A Real-Time Privacy Amplification Scheme in Quantum Key Distribution

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Abstract..QKD (Quantum Key Distribution) technology, based on the laws of physics, can create an unconditionally secure key between communication parties. In recent years, researchers draw more and more attention to the QKD technology. Privac amplification is a very significant procedure in QKD system. In this paper, we propose the real-time privacy amplification (RTPA) scheme which converts the weak secret string W to a uniform key that is fully secret from Eve. We implement RTPA scheme based on CLIP (Cvqkd Ldpc experImental Platform) which is connected to the real quantum communication systems. Experimental results show that, our proposed RTPA scheme is very efficient when the bit error rate of quantum channel is lower than 0.06.

Keywords: QKD, privacy amplification, security

1 Introduction

Quantum Key Distribution (QKD) [1, 2] is technology for solving the key distribution problem. QKD system, based on the laws of physics, rather than the computational complexity of the mathematical problems assumed by current cryptography methods, can create an unconditionally secure key between communication parties. These keys are generated over unsecured channels, where may exist an active computationally unbounded adversary Eve.

After the procedure of information reconciliation [3], Alice and Bob have own almost uniform keys with comparative low BER (Bit Error Rate). But Eve may have partial knowledge about the keys by eavesdropping or other ways.

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Therefore, in order to gain the absolutely security keys, we must ensure the keys are privacy amplified. Privacy amplification (PA) [4] is a technology, through a public channel, to improve the information confidentiality. Privacy amplification converts the weak secret string W to a uniform key that is fully secret from Eve.

Privacy amplification technology typically applies a random and public hash function to shorten the weak secret key and reduce the amount of information obtained by Eve as much as possible. By sacrificing partially key information of Alice and Bob，privacy amplification makes the knowledge obtained by Eve been meaningless.

Though the majority researches (such as [5, 6]) about privacy amplification focusing on the theoretical study and proof of security, implementing an efficient privacy amplification scheme in QKD system has been more and more significant.

In this paper, we propose a real-time privacy amplification scheme (RTPA) and implement RTPA in CLIP system [7], which is connected to the quantum communication system. After extensive experiments, the performance and the detail analysis are described in Section III. Experimental results show the efficiency of our proposed RTPA scheme for generating unconditional security keys in quantum cryptography.

2 The Proposed RTPA Scheme

2.1 Privacy Amplification Protocol

After analyzed and researched the classical and quantum privacy amplification theoretical study in [4, 8, 9, 10] and etc., we approach the RTPA protocol (Real-time Privacy Amplification Protocol).

We assume that the key information of Alice and Bob after information reconciliation is W and its length is N, the length of key information used for reconciliation, confirmation and etc. is K, the length of key information may obtained by Eve is T, the security parameter is S, and the final key length is R. We describe the RTPA protocol as follows:

- Alice and Bob select the security parameter S, according to the quantum key state, key length N and other information;
- Alice generates the description information about hash function randomly, the seed string $Seed$ and the shift string $Shift$. $Seed$ and $Shift$ send to Bob through the public channel;
- Alice and Bob construct the hash function f, $f \in F$, $F:\{0,1\}^N \to \{0,1\}^R$, $R = N - T - K - S$;
- Alice and Bob gain the final key y, $y = f(W)$.

In RTPA protocol, hash function f is randomly chosen from class H_3 of universal$_2$. The hash function is described by Toeplitz matrix construction method [11, 12, 13]. After applying the privacy amplification procedure, the final key is unconditionally safe to Eve.

2.2 The RTPA Scheme

The RTPA scheme mainly consists three parts: Hash function construction, data communication and privacy amplification. The architecture of RTPA is shown in Figure 1.

Hash function construction

As shown in Figure 1, the parameter controller carries out the security parameter S and controls the generation of $Shift$ and $Seed$ based on the quantum channel states. Then, the hash function construction module constructs Toeplitz hash function scaling $N \times R$.

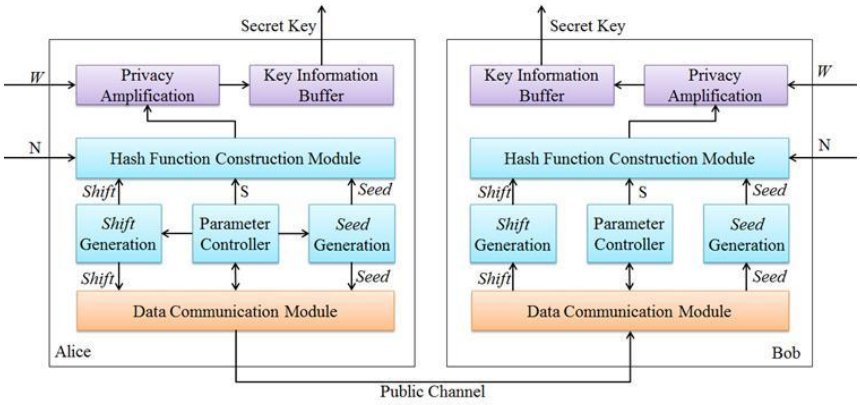

Fig. 1. The Architecture of RTPA

- **Data communication**
  In this part, Alice sends $Shift$ and $Seed$, the description information of hash function , to Bob through a public channel.
- **Privacy amplification**
  Privacy amplification is applied to convert W to an absolutely secret key with length R . These keys used for quantum cryptography are stored in the Key information buffer.

