

Minimum density of identifying codes of king grids

Rennan Dantas, Rudini Sampaio, Frédéric Havet

► **To cite this version:**

Rennan Dantas, Rudini Sampaio, Frédéric Havet. Minimum density of identifying codes of king grids. *Electronic Notes in Discrete Mathematics*, Elsevier, 2017, 62, pp.51 - 56. <10.1016/j.endm.2017.10.010>. <hal-01634305>

HAL Id: hal-01634305

<https://hal.inria.fr/hal-01634305>

Submitted on 13 Nov 2017

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

Minimum density of identifying codes of king grids

Rennan Dantas, Rudini M. Sampaio¹

Universidade Federal do Ceará, Fortaleza, Brazil

Frédéric Havet¹

Université Côte d'Azur, CNRS, I3S, INRIA, France

Abstract

A set $C \subseteq V(G)$ is an *identifying code* in a graph G if for all $v \in V(G)$, $C[v] \neq \emptyset$, and for all distinct $u, v \in V(G)$, $C[u] \neq C[v]$, where $C[v] = N[v] \cap C$ and $N[v]$ denotes the closed neighbourhood of v in G . The minimum density of an identifying code in G is denoted by $d^*(G)$. In this paper, we study the density of king grids which are strong product of two paths. We show that for every king grid G , $d^*(G) \geq 2/9$. In addition, we show this bound is attained only for king grids which are strong products of two infinite paths. Given $k \geq 3$, we denote by \mathcal{K}_k the (infinite) king strip with k rows. We prove that $d^*(\mathcal{K}_3) = 1/3$, $d^*(\mathcal{K}_4) = 5/16$, $d^*(\mathcal{K}_5) = 4/15$ and $d^*(\mathcal{K}_6) = 5/18$. We also prove that $\frac{2}{9} + \frac{8}{81k} \leq d^*(\mathcal{K}_k) \leq \frac{2}{9} + \frac{4}{9k}$ for every $k \geq 7$.

Keywords: Identifying code, King grid, Discharging Method.

1 Introduction

Let G be a graph. The *neighbourhood* of a vertex v of G , denoted by $N(v)$, is the set of vertices adjacent to v in G , and the *closed neighbourhood* of v is the set $N[v] = N(v) \cup \{v\}$. Given a set $C \subseteq V(G)$, let $C[v] = N[v] \cap C$. We say that C is an *identifying code* of G if $C[v] \neq \emptyset$ for all $v \in V(G)$, and $C[u] \neq C[v]$ for all distinct $u, v \in V(G)$. Clearly, a graph G has an identifying code if and only if it contains no *twins* (vertices $u, v \in V(G)$ with $N[u] = N[v]$).

Let G be a (finite or infinite) graph with bounded maximum degree. For any non-negative integer r and vertex v , we denote by $B_r(v)$ the ball of radius r in G centered at v , that is $B_r(v) = \{x \mid \text{dist}(v, x) \leq r\}$. For any set of vertices $C \subseteq V(G)$, the *density* of C in G , denoted by $d(C, G)$, is defined by

$$d(C, G) = \limsup_{r \rightarrow +\infty} \frac{|C \cap B_r(v_0)|}{|B_r(v_0)|},$$

¹ {rennan,rudini}@lia.ufc.br, frederic.havet@cnrs.fr

where v_0 is an arbitrary vertex in G . The infimum of the density of an identifying code in G is denoted by $d^*(G)$. Observe that if G is finite, then $d^*(G) = |C^*|/|V(G)|$, where C^* is a minimum-size identifying code of G .

The problem of finding low-density identifying codes was introduced in [12] in relation to fault diagnosis in arrays of processors. Particular interest was dedicated to grids as many processor networks have a grid topology. Many results have been obtained on square grids [4,1,9,2,11], triangular grids [12,10], and hexagonal grids [5,7,8]. In this paper, we study *king grids*, which are strong products of two paths. The *strong product* of two graphs G and H , denoted by $G \boxtimes H$, is the graph with vertex set $V(G) \times V(H)$ and edge set :

$$\begin{aligned} E(G \boxtimes H) = & \{(a, b)(a, b') \mid a \in V(G) \text{ and } bb' \in E(H)\} \\ & \cup \{(a, b)(a', b) \mid aa' \in E(G) \text{ and } b \in V(H)\} \\ & \cup \{(a, b)(a', b') \mid aa' \in E(G) \text{ and } bb' \in E(H)\}. \end{aligned}$$

