Hamilton Circuits in the Directed Butterfly Network

Jean-Claude Bermond, Eric Darrot, Olivier Delmas, Stéphane Pérennes

INSTITUT NATIONAL DE RECHERCHE EN INFORMATIQUE ET EN AUTOMATIQUE

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Abstract: in this paper, we prove that the wrapped Butterfly digraph $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ of degree d and dimension n contains at least $d-1$ arc-disjoint Hamilton circuits, answering a conjecture of D. Barth. We also conjecture that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ can be decomposed into d Hamilton circuits, except for $d=2$ $n=2$, $d=2$ $n=3$ and $d=3$ $n=2$. We show that it suffices to prove the conjecture for d prime and $n=2$. Then, we give such a Hamilton decomposition for all primes less than 12000 by a clever computer search, and so, as corollary, we have a Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ for any d divisible by a number q, with $4 \leq q \leq 12000$.

Key-words: Butterfly graph, graph theory, Hamiltonism, Hamilton decomposition, Hamilton cycle, Hamilton circuit, perfect matching.

(Résumé : tsvp)

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Email : `{bermond, darrot, delmas, sp}@unice.fr`

Unité de recherche INRIA Sophia-Antipolis
2004 route des Lucioles, BP 93, 06902 SOPHIA-ANTIPOLIS Cedex (France)
Téléphone : (33) 93 65 77 77 – Télécopie : (33) 93 65 77 65

Circuits Hamiltoniens dans le réseau Butterfly orienté

Résumé : dans cet article, nous prouvons que le graphe Butterfly rebouclé (dans sa version orienté) $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ de degré d et de dimension n contient au moins $d-1$ circuits Hamiltoniens arc-disjoints, répondant à une conjecture de D. Barth. Nous conjecturons aussi que $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ peut être décomposé en d circuits Hamiltoniens, sauf pour $d = 2$ $n = 2$, $d = 2$ $n = 3$ et $d = 3$ $n = 2$. Nous montrons qu'il suffit de prouver cette conjecture pour d premier et $n = 2$. Puis nous donnons une telle décomposition pour tous les premiers jusqu'à 12000 grâce à une subtilité de programmation, induisant comme corollaire que $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ admet une décomposition Hamiltonienne pour tout d divisible par un nombre q tel que $4 \leq q \leq 12000$.

Mots-clé : graphe Butterfly, théorie des graphes, Hamiltonisme, décomposition Hamiltonienne, cycle Hamiltonien, circuit Hamiltonien, couplage parfait.

1 Introduction and notations

1.1 Butterfly networks

Many interconnection networks have been proposed as suitable topologies for parallel computers. Among them the *Butterfly networks* have received particular attention, due to their interesting structure.

First, we have to warn the reader that under the name *Butterfly* and with the same notation, different networks are described. Indeed, if some authors consider the *Butterfly networks* as multistage networks used to route permutations, others consider them as point-to-point networks. In what follows, we will call the multistage version *Butterfly* and we will use Leigthon's terminology [13], namely *wrapped Butterfly*, for the point-to-point version. Furthermore, these networks can be considered either as undirected or directed. To be complete, we recall that some authors consider only *binary Butterfly networks* that is the restricted class of networks obtained when the out-degree is 2 (directed case) or 4 (undirected case).

In this article, we will use the following definitions and notation. $\mathbb{Z}_q$ will denote the set of integers modulo q. (For definitions not given here see [15]).

Definition 1.1 *The* **Butterfly digraph** *of degree d and dimension n, denoted $\vec{\mathcal{BF}}(d,n)$ has as vertices the couples (x,l), where x is an element of $\mathbb{Z}_d^n$ that is a word $x_{n-1}x_{n-2}\cdots x_1x_0$ where the letters belong to $\mathbb{Z}_d$ and $0 \le l \le n$ (l is called the level). For $l < n$, a vertex $(x_{n-1}x_{n-2}\cdots x_1x_0, l)$ is joined by an arc to the d vertices $(x_{n-1}\cdots x_{l+1}, \alpha, x_{l-1}\cdots x_0, l+1)$ where α is any element of $\mathbb{Z}_d$.*

$\vec{\mathcal{BF}}(d,n)$ has $(n+1)d^n$ vertices. Each vertex in level $l < n$ has out-degree d. This digraph is not strongly connected. It is mainly used as a multistage interconnection network (the levels corresponding to the stages) in order to route some one-to-one mapping of d^n inputs (nodes at level 0) to d^n outputs (nodes at level n).

The underlying undirected graph obtained by forgetting the orientation will be denoted $\mathcal{BF}(d,n)$.

Figure (**1**) shows simultaneously $\mathcal{BF}(3,2)$ and $\vec{\mathcal{BF}}(3,2)$. The orientation on $\vec{\mathcal{BF}}(d,n)$ being obtained by directing the edges from left to right.

Note that $\vec{\mathcal{BF}}(d,n)$ is often represented (for example in [13, 15]) in an opposite way to our drawing as the authors denote the nodes $(x_0x_1\cdots x_{n-1})$. We have chosen the representation which emphasises the most on the recursive decomposition of $\vec{\mathcal{BF}}(d,n)$ and provides us an easy representation of our inductive construction (see section (**3**)).

Definition 1.2 *The* **wrapped Butterfly digraph** *$\mathcal{W}\vec{\mathcal{BF}}(d,n)$ is obtained from $\vec{\mathcal{BF}}(d,n)$ by identifying the vertices of the last and first level namely (x,n) with $(x,0)$. In other words the vertices are the couples (x,l) where x is an element of $\mathbb{Z}_d^n$ that is a word $x_{n-1}x_{n-2}\cdots x_1x_0$ where the letters belong to $\mathbb{Z}_d$ and $l \in \mathbb{Z}_n$ (l is called the level). For any l, a vertex $(x_{n-1}x_{n-2}\cdots x_1x_0, l)$ is joined by an arc to the d vertices $(x_{n-1}\cdots x_{l+1}, \alpha, x_{l-1}\cdots x_0,\ l+1)$ where α is any element of $\mathbb{Z}_d$ (and where $l+1$ has to be taken modulo n).*

Usually to represent the wrapped Butterfly (di)graph we use the representation of $\vec{\mathcal{BF}}(d,n)$ by repeating at the end level 0. Hence the reader has to remind that levels 0 and n are identified for $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$. $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ is a is d-regular digraph with nd^n vertices; its diameter is $2n-1$. The underlying wrapped Butterfly network will be denoted $\mathcal{WBF}(d,n)$; it is regular of degree $2d$, with diameter $\lfloor \frac{3n}{2} \rfloor$.

