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Abstract Graph clustering is one of the key techniques to understand the structures present in the graph data. In addition to cluster detection, the identification of hubs and outliers is also a critical task as it plays an important role in the analysis of graph data. Recently, several graph clustering algorithms have been proposed and used in many application domains such as biological network analysis, recommendation systems and community detection. Most of these algorithms are based on the structural clustering algorithm SCAN. Yet, SCAN algorithm has been designed for small graphs, without significant support to deal with big and dynamic graphs. In this paper, we propose DISCAN, a novel distributed and incremental graph clustering algorithm based on SCAN. We present an implementation of DISCAN on top of BLADYG framework, and experimentally show the efficiency of DISCAN in both large and dynamic networks.

Keywords Graph processing · Structural graph clustering · Big Graph Analysis · Community detection · Dynamic graph clustering

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1 Introduction

Recently, the graph model has risen one of the most used data models in several applications like social networks [26], road maps [8], pattern mining [17] and bioinformatics [36]. For instance, a recent ranking shows that the popularity of graph databases model increased up to around 500% in the last years [21]. It has shown an optimum data model that allows to represent easily many relationships and facilitates the exploration of data. Taking social networks as an example, the graph model organizes data elements into a set of vertices representing the members, and a list of edges to materialize the relationships between vertices. Also, the last years featured a big data explosion especially in graph-based social networks. As an example, Facebook in 2013 had over 874 million monthly users [5]. This proliferation of a huge amount of data and the massiveness of graphs introduce additional factors to the renewed popularity of graph analytics [21]. Consequently, new applications and use cases have been mentioned in the literature. One important analysis technique in graph mining and graph analysis fields is graph clustering. Graph clustering helps identifying tightly connected regions within a graph [38,19]. It has been used to solve various problems such as discovering communities in social networks and detecting protein complexes in Protein–Protein Interaction (PPI) networks [36].

It is important to mention that in some applications, the used graph could be very large. In such situation, the considered graph could be partitioned into several sub-graphs and the computation is performed in a parallel/distributed way [12,2]. In this context, graph clustering algorithms are used to ensure the partitioning of the graph into several parts [1,37].

In this paper, we propose DISCAN: a novel graph clustering algorithm in the context of large and dynamic graphs. The main contributions of this work are summarized as follows:

- We propose a distributed and incremental graph clustering algorithm of both large and dynamic graphs (DISCAN). The proposed method is exact and allows discovering the same clusters that are discovered by SCAN, the reference algorithm for structural graph clustering [35].
- We conduct an extensive experimental study with large graphs, to evaluate the scalability of DISCAN. We compare DISCAN with four existing structural graph clustering algorithms.

The rest of the paper is organized as follows. After introducing graph clustering in Section 1, we present a brief overview of related work in Section 2. In Section 3, we present the basic concepts related to the structural graph clustering. In Section 4, we present our proposed algorithms for large and dynamic graph clustering. In Section 5, we provide the experimental results. The last section is devoted to the conclusion and the future work.

2 Related work

Several graph clustering algorithms have been proposed in the literature. Taking as examples, modularity-based approaches [23] that represent an optimization solution of the modularity measure for each partitioning schema (generated randomly or according to a heuristic function) [23]. The Louvain method [4] represents one of the clustering algorithms based on modularity, that initially generates a random clustering. After that, it starts to change every time a vertex from a cluster to another until getting a maximum modularity. Despite the fact that it gives very connected clusters in the larger graphs, the modularity measure cannot capture the small clusters [22]. Graph partitioning [6] and min-cut [14] are other methods used for the graph clustering which consist of splitting a graph into sub-graphs while optimizing the cut edge during the partitioning. The spectral clustering [33] is based on the graph density. It represents an input graph with a matrix and transforms this matrix so that to apply the basic clustering algorithm, like k-means [20].
A sampling-based distributed graph clustering method has been proposed in [30]. The proposed algorithm is based on the density of the graph to produce a set of sparse subgraphs. Graph embedding has also been used in the graph clustering. In [18], the authors discussed the use of the graph embedding technique to combine the structural and attributed similarity over the graph clustering. The above methods provide, as output, a list of clusters which are not really sufficient to understand the graph behavior. To address this issue, the Structural Clustering Algorithm for Networks (SCAN) was proposed in [35] aiming, not only to identify the clusters in a graph, but also to provide additional informations like hubs (vertices between one or more clusters) and outliers (vertices that do not belong to any cluster). These additional pieces of information can be used to detect vertices that can be considered as noise and also vertices that can be considered as bridges between clusters. The functional principle of SCAN is based on graph topology. It consists of grouping vertices that share the maximum number of neighbors. Moreover, it computes the similarity between all the edges of the graph in order to perform the clustering. The similarity computation step in SCAN is linear according to the number of edges, which degrades SCAN performance especially in case of large graphs. Structural graph clustering is one of the most effective clustering methods for differentiating the various types of vertices in a graph. In the litturature, several works have been proposed for the strucural graph clustering to overcome the drawback of SCAN. In [28], Shiokawa et al. proposed an extension of the basic SCAN algorithm, namely SCAN++. The proposed algorithm aims to introduce a new data structureof directly two-hop-away reachable node set (DTAR). This new data structure is the set of two-hop-away nodes from a given node that are likely to be in the same cluster as the given node. SCAN++ could save many structural similarity operations, since it avoids several computations of structural similarity by vertices that are shared between the neighbors of a vertex and its two-hop-away vertices.

In the same way, the authors in [9] suppose that the identification of core vertices represents an essential and expensive task in SCAN. Based on this assumption, they proposed a pruning method for identifying the core vertices after a pruning step, which aims to avoid a high number of structural similarity computations. To improve the performance and robustness of the basic SCAN, an algorithm named LinkSCAN* has been proposed in [24]. LinkSCAN* is based on a sampling method, which is applied on the edges of a given graph. This sampling aims to reduce the number of structural similarity operations that should be executed. However, LinkSCAN* provides approximate results.

Other works have proposed parallel implementations of SCAN algorithm [31] [25] [10] [29] [32]. In [31], the authors proposed an approach based on openMP library [7]. The author's aims were to ensure a parallel computation of the structural similarity and to show the impact of parallelism on the response time. Their method was proven to be faster than the basic SCAN algorithm. Other works have focused on the problem of dynamic graph clustering. In [25] and [10], the authors have extended SCAN algorithm to deal with the addition or removal of nodes and edges. Authors in [29], used the graphical processing unit (GPU) whose purpose is to parallelize the processing and to benefit from the high number of processing slots in the GPU which increase the degree of parallelism.
On the other hand, an index approach is proposed in [32] where the authors have emphasized the impact of the two parameters μ and ϵ on the running time of the structural graph clustering. Using the index approach a significant improvement has been shown b in terms of clustering running time. Again, the two sensitive parameters μ and ϵ have been taken into consideration by the authors in [34]. The main goal of the latter consists of estimate the best μ and ϵ values in order to enhance the accuracy and efficiency of the clustering results.

Most of the above presented works suffer from two major problems: (1) they do not deal with big graphs and (2) they do not consider already distributed/partitioned graphs.

Through Table 1, we have summarized the discussed approaches according to some features.
As shown in Table. 1, most proposed algorithms allow parallel processing but could not deal with very large graphs. It is also clear that none of the studied approaches allow distributed computing. In addition, these approaches are unable to process large and dynamic graphs. Also in some applications, like social network, graphs are distributed by nature. Thereby, using the discussed algorithms, it should aggregate in one machine all partitions of a distributed graph in order to run the graph clustering. based on this limits, in this work, we tackle the problem of large and dynamic graph clustering in a distributed setting.