3 Experimental Results And Analysis

The RTPA scheme is implemented in CLIP [7] which is connected to the real quantum communication system. The experimental environment is shown in Figure 2.

We conducted extensive experiments to evaluate the performance of RTPA. We analyzed the privacy amplification overhead, average bit error rate (avBER) of key information.

3.1 Privacy Amplification Overhead

Various hash function constructed for different input key lengths, will lead to different time overhead per privacy amplification process. While the input key length should be long enough in order to gain an absolutely security key, the privacy amplification overhead will be very high. In this experiment, we test the privacy amplification overhead of different hash function scale. The result is shown in Figure 3.

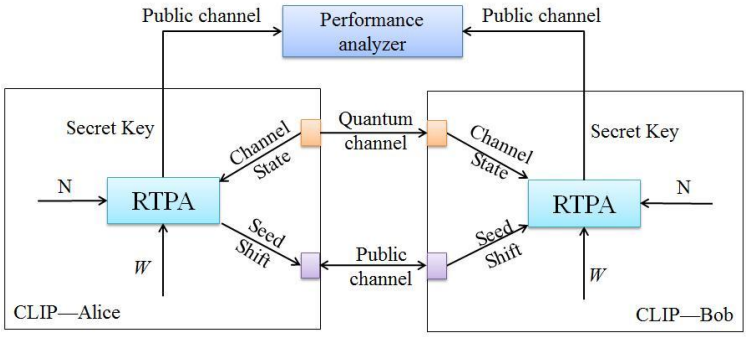

Fig. 2. The experimental environment of RTPA

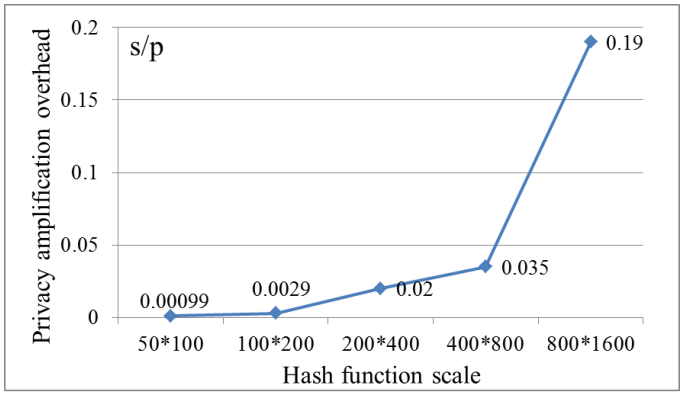

Fig. 3. The privacy amplification overhead of various hash function scales

For example, it will cost 0.19s/p (second per amplification process) when converting a key from 1600 bits to 800bits. Though costing 9.5 times overhead than the scale of 200*400, the security of keys is enhanced by thousands of times. The hash function scale should be balanced between the security demands and the time overhead.

3.2 Average Bit Error Rate

After the procedure of information reconciliation, Alice and Bob have own almost uniform keys with comparative low BER (Bit Error Rate). When applying hash func-

tion to these keys, it may generate quite different strings for Alice and Bob. Therefore, we test the average Bit Error Rate for the final keys with different quantum channel Bit Error Rates.

As it shown in Figure 4, privacy amplification can work effectively when the BER of quantum channel is lower than 0.06. When the BER of quantum channel ranges from 0.06 to 0.10, the information reconciliation procedure still works effectively, the BER after information reconciliation is close to zero, but the BER after privacy amplification is very high. And it doesn't meaningless when the quantum channel BER is higher than 0.10.

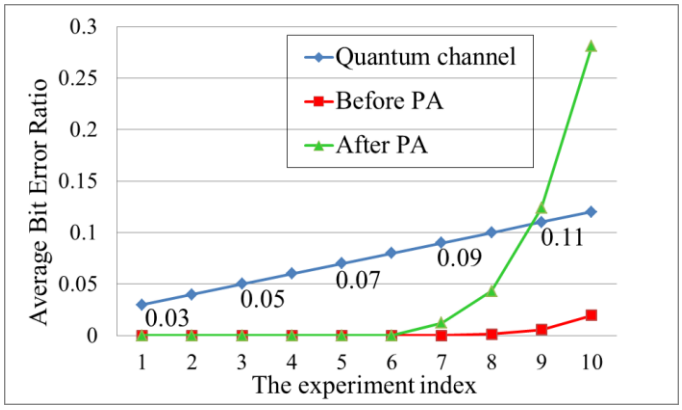

Fig. 4. The Average Bit Error Ratios with different scenes

4 Conclusion

In this paper, we approached the privacy amplification protocol and proposed the RTPA scheme, a real-time quantum privacy amplification procedure in QKD systems. To evaluate the performance of RTPA, we built a prototype QKD system based on CLIP [7]. Experimental results showed the efficiency of our proposed RTPA scheme when the bit error rate of quantum channel is lower than 0.06. The results showed that the performance of RTPA is greatly affected by the quantum channel BER and the information reconciliation. In order to gain an efficient performance, we must enhance the performance of information reconciliation to gain a low BER of key information before privacy amplification.

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