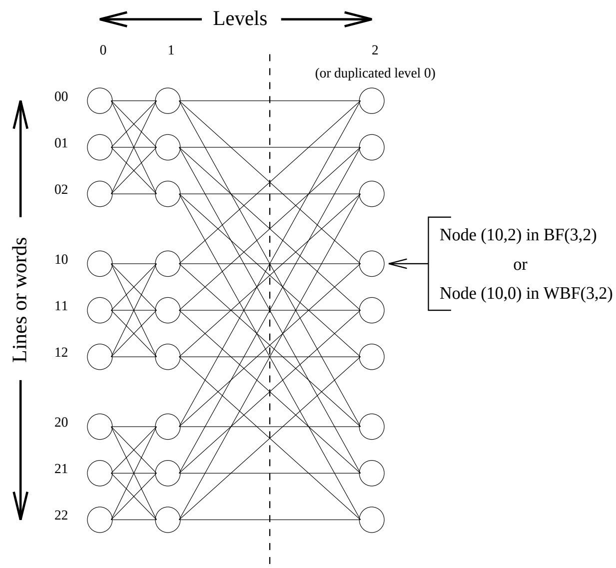

Figure 1: *The graphs* $\mathcal{BF}(3,2)$ *(multistage version) with 3 levels, or* $\mathcal{WBF}(3,2)$ *(point-to-point version) by duplicating level* 0. *For digraphs* $\vec{\mathcal{BF}}(3,2)$ *or* $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(3,2)$ *the edges must be directed into arcs from left to right .*

1.2 Other definitions and general results

- $\mathcal{K}_d$ will denote the complete graph on d vertices,
- $\mathcal{K}_{d,d}$ will denote the complete bipartite graph where each set of the bipartition has size d,
- G^* will denote the symmetric digraph obtained from the graph G by replacing each edge by two opposite arcs. In particular $\mathcal{K}_d^*$ *(resp.* $\mathcal{K}_{d,d}^*$*)* will denote the complete symmetric *(resp. bipartite)* digraph on d *(resp.* $d \times d$*)* vertices,
- $\mathcal{K}_d^+$ will denote the complete symmetric digraph with a loop on each vertex,
- a circuit or directed cycle of length n will be denoted $\vec{C}_n$, and a dipath of length n $\vec{P}_n$.

- $\vec{\mathcal{K}}_{d,d}$ will denote the directed digraph obtained from $\mathcal{K}_{d,d}$ by orienting each edge from the left part to the right part.

Definition 1.3 *(see [15]) Let G be a directed graph. The* **line digraph** $L(G)$ *of G is the directed graph whose vertices are the arcs of G and whose arcs are defined as follows: there is an arc from a vertex e to a vertex f in $L(G)$ if and only if, in G, the initial vertex of f is the end vertex of e.*

Note that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,1)$ is nothing else than $\mathcal{K}_d^+$ and that $\vec{\mathcal{B}\mathcal{F}}(d,1)$ is $\vec{\mathcal{K}}_{d,d}$. We will see in section (**4**) (corollary (**4.1**)) that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2)$ is the line digraph of $\mathcal{K}^*_{d,d}$.

Definition 1.4 *A* **1-difactor** *of a digraph G is a spanning subgraph of G with in and out-degree 1. It corresponds to a partition of the vertices of G into circuits.*

Definition 1.5 *A* **Hamilton cycle (resp. circuit)** *of a graph* (resp. digraph) *is a cycle* (resp. circuit) *which contains every vertex exactly once.*

Definition 1.6 *We will say that a graph* (resp. digraph) *has a* **Hamilton decomposition** *or* **can be decomposed into Hamilton cycles (resp. circuit)** *if its edges* (resp. arcs) *can be partitioned into Hamilton cycles* (resp. circuits).

Remark 1 A Hamilton circuit is a connected 1-difactor.

The existence of one and if possible many edge(arc)-disjoint Hamilton cycles (circuits) in a network can provide advantage for algorithms that make use of a ring structure. Furthermore, the existence of a Hamilton decomposition allows also the message traffic to be evenly distributed across the network. Various results have been obtained on the existence of Hamilton cycles in classical networks (see for example the survey [2, 11]). For example it is well-known that any Cayley graph on an abelian group is hamiltonian. Furthermore it has been conjectured by Alspach [1] that:

Conjecture 1 (Alspach) *Every connected Cayley graph on an abelian group has a Hamilton decomposition.*

This conjecture has been verified for all connected 4-regular graphs on abelian groups in [9]. That includes in particular the toroidal meshes (grids). For the hypercube it is also known that $\mathcal{H}(2d)$ can be decomposed into d Hamilton cycles (see [3, 2]).

Concerning line digraphs it has been shown in [12] that d-regular line digraphs always admit $\lfloor \frac{d}{2} \rfloor$ Hamilton circuits. In the case of de Bruijn or Kautz digraphs which are the simplest line digraphs, partial results have been obtained successively in [6, 14], and quasi optimal results have been obtained for undirected de Bruijn and Kautz graphs [4].

1.3 Results for the Butterfly networks

The wrapped Butterfly (di)graph is actually a Cayley graph (on a non abelian group) and a line digraph. So the decomposition into Hamilton cycles *(resp. circuits)* of this graph *(resp. digraph)* has received some attention. First it is well-known that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ is hamiltonian (see [13, page 465] for a proof in the case $d = 2$). In [7], Barth and Raspaud proved that $\mathcal{W}\mathcal{B}\mathcal{F}(2,n)$ has a Hamilton decomposition answering a conjecture of J. Rowley and D. Sotteau (private communication).

Theorem 1.1 (Barth, Raspaud) *$\mathcal{WBF}(2,n)$ can be decomposed into 2 Hamilton cycles.*

They also gave the following conjecture.

Conjecture 2 (Barth, Raspaud) *For $n \geq 2$, $\mathcal{WBF}(d,n)$ can be decomposed into d Hamilton cycles.*

In his thesis [5] , Barth also stated the following conjecture for the directed case:

Conjecture 3 (Barth) *For $n \geq 2$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ contains $d-1$ arc-disjoint Hamilton circuits.*

Recall that for $n = 1$ $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,1)$ is nothing else than $\mathcal{K}_d^+$ which itself is the arc-disjoint sum of $\mathcal{K}_d^*$ and loops. So conjecture (**3**) can be seen as an extension of a theorem of Tillson [17].

Theorem 1.2 (Tillson) *The complete symmetric digraph $\mathcal{K}_d^*$ can be decomposed into $d-1$ Hamilton circuits except for $d = 4$ and 6.*

In this paper we focus mainly on the decomposition of the wrapped Butterfly digraph $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$. Our main result implies that the number of arc-disjoint Hamilton circuits contained in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ can only increase when n increases.

Proposition 1.1 *If $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ contains p arc-disjoint Hamilton circuits, then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n')$ contains also at least p arc-disjoint Hamilton circuits, for any $n' \geq n$.*

This proposition with Tillson's theorem and a special study for $d = 4$ and 6, implies conjecture (**3**).

Theorem 1.3 *For $n \geq 2$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ contains $d-1$ arc-disjoint Hamilton circuits.*

Furthermore it appears that, except for three cases, for all small values of d, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ can be decomposed into d Hamilton circuits, so we conjecture that:

Conjecture 4 *For $n \geq 2$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ can be decomposed into d Hamilton circuits, except for ($d = 2$, $n = 2$ or 3) and ($d = 3$, $n = 2$).*

By the proposition (**1.1**) it suffices to prove the conjecture for $n = 2$. Using results of section (**4**) on conjunction of graphs, we have been able to reduce the study to prime degrees. So conjecture (**4**) would follow from conjecture (**5**).

Conjecture 5 *For any prime number $p > 3$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p,2)$ can be decomposed into p Hamilton circuits.*

With a clever computer search, we have been able to prove conjecture (**5**) for any prime less than 12000, leading to the following statement:

Theorem 1.4 *If d is divisible by any number q, $4 \leq q \leq 12000$ then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2)$ and consequently $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ has a Hamilton decomposition.*

Finally the methods used in this paper can be applied with other ideas to the undirected case and lead us to prove conjecture (**2**) in a forthcoming paper [8].

Theorem 1.5 *For $n \geq 2$, $\mathcal{WBF}(d,n)$ can be decomposed into d Hamilton cycles.*

2 Circuits and Permutations

2.1 More definitions

First, we will show that the existence of k arc-disjoint Hamilton circuits in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$, is equivalent to the ability to route k *compatible cyclic realisable* permutations between levels 0 and n in $\vec{\mathcal{BF}}(d,n)$. For this purpose we need some specific definitions.