Table 1: Comparative study on existing graph clustering methods

Approach	**Parallel**	**Distributed**	**Processing model**	**Graph partitioning**	**Main contributions**
SCAN [35]	No	No	Sequential	No	Basic implementation of structural graph clustering
SCAN++ [28]	No	No	Sequential	No	Reducing the number of similarity computations
AnySCAN [39]	Yes	No	Parallel	No	Parallelizing SCAN
pSCAN [9]	Yes	No	Parallel	No	Reducing the number of similarity computations
Index-based SCAN [32]	No	No	Sequential	No	Interactive SCAN
ppSCAN [11]	Yes	No	Parallel	No	Parallel version of pSCAN
SCAN based on GPU [29]	Yes	No	Parallel	No	A GPU-based version of SCAN
pm-SCAN [27]	Yes	No	Parallel	Yes	Graph partitioning to reduce the I/O costs

3 Background

3.1 Structural graph clustering: Basic concepts

Graph clustering consists in dividing a graph into several partitions or subgraphs. As with other clustering techniques, we must use one or more metrics to measure the similarity between two vertices or partitions in the graph. In the structural clustering technique, the graph structure or topology is splitted into a set of subgraphs that are relatively distant and a set of vertices that are strongly connected.

As a well-known algorithm for structural graph clustering, SCAN algorithm uses the structural similarity between vertices to perform the clustering. In addition, it provides other pieces of information like outliers and bridges. In what follows, we first present an overview of graphs and graph clustering with the SCAN algorithm.
Consider a graph $G = \{V, E\}$, where V is a set of vertices, and E is a set of edges between vertices. Each of those elements can represent a real property in real-word applications. Let u and v be two vertices in V. We denote by (u,v) an edge between u and v; u (resp. v) is said to be a neighbor of v (resp. u).
In the following, we extend some basic definitions of structural graph clustering, which will be used in our proposed algorithm.

Definition 1 (Structural neighborhood) The structural neighborhood of a vertex v,is denoted by $N(v)$, and represents all the neighbors of a given vertex $v \in V$, including the vertex v:

$$N(v) = \{u \in V | (v, u) \in E\} \cup \{v\} \tag{1}$$

Definition 2 (Structural similarity) The structural similarity between each pair of vertices (u, v) in E represents a number of shared structural neighbors between u and v. It is defined by $\sigma(u, v)$.

$$\sigma(u, v) = \frac{|N(u)| \cap |N(v)|}{\sqrt{|N(u)| \,.\, |N(v)|}} \tag{2}$$

After calculating the structural similarity with Eq. (2), SCAN uses two parameters to detect the core vertices in a given graph G.

Definition 3 (ϵ-neighborhood) Each vertex has a set of structural neighbors, like it is mentioned in Definition 1. To group one vertex and its neighbors in the same cluster, they must have a strong connection (denoted by ϵ-neighborhood) between them. SCAN uses a ϵ threshold and Eq. (3) to filter, for each vertex, its strongest connections. The ϵ-neighborhood is defined as follows.

$$N_\epsilon(u) = \{N(u) | \sigma(u, v) \geq \epsilon\} \tag{3}$$

The ϵ threshold $0 < \epsilon \leqslant 1$ shows to what extent two vertices u and v are connected based on the shared structural neighbors. In addition, it represents a metric with which the most important vertices, also called *core* vertices, are detected.

Definition 4 (Core) Core vertices detection is a fundamental step in SCAN algorithm. It consists of finding the dominant vertices in a given graph G. This step allows to build the set of clusters or a clustering mapping. A core vertex v is a vertex which has a sufficient number of neighbors strongly connected with it $N_\epsilon(v)$. We use μ as a minimum number of strong connected neighbors (see Definition 3). A core vertex is modeled as follows:

$$V_c = \{v \mid |N_\epsilon(v)| \geq \mu\} \tag{4}$$

Definition 5 (Border) Let v_c be a core vertex. v_c has two lists of structural neighbors: (1) weak connected neighbors to v_c, also called noise vertices $(N(V_c) \setminus N_\epsilon(V_c))$, and (2) strong connected neighbors called *reachable structured neighbors*. In our work, reachable structured neighbors are called border vertices. $N_\epsilon(V_c)$ represents the border vertices of a core vertex V_c.

Once the nodes and their borders are determined, it is straightforward to start a clustering step. To do so, we use the following definition:

Definition 6 (Cluster) A cluster C ($|C| \geq 1$) is a nonempty subset of vertices, where its construction is based on the set of core vertices and their border vertices. The main steps of clusters' building algorithm are the following: first, randomly chose a core from the cores' list and create a cluster C, then push the core and its borders into the same cluster. At the same time, the algorithm checks if the list of borders has a core vertex. Then, it inserts their borders into the same cluster and it applies this step recursively until adding all the borders of the connected cores. Finally, the algorithm chooses other core vertices and applies the same previous steps until checking all core vertices.

Among the fundamental information returned by SCAN, compared to other graph clustering algorithms, we mention bridge and outlier information, defined as below:

Definition 7 (Bridge and Outlier) The clustering step aggregates the core vertices and their borders into a set of clusters. However, some vertices do not have strong connections with a core vertex, which does not give the possibility to join any cluster. In this context, SCAN algorithm classifies these vertices into two families: bridges and outliers.
A vertex v, that is not part of any cluster and has at least two neighbors in different clusters, is called bridge. Otherwise it is considered as an outlier.

Algorithm 1: Basic SCAN algorithm

```
   Input  : A Graph G and parameters (μ ,ϵ)
   Output: ℂ, 𝔹, 𝕆
 1 foreach (u, v) ∈ E do
 2 |   Compute σ(u, v)
 3 end
 4 Core ← ∅
 5 foreach u ∈ V do
 6 |   if (|N_ϵ(u)| ≥ μ) then
 7 |   Core = Core ∪ {u} ;
 8 end
 9 Cluster ← ∅
10 foreach unprocessed core vertex u ∈ Core do
11 |   Cluster= ← {u}
12 |   Mark u as processed
13 |   foreach unprocessed border of vertex v ∈N_ϵ(u) do
14 |   |   Cluster= ←Cluster ∪ {v}
15 |   |   Mark v as processed
16 |   end
17 end
```

3.2 Running example

In this section, we explain through a running example, how SCAN algorithm works. Consider a graph G presented in Figure 1 and the parameters $\epsilon = 0.7$

and $\mu = 3$. In the first step, lines 1-3 of Algorithm 17 use Eq. (2) to compute the structural similarity for each edge $e \in \mathbb{E}$. Then, Eq. (3) (lines 5-8) is used to define the core vertices (see gray vertices in Fig. 1). After that, we proceed to the clustering step, then we apply Definition 6 (lines 11-16) to build the clustering schema. In our example we have four core vertices: 0,2,9 and 10. Each core has a list of border vertices (the neighbors that have strong connections with a core). In our example, vertex number 2 is a core. This later has the vertices number 1, 4, 5, 3 and 0 as the list of borders since they have strong connections with the core. The core and its borders build a cluster, and if a border is a core we join all its borders into the cluster. Like in our example, the vertex 2 is a core. Hence, we join all its borders (1, 3 ,5 ,4 including 0 as being a core vertex). In this case, if the vertex 3 is not a border of vertex 2 and it is a border of vertex 0, then vertex 3 should belong to the same cluster of vertex 2 (which is a core vertex), since it is a reachable border of vertex 2. The last step of SCAN algorithm consists in defining the bridges and the outliers. As shown in Fig. 1, we have two clusters. The first one is composed of vertices 1,2,3,4, and 5, whereas the second cluster is composed of nodes 8,9,10, and 11. The remaining vertices (6 and 7) must be categorized as outliers and bridges according to the Definition 7. In our example, vertex 7 has two connections with two different clusters, and vertex 6 has only one connection with one cluster which makes vertex 7 a bridge and vertex 6 an outlier.