The *two-way infinite path*, denoted by $P_{\mathbb{Z}}$, is the graph with vertex set \mathbb{Z} and edge set $\{\{i, i + 1\} \mid i \in \mathbb{Z}\}$, and the *one-way infinite path*, denoted by $P_{\mathbb{N}}$, is the graph with vertex set \mathbb{N} and edge set $\{\{i, i + 1\} \mid i \in \mathbb{N}\}$. A *path* is a connected subgraph of $P_{\mathbb{Z}}$. For every positive integer k , P_k is the subgraph of $P_{\mathbb{Z}}$ induced by $\{1, 2, \dots, k\}$. A *king grid* is the strong product of two (finite or infinite) paths. The *plane king grid* is $\mathcal{G}_K = P_{\mathbb{Z}} \boxtimes P_{\mathbb{Z}}$, the *half-plane king grid* is $\mathcal{H}_K = P_{\mathbb{Z}} \boxtimes P_{\mathbb{N}}$, the *quater-plane king grid* is $\mathcal{Q}_K = P_{\mathbb{N}} \boxtimes P_{\mathbb{N}}$, and the *king strip of height k* is $\mathcal{K}_k = P_{\mathbb{Z}} \boxtimes P_k$.

In 2002, Charon et al. [3] proved that $d^*(\mathcal{G}_K)$ is $2/9$. They provided the tile depicted in Figure 1, which generates a periodic tiling of the plane with periods $(0, 6)$ and $(6, 0)$, yielding an identifying code C_{∞} of the bidimensional infinite king grid with density $\frac{2}{9}$.

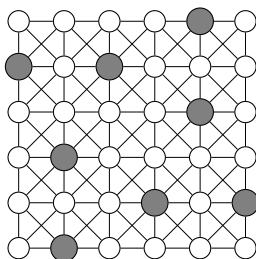


Fig. 1. Tile generating an optimal identifying code of the bidimensional infinite grid. Black vertices are those of the code.

In this paper, using the Discharging Method (see Section 3 of [10] for a detailed presentation of this technique for identifying codes), we provide the following tight general lower bound on the minimum density of identifying

codes of king grids.

Theorem 1.1 *If G is a (finite or infinite) king grid, then $d^*(G) \geq \frac{2}{9}$.*

Keeping on, we prove the following.

Theorem 1.2 *If G is a finite king grid, then $d^*(G) > \frac{2}{9}$.*

Finally, we give some bounds for king strips. Pushing further the proof of Theorem 1.1, we prove the following.

Theorem 1.3 *For every $k \geq 6$, $d^*(\mathcal{K}_k) \geq \frac{2}{9} + \frac{8}{81k}$.*

Modifying C_∞ , we construct identifying codes of \mathcal{K}_k yielding the following upper bounds.

Theorem 1.4 *For every $k \geq 5$,*

$$d^*(\mathcal{K}_k) \leq \begin{cases} \frac{2}{9} + \frac{6}{18k}, & \text{if } k \equiv 0 \pmod{3}, \\ \frac{2}{9} + \frac{8}{18k}, & \text{if } k \equiv 1 \pmod{3}, \\ \frac{2}{9} + \frac{7}{18k}, & \text{if } k \equiv 2 \pmod{3}. \end{cases}$$

Finally, we show some identifying codes of \mathcal{K}_3 , \mathcal{K}_4 , \mathcal{K}_5 and \mathcal{K}_6 (see Figures 2, 3, 4, and 5.) and prove that they are optimal. This yields the following.

Theorem 1.5 $d^*(\mathcal{K}_3) = 1/3 = 0.333\dots$ $d^*(\mathcal{K}_4) = 5/16 = 0.3125$
 $d^*(\mathcal{K}_5) = 4/15 = 0.2666\dots$ $d^*(\mathcal{K}_6) = 5/18 = 0.2777\dots$

Clearly $d^*(\mathcal{K}_1) = 1/2$ (as $\mathcal{K}_1 = \mathcal{S}_1 = P_{\mathbb{Z}}$) and \mathcal{K}_2 has no identifying code because it has twins. All these results imply that \mathcal{G}_K , \mathcal{H}_K and \mathcal{Q}_K are the unique king grids having an identifying code with density $2/9$. (One can easily derive from C_∞ identifying codes with density $2/9$ of \mathcal{H}_K and \mathcal{Q}_K).