In the whole paper, π will always denote a permutation of $\mathbb{Z}_d^n$ which associates to x, $\pi(x)$. It follows from the definition of $\vec{\mathcal{BF}}(d,n)$ that there exists a unique dipath connecting a vertex $(x,0)$ to a vertex (y,n). So we can associate to a permutation π of $\mathbb{Z}_d^n$ a set of dipaths in $\vec{\mathcal{BF}}(d,n)$ connecting vertex $(x,0)$ to vertex $(\pi(x),n)$ for any x in $\mathbb{Z}_d^n$.

The **composition** $\pi \cdot \pi'$ **of two permutations** π and π' is the permutation which associate to the element a the element $\pi(\pi'(a))$.

Definition 2.1 *A permutation π is* **realisable** *in $\vec{\mathcal{BF}}(d,n)$ or equivalently $\vec{\mathcal{BF}}(d,n)$ realizes the permutation π if the d^n associated dipaths from the inputs to the outputs are vertex-disjoint.*

Definition 2.2 *A permutation π is* **cyclic** *if all the elements $\pi^i(x)$ are distinct for $0 \leq i < d^n$.*

Definition 2.3 *A set of k permutations $\pi_0, \pi_1, \cdots, \pi_{k-1}$ realisable in $\vec{\mathcal{BF}}(d,n)$ is* **compatible** *if the kd^n dipaths from $(x,0)$ to $(\pi_j(x),n)$ for x in $\mathbb{Z}_d^n$ and $0 \leq j \leq k-1$ are arc-disjoint. We will also say that $\vec{\mathcal{BF}}(d,n)$ realises k compatible permutations.*

2.2 Hamilton circuits and permutations

We are now ready to prove that there is an immediate connection between the existence of compatible cyclic realisable permutations in $\vec{\mathcal{BF}}(d,n)$ and that of Hamilton circuits in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$.

Lemma 2.1 *$\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ contains k arc-disjoint Hamilton circuits if and only if $\vec{\mathcal{BF}}(d,n)$ realizes k compatible cyclic permutations.*

Proof. Let π be a cyclic permutation realisable in $\vec{\mathcal{BF}}(d,n)$ and let P_i be the unique dipath joining $(\pi^i(x),0)$ to $(\pi^{i+1}(x),n)$. Let P_i' be the dipath of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ obtained from P_i by identifying $(\pi^{i+1}(x),n)$ with $(\pi^{i+1}(x),0)$. The concatenation of the dipaths P_i' forms a Hamilton circuit of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ because, π being cyclic, the $\pi^i(x)$ span the set of vertices of level 0. Conversely, if H is a Hamilton circuit of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$, let $(x_0,0),\cdots,(x_i,0),\cdots,(x_{d^n-1},0)$ be the vertices we meet successively on level 0 by following the cycle H. Let us consider the permutation defined by $\pi(x_i) = x_{i+1}$. As H is a Hamilton circuit, all the x_i's are distinct and all the inside dipaths are vertex-disjoint, so π is a cyclic realisable permutation in $\vec{\mathcal{BF}}(d,n)$. Finally by the equivalence above, arc-disjoint Hamilton circuits correspond to compatible cyclic realisable permutations (see Figure (**2**) for an example). $\square$

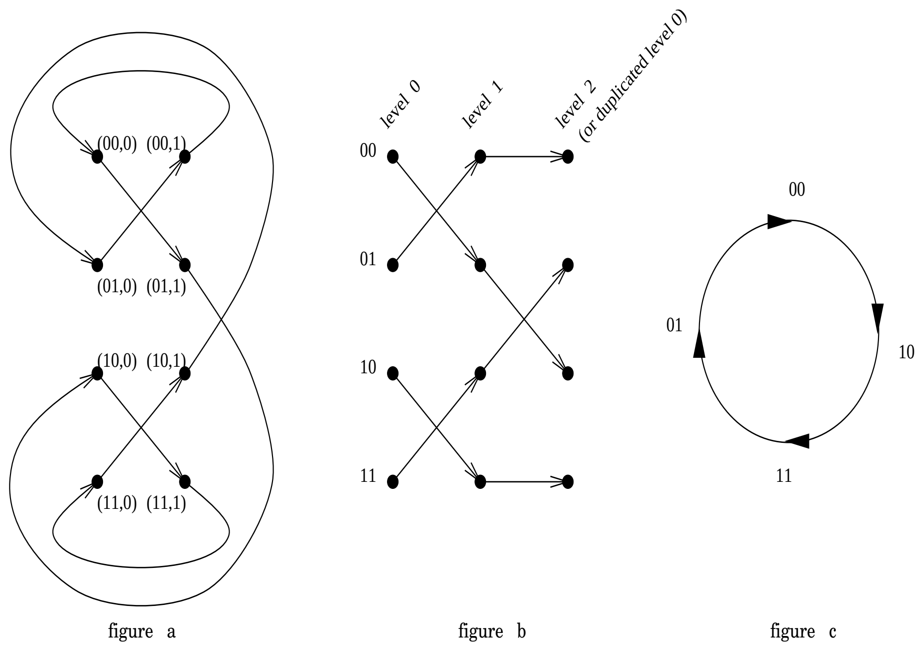

Figure 2: *A Hamilton circuit of* $\mathcal{W}\vec{\mathcal{B}\mathcal{F}}(2,2)$ *(figure a) or equivalently the associated permutation realisable in* $\vec{\mathcal{B}\mathcal{F}}(2,2)$ *(figure b) and the cyclic permutation which is used (figure c).*

3 Recursive construction

3.1 Recursive decomposition of $\vec{\mathcal{B}\mathcal{F}}(d,n)$

The permutation network $\vec{\mathcal{B}\mathcal{F}}(d,n)$ has a simple recursive property: the $n+1$ first levels of $\vec{\mathcal{B}\mathcal{F}}(d,n+1)$ form d vertex-disjoint subgraphs isomorphic to $\vec{\mathcal{B}\mathcal{F}}(d,n)$. We shall call them left Butterflies. If the vectors of $\mathbb{Z}_d^{n+1}$ are denoted $y=(ax)\in\mathbb{Z}_d\times\mathbb{Z}_d^n$, then each left Butterfly connects the set of vertices having the same left part a. So we will label a left Butterfly by $\mathcal{B}_{left}(a)$. In the same way, the two last levels of $\vec{\mathcal{B}\mathcal{F}}(d,n+1)$ are built with d^n disjoint subgraphs isomorphic to $\vec{\mathcal{B}\mathcal{F}}(d,1)=\vec{\mathcal{K}}_{d,d}$ that we shall call right Butterflies; each right Butterfly connects all the vertices having the same right part x and we will label it by $\vec{\mathcal{K}}_{d,d}(x)$.

We can summarize the situation as follows:

- vertices of $\vec{\mathcal{B}\mathcal{F}}(d,n+1)$ are denoted (ax,l),
- the left Butterfly labeled by $a\in\mathbb{Z}_d$ is formed by the vertices $a*$ of the $n+1$ first levels. It will be denoted $\mathcal{B}_{left}(a)$,
- the right Butterfly with label $x\in\mathbb{Z}_d^n$ is formed by the vertices $*x$ of the 2 last levels. It will be denoted $\vec{\mathcal{K}}_{d,d}(x)$.

Remark 2 In $\vec{\mathcal{BF}}(d, n+1)$, vertices of level $n+1$ are shared by the left and right Butterflies, the outputs of the left Butterflies being considered as the inputs of the right Butterflies. Moreover all the subgraphs defined above are arc-disjoint.

Figure (**3**) displays such a recursive decomposition.