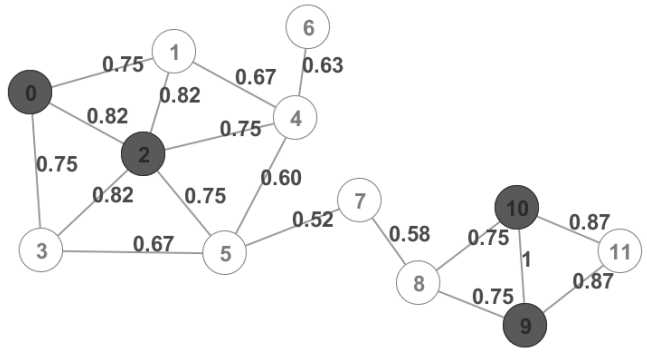

Fig. 1: Running example ($\epsilon = 0.7$ and $\mu = 3$).

4 DISCAN: A Distributed and Incremental Algorithm for Large-Scale Graph Clustering

In this section, we introduce DISCAN: a new distributed and incremental algorithm for structural graph clustering. Our proposed approach is based on a master/slave architecture and is implemented on top of BLADYG framework [3]. This latter is a distributed graph processing framework in which

the slaves are responsible for the execution of a specific computation and the master machine coordinates between all the slaves. The input data must be divided into sets of chunks (subgraphs in our case). Each chunck/partition is assigned to a worker, which performs a specific computation. The master machine orchestrates the workers' execution. In the following, we present the main three steps of DISCAN: (1) graph partitioning, (2) distributed graph clustering named DSCAN and (3) incremental graph clustering.

4.1 Distributed graph partitioning

In this step, we split the input graph G into several small partitions $P_1, P_2, ..P_n$, while keeping data consistency (graph structure). To ensure the consistency property while dividing the input graph, we must identify a list of cuts edges in order to have a global view of G. Usually, the graph partitioning problem is categorized under the family of NP-hard problems, that need to parse all the combinations in order to have the best partitioning result [13]. For this reason, we proposed an approximation and a distributed partitioning algorithm as a preliminary step for our distributed clustering algorithm. Algorithm 2 shows that, at the beginning, the master machine divides equitably an input graph file into sub-files according to the number of edges, and sends the sub-files to all the workers. Secondly, each worker gets a list of edges and vertices from its sub-file. Thereafter, it sends its list of vertices to all workers, in order to determine the frontier vertices so that to get the cuts edges. Afterwards, when each worker receives a list of vertices from its neighbor workers, it determines the vertices that belong to the current worker (partition). Consequently, these vertices are considered as frontier vertices in their partitions. In the last step, when each worker could determine the frontier vertices, it starts to fix the cuts edges, i.e. edges that have a frontier vertex.

4.2 Initial distributed graph clustering

Initial graph clustering (DSCAN) has two main steps: (1) local clustering step and (2) merging step.

Step 1: Local clustering. As presented in Algorithm 3, the input graph G is splitted into multiple subgraphs/partitions ($\mathbb{P}$), each one is assigned to a worker machine. The partitioning step, as mentioned in Section 4.1, is performed according to the α parameter, which refers to the number of worker machines (line 1). To overcome the loss of information during the partitioning step (edges connecting nodes in different partitions), frontier vertices are duplicated into neighboring partitions (line 5-8). Subsequently, for each partition P_i, a local clustering is performed (lines 9-14) on each worker machine. Fig. 2 shows a demonstrative example of the duplication step. The demonstrative example describes how the graph consistency will be ensured during the partitioning step.

Algorithm 2: Distributed partitioning

```
   Input  : Graph file GF as a text file, parameter (NP number of partitions)
   Output: ℙ set of partitions
 1 BLADYG initialization according to the NP parameter
 2 Master machine: split GF into a set of sub-files 𝔾𝔽
 3 foreach GF_i ∈ 𝔾𝔽 do
 4 |   Assign GF_i to worker W_i
 5 end
   /* In parallel                                                    */
 6 foreach Worker W_i ∈ 𝕎 do
 7 |   Get list of vertices and edges from GF_i
 8 |   Find the list of frontier vertices from the neighbor workers
 9 end
   /* In parallel                                                    */
10 foreach Worker W_i ∈ 𝕎 do
11 |   foreach Edge ∈ 𝔼 do
12 |   |   (a,b)=Edge
13 |   |   Let P_i^f the frontier vertices
14 |   |   if (a ∈ P_i^f ∪ b ∈ P_i^f ) then
15 |   |   |   Set the edge E as a cut edge
16 |   |   end
17 |   end
18 end
```

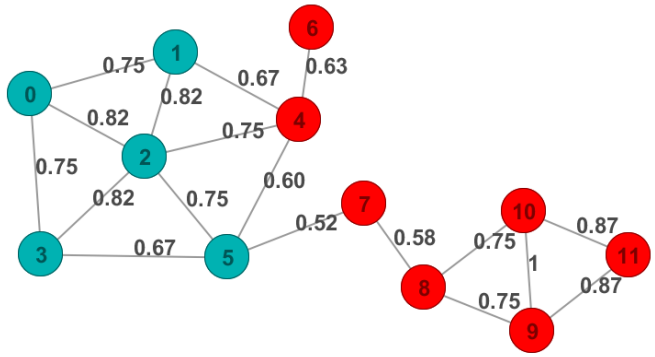

Fig. 2: Illustrative example

Assume that a graph G is partitioned into two partitions $P1$ (vertices in blue) and $P2$ (vertices in red), like it is depicted in Fig. 2, and each partition has a set of vertices connected with other partitions. We call them frontier vertices of a given partition P, and they are denoted by V_P^f. For example, V_{P1}^f={1,2,5} and V_{P2}^f={4,7}. Each $v \in V_P^f$ has a set of internal neighbors and external neighbors. For example, vertex 4 is a vertex of the partition $P2$. It has the vertex 6 as internal and the vertices 1,2 and 5 as external neighbors. Computing the structural similarity of a frontier vertex u considers that all the neighbors of u belong to the same partition. Thereby, we duplicate all frontier vertices in partition P to all neighbor partitions and we call them

Algorithm 3: DSCAN

Input : Graph G , parameters (μ, ϵ, α)
Output: $\mathbb{C}$lusters, $\mathbb{B}$ridges, $\mathbb{O}$utliers

/* Divide G into subgraphs $\mathbb{G}$ = $\{G_1 G_2 .. G_\alpha\}$ according to parameter α */
$\mathbb{P} \longleftarrow$ ***Partition*** (G,α)
BLADYG initialization
/* In parallel */
foreach $P_i \in \mathbb{P}$ **do**
  Assign P_i *to* W_i
end
/* In parallel */
/* Step 1: Local clustering */
foreach $Worker\ W_i \in \mathbb{W}$ **do**
  Let P_i *the current partition*
  Find the frontier vertices in P_i *and duplicate them into neighbor partitions*
end
/* In parallel */
foreach $Worker \in \mathbb{W}$ **do**
  Compute the structural similarity of a partition P_i *using* V^f *list*
  Retrieve local Cores and Borders in P_i
  Build local clusters in P_i
  Find local Bridges and Outliers in P_i
end
/* Step 2: Merging */
All workers exchange theirs local clusters between them; using Worker2Worker message
foreach $Worker\ W_i \in \mathbb{W}$ **do**
  if $(C_1 \cap C_2 \cap .. \cap C_\alpha = \mathbb{V}$; *and* $\exists V_i \in \mathbb{C}ore)$ **then**
    $C \leftarrow$ ***Merge***$(C_1, C_2, ..C_\alpha)$
    Send C *to the master*
  end
  else
    Send local clusters to the master
  end
  for V_i ***IN*** $\mathbb{O}utliers$ **do**
    if $(V_i \in \mathbb{C}ore \cup \mathbb{B}order \cup \mathbb{B}ridges\ )$ **then**
      Remove V_i *from the list of* $\mathbb{O}utliers$
    end
  end
  for V_i ***IN*** $\mathbb{O}utliers$ **do**
    $Nb_{Connections} \leftarrow 0$
    for C_i ***IN*** $\mathbb{C}lusters$ **do**
      if $(N(V_i) \cap C \neq \varnothing)$ **then**
        $Nb_{Connections}++$
      end
    end
    if $(\ Nb_{Connections} \geq 2\ )$ **then**
      Add V_i *to* $\mathbb{B}ridges$
      Remove V_i *from* $\mathbb{O}utliers$
    end
  end
end
Send $\mathbb{C}lusters$,$\mathbb{B}ridges$,$\mathbb{O}utliers$ *to the master using Worker2Master message*

external vertices. For example in our running example, $P1$ has frontier vertices V^f_{P1}={1,2,5} and $P2$ is the neighboring partition. Thus, we must duplicate V^f_{P1} into $P2$ to ensure the accuracy of structural similarity of V^f_{P2} (see Fig. 4).

