$P_1$ is a set of match-equations with $\Sigma_1$ -rooted left-hand sides
<b>Solve<sub>2</sub></b> :	$\frac{P_2 \cup P}{A_2(P_2) \cup P}$	$P_2$ is a set of match-equations with 2-pure left-hand sides
<b>VA</b> :	$\frac{\{s[u] \leq^? t\} \cup P}{\{s[X] \leq^? t, u =^? X\} \cup P}$	if $s$ is $\Sigma_2$ -rooted, $u$ is an alien subterm
<b>RemEq</b> :	$\frac{\{t =^? t'\} \cup P}{\{t \leq^? t'\} \cup P}$	if $t'$ is ground
<b>Rep</b> :	$\frac{\{Y \leq^? u, t =^? t'[Y]\} \cup P}{\{Y \leq^? u, t =^? t'[u]\} \cup P}$	
<b>Merge</b> :	$\frac{\{X \leq^? t, X \leq^? s\} \cup P}{\{X \leq^? t\} \cup P}$	if $s =_{E_1 \cup E_2} t$
<b>Clash</b> :	$\frac{\{X \leq^? t, X \leq^? s\} \cup P}{Fail}$	if $s \neq_{E_1 \cup E_2} t$

**Fig. 2.**  $\mathfrak{C}_M$ : inference system for the combination of matching

1. If there is some equation  $t =^? t'$  in  $P$ , we have two possible cases. First, if  $t'$  is ground, then **RemEq** applies. Second, if  $t'$  contains some ground-solved variable, then a match-equation containing some fresh variable must also occur in  $P$ . If this match-equation is solved, then **Rep** applies. Otherwise, just like in the third case below, either **Solve<sub>1</sub>** or **Solve<sub>2</sub>** or **VA** applies on this unsolved match-equation.
2. If there are  $X \leq^? t$  and  $X \leq^? s$  in  $P$ , then either **Merge** or **Clash** can be applied.
3. If there is some match-equation  $s \leq^? t$  in  $P$  where  $s$  is not a variable, then either **Solve<sub>1</sub>** or **Solve<sub>2</sub>** or **VA** can be applied.

□

To show the soundness and completeness of  $\mathfrak{C}_M$ , it remains to show that  $\mathfrak{C}_M$  terminates for any input.

**Lemma 11.** *Let  $P$  be any input matching problem. Any repeated application of rules in  $\mathfrak{C}_M$  on  $P$  terminates.*

*Proof.* Consider the complexity measure  $ms$  introduced for Lemma 9 plus the following ones:

- $m_1$ : number of equations (denoted by  $=^?$ )
- $m_2$ : number of  $\Sigma_2$ -rooted match-equations
- $m_3$ : number of match-equations
- $m_4$ : number of variables occurring in equations (denoted by  $=^?$ )

Then, the termination of  $\mathfrak{C}_{\mathcal{M}}$  can be obtained by considering a lexicographic combination of these complexity measures, more precisely the one given by the tuple  $(ms, m_1, m_2, m_3, m_4)$ , as shown in the table below.

	$ms$	$m_1$	$m_2$	$m_3$	$m_4$
<b>Solve<sub>1</sub></b>	↓				
<b>VA</b>	↓				
<b>RemEq</b>	↓=	↓			
<b>Solve<sub>2</sub></b>	↓=	↓=	↓		
<b>Merge/Clash</b>	↓=	↓=	↓=	↓	
<b>Rep</b>	↓=	↓=	↓=	↓=	↓

□

**Theorem 2.** *For any finite syntactic hierarchical combination  $(E_1, E_2)$ , there exists a combined  $E_1 \cup E_2$ -matching algorithm (Fig. 1 and Fig. 2), provided that an  $E_2$ -matching algorithm is known.*

*Proof.* Direct consequence of Lemmas 10 and 11. □

### 5.3 Example: Matching in a Theory of Distributive Exponentiation

As an example of the above combination method for matching algorithms we consider in this section a theory of distributive exponentiation, namely  $\mathcal{E}_{AC}$  for which a rule-based unification algorithm is presented in [9]. We recall that  $\mathcal{E}_{AC} = (E_1, E_2)$ , where  $E_2 = AC(\otimes) = \{(x \otimes y) \otimes z = x \otimes (y \otimes z), x \otimes y = y \otimes x\}$  and  $E_1$  is given by the following convergent rewrite system modulo  $E_2$ :

$$R_1 = \begin{cases} \exp(\exp(x, y), z) \rightarrow \exp(x, y \otimes z) \\ \exp(x * y, z) \rightarrow \exp(x, z) * \exp(y, z) \end{cases}$$