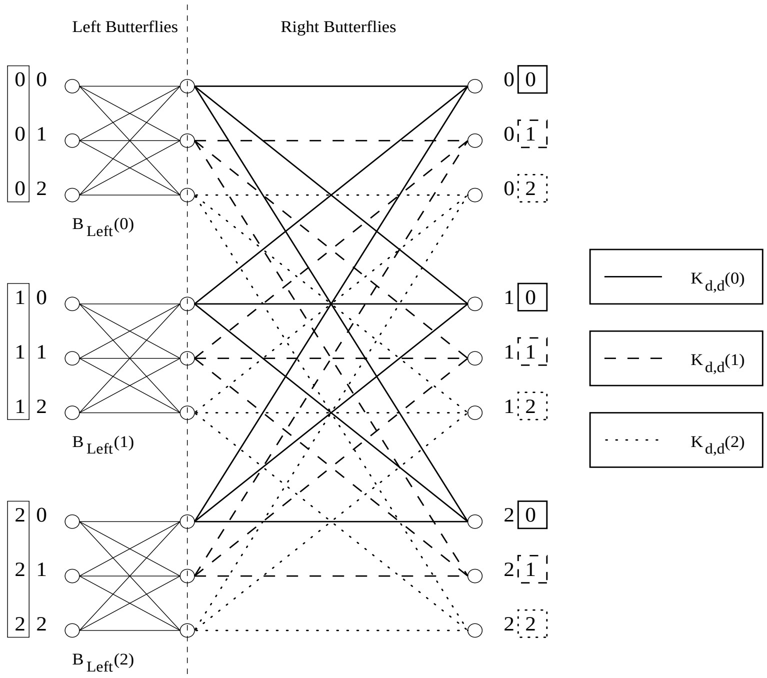

Figure 3: *The recursive decomposition of $\mathcal{BF}(3,2)$. To obtain the directed version $\vec{\mathcal{BF}}(3,2)$ the edges must be directed into arcs from left to right. The vertices are denoted $y = (ax) \in \mathbb{Z}_3 \times \mathbb{Z}_3^1$. In $\vec{\mathcal{BF}}(3,2)$ the 2 first levels form 3 vertex-disjoint subgraphs, each one isomorphic to $\vec{\mathcal{BF}}(3, 2-1)$. These 3 subgraphs are labeled $\mathcal{B}_{left}(a)$. In the same way, the 2 last levels of $\vec{\mathcal{BF}}(3,2)$ are built with 3^1 disjoint subgraphs isomorphic to $\vec{\mathcal{BF}}(3,1) = \vec{\mathcal{K}}_{d,d}$ labeled $\vec{\mathcal{K}}_{d,d}(x)$.*

3.2 Iterative Construction

We will now give a simple construction which enables us to construct p compatible cyclic realisable permutations in $\vec{\mathcal{BF}}(d, n+1)$ from p compatible cyclic realisable permutations in $\vec{\mathcal{BF}}(d, n)$.
In what follows M_x will denote a permutation realisable in the right Butterfly $\vec{\mathcal{K}}_{d,d}(x)$. To M_x is associated a perfect matching in $\vec{\mathcal{K}}_{d,d}(x)$; the arcs in the perfect matching join ax to $M_x(a)x$.

Definition 3.1 *Let $\mathcal{M}_j = (M_{x,j},\, x \in \mathbb{Z}_d^n)$, $0 \le j \le p-1$ be p families of permutations realisable in $\vec{\mathcal{K}}_{d,d}(x)$, $x \in \mathbb{Z}_d^n$. The* **families** *$\mathcal{M}_j$ are* **compatible***, if for every x the set of permutations $M_{x,j}$, $0 \le j \le p-1$, is compatible (in other words they correspond to arc-disjoint matchings of $\vec{\mathcal{K}}_{d,d}(x)$).*

Definition 3.2 *A family of permutations satisfies the* **cyclic property** *if the composition of the permutations of the family is a cyclic permutation for any order of the composition.*

Lemma 3.1 *There exist d compatible families $\mathcal{M}_j = (M_{x,j},\ x \in \mathbb{Z}_d^n)$, $0 \le j \le d-1$ of permutations realisable in $\vec{\mathcal{K}}_{d,d}(x)$, $x \in \mathbb{Z}_d^n$ such that for each j, the family $\mathcal{M}_j$ has the cyclic property.*

Proof. Let $M_{x,j}$ be the permutation of $\mathbb{Z}_d$ defined by $M_{x,j}(a) = a+j$ if $x \neq 0$ and $M_{0,j}(a) = a+j+1$. Clearly, $M_{x,j}$ is realisable in $\vec{\mathcal{K}}_{d,d}(x)$. Furthermore for $j \neq j'$, the matchings associated to $M_{x,j}$ and $M_{x,j'}$ are arc-disjoint; so for a given x the permutations $M_{x,j}$ are compatible. Finally for a given j the composition of the d^n permutations $M_{x,j}$ taken in any order is the permutation which associates to a the element $a + (d^n-1)j + j + 1 = a + d^n j + 1 = a+1$ which is clearly cyclic. □

Lemma 3.2 (Inductive lemma) *Let π be a cyclic permutation realisable in $\vec{\mathcal{BF}}(d,n)$ and let $\mathcal{M} = (M_x,\, x \in \mathbb{Z}_d^n)$ be a family of realisable permutations satisfying the cyclic property. Then the permutation $f_{(\pi,\mathcal{M})}$ of $\mathbb{Z}_d^{n+1}$ defined by $f_{(\pi,\mathcal{M})}(ax) = b\pi(x)$ where $b = M_{\pi(x)}(a)$, is a cyclic permutation realisable in $\vec{\mathcal{BF}}(d,n+1)$.*

Proof. The dipath associated to $f_{(\pi,\mathcal{M})}$ from $(ax,0)$ to $(b\pi(x),n+1)$ consists of the dipath from $(ax,0)$ to $(a\pi(x),n)$ in $\mathcal{B}_{left}(a)$ associated to the permutation π of $\mathcal{B}_{left}(a)$ (which is isomorphic to $\vec{\mathcal{BF}}(d,n)$) and then of the arc joining $(a\pi(x),n)$ to $(b\pi(x),n+1)$ in $\vec{\mathcal{K}}_{d,d}(\pi(x))$ defined by the matching $M_{\pi(x)}$, that is $b = M_{\pi(x)}(a)$. We claim that the dipaths joining two distinct inputs $(ax,0)$ and $(a'x',0)$ to their outputs are vertex-disjoint and so $f_{(\pi,\mathcal{M})}$ is realisable. Indeed, if $a \neq a'$ their first parts are in two different $\mathcal{B}_{left}(a)$ and $\mathcal{B}_{left}(a')$ and the last arcs are disjoint as either $x \neq x'$ or $x = x'$ and $M_{\pi(x)}$ is realisable in $\vec{\mathcal{K}}_{d,d}(\pi(x))$. If $a = a'$, π being realisable the first dipaths are vertex-disjoint and as $x \neq x'$ the last arcs belong to two different $\vec{\mathcal{K}}_{d,d}$. Finally it remains to show that $f_{(\pi,\mathcal{M})}$ is cyclic. To prove that it suffices to show that $f^i_{(\pi,\mathcal{M})}(ax) \neq ax$ for $1 \le i \le d^{n+1}-1$. Suppose $f^i_{(\pi,\mathcal{M})}(ax) = ax$. By definition $f^i_{(\pi,\mathcal{M})}(ax) = a'\pi^i(x)$. So $\pi^i(x) = x$, which implies that $i = kd^n$. But for $i = d^n$, a' is the image of a by the composition of the d^n elements of $\mathcal{M}$ in some order and as $\mathcal{M}$ has the cyclic property for $i = d^n$, $a' = \sigma(a)$ where σ is a cyclic permutation. So $f^{kd^n}_{(\pi,\mathcal{M})}(ax) = \sigma^k(a)x \neq ax$ for $1 \le k < d$. So for any i $1 \le i \le d^{n+1}-1$, $f^i_{(\pi,\mathcal{M})}(ax) \neq ax$. □

