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Improve type generic programming (slides)

Jens Gustedt

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Improve type generic programming

ISO/IEC JTC 1/SC 22/WG14 N2890

WG21 P2304

Jens Gustedt

INRIA – Camus

ICube – ICPS

Université de Strasbourg

inria
informatiques mathématiques

ICUBE



<https://gustedt.gitlabpages.inria.fr/modern-c/>

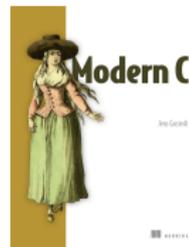


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- 4 Missing features
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Example

"Five" point tangent evaluation for approximation of a derivative:

$$(-\text{func}(x+2*h) + 8*\text{func}(x+h) - 8*\text{func}(x-h) + \text{func}(x-2*h)) / (12*h)$$

With a macro?

```
// WARNING: multiple argument evaluation
#define TANGENT5(FUNC, X, H) \
    (- FUNC((X)+2*(H)) + 8*FUNC((X)+ (H)) \
     - 8*FUNC((X)- (H)) + FUNC((X)-2*(H))) / (12*(H))
```

How to create an interface that is simple and safe to use?

With a lambda:

```
auto const λ5 = [](double x, double h, double (*func)(double)) {
    return (-func(x+2*h)+8*func(x+h)-8*func(x-h)+func(x-2*h)) / (12*h);
};
// Can be used in function call or for a function pointer
double (*fp)(double, double, double (*)(double)) = λ5;
```

Example

or so:

```
// freeze  $\varepsilon$  to  $\delta$  and have the function parameter dependent
auto const  $\lambda 5\varepsilon = [\delta = \varepsilon]$  (double x, typeof(x) func(typeof(x))) {
    double h =  $\delta * x$ ;
    return (-func(x+2*h)+8*func(x+h)-8*func(x-h)+func(x-2*h)) / (12*h);
};
```

or so:

```
// also freeze a function, and have the parameter dependent
auto const  $\lambda 5\varepsilon\_func = [\delta = \varepsilon, func = f]$  (typeof(func(0)) x) {
    auto h =  $\varepsilon * x$ ;
    return (-func(x+2*h)+8*func(x+h)-8*func(x-h)+func(x-2*h)) / (12*h);
};
```



Policy

- *extend* the standard
 - valid code remains valid
 - new feature integrates syntactically and semantically
- fix as much requirements as possible through constraints
 - specific syntax
 - explicit constraints
- avoid new undefined behavior
 - only, if property is not (or hardly) detectable at translation time
 - or we leave design space to implementations
- don't mess with ABI
 - no changes
 - no extensions

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A leveled specification: type inference

Type inference from identifiers, value expressions and type expressions

See JeanHeyd's paper [N2899](#) on `typeof`

`typeof`: keep qualifiers and `_Atomic`
`remove_qual`: remove qualifiers

[N2891](#) Type inference for variable definitions and function return
`auto` feature

- type is inferred from initializer or `return`
- conversion: lvalue, array-to-pointer, function-to-pointer
- lose qualifiers and `_Atomic`
- as-if `auto` were replaced by a `typeof` expression
- no new types in an `auto` declaration!

A leveled specification: lambdas

N2892 Basic lambdas for C

- primary use: function call expressions
- no default captures, all captures are explicit
- local *identifiers* are captured by `&id` notation
- *values* are captured by `id=expr`
- modest syntax ambiguity
- conversion: no captures \rightarrow function pointer

N2893 Options for lambdas

- `&` for default identifier captures migration path for gcc
- `=` for default *shadow* captures clang
- `id` for explicit shadow captures
- more syntax ambiguity

A leveled specification: usage patterns for lambdas

N2894 Type-generic lambdas

- **auto** parameters, types are inferred from
 - function call
 - function pointer conversion

N2862 Function Pointer Types for Pairing Code and Data

Lead by Martin Uecker

- main feature is independent of lambdas:
 - provide API for existing ABI
 - additional context pointer to function calls
- option to integrate lambdas

N2895 A simple defer feature for C

- use lambdas as syntax to describe **defer**

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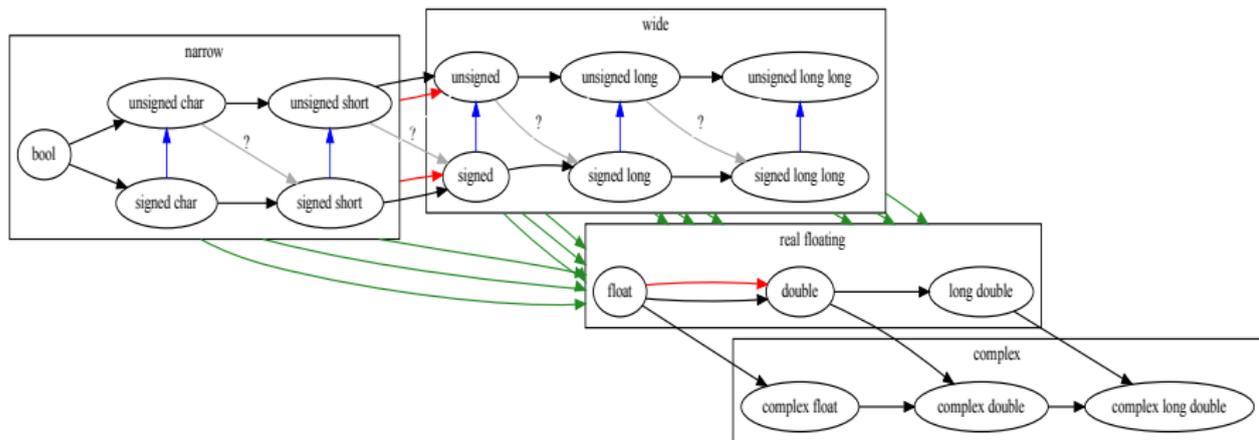
Existing type-generic features in C