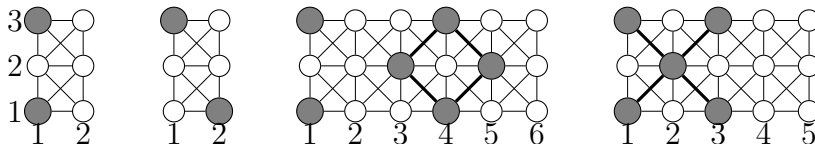


Fig. 2. Four tiles generating optimal identifying codes of \mathcal{K}_3 (density $1/3$)

2 Sketches of proofs

Sketch of proof of Theorem 1.1. Let G be a king grid and C an identifying code of G . We shall prove that $d(C, G) \geq 2/9$. For this, we use the Discharging

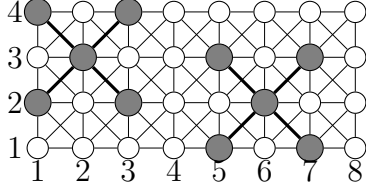


Fig. 3. Tile generating an optimal identifying code of \mathcal{K}_4 (density $5/16$)

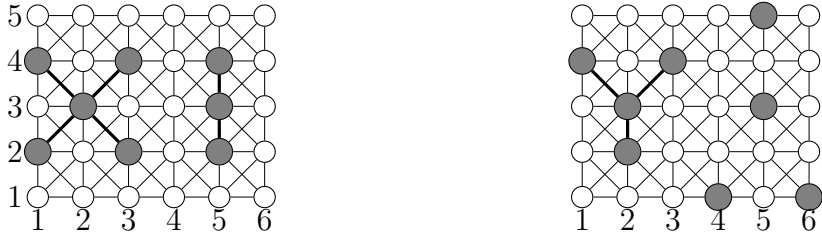


Fig. 4. Two tiles generating optimal identifying codes of \mathcal{K}_5 (density $4/15$)



Fig. 5. Two tiles generating optimal identifying codes of \mathcal{K}_6 (density $5/18$)

Method. The initial charge of a vertex v is 1 if $v \in C$ and 0 otherwise. We then apply some local discharging rules. We shall prove that the final charge of every vertex in C is at least $2/9$. This would imply the result.

We set $U = V(G) \setminus C$. Given $X \subseteq V(G)$ and $1 \leq i \leq 9$, we denote by X_i (resp. $X_{\geq i}$) the set of vertices in X having exactly i vertices (resp. at least i vertices) in their identifier. An X -vertex is a vertex in X . A vertex is *full* if its eight neighbours in \mathcal{G}_K are in G ; otherwise it is a *side vertex*.

We first establish the following properties of C .

- (i) Two C_2 -vertices are not adjacent.
- (ii) Every C -vertex has at most one neighbour in U_1 .
- (iii) Every full C_2 -vertex has at least three neighbours in $U_{\geq 3}$.
- (iv) Every full C_3 -vertex has a neighbour in $U_{\geq 3}$.

- (v) Every C_1 -vertex (a, b) has no neighbour in U_1 and at most six neighbours in U_2 . Furthermore, if it has six neighbours in U_2 , then either $\{(a-1, b-2), (a-2, b-1), (a+2, b+1), (a+1, b+2)\} \subseteq C$ or $\{(a+1, b-2), (a+2, b-1), (a-2, b+1), (a-1, b+2)\} \subseteq C$.

A *defective vertex* is a vertex in C_1 with six neighbours in U_2 . Let $v = (a, b)$ be a defective vertex. The *team* of v is a set among $\{(a-1, b-2), (a-2, b-1), (a+2, b+1), (a+1, b+2)\}$ and $\{(a+1, b-2), (a+2, b-1), (a-2, b+1), (a-1, b+2)\}$ which is included in C . By Property (v), the team exists. Moreover, by Property (i), at least two vertices of the team are in $C_{\geq 3}$. Those vertices are the *partners* of v .

We apply the following discharging rules.

- (R1) Every C -vertex sends $\frac{2}{9i}$ to each of its neighbours in U_i .
(R2) Every defective vertex receives $\frac{1}{54}$ from each of its partners.

Using the above properties, we then prove that the final charge of every vertex v is at least $2/9$. \square

Sketch of proof of Theorem 1.2. We only need to prove that, at the end of the proof of Theorem 1.1, one vertex has final charge greater than $2/9$. To do so we shall prove that there is a side C -vertex or a $C_{\geq 3}$ -vertex and check that such a vertex has final charge at least $\frac{2}{9} + \frac{1}{27}$. \square

Sketch of proof of Theorem 1.3. Using the Discharging Method, we prove that in average, for every column, there is an extra charge of at least $\frac{4}{81}$ on the three top vertices and an extra charge of at least $\frac{4}{81}$ on the three top vertices. \square