Corollary 3.1 *If there exist p compatible cyclic realisable permutations in $\vec{\mathcal{BF}}(d,n)$, then there exist p compatible cyclic realisable permutations in $\vec{\mathcal{BF}}(d,n+1)$.*

Proof. Let π_j, $0 \le j \le p-1$, be p compatible cyclic realisable permutations of $\vec{\mathcal{BF}}(d,n)$ and let $\mathcal{M}_j = (M_{x,j},\ x \in \mathbb{Z}_d^n)$, $0 \le j \le p-1$, be p compatible families of realisable permutations each satisfying the cyclic property. Such p families exist by lemma (**3.1**) as $p \le d$. By the above proposition for each j the $f_{(\pi_j, \mathcal{M}_j)}$ are cyclic realisable permutations of $\vec{\mathcal{BF}}(d, n+1)$ and they are compatible as both the π_j and the $\mathcal{M}_j$ are compatible (all the dipaths associated are arc-disjoint). □

Now, we are ready to prove our main proposition stated in the introduction.

Proposition 3.1 (Main) *If $\mathcal{W}\vec{\mathcal{BF}}(d,n)$ contains p arc-disjoint Hamilton circuits, then $\mathcal{W}\vec{\mathcal{BF}}(d,n')$ contains also at least p arc-disjoint Hamilton circuits, for any $n' \ge n$.*

Proof. The result for $n' = n+1$ follows from corollary (**3.1**) and lemma (**2.1**). A recursive application of this property gives the above proposition. □

For an example of the construction see Figure (**4**).

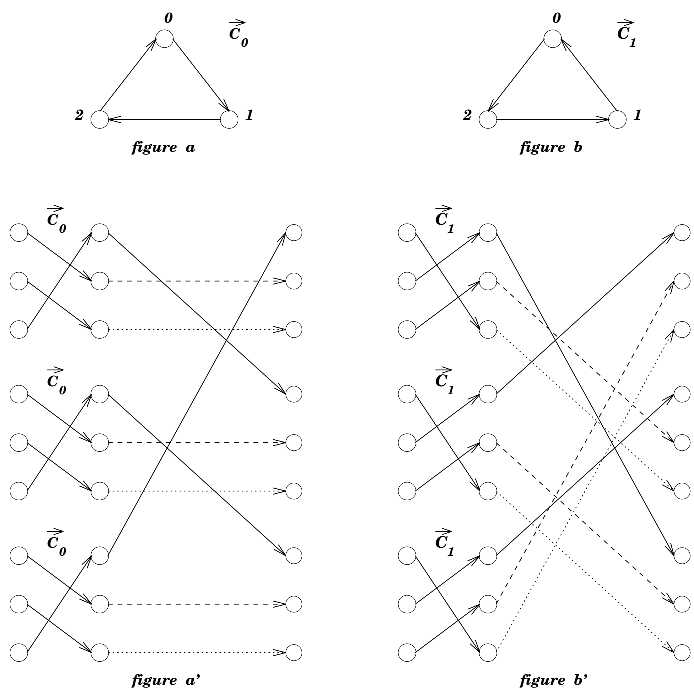

Figure 4: *Two Hamilton circuits (figures a' and b') of $\vec{\mathcal{BF}}(3,2)$ are obtained from two Hamilton circuits of $\mathcal{W}\vec{\mathcal{BF}}(3,1) = \mathcal{K}_3^+$ by the construction of lemma (**3.2**). Figures (a) and (b) show two arc-disjoint Hamilton circuits of $\mathcal{K}_3^+$, $(\vec{C}_0 \mid x \to x+1 \pmod 3)$ and $(\vec{C}_1 \mid x \to x+2 \pmod 3)$. We used the families $\mathcal{M}_0$ (figure a') and $\mathcal{M}_1$ (figure b') defined in the proof of lemma (**3.1**).*

3.3 Consequences

Corollary 3.2 *$\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,n)$ can be decomposed into two Hamilton circuits as soon as $n \geq 4$. For $1 \leq n \leq 3$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,n)$ admits only one Hamilton circuit.*

Proof. A computer search has given a decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,4)$ into two arc-disjoint Hamilton circuits. Therefore by proposition (**3.1**) $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,n)$ has a Hamilton decomposition, for any $n \geq 4$. For $1 \leq n \leq 3$, an exhaustive computer search shows that there cannot exist two arc-disjoint Hamilton circuits. □

Corollary 3.3 *$\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(3,n)$ can be decomposed into three Hamilton circuits as soon as $n \geq 3$. For $1 \leq n \leq 2$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(3,n)$ admits only two arc-disjoint Hamilton circuits.*

Proof. That follows from the existence of a Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(3,3)$ obtained by computer (figures of decompositions available on demand). □

Now we are able to prove Barth's conjecture (**3**).

Theorem 3.3 *For $n \geq 2$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ contains $d-1$ arc-disjoint Hamilton circuits.*

Proof. By Tillson's decomposition ((**1.2**)), for $d \neq 4$, and $d \neq 6$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,1) = \mathcal{K}_d^+$ contains $d-1$ arc-disjoint Hamilton circuits. So by proposition (**3.1**), for $d \neq 4$ and $d \neq 6$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ contains at least $d-1$ arc-disjoint Hamilton circuits. For $d = 4$ *(resp. 6)*, we have found by computer search 4 *(resp. 6)* arc-disjoint Hamilton circuits in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(4,2)$ *(resp.* $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(6,2)$*)*. So by proposition (**3.1**), $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(4,n)$ *(resp.* $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(6,n)$*)* contains 4 *(resp. 6)* arc-disjoint Hamilton circuits. □

As seen in the proof above for $d = 4$, 6 there exists for $n \geq 2$ a Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$. These results and those of the next section lead us to propose the following conjecture, which would completely close the study of the Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$.

Conjecture 6 *For $d \geq 4$ and $n \geq 2$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ can be decomposed into Hamilton circuits.*

By proposition (**3.1**) it suffices to prove the conjecture for $n = 2$ or equivalently, as $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2) = L(\mathcal{K}^*_{d,d})$ (see corollary (**4.1**)) that $\mathcal{K}^*_{d,d}$ admits d compatible eulerian tours (see [12]).

4 Decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2)$ into Hamilton circuits

4.1 Line digraphs and conjunction

We need some more definitions and results concerning *conjunction*, *line digraphs* and *de Bruijn digraphs*.

Definition 4.1 *(see [2])*

1. *The* **conjunction** $G_1 \cdot G_2$ *of two digraphs* $G_1 = (V_1, E_1)$ *and* $G_2 = (V_2, E_2)$ *is the digraph with vertex-set* $V_1 \times V_2$ *and an arc joining* (u_1, u_2) *to* (v_1, v_2) *if and only if there is an arc joining* u_1 *to* v_1 *in* G_1 *and an arc joining* u_2 *to* v_2 *in* G_2.

2. *The* **edge-disjoint union (sum)** *of two digraphs* A, B *on the same set of vertices is denoted* $A \oplus B$.

3. cG *will denote the digraph made of* c *disjoint copies of* G.

For example $\mathcal{K}^*_{d,d} = \mathcal{K}^+_d \cdot \vec{C}_2$ and $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,4) = A \oplus B$ where A and B are the two arc-disjoint Hamilton circuits of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,4)$.