After that, when we apply a local clustering on $P1$ and $P2$, we will find several conflicts such as the vertex $v \in V^f_{P1}$ is a core vertex in $P1$, and an outlier in $P2$. These conflicts should be avoided in the merging step.

Step 2: Merging. The distribution of similarity computation and the local clustering step can improve the response time of DSCAN, compared to the basic sequential algorithm. However, we should take into consideration the exactness of the returned results compared to those of the basic SCAN. To ensure the same result of basic SCAN, we defined a set of scenarios regarding the merging step. These scenarios will repetitively be applied to every two partitions of G, until combining all the partitions (see Algorithm 3, lines 16-37). For each pair of partitions P_i and P_j, a merging function is executed to combine the local results from P_i and P_j and store them in global variables like clusters, borders, bridges and outliers. Algorithm 3 also achieves several scenarios (Lemmas 1, 2 and 3) to solve the encountered conflicts, mentioned below:

Lemma 1 ***(Merging local clusters)*** *Let $\mathbb{C}_1$ and $\mathbb{C}_2$ two sets of local clusters in different partitions P_1 and P_2, respectively. $\exists c_1 \in \mathbb{C}_1$ and $\exists c_2 \in \mathbb{C}_2$, $Core(c_1) \cap Core(c_2) \neq \emptyset$.*

Proof Let C_i be a cluster that groups a set of border and core vertices. If C_i shares at least one core vertex c with another cluster C_j, then c has a set of borders in C_i and C_j, and all the vertices in C_i and C_j are reachable from c. Hence, C_i and C_j should be merged into the same cluster.

Lemma 2 ***(Outlier to Bridge)*** *Let $\mathbb{C}_1$ and $\mathbb{C}_2$ two sets of local clusters in partitions P_1 and P_2, respectively. $\exists$ C_1 and C_2 two clusters that belong to the two different sets of clusters $\mathbb{C}_1$ and $\mathbb{C}_2$. In addition, $\exists$ o an outlier in both partitions P_1 and P_2, with $N(o) \cap c_1 \neq \emptyset$ and $N(o) \cap c_2 \neq \emptyset$*

Proof If C_i and C_j $(i \neq j)$ share an outlier o, this means that o is weakly connected with the two clusters C_i and C_j. Hence, according to the Definition 7, o should be changed to a bridge vertex.

Lemma 3 ***(Bridge to Outlier)*** *Let c_1 and c_2 two local clusters in the partitions P_1 and P_2, respectively.*
$\exists$ a bridge b according to only two clusters c_1 and c_2, when c_1 and c_2 two clusters that will be merged into one cluster, b should be changed into an outlier.

Proof Let C_i and C_j two clusters that have a set of vertices (borders or cores), and a set of bridges with other local clusters, and $\exists$ b a bridge vertex according to the two clusters C_i and C_j only, where $i \neq j$. In the merging step and according to Lemma 1, if one or several clusters share at least a core vertex, then they will be merged into a single cluster. In this case, $(C_i,C_j) \Rightarrow C$ which makes b be weakly connected with only one cluster C, then according to Definition 7, b should change its status from bridge to outlier.

For instance looking at lines 16 to 22, we have focused on the shared cores between two clusters and the case when they share at least one core. According to Lemma 1, we should merge them into one single cluster. Subsequently, in lines 23-25, we verify for each outlier if it does not belong to some sets of cores,

borders or bridges. In this case, we must remove it from the list of outliers. Otherwise, a vertex v should be changed as a bridge if it is classified as an outlier in the two clusters C_i and C_j that are not in the same partition, and if it has two connections with different clusters in the merging step.

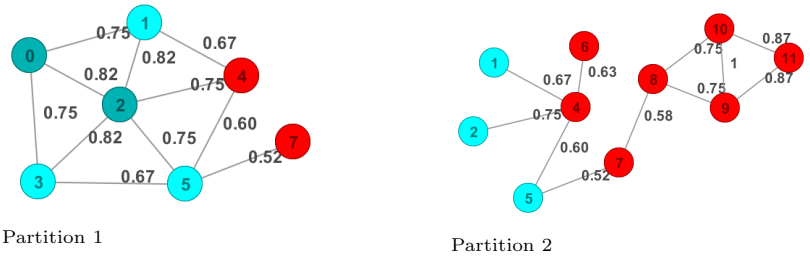

Fig. 4: The partitioned graph G used in the running example in Section 3

4.2.1 Illustrative example

Fig. 4 shows a demonstrative example of a graph clustering using DSCAN. In this example, we use the same parameters (ϵ and μ) of the running example in Section 3, and two partitions ($P1$ and $P2$). In the first step of DSCAN, the input graph G is divided into two partitions $P1$ and $P1$ as it presented in Fig. 4, where the blue vertices represent the first partition and the red vertices represent the second partition. In the second step, DSCAN duplicates the frontier vertices in each pair of adjacent partitions. For example, the blue vertices 1,2 and 5 are duplicated in partition $P2$ since they represent frontier vertices in their partition. In the same way, the vertices 4 and 7 are duplicated in the partition $P2$. In the third step, all the workers perform the similarity computation, check the status for each vertex and build the local clusters, as it is shown in Fig 4. The last step of DSCAN consists of combining all local results returned by each worker. As shown in Fig. 4, there are some conflicts in terms of node status. For example, vertex 2 is a core vertex in $P1$ whereas it is a noise (outlier) vertex in $P2$. In the same way, vertex 4 is a border in $P1$ and a noise in $P2$. Then, after building the local clusters, $P1$ has one cluster (1,2,3,4, and 5) in which vertex 7 is a noise vertex in $P2$. As for $P2$, it groups the vertices 8,9, 10 and 11 as a cluster and the remaining vertices (4,5,1,2,7 and 6) are considered as noise vertices including vertex 7. This latter is a bridge according to basic SCAN (see running example in Section 3)). In the merging step, DSCAN considers that the vertex 7 has two weak connections with two different clusters. Thus vertex 7 is marked as a bridge.

4.2.2 Time complexity analysis of DSCAN

Let $\mathbb{E}$ be the set of all edges (internal and cut edges), $\mathbb{V}$ be the set of all vertices(internal and external) and the set of core vertices is $\mathbb{V}_c$. The time complexity of DSCAN is like basic SCAN, and can be divided into three step, (i) the structural neighborhood step, when DSCAN according to Definition 1 aggregates for each vertex its list of neighbors. Thus, the complexity is of the order of $\mathcal{O}((v_i, v_j)_{i,j=1}^{|V|}, (u_i, u_j) \in \mathbb{E})$=$\mathcal{O}(|\mathbb{E}|)$.(ii) the time complexity of the computation of structural similarity which is defined by Definition 2. In this step, DSCAN enumerates for each edge (u, v) the set of common neighbors. Therefore, the time complexity is $\mathcal{O} = (min(|N(u)|, |N(v)|).|\mathbb{E}|)$. (iii) DSCAN perfoms the clustering step using the μ in order filter the core vertices, and groups those which have strong connections into a same cluster. Like this, the complexity is $\mathcal{O} = (|V_c| - |(v_i, v_j)|_{i,j=1}^{|V|}, v_i, v_j \in V_c$ and $\sigma(v_i, vj) \geq \epsilon)$

4.3 Dynamic and distributed graph maintenance

Consider a large and dynamic graph G, which will undergo several changes over time such as adding or deleting edges or/and vertices. When a clustering will be performed, the classic approaches re-run the clustering from scratch, which is obviously very expensive especially in large graphs. In this work, we propose an incremental and distributed algorithm for large and dynamic graph clustering. The main idea of our approach is to determine in each update the vertices and the edges concerned by this update, and according to that, we apply the new updates on the old clustering results.
Assume that $G = (V,E)$ is going to undergo several updates U over time. U can be adding/deleting vertices or edges from/to G. These changes can affect the similarity values and some vertex status (borders, core, etc). Consequently, the clustering schema may be affected in several situations.