Sketch of proof of Theorem 1.5. The b th row of \mathcal{K}_k is $R_b = \{(a, b) \mid a \in \mathbb{Z}\}$. We have $d(C, \mathcal{K}_k) = \frac{1}{k} \sum_{i=1}^k d(C, R_i)$. We show that if C is an identifying code of \mathcal{K}_k ($k \geq 3$), then $d(C, R_1) + d(C, R_2) \geq 1/2$, $d(C, R_k) + d(C, R_{k-1}) \geq 1/2$, $d(C, R_3) \geq 1/3$ and $d(C, R_{k-2}) \geq 1/3$. One easily derives that if C is an identifying code of \mathcal{K}_5 (resp. \mathcal{K}_6), then $d(C, \mathcal{K}_5) \geq 4/15$. (resp. $d(C, \mathcal{K}_6) \leq 5/18$.)

To prove lower bounds on $d^*(\mathcal{K}_k)$ for $k \in \{3, 4\}$, we use the Discharging Method on the columns $Q_a = \{(a, b) \mid 1 \leq b \leq k\}$. Let C be an identifying code of \mathcal{K}_k . We set the initial charge of every integer $a \in \mathbb{Z}$ to $\text{chrg}_0(a) = |Q_a \cap C|$. We say that $a \in \mathbb{Z}$ is *satisfied* if its charge is least q_k and *unsatisfied* otherwise, where $q_3 = 1$ and $q_4 = 5/4$. We apply five discharging rules, Rule i for $i = 1$ to 5 one after another. We denote by $\text{chrg}_i(a)$ the charge of a after applying Rule i .

Rule i : every unsatisfied $a \in \mathbb{Z}$ receives $\min\{\text{chrg}_{i-1}(a-i) - q_k, q_k - \text{chrg}_{i-1}(a)\}$

from $a - i$, if $a - i$ is satisfied (before Rule i).

Finally, we prove that, after these rules, every integer $a \in \mathbb{Z}$ is satisfied. This implies $d(C, \mathcal{K}_k) \geq q_k/k$. \square

Acknowledgments This research was supported by FAPERJ, Capes (Project STIC-AmSud 2012), CNPq (Project Universal/Bolsa de Produtividade), ANR (Contract STINT ANR-13-BS02-0007), and the FUNCAP/CNRS project GA-IATO INC-0083-00047.01.00/13.

References

- [1] Y. Ben-Haim and S. Litsyn. Exact minimum density of codes identifying vertices in the square grid. *SIAM J. Discrete Math.* 19: 69–82, 2005.
- [2] M. Bouznif, F. Havet, M. Preissman. Minimum-density identifying codes in square grids. *AAIM 2016. Lect. Notes Computer Science*, 9778: 77–88, 2016.
- [3] I. Charon, O. Hudry and A. Lobstein. Identifying Codes with Small Radius in Some Infinite Regular Graphs. *Elec. J. Combinatorics* 9: R11, 2002.
- [4] G. Cohen, S. Gravier, I. Honkala, A. Lobstein, M. Mollard, C. Payan, and G. Zémor. Improved identifying codes for the grid, *Comment to [6]*.
- [5] G. Cohen, I. Honkala, A. Lobstein, and G. Zemor. Bounds for Codes Identifying Vertices in the Hexagonal Grid. *SIAM J. Discrete Math.* 13: 492–504, 2000.
- [6] G. Cohen, I. Honkala, A. Lobstein, and G. Zémor. New bounds for codes identifying vertices in graphs. *Elec. J. Combinatorics* 6(1): R19, 1999.
- [7] D.Cranston, G.Yu, A new lower bound on the density of vertex identifying codes for the infinite hexagonal grid. *Elec. J. Combinatorics* 16: R113, 2009.
- [8] A. Cukierman, G. Yu. New bounds on minimum density of an identifying code for the infinite hexagonal graph grid. *Disc. App. Math.* 161: 2910–2924, 2013.
- [9] M. Daniel, S. Gravier, and J. Moncel. Identifying codes in some subgraphs of the square lattice. *Theoretical Computer Science* 319: 411–421, 2004.
- [10] R. Dantas, F. Havet, R. Sampaio. Identifying codes for infinite triangular grids with finite number of rows. *Discrete Math.*, doi 10.1016/j.disc.2017.02.015
- [11] M. Jiang. Periodicity of identifying codes in strips. *arXiv:1607.03848 [cs.DM]*
- [12] M. Karpovsky, K. Chakrabarty, and L. B. Levitin. On a new class of codes for identifying vertices in graphs. *IEEE Trans. Inform. Theory* 44: 599–611, 1998.