Properties 4.1

$$F \cdot G = G \cdot F$$

$$L(F \cdot G) = L(F) \cdot L(G)$$

$$(F \cdot G) \cdot H = F \cdot (G \cdot H) = F \cdot G \cdot H$$

$$(A \oplus B) \oplus C = A \oplus (B \oplus C) = A \oplus B \oplus C$$

$$(A \oplus B) \cdot F = (A \cdot F) \oplus (B \cdot F)$$

Proof. These results are clear from the definitions. □

There is a very strong connection between the *de Bruijn* digraph and the wrapped Butterfly digraph. We recall the definition of the de Bruijn digraph:

Definition 4.2 *The* **de Bruijn digraph** $\vec{\mathcal{B}}(d,n)$ *of out-degree* d *and diameter* n *has as vertices the words of length* n *on an alphabet of* d *letters. Vertex* $x_0 \dots x_{n-1}$ *is joined by an arc to the vertices* $x_1 \dots x_{n-1}\alpha$ *where* α *is any letter from the alphabet.*

Proposition 4.2

$$\vec{\mathcal{B}}(d,n) = L^{n-1}(\mathcal{K}^+_d) \tag{1}$$

$$\vec{\mathcal{B}}(d_1 d_2, n) = \vec{\mathcal{B}}(d_1, n) \cdot \vec{\mathcal{B}}(d_2, n) \tag{2}$$

$$\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n) = \vec{\mathcal{B}}(d,n) \cdot \vec{C}_n = L^{n-1}(\mathcal{K}^+_d \cdot \vec{C}_n) \tag{3}$$

$$\vec{\mathcal{B}}\mathcal{F}(d,n) = \vec{\mathcal{B}}(d,n) \cdot \vec{P}_n \tag{4}$$

$$\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1 d_2, n) = \mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1, n) \cdot \vec{\mathcal{B}}(d_2, n) \tag{5}$$

Proof. Equality (**1**) is well known, and even sometimes considered as the proper definition of de Bruijn digraphs (see [10, 15]).

Result (**2**) can be found in [16] and can be proved as follows: $\vec{\mathcal{B}}(d_1d_2,n) = L^{n-1}(\mathcal{K}^+_{d_1d_2})$ from (**1**), but $\mathcal{K}^+_{d_1d_2} = \mathcal{K}^+_{d_1} \cdot \mathcal{K}^+_{d_2}$. By properties (**4.1**), $L^{n-1}(\mathcal{K}^+_{d_1} \cdot \mathcal{K}^+_{d_2}) = L^{n-1}(\mathcal{K}^+_{d_1}) \cdot L^{n-1}(\mathcal{K}^+_{d_2})$ which is indeed $\vec{\mathcal{B}}(d_1,n) \cdot \vec{\mathcal{B}}(d_2,n)$.

Result (**3**) is implicit in different papers. It can be obtained by considering the following isomorphism from $\vec{\mathcal{B}}(d,n) \cdot \vec{C}_n$ to $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$. To the vertex (x,l) in $\vec{\mathcal{B}}(d,n) \cdot \vec{C}_n$ with $x = x_0x_1\cdots x_{n-1}$, and $l \in \mathbb{Z}_n$, we associate the vertex $\phi((x,l)) = (x',l)$ in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$, where $x' = x'_{n-1}x'_{n-2}\cdots x'_0$ and $x'_i = x_{i-l}$. By definitions (**4.1-1**) and (**4.2**), the out-neighbors of (x,l) in $\vec{\mathcal{B}}(d,n) \cdot \vec{C}_n$ are the vertices $(y,l+1)$ with $y = y_0y_1\cdots y_{n-1}$ such that $y_i = x_{i+1}$ for $i \neq n-1$, and $y_{n-1} = \alpha$, α being any letter from the alphabet. The associated vertices in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ are $\phi((y,l+1)) = (y',l+1)$ where $y' = y'_{n-1}y'_{n-2}\cdots y'_0$ and $y'_i = y_{i-l-1}$. For $i-l-1 \neq n-1$, or equivalently $i \neq l$, $y'_i = x_{i-l} = x'_i$, and for $i = l$, $y'_i = \alpha$. So by definition (**1.2**), the vertices $(y',l+1)$ are exactly the out-neighbors of (x',l) in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$. The second part of the equality is due to the fact that $L^{n-1}(\vec{C}_n) = \vec{C}_n$, hence $L^{n-1}(\mathcal{K}^+_d) \cdot \vec{C}_n = L^{n-1}(\mathcal{K}^+_d) \cdot L^{n-1}(\vec{C}_n) = L^{n-1}(\mathcal{K}^+_d \cdot \vec{C}_n)$.

An example is displayed in Figure (**5**). Result (**4**) is equivalent to (**3**).

The last equality follows directly from (**2**) and (**3**). □

Corollary 4.1 $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2) = L(\mathcal{K}^*_{d,d})$.

Proof. Follows from proposition (**4.2**) equality (**3**) for $n = 2$. □

Lemma 4.1 *When r and s are relatively prime, $\vec{C}_{qs} \cdot \vec{C}_{qr} = q\vec{C}_{qsr}$.*

Proof. $\vec{C}_{qs} \cdot \vec{C}_{qr}$ is a regular digraph with in and out-degree 1. So it is the union of circuits. Starting from a vertex (u,v) we find at distance i the vertex $(u+i, v+i)$ where $u+i$ *(resp.* $v+i$*)* has to be taken modulo qs *(resp.* qr*)*. So the length of any circuit is the smallest common multiple of qs and qr, that is qrs as r and s are relatively prime. As the number of vertices in the digraph is q^2rs there are q such cycles. □

Proposition 4.3 $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1,n) \cdot \mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_2,n) = n\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n)$

Proof. Let $G = \mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1,n) \cdot \mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_2,n)$. By property (**3**) of proposition (**4.2**), we have $G = (\vec{\mathcal{B}}(d_1,n) \cdot \vec{C}_n) \cdot (\vec{\mathcal{B}}(d_2,n) \cdot \vec{C}_n)$. As $\vec{C}_n \cdot \vec{C}_n = n\vec{C}_n$ (from lemma (**4.1**) with $q = n$ and $s = r = 1$), then we obtain:

$$G = \vec{\mathcal{B}}(d_1,n) \cdot \vec{\mathcal{B}}(d_2,n) \cdot (n\vec{C}_n) = n(\vec{\mathcal{B}}(d_1,n) \cdot \vec{\mathcal{B}}(d_2,n) \cdot \vec{C}_n) = n\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n).$$

□

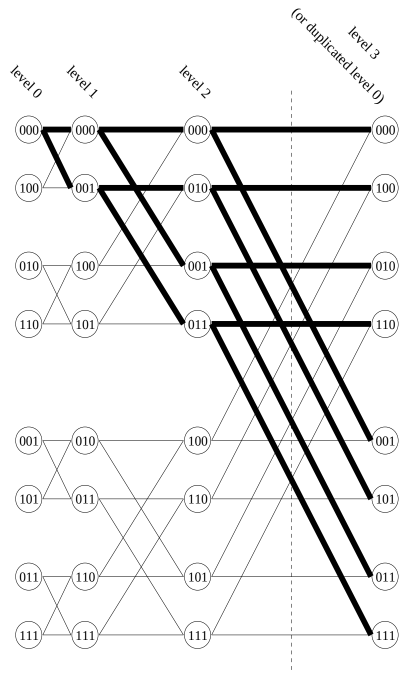

Figure 5: *The graph $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,3)$ as a conjunction of $\vec{\mathcal{B}}(2,3)$ with $\vec{C}_3$.*

Corollary 4.2 *If d_1 and d_2 are relatively prime, and if $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1,n)$* (resp. $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_2,n)$) *admit a_1* (resp. a_2) *arc-disjoint Hamilton circuits, then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n)$ admit a_1a_2 arc-disjoint Hamilton circuits.*

Proof. Let $\vec{C}_{nd_1^n}$ *(resp. $\vec{C}_{nd_2^n}$)* be a Hamilton circuit in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1,n)$ *(resp. $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_2,n)$).* From lemma (**4.1**), the conjunction $\vec{C}_{nd_1^n}\cdot\vec{C}_{nd_2^n}$ is a set of n circuits of length $nd_1^nd_2^n$.