**Lemma 12.**  $\mathcal{E}_{AC}$  is a finite syntactic hierarchical combination  $(E_1, AC(\otimes))$  where  $E_1$  admits the following resolvent presentation:

$$S_1 = \begin{cases} \exp(\exp(x, y), z) = \exp(x, y \otimes z) \\ \exp(x * y, z) = \exp(x, z) * \exp(y, z) \end{cases}$$

Note that Restriction 4(2) is addressed trivially since there are well-known  $AC$ -matching algorithms. Therefore, the main task is to instantiate the inference system  $\mathcal{M}$  (Fig. 1). This leads to the mutation rules for  $\mathcal{E}_{AC}$ -matching shown in Fig. 3.

*Example 2.* Consider the equational theory  $\mathcal{E}_{AC}$  and the matching problem

$$\exp(X, V \otimes c_1) \stackrel{?}{\leq} \exp(b, c_1 \otimes c_2 \otimes c_3)$$

The combination algorithm  $\mathfrak{C}_{\mathcal{M}}$  works as follows with this input. First, **Matching Decomposition** is applied and leads to  $\{X \stackrel{?}{\leq} b, V \otimes c_1 \stackrel{?}{\leq} c_1 \otimes c_2 \otimes c_3\}$ . Then **Solve<sub>2</sub>** applies and provides a first solved form  $\{X \stackrel{?}{\leq} b, V \stackrel{?}{\leq} c_2 \otimes c_3\}$ .

Another possibility is to apply the first **Mutate** rule from Fig. 3, yielding  $\{Y \otimes V \otimes c_1 \leq^? c_1 \otimes c_2 \otimes c_3, X =^? exp(b, Y)\}$ . By **Solve**<sub>2</sub>, the above 2-pure match-equation has two solutions. The first solution is  $\{Y \leq^? c_2, V \leq^? c_3\}$ . After **Rep** and **RemEq**, we obtain a new solved form:  $\{V \leq^? c_3, X \leq^? exp(b, c_2)\}$ . Similarly, for the second solution  $\{Y \leq^? c_3, V \leq^? c_2\}$ , we get the solved form  $\{V \leq^? c_2, X \leq^? exp(b, c_3)\}$ . The other **Mutate** rules lead to a failure thanks the application of **Matching Clash**.

<b>Mutate</b>
$\{exp(s_1, s_2) \leq^? exp(t_1, t_2)\} \cup P \longrightarrow \{Y \otimes s_2 \leq^? t_2, s_1 =^? exp(t_1, Y)\} \cup P$
$\{exp(s_1, s_2) \leq^? exp(t_1, t_2)\} \cup P \longrightarrow \{exp(s_1, Y) \leq^? t_1, s_2 =^? Y \otimes t_2\} \cup P$
$\{exp(s_1, s_2) \leq^? t_1 * t_2\} \cup P$
$\longrightarrow \{exp(X, s_2) \leq^? t_1, exp(Y, s_2) \leq^? t_2, s_1 =^? X * Y\} \cup P$
$\{s_1 * s_2 \leq^? exp(t_1, t_2)\} \cup P$
$\longrightarrow \{X * Y \leq^? t_1, s_1 =^? exp(X, t_2), s_2 =^? exp(Y, t_2)\} \cup P$

**Fig. 3.** Mutation rules for  $\mathcal{E}_{AC}$ -matching

## 6 Conclusion

We have presented a collection of new results about our hierarchical combination approach for solving unification problems. First we defined a simpler reformulation of the combination method, which is sufficient for the problems we focus on in this paper. Our application to shallow extensions complement and improve our earlier work on hierarchical combination presented in [7,9]. Hierarchical combination requires a solver  $A_1$ , taking in account the axioms of  $E_1$ , to produce partial solved forms. Although a general sound and complete method is available to construct  $A_1$ , the problem of termination still remains. We solve this problem for shallow theories by showing how to exploit unification algorithms known for them. Second, we have shown another combination method for the matching problem in finite syntactic extensions. Future work includes applying the general method developed here to partly ground unification problems and finding conditions which allow us to combine unification algorithms of larger classes of equational theories.

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