Definition 8 Each vertex $u_i \in U$ generates modifications of several structural similarities, which is defined as the affected edges E_A.These vertices also change the status of several vertices and are noted as affected vertices V_A.

Definition 9 (Affected edges) When adding a vertex V_{new} associated with an existing vertex V_{old}, new edge (V_{new}, V_{old}) will be created and a set of edges denoted E_A will be marked as affected edges and should be updated. Note that $e \in E_A$ and $e = (u, v)$, where $u \in V_{old}$ and $v = v_{new}$.

In the case of removing a vertex V_{TBD}, a set of edges noted as E_{TBD} should be deleted from G. Therefore, a V_c of vertices is concerned by this update. Here $V_c = (u, v) \in E_{TBD} \setminus V_{TBD}$ and, according to V_c, the affected edges are $E_A = E_i \subset G$ and $E_i = (u, v)$ / {u,v} $\in V_c$.

Also when adding or deleting an edge (v_i, v_j), the list of neighbors of both vertices v_i and v_j will be changed which can produce several other edges. The affected edges update their structural similarities. E_A is a set of edges (u, v) $\in E$ and $v_i = u$ or $v_j = v$

Definition 10 (Affected Vertices) For each update on G, some structural similarities can be changed. Thus, some vertices change their status (border, core vertices). These vertices are defined as affected vertices V_A. V_A can be divided into two sub-sets: (i) directly affected V_{AD}, which update their status depending directly on E_A, and are defined as V_{AD}, (ii) indirectly affected V_{ID} which depend on the status of V_{AD}. For example, when a core vertex changes his status to outlier or border, then all these border vertices will become affected indirectly according to this update. Let E_A=(v_1, v_2) be following an update V_{AD} =$V_{AD} \cup v_1 \cup v_2$, while V_{AI} are all neighbors of V_{AD}. Other vertices can be affected indirectly with several updates, and are dependent on the changed clusters. These vertices are defined the affected bridges and noted as V_{AB}. V_{AB} is the list of outliers or bridges which are connected with affected clusters described in the Definition 11.

Definition 11 (Affected clusters) The affected vertices in several situations can affect the clustering schema. The Affected (concerned) clusters represent the sub set of clusters which contain one or more affected vertices, denoted by $C_{affected}$. $C_{affected}$=$c \in \mathbb{C}$, $\exists v \in (v_{AD} \cup v_{ID})$ and $v \in c$.

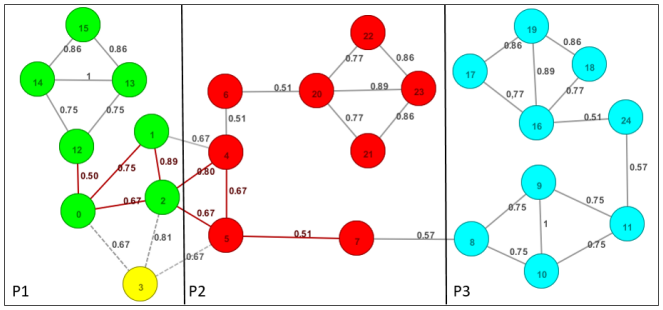

Fig. 5: A partitioned graph G (P1, P2, P3), affected vertices, edges when deleting vertex 3

Lemma 4 ***Change of structural similarity of affected edges*** *E_A are mainly depend on the Definition 1. The similarity value is based on the neighbors of both vertices of an edge. Therefore, each update (adding/deleting an edge/vertex) can change the list of neighbors of some vertices. As a result, some edges must change their structural similarity.*

Lemma 5 ***Change in the status of affected vertex*** *According to the Definition 4 and the Definition 3, for some updates in the structural similarity, when the similarity value exceeds or decreases the threshold ϵ the status of this edge maybe changed to strong or weakly connection. Thereby, this change can affect the status of each vertex $v \in V_{AD}$. Consequently, according to 5 the V_DA can change other vertices status indirectly which are defined as indirectly*

affected vertices V_{IA}. Rationally, those changes can affect the mapping of the clusters, and several clusters may be changed. Thereby, according to Definition 7, $\exists$ a vertex $v \in$ borders list where v only depends on two clusters c_1 and c_2. When c_1 and c_2 will be merged into one cluster c_{new}, v will be weakly connected with only c_{new}. In this case and according to Definition 7 v will be an outlier vertex. In the same way, when we have a cluster c and v is an outlier to c with two weak connections related to v_1 and $v2$, after some updates c becomes splited to two clusters c_1 and c_2 when v_1 and v_2 are respectively in c_1 and $C2$. In this case, v become a bridge between the new clusters c_1 and $C2$ according to the connection of v with v_1 and v_2.

Fig. 5 shows an example of a graph G with 3 partitions forwhich we deleted the vertex 3 (yellow vertex). Then, all affected edges E_A (red edges) change their structural similarities since the vertices 0,2 and 5 (old neighbors of the deleted vertex 3) lost the vertex 3 as a neighbor, and all the edges associated with the vertex (3) are removed from G. This new update affects only a subset of vertices in G. The example shows that the vertex 2 ($v_2 \in v_A$) lost its core status, whereas vertex 9 keeps its core status ($v_9 \notin v_A$). Thus, E_A can change several vertices ' status directly (v_{AD}) or indirectly (v_{AI}). Like it is shown in our example, vertex 2 $\in v_{AD}$, it lost on strong connection with the removed vertex which changed to an outleir in this case. In the same way, v_{AD} can change the status of v_{AI}. Vertices 0,1,4 and 5 were considered as border vertices according to vertex 2 (in the initial graph), and in our case, when the vertex 2 lost the core status, as a result v_{AI} are changed to outliers. According to Definition 11, v_{AD} and v_{AI} can change the old clustering schema and our example confirm that. Partition 1 of G groups two clusters. The first one is C_1 = { 12,13,14} and the second one is C_2 ={0,1,2,3,4,5} (4 and 5 are external vertices). Thus, $C1$ and $C2$ are affected because $C1$ contains the vertex 12 as an affected vertex, and in the $C2$, all its vertices are affected because of the removed vertex. These clusters can affect in the last step other vertices (bridge or outlier vertices). Like it is depicted in Fig. 5 C_2 should be removed since it does not have any core vertex, and the vertices 6 and 7 were considered as bridges with $C2$. Thereby, these bridges are susceptible to change their status.

Algorithm 4 describes the main steps of the proposed DISCAN. DISCAN provides real-time graph maintenance and a micro-batch clustering. Moreover, it performs the new clustering after several changes, in order to optimize the incremental clustering. We present below the main steps of DISCAN:

Step 1: Graph maintenance. In this step, DISCAN runs the graph maintenance in real time. In each update U, it (i) updates the graph G structure, (ii) checks the affected edges E_A, and (iii) recomputes the similarities of E_A (lines4-12 of the pseud-code 4). Then, it gets V_A according to E_A in order to memorize them in a global variable. Like discussed in Definition 11, V_A can affect several clusters in the current partition or in other partitions. Thereby, each worker shares its V_A with all workers so that each one determines the affected clusters and saves them. Once the affected clusters are fixed, they affect some bridges or outliers vertices. So the affected bridges should be added

to the list of affected vertices to be checked in the next step.