- operators
- promotions and conversions
- macros
- variadic functions
- function pointers
- **void** pointers
- type-generic C library functions
- **_Generic** primary expressions

Operators

- most binary operators, have the same type for both operands
- bit-wise operators are defined for
 - wide integer types
- additionally for multiplicative operators
 - real floating point types
 - complex types
- additionally for additive operators
 - object pointers

promotions and conversions



- implicit conversion
- promotion and default argument conversion
- default arithmetic conversion

Macros

- Macros for type-generic expressions (see intro above)
 - no local variables
 - dangerous because of multiple evaluation of arguments
- Macros placeable as statements
 - weird conventions, usually

<code>/* Macro */</code>		<code>/* Type-generic lambda */</code>
<code>#define myfeature(X)</code>	<code>\</code>	<code>#define myfeature</code>
<code>do {</code>	<code>\</code>	<code>[] (auto x) {</code>
<code> typeGuess x = (X);</code>	<code>\</code>	<code> /* do something with x */</code>
<code> /* do something with x */</code>	<code>\</code>	<code>}</code>
<code>} while(false)</code>		
<code>/* Macro usage, conversion? */</code>		<code>/* Lambda usage, type safe */</code>
<code>myfeature(42);</code>		<code>myfeature(42);</code>

- Macros for declarations and definitions

Type oblivion

the user has all the burden

- Variadic functions
 - weird default conversions
 - weird library support (`va_list` a reference type?)
 - intrinsically unsafe
 - rarely used for new code
- `void*` pointers
 - unhuman effort has to be made to keep all the types correct
 - not even used by variadic functions
- function pointers
 - used with `void*` parameters for type-genericity (`bsearch`, `qsort`)



Automatic type deduction

- type-generic C library functions
 - `<tgmath.h>`
 - `<stdatomic.h>`
- **`_Generic`** primary expressions
 - difficult to extend
 - mostly restricted to function-like macros
 - not widely used

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Missing features

- temporary variables
 - temporary objects with a name
- controlled encapsulation
 - don't steal information from the surrounding scopes
 - don't pollute the surrounding scopes
- controlled constant propagation
 - control exactly what information is considered constant

Missing features

- automatic instantiation of function pointers
 - missing for `<tgmath.h>`
- automatic instantiation of specializations
 - works well with controlled constant propagation
- direct type inference
 - avoid guessing or forcing a type
 - avoid implicit conversions

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Identifiers of surrounding scopes

use of identifiers distinguishes

- visibility by scope
- access
 - no linkage (same function, relative addressing)
 - internal linkage (same TU, global addressing)
 - external linkage (same program, linktime resolution)

Identifiers of surrounding scopes

What does an identifier mean in a local function?

- local function, short lifetime, multiple instances
 - when is the definition evaluated?
- when is an outer identifier evaluated?
 - evaluation of function definition or lambda expression
 - function call
- how much evaluation?
 - remains lvalue
 - lvalue conversion
 - promotion

Lambdas: design space and terminology

The design space for captures and closures

access to automatic variables

- evaluate *expression* when seeing lambda, *value capture*
 - rvalue (no address)
 - unmutable value (**const** qualified, addressable)
 - mutable value (C++ keyword `mutable`)
- evaluate *variable* when seeing lambda, *shadow capture*
 - same possible differentiation as above
- evaluate when calling lambda
 - *identifier capture*

Terminology

- *function literal* \Leftrightarrow anonymous function, no captures
- *closure* \Leftrightarrow any capture

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Type inference

in C implementations and in other related programming languages

- **auto** type inference (`__auto_type__`)
- the **typeof** feature
- the **decltype** feature

Lambdas: existing extensions

- Objective C's blocks clang
 - shadow capture per default
 - managed memory for identifier captures
- Statement expressions gcc, clang, intel, ibm xl, open64
 - weird specification of the effective value
 - all captures are identifier captures
- Nested functions gcc
 - all captures are identifier captures
 - separation of definition and call
- C++ lambdas since C++11
 - capture model chosen by user
 - [un]mutable value or shadow captures
 - identifier and *alias* captures
 - defaults `&` and `=`