As $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n)$ has $nd_1^nd_2^n$ vertices, the 1-difactor $\vec{C}_{nd_1^n}\cdot\vec{C}_{nd_2^n}$ consists of n circuits each one being a Hamilton circuit of a connected component of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1,n)\cdot\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_2,n)$ isomorphic to $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n)$. So, the conjunction of one Hamilton circuit of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1,n)$ with one Hamilton circuit of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_2,n)$ provides one Hamilton circuit in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n)$. Applying this results to the a_1a_2 different couples of circuits, provides a_1a_2 arc-disjoint Hamilton circuits in $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d_1d_2,n)$. $\square$

Given a Hamilton decomposition of the digraphs $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p^i,n)$ with p prime, corollary (**4.2**) enables to construct a Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$, so it is enough to prove conjecture (**6**) for d being a power of a prime.

4.2 Reduction to the case where p is prime

We would like to prove that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2) = \vec{\mathcal{B}}(d,2)\cdot\vec{C}_2$ has a Hamilton decomposition. That appears quite difficult. However, we will prove that for $n \geq 3$, $\vec{\mathcal{B}}(d,2)\cdot\vec{C}_n$ has a Hamilton decomposition. Such a decomposition will then be sufficient to reduce the problem to case of prime degrees.

Lemma 4.2 *For any number $n \geq 3$ and any prime p, $\vec{\mathcal{B}}(p,2)\cdot\vec{C}_n$ can be decomposed into p Hamilton circuits.*

Proof. Let the nodes of $\vec{\mathcal{B}}(p,2)\cdot\vec{C}_n$ be labeled (xy,l) with $x \in \mathbb{Z}_d$, $y \in \mathbb{Z}_d$ and $l \in \mathbb{Z}_n$. The digraph $\vec{\mathcal{B}}(p,2)\cdot\vec{C}_n$ is similar to the wrapped butterfly digraph, so we can also define a multistage network obtained by duplicating the level 0 into a level n. Formaly this multistage network is $\vec{\mathcal{B}}(p,2)\cdot\vec{P}_n$ where $\vec{P}_n$ is a directed path of length n (.i.e. with $n+1$ vertices); its vertices will be labeled (xy,l) with $x \in \mathbb{Z}_d$, $y \in \mathbb{Z}_d$ and $l \in \{0,1,\cdots,n\}$.

Like in section (**2**), we can define a notion of realisable permutation in $\vec{\mathcal{B}}(p,2)\cdot\vec{P}_n$ except that now there is more than one dipath connecting $(xy,0)$ to $(\pi(xy),n)$. We will say that $\vec{\mathcal{B}}(p,2)\cdot\vec{P}_n$ realizes k compatible permutations of $\mathbb{Z}_p^2$ $\pi_0,\pi_1,\ldots,\pi_{k-1}$ if there exist kp^2 dipaths $P_j(xy)$, $0 \leq j \leq k-1$, $xy \in \mathbb{Z}_p^2$, where $P_j(xy)$ connects $(xy,0)$ to $(\pi_j(xy),n)$ in $\vec{\mathcal{B}}(p,2)\cdot\vec{P}_n$ satisfying the following properties. For a given j the p^2 dipaths $P_j(xy)$ are vertex-disjoint (realisability's property) and all the kp^2 dipaths $P_j(xy)$ are arc-disjoint (compatibility's property).

Using the same argument than in lemma 2.1, we can state that $\vec{\mathcal{B}}(p,2)\cdot\vec{C}_n$ can be decomposed into p Hamilton circuits if and only if $\vec{\mathcal{B}}(p,2)\cdot\vec{P}_n$ realizes p compatible cyclic permutations.

We will show by induction that $\vec{\mathcal{B}}(p,2)\cdot\vec{P}_n$ realizes p compatible cyclic permutations; more exactly we will prove that if the property is true for n, it is also true for $n+3$. First we give for $n = 3, 4$ and 5 the dipaths $P_j(xy)$ associated to compatible cyclic permutations.

In all the dipaths that we consider, a vertex (xy,l) is followed by a vertex $(yx', l+1)$ with $x' = g_l(x,y,j) = ax + f(y) + cj$, where a, f and c depend on the level l.

- For any j, the dipaths $P_j(xy)$ are vertex-disjoint if and only if the functions $x \rightarrow g_l(x,y,j)$ are for given l and j one to one mappings. As p is prime that is satisfied if and only if $a \neq 0$.

- The dipaths $P_j(xy)$ are arc-disjoint if and only if the functions $j \rightarrow g_l(x,y,j)$ are for a given l one to one mappings. That is satisfied if and only if $c \neq 0$.

In order to simplify the notations we omit in the labels of the vertices the values of the levels; as a vertex of level l is always followed by a vertex of level $l+1$.

4.2.1 Initial constructions

Let δ_0 denote the function of $\mathbb{Z}_p$ into $\{0,1\}$:

$$\delta_0(x) = \begin{cases} 1 & \text{if } x = 0 \\ 0 & \text{if } x \neq 0 \end{cases}$$

$n = 3$

$$P_j(xy) = \quad xy \quad y(x+y+j) \quad (x+y+j)(x+1) \quad (x+1)(y+\delta_0(x+1))$$

$n = 4$

$$P_j(xy) = \quad xy \quad y(x+j) \quad (x+j)(y+j) \quad (y+j)(x+1) \quad (x+1)(y+\delta_0(x+1))$$

Figure (**6**) shows one decomposition of $\vec{\mathcal{B}}(3,2)\cdot\vec{C}_4$ into circuits. To produce a clearer figure, vertices ab on odd levels are ranked lexicographically and those on even levels in the following order: $ab < a'b'$ if $b < b'$ or $b = b'$ and $a < a'$.

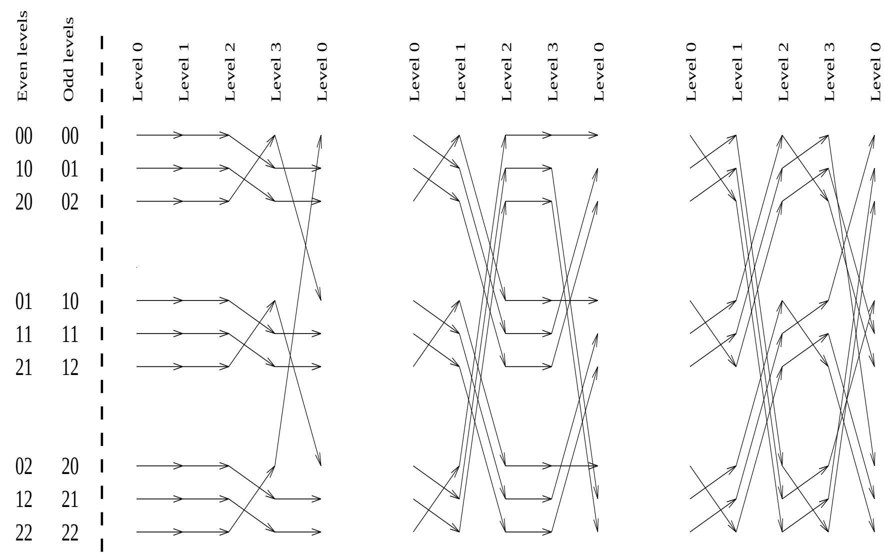

Figure 6: *A decomposition of $\vec{\mathcal{B}}(3,2)\cdot\vec{C}_4$, presented with a special ranking of the vertices.*