Step 2: Incremental local clustering. This step is performed in micro-batch processing mode, after a number of updates or using a window time. In both methods, DISCAN runs the some scenarios. These latter consist of two choices for a user who selects one of them according to the number of changes per seconds or according to other needs such as the graph size, graph evolution. Therefore, in each batch, DISCAN first checks the core vertices. The checking is performed only on the affected vertices V_A, where in general the $|V_A| \ll |V|$.

Secondly, DISCAN checks the border vertices like in the first clustering but using only E_A, since the border vertices are dependent on the strong connections with core vertices. For this reason, DISCAN checks the core vertices according to E_A only.

After that, the updated cores and borders vertices will change the clustering schema depending of the affected vertices (cores, borders) like it is described in Definition 11. The affected clusters feat are four possible cases owing to each update of a graph: (1) split one cluster to small clusters, (2) remove an existing cluster, (3) build new clusters, (4) update the existing clusters and (5) merge two or more clusters. Consider set of cores and theirs borders, and the core vertices have strong connections beteen them. After the removal of a strong connection between two cores (c_1,c_2) and (c_1,c_2), then the clusters should be splited into sub-clusters depending on c_1 and c_2. For the case (2), each cluster is built on a list of cores and theirs borders. If one cluster does have any core vertex it should be removed. Sometimes, a new update makes an outlier vertex to a core vertex c_{new}. If this vertex has a strong connection with any other old core c, it joins the cluster of c, like in the case (4). If c_{new} does have any strong connections with any other core it builds a new cluster, like the case (3). In the last case (5), if there exist a new strong connection between two cores in two different clusters, then we merge these latter. In order to deal with all these cases, we remove all the affected clusters and re-build a new clustering in a same way of the first clusetring presented in Definition 6 and using only the affected vertices (cores and borders vertices). The new clustering is performed only on the changed vertices which will be very faster compared to the first clustering.

Finally, DISCAN uses the new clusters and the remaining affected vertices to check if they represent new border vertices.

Step 3: Merge the new updates. After each batch, each worker starts to merge the new updates (see Algorithm 4 lines 19-23). The master machine gets all the affected vertices from all workers in order to facilitate the combination of the new changes. Here some clusters should be verified. In each batch, the master keeps only the past unchanged clusters and requests all workers to get the changed clusters. Thereby, all workers combine the updated clusters eventually including new clusters. In the merging step, DISCAN uses the same scenarios as in the first clustering (see Lemma 1). If at most one cluster shares at least one core vertex, they can be merged into one cluster. In the next step, after getting the new clusters, the master machine requests all workers

to get theirs borders vertices. Then, each worker filters only the affected vertices which will belong to its local bridges, and sends them to the master. The rest of affected vertices will be added to the global list of outliers. Finally, DISCAN initializes the global affected vertices and global affected clusters to empty lists which will to be used in the next batch.

Algorithm 4: DISCAN: Incremental DSCAN

Input : Initial graph G as a text file, parameter (NP number of partitions), a new update U
Output: $\mathbb{C}, \mathbb{B}, \mathbb{O}$

```
/* Global affected vertices and clusters in each worker                    */
```
$\mathbb{G}lobalAffectedVertices = \emptyset$
$\mathbb{G}lobalAffectedClusters = \emptyset$
```
/* Step 1: Graph maintenance                                               */
/* Update the initial graph in each new update and get the affected vertices
   and the affected clusters                                               */
```
Master machine: send a new update u to all workers
```
/* In parallel                                                             */
```
foreach $Worker_i$ $w_i \in \mathbb{W}$ **do**
    Update the current partition according to U
    Get E_A
    Recompute the similarities of E_A
    Get the immediately affected vertices V_I
    Share the immediately affected vertices V_I *with all workers*
    Get the indirectly affected vertices V_{Ind}, *and the affected clusters* C_A
end
```
/* Step 2: Local clustering schema maintenance                             */
/* Update the old local clustering schema according to the global affected
   vertices                                                                */
```
foreach $Worker_i$ $w_i \in \mathbb{W}$ **do**
    Check Core vertices from GlobalAffectedVertices
    Check Border vertices from GlobalAffectedVertices
    Check affected clusters
    Check affected bridges
end
```
/* Step 3: Merge the new updates                                           */
```
foreach $Worker_i$ $w_i \in \mathbb{W}$ **do**
    Merge all affected clusters from all workers
    Check new Bridge vertices according to the new clusters
    Check the remaining outlier vertices
end
```
/* Reset the global affected vertices and clusters in each worker          */
```
$\mathbb{G}lobalAffectedVertices = \emptyset$
$\mathbb{G}lobalAffectedClusters = \emptyset$
Send $\mathbb{C}$,$\mathbb{B}$,$\mathbb{O}$ *to master using Worker2Master message*

4.3.1 Time complexity analysis of DISCAN

The time complexity of DISCAN mainly depends on the affected edges and the affected vertices in each new update. Moreover, this complexity is very low

compared to the initial complexity of DSCAN. Since the number of affected edges $|E_A| \ll |E|$ and the number of affected vertices $|V_A| \ll |V|$. In DISCAN, the time complexity is based on three steps: (1) The graph maintenance, the complexity in this step is about $\mathcal{O} = |E|/n$ in worst case and $\mathcal{O} = 1$ in best case, where n is the number of partitions. (2) In the similarity computation the complexity is $\mathcal{O} = (min(|N(u)|, |N(v)|).|\mathbb{E}_A|)$, and (3) also the clustering step manly depends on V_A. Thereby the incremental clustering complexity is about $\mathcal{O} = (|V_c \in V_A| - |(v_i, v_j)|_{i,j=1}^{|V_A|}, v_i, v_j \in V_c$ and $\sigma(v_i, vj) \geq \epsilon)$.

5 Experiments

In this section, we present our experimental study and we evaluate the effectiveness and efficiency of our proposed algorithm for structural clustering of large and dynamic distributed graphs. We performed some experiments like the impact of graph size on DISCAN in the case of static graphs and the number of updates in a dynamic graph setting. We also evaluated some features of DISCAN.

5.1 Experimental protocol

In the first experiment part, we compared DISCAN with four existing structural graph clustering algorithms in the case of static graphs. Then, in the second experiment is cared out on dynamic graphs, and we performed some other experience in order to evaluate some features of DISCAN such as the impact of number of update on the running time.

1. Basic SCAN[1].
2. pSCAN: a pruning SCAN algorithm[2].
3. AnyScan: a parallel implementation of basic SCAN using OpenMP library[3].
4. ppSCAN: a pruning and parallel SCAN implementation[4].

The above mentioned algorithms are implemented with C language. Thus, we used the GCC/GNU compiler to build their binary versions. The compared algorithms are divided into two categories: (1) centralized algorithms such as SCAN and pSCAN, and (2) parallel algorithms such as AnyScan and ppSCAN. To run both centralized and parallel algorithms, we used a *T3.2xlarge* virtual machine on Amazon EC2. This machine is equipped with an 8 vCPU Intel Skylake CPUs at 2.5 GHz and 32 GB of main memory on a Ubuntu 16.04 server distribution. In order to evaluate DISCAN, we used a cluster of 10 machines,

[1] https://github.com/eXascaleInfolab/pSCANdeploymet
[2] https://github.com/RapidsAtHKUST/ppSCAN/tree/master/SCANVariants/anySCAN
[3] https://github.com/RapidsAtHKUST/ppSCAN/tree/master/SCANVariants/anySCAN
[4] https://github.com/RapidsAtHKUST/ppSCAN/tree/master/ppSCAN-release

Table 2: Graph datasets

Dataset name	Number of vertices	Number of edges	Diameter	Avg. CC
G1: California road network	1 965 206	2 766 607	849	0.04
G2: Youtube	1 134 890	2 987 624	20	0.08
G3: Orkut	3 072 441	117 185 083	9	0.16
G4: LiveJournal	3 997 962	34 681 189	17	0.28
G5: Friendster	65 608 366	1 806 067 135	32	0.16

each of them is equipped with a 4Ghz CPU, 8 GB of main memory and operating with Linux Ubuntu 16.04. This link https://discan.yo.fr/DISCAN.html provides some details about the configuration, deployment of DISCAN, and also a user guide is presented.

