

L2-Orthogonal ST-Code Design for Multi-h CPM with fast Decoding

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Université Nice Sophia Antipolis Abstract \oplus CPM: favorable trade-off between power and bandwidth efficiency. \oplus Multi-*h* CPM: generalization to further decrease the need for bandwidth. \ominus Difficult decoding in multi-path environments with no diversity. How to overcome these limitations: \oplus **IDEA:** To combine CPM with Space-Time Block Coding (STBC). \rightarrow Non trivial extension to L^2 -orthogonal Space-Time codes provides full diversity and better spectral compactness.[1][2] \rightarrow Decoding complexity greatly decreased.[1][2] The System Model *M*-ary multi-*h* antennas. $\mathbf{r}(t,\mathbf{d})_{(L_r imes 1)}$, $\mathbf{s}(t,\mathbf{d})_{(L_t imes 1)}$ $\mathbf{H}_{(L_r imes L_t)}$ with elements $h_{n,m}$ $\mathbf{n}(t)$ AWGN. The baseband general form [3] $s(t, \mathbf{d}) = \sqrt{\frac{E_s}{T}} \exp(j\phi(t, \mathbf{d}))$ The information-carrying phase function pulse q(t), K-1 $\phi(t, \mathbf{d}) = 2\pi \sum_{k=1}^{\infty} h_{[k]} d_k q(t - (k - 1)T)$ $q(t) = \left\{ \ \bar{q}(t) = \right\}$ modulation indices h_1, \ldots, h_H , cycle in time with period *H* as: [k] = mod(k, H) + 1the symbol index is given by $k = \lfloor t/T \rfloor$ $h_{[k]}$ quotient between two relative prime integers. $H_k = (h_1, h_2, \dots, h_k) = \left(\frac{p_1}{a}, \frac{p_2}{a}, \dots, \frac{p_k}{a}\right)$ L^2 Orthogonality Based on $\|\mathbf{s}\|_{L^2}^2 = \int \mathbf{s}(t) \mathbf{s}^H(t) dt$, each block fulfills the condition Time Slots defined within a Block Blocks defined within a Frame T_x l_E=1...L_E $l_F L_B T$ $\mathbf{s}(t$ $(l_f-1)L_BT$ **⊦**----, 1...L_BT Block

L2-Orthogonal ST-Code Design for Multi-H CPM with fast Decoding

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The phase continuity is ensured by an antenna dependent phase memory:

$$\theta_{m,k} = \theta_{m,k-1} + \frac{n_{\lfloor k-\gamma \rfloor}}{2} d_{m,k-\gamma}$$

 $v_{\gamma} + c_{m,k}(T) - c_{m,k}(0).$ $-\theta_{m',k} - \sum_{i=k-\alpha+1}^{k} h_{[i]} d_{m',i} q(t - (i-1)T) - c_{m',k}(t)]) dt = 0$ 2 d - - d(t-1)T < t < kT $s_m(t, \mathbf{d}) = s(t, \mathbf{d}) \exp(j2c_m^{lin}(t))$ L^2 Decoding $3L_BT$ $L_B T$ $2L_BT$ inter-block $\Theta(1)$ trellis M^{L_t} pps $\Theta(pM^{\gamma})$ $2L_BT$ inner-block trellis $\Theta(pM^{\gamma-1})$

$$\sum_{k=1}^{L_B} \int_{(k-1)T}^{kT} \exp(j2\pi [\theta_{m,k} + \sum_{i=k-\gamma+1}^k h_{[i]} d_{m,i} q(t - (i-1)T) + c_{m,k}(t) - q_{m,k}(t) - q$$

1.
$$c_{m,k}(t) = c_{m,k'}(t)$$

$$d_{m,k} = d_{m',k}$$

Auto-correlation coefficients cancel and cross-correlation coefficients between antennas put to 0 To ease the design, two assumptions are introduced, For any arbitrary number of transmit antennas, we introduce correction functions $c_{m,k}$ as: For $m = 1, \ldots, L_t$ and the transmitted signal takes the form of

$$c_m^{lin}(t) = \frac{m-1}{L_B T} t \quad \text{for } (k$$



FIG. 1: Inner-block phase transition for $h_1 = \frac{1}{4}$ and $h_2 = \frac{3}{4}$,

Classical correlation based multi-*h* detector expressed blockwise:

$$D(\mathbf{d}, \tilde{\mathbf{d}} | \Theta(k)) = \sum_{k=1}^{L_F} \sum_{n=1}^{N_s k} \int_{kTL_B}^{(k+1)TL_B} Re\left\{ r(t, \mathbf{d}_{1_{[n]}}, \dots, \mathbf{d}_{L_B[n]}) \cdot \left(\sum_{m=1}^{L_t} h_m^* s_m^*(t, \theta, \tilde{\mathbf{d}}_{1_{[n]}}, \dots, \tilde{\mathbf{d}}_{L_B[n]}) \right) \right\} dt$$

\ni Impractical:

 $pM^{(L_tL_BN_s+\gamma-1)}$ matched filters of length N_sL_BT employed $pM^{\gamma-1}$ times. pM^{L_B-1} states and $M^{L_BL_T}$ symbols to decode with N_s samples each. **Orthogonality:** the cross correlation terms are canceled out with blockwise decoding.

 $D_B(\mathbf{d}, \tilde{\mathbf{d}} | \Theta(k)) = \sum_{m=1}^{L_T} \int_{0}^{t_T} Re\{r(t, \mathbf{d}) h_m^* c_m^*(t) s^*(t, \tilde{\mathbf{d}})\} dt.$

To get decoding with complexity growing linearly w.r.t number of transmit antennas, we introduce

$$x(t,\mathbf{d}) = r(t,\mathbf{d})\sum_{m=1}^{L_T}$$

A simplified expression for a classical multi-*h* decoder [4] then is given by

$$D(\mathbf{d}, \tilde{\mathbf{d}} | \Theta(k)) = \int_{(l-1)T}^{lT} Re\{x(t, t)\}$$

MIMO system of
$$L_t \times L_r$$

 $\mathbf{r}(t, \mathbf{d}) = \mathbf{Hs}(t, \mathbf{d}) + \mathbf{n}(t)$

The vector of transmitted signals:

 $\mathbf{s}(t, \mathbf{d}) = \left[\mathbf{s}_l(t, \mathbf{d}) \dots \mathbf{s}_{L_t}(t, \mathbf{d})\right]^T$

The phase continuity is ensured by the phase

$$\begin{cases} 0 & t < 0\\ \bar{q}(t) & elsewhere\\ 1/2 & t \ge \gamma T \end{cases}$$

 $ar{q}(0)=0$, $ar{q}(\gamma T)=1/2$ and $\lim_{t
ightarrow au}ar{q}(t)=ar{q}(au)$ for $0 \leq \tau < \gamma T$. γ is the overlapping factor and



$$(t, \mathbf{d})\mathbf{s}^{H}(t, \mathbf{d})dt =$$

M pps, $D(\mathbf{d}, \mathbf{\tilde{d}})$

FIG. 2: Simplified detection with k = 1 (ppspaths per state).

 $h_m^* c_m^*(t).$

 $(t, \mathbf{d}) \cdot s(t, \mathbf{\tilde{d}})_{|\phi_l|} \} dt$

minimum distance as



References

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