5.2 Experimental data

For all test cases of static graphs, we used real-world graphs (see Table 2) obtained from the Stanford Network Analysis Project (SNAP) [5]. Then, in the dynamic graphs, we generate a set of new edges for each used graph in the set of static datasets (G2,G3,G4 and G5). The new updates are grouped as a set of stackor batches. In each one, we generated five batches with different sizes that vary from 2000 to 10000 new edges.

5.3 BLADYG framework

BLADYG is a distributed and parallel graph processing framework that runs on a commodity hardware. The architecture of BLADYG is based on a master/slaves topology. BLADYG starts by reading the input graph from many different sources, which can be local or distributed files such as Hadoop Distributed File System (HDFS) and Amazon Simple Storage Service (Amazon S3). The communication model used by BLADYG is the message passing technique, which consists in sending messages explicitly from one component to another in order to get or send useful data during the graph processing. In the same way, BLADYG defines two types of messages: (1) worker-to-worker messages, and (2) master-to-worker messages. BLADYG allows its users to implement their own partitioning techniques.

[5] https://snap.stanford.edu/data/

5.4 Experimental results

Speedup of initial distributed graph clustering. DISCAN is designed to deal with dynamic graphs, however, it start in first step by clustering of static graph (initial clustering). Then, in this experiment we evaluated the speedup of DISCAN in a static graphs compared to the basic SCAN and its variants presented in Section 5.1. The compared algorithms use different graph representations. In fact, AnyScan and SCAN implementations use the adjacency list representation [15], whereas both pSCAN and ppSCAN use the Compressed Sparse Row (CSR) format [16]. In our proposed algorithm, we used an edge list format, in which each line represents one edge of the graph. The incompatibility of the graph representations poses an additional transformation cost while evaluating the studied methods. For example, the transformation of the live journal dataset from edge list to adjacency list takes around 100 seconds using one machine equiped with an 8 vCPU and 32 GB of main memory.

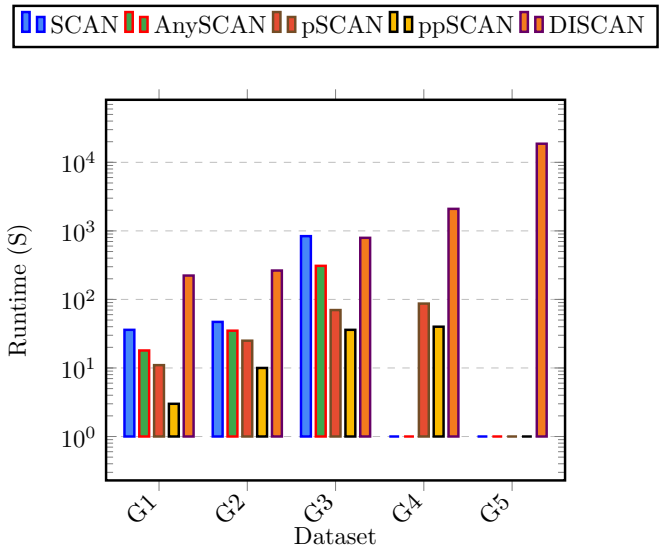

Fig. 6: Impact of the graph size on the processing time of SCAN variants and DSCAN

As shown in Fig. 6, our approach is slower than the other algorithms in the case of small graphs (G1,G2,G3) and there was a very large gap between DISCAN and the other algorithms especially with pSCAN and ppSCAN. On the other hand, this gap become reduced when the graph size increases (case of G4 dataset). The plots bars in Fig. 6 shows a gap of 12x between DISCAN and basic SCAN with the G1 dataset and 2x only with the G4 dataset. We notice that the gap between DISCAN and ppSCAN depends mainly on the size of the used dataset. For example, with the G1 dataset, the gap between DISCAN

and ppScan is about 20x, while this gap is reduced to 11x for the G4 dataset. This can be explained by reach the pruning step of pSCAN, which exempts several similarity computings during the clustering step. It is important to mention that DISCAN is a distributed implementation of SCAN and the other studied algorithms are centralized. This leads to additional costs related to data distribution, synchronization and communication. Fig 6 also shows that with the modest hardware configurations, only DISCAN can scale with large graphs like the G5 dataset.

SpeedUp on incremental clustering. The main goal of this experiment is to compare DISCAN, with other algorithms, since there is no incremental algorithm for the structural graph clustering, we evaluated the speedup of DISCAN compared to the fastest existing algorithm (pSCAN and ppSCAN) which are presented in the previous section. Fig. 7 shows the running time of the tested algorithms with different graphs. In the first time, we start by running all the algorithms with the used graphs, and in each moment, we added a new batch of updates in different sizes. In almost used graphs, DISCAN is slow compared to other algorithms, except for pSCAN which does not support the G5 graph. Despite that pSCAN and ppSCAN re-execute the clustering from scratch, they are faster than DISCAN in case of G2 and have the same response time in the case of G3. In the case of the G4, DISCAN is initially better than pSCAN and ppSCAN. Initially, ppSCAN and pSCAN are faster than DISCAN, but the gap in running time begins to narrow between G2 and G4. Then, when we add a new batch to the initial graphs, DISCAN becomes faster than pSCAN and ppSCAN. In the case of G3, Fig. 7 shows a gap of 2X between DISCAN and pSCAN and almost the same running time of DISCAN and ppSCAN. Furthermore, the gap between DISCAN and pSCAN starts to increase, to finally reach 3X, and 10% between DISCAN and ppSCAN.

DISCAN scalabilily. Fig. 8 shows the scalability of DISCAN w.r.t. the number of workers, with G2 and G3 datasets and for different update batches using the default parameters ϵ=0.5, μ=3.

Overall, the number of workers affects the running time in terms of size of the used dataset and the number of updates. In Fig. 8 with G2 dataset, all curves have almost the same look. They decrease for 2 to 4 workers, but with varying the degrees of improvement depending on the number of updates. These varying improvements are between 10% and 50%, respectively for 2000 and 10000 updates. Afterwards, the curves grow until 8 workers, then comes down. In the G3 dataset, the scalability becomes very high. The response time decreases depending on the number of workers. This improvement is about 30% and 40% when related to the number of updates. This behavior can be attributed to the number of affected vertices and clusters. When this number becomes small, the needed resources (e.g.,workers) should be small. Nevertheless, using many workers, DISCAN uses a lot of communication which increases the processing time.

Impact of the batch size. The main goal of this experiment is to evaluate the impact of batch size on the response time. Initially, we run DISCAN with different datasets (G2, G3, G4 and G5) and different batch sizes (between

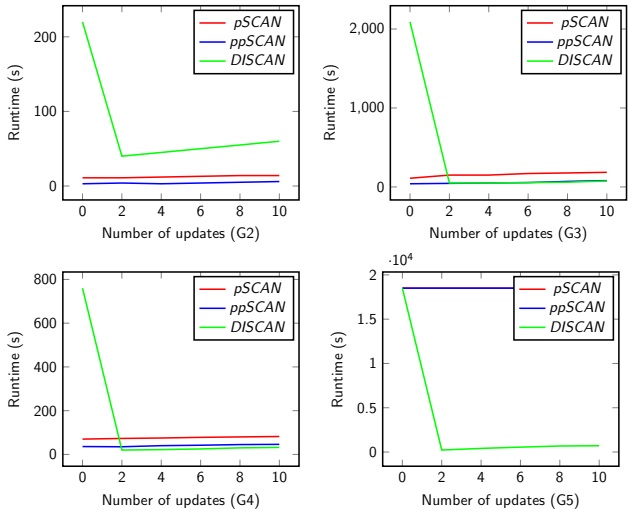

Fig. 7: Impact of the number of updates on the processing time of the pSCAN, ppSCAN and DISCAN (ϵ=0.5, μ=3)

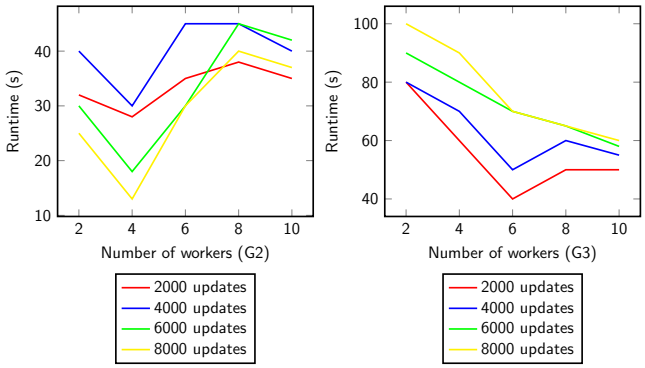

Fig. 8: Impact of the workers' number on the processing time of DISCAN (ϵ=0.5, μ=3)

2000 en 1000) on a cluster of 10 machines and using the default parameters ϵ=0.5, μ=3. Fig. 9 shows that the running time increases for all curves. This is explained by the initial graph size, as we noticed in Fig. 7, but also in function of the batch size. We can also notice that the increase depends on the size of the initial graph and the batch size. The increase rate is about 5%, 50% and 150%, respectively for G2, G4 and G5. This can mainly be explained by the graph size, since for each update DISCAN in the graph maintenance step check all affected vertices and edges in order to update the affected similarities. Therefore, this checking depends on the graph size. On the other hand, in the reclustering step, DISCAN first performs the clustering

on the affected vertices. Then, it merges all the affected clusters in the master machine. Thus, the running time of the merging step depends on the number of affected edges because we must check all affected clusters. The affected clusters' checking requires using the communications between the master and other workers. In this way, when the batch is small, i.e. few affected vertices, a little communication should be used.

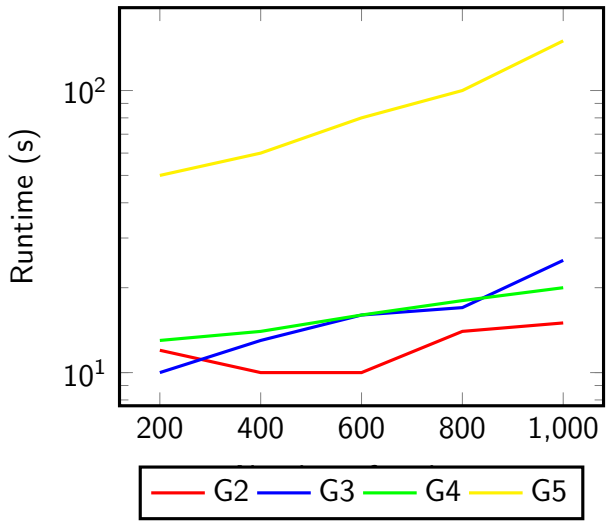

Fig. 9: Impact of the batch size on the running time of DISCAN (ϵ=0.5, μ=3)

Impact of the vertex type (internal vs external). Each update on an existing graph maybe in internal or frontier (external) vertices. With, this experiment we show the difference in terms of the running time between the updates in internal vertices and external vertices. We generated for G2,G3,G4 and G5 several internal and external batches of adding new edges. After that, we run DISCAN using them in order to compare their response times. As depicted in the Fig. 10 the updates on the frontier vertices are very expensive with regard to running time. This latters' gap trends to around 50% between internal and external vertices updates. This can be explained by the graph' maintenance step performed by DISCAN, that consumes a lot of processing time in the case of external vertices, compared to internal vertices. In fact, in the internal vertices all maintenance operations are done locally (in the same partition), whereas the graph maintenance is done in several partitions, which requires additional costs. These latter's are mainly related to the duplication of the frontier partitions and obviously the communication between all the affected workers.

DISCAN Steps. We notice, from the previous experiment, that the response

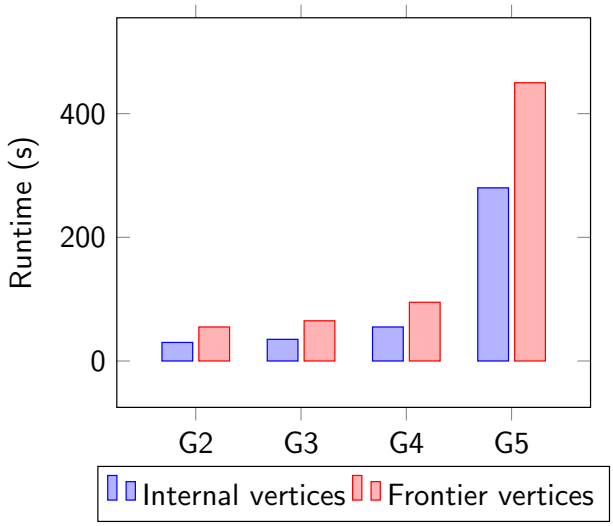

Fig. 10: Impact of the vertices' type (internal or external) on the processing time (ϵ=0.5, μ=3)

time depends on the type of vertices that can be internal or external. Especially, for the external vertices the response time is slow because graph ' maintenance is expensive, compared to the case of internal vertices.

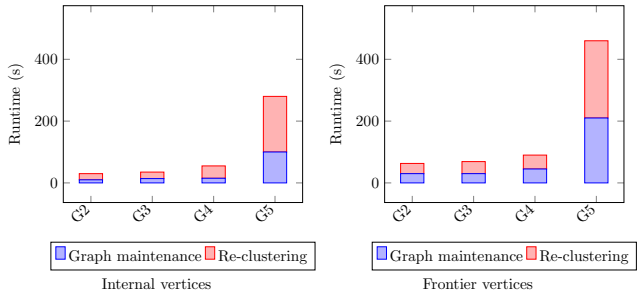

Fig. 11: Performance of DISCAN steps (ϵ=0.5, μ=3)

Fig. 11 clearly shows a span of the difference between graph maintenance and reclustering steps for both the internal and external vertices. The reclustering step takes about 30% of the global running time in case of internal vertices, while this rate is about 50% in case of frontiers updates. On the other hand, the difference between internal and external updates in terms of running time during the reclustering is about 3X. This is understandable because

the external updates need additional computing time like the duplication of frontier vertices and the graph maintenance according to the neighboring partitions. The additional treatment sometimes requires several communications with other workers.

6 Conclusion

In this paper, we proposed DISCAN, a distributed and incremental algorithm for big dynamic graph clustering based on the structural similarity. DISCAN is build on top of based on a distributed and master/slaves architecture which makes it scalable and works on the community of modest machines. The proposed algorithm is able to deal with any graph size and is scalable with a large number of machines, in a parallel way. We have presented the main functions of DISCAN starting from the partitioning to the combining of intermediate results for each worker. Also, we have performed an extensive experimentation about our proposed algorithm, compared with other ones. The experiments have shown that DISCAN featured an horizontal scalability that is not guaranteed with other algorithms. In our future works, we will improve the graph partitioning step of DISCAN. Then, in the partitioning step, we plan to use the density feature during the graph partitioning, whereas for the dynamic graph clustering, we plan to make DISCAN support big evolving graphs in order to check all the snapshots of a graph. Having this hand, we can evaluate and follow the evolution of each cluster and other vertex types (border, outlier, and bridge).

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