$n = 5$ **and** $p \neq 2$

$$P_j(xy) = \quad xy \quad y(x+y+j) \quad (x+y+j)(x+2j) \quad (x+2j)y \quad y(x+1) \quad (x+1)(y+j+\delta_0(x+1))$$

$n = 5$ **and** $p = 2$

$$P_j(xy) = \quad xy \quad y(x+y+j+1) \quad (x+y+j+1)(y+j) \quad (y+j)(x+j+1) \quad (x+j+1)y \quad y(x+1)$$

In all the cases one can easily verify that the functions $g_l(x, y, j)$ are of the form $ax + f(y) + cj$ with $a \neq 0$ and $c \neq 0$. For example, in the construction for $n = 3$, the functions implicitely defined are:

$$\begin{array}{ccc} (X,Y) & \stackrel{X'=aX+f(Y)+cj}{\longrightarrow} & (Y,X') \\ (x,y) & \stackrel{X'=X+Y+j}{\longrightarrow} & (y,x+y+j) \\ (y,x+y+j) & \stackrel{X'=Y-X-j+1}{\longrightarrow} & (x+y+j,x+1) \\ (x+y+j,x+1) & \stackrel{X'=X-Y-j+1+\delta_0(Y)}{\longrightarrow} & (x+1,y+\delta_0(x+1)) \end{array}$$

To complete the proof it remains to note that in the three first cases the permutation induced by the construction $\pi(xy) = (x + 1)(y + cj + \delta_0(x + 1))$ is cyclic, and that in the case $n = 5$, $p = 2$ $\pi(xy) = y(x + 1)$ is also cyclic as $p = 2$.

4.2.2 Induction step

The induction step follows from two facts. First, it can be easily seen that $\vec{\mathcal{B}}(p, 2) \cdot \vec{P}_{n+m}$ realizes p compatible permutations π_j, $0 \leq j \leq p - 1$, if and only if there exist two sets of permutations π'_j and π''_j, $0 \leq j \leq p - 1$, such that:

- for $0 \leq j \leq p - 1$, $\pi_j = \pi'_j \pi''_j$,
- $\vec{\mathcal{B}}(p, 2) \cdot \vec{P}_n$ realizes the p compatible permutations π'_j, $0 \leq j \leq p - 1$,
- $\vec{\mathcal{B}}(p, 2) \cdot \vec{P}_m$ realizes the p compatible permutations π''_j, $0 \leq j \leq p - 1$.

Secondly, $\vec{\mathcal{B}}(p, 2) \cdot \vec{P}_3$ realizes p compatible permutations π_j, $0 \leq j \leq p - 1$, such that each $\pi_j = e$ is the identity permutation. Indeed, let us consider the dipaths:

$$P_j(xy) = \quad xy \quad y(x + y + j) \quad (x + y + j)x \quad xy$$

Once again a vertex XY of level l is joined to a vertex YX' of level $l+1$ with $X' = g_l(X, Y, j) = aX + f(Y) + cj$ with $a \neq 0$ and $c \neq 0$.

So if $\vec{\mathcal{B}}(p, 2) \cdot \vec{P}_n$ realizes p compatible permutations π'_j then $\vec{\mathcal{B}}(p, 2) \cdot \vec{P}_{n+3}$ realizes the same compatible permutations.

So we can conclude by induction that $\vec{\mathcal{B}}(p, 2) \cdot \vec{C}_n$ can be decomposed into p Hamilton circuits for any number $n \geq 3$. $\square$

Theorem 4.3 *If a digraph G with at least 3 vertices, contains k arc-disjoint Hamilton circuits, then $\vec{\mathcal{B}}(d, 2) \cdot G$ contains dk arc-disjoint Hamilton circuits.*

Proof. First we prove the result for d prime. by hypothesis $G \supset \bigoplus_{0 \leq i < k-1} \vec{C}^i_l$, where l being the number of vertices of G is greater than 3. Hence

$$\vec{\mathcal{B}}(d, 2) \cdot G \supset \vec{\mathcal{B}}(d, 2) \cdot \bigoplus_{0 \leq i < k-1} \vec{C}^i_l = \bigoplus_{0 \leq i < k-1} \vec{\mathcal{B}}(d, 2) \cdot \vec{C}^i_l$$

From lemma (**4.2**) $\vec{\mathcal{B}}(d,2)\cdot\vec{C}_l = \bigoplus_{0\le j\le d-1}\vec{C}^{j}_{d^2l}$. So, $\vec{\mathcal{B}}(d,2)\cdot G \supset \bigoplus_{0\le i<k-1}\bigoplus_{0\le j\le d-1}\vec{C}^{i,j}_{d^2l}$ which gives $\vec{\mathcal{B}}(d,2)\cdot G \supset \bigoplus_{0\le m\le kd-1}\vec{C}^{m}_{d^2l}$.

Suppose now that the result holds for all integers strictly less than d. If d is prime we just proved the result. Otherwise $d = d_1 p$ where p is a prime and $d_1 < d$. By proposition (**4.2-2**), $\vec{\mathcal{B}}(d,2) = \vec{\mathcal{B}}(p,2)\cdot\vec{\mathcal{B}}(d_1,2)$. As a consequence $\vec{\mathcal{B}}(d,2)\cdot G = \vec{\mathcal{B}}(p,2)\cdot(\vec{\mathcal{B}}(d_1,2)\cdot G)$. By induction $G' = \vec{\mathcal{B}}(d_1,2)\cdot G$ contains at least d_1k arc-disjoint Hamilton circuits. Moreover p being prime $G = \vec{\mathcal{B}}(p,2)\cdot G'$ will contain $pd_1k = dk$ arc-disjoint Hamilton circuits. □

When G can be decomposed into Hamilton circuit the above theorem can be restated as:

Theorem 4.4 *If G with more than 3 vertices can be decomposed into Hamilton circuits, then $\vec{\mathcal{B}}(d,2)\cdot G$ can also be decomposed into Hamilton circuits.*

Corollary 4.3 *If $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2)$ with $d \neq 1$ can be decomposed into Hamilton circuits, then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(qd,2)$ can also be decomposed into Hamilton circuits for any integer q.*

Proof. Just apply theorem (**4.4**) to $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(qd,2)$ which is by proposition (**4.2**) $\vec{\mathcal{B}}(q,2)\cdot\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2)$. Note that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(1,2)$ has only 2 vertices so the corollary cannot be applied for $d = 1$. □

Example 1 As $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(4,2)$ has a Hamilton decomposition then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(4q,2)$ has a Hamilton decomposition for any integer q. In particular $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2^i,2)$ has a Hamilton decomposition for $i \ge 2$.

Corollary 4.4 *To prove conjecture (**6**) it suffices to prove that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p,2)$ has a Hamilton decomposition for any prime $p \ge 5$.*

Proof. Let d be a non prime number. If d has a prime factor p different from 2 or 3, by corollary (**4.3**), it suffices to prove the conjecture for $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p,2)$. If $d \ge 4$ has only prime factors equal to 2 or 3, then $d = 2^i3^j$ with $i + j \ge 2$. A computer search shows that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(4,2)$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(6,2)$ and $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(9,2)$ have a Hamilton decomposition. So, according to corollary (**4.3**), $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2^i3^j,2)$ with $i + j \ge 2$ have a Hamilton decomposition. □

Remark 3 Although it is not the purpose of this article, proposition (**4.3**) can be used to improve results on the decomposition of de Bruijn digraphs into Hamilton circuits,

Proposition 4.4

- *If p is the greatest prime dividing d, then $\vec{\mathcal{B}}(d,2)$ contains $\frac{p-1}{p}d$ Hamilton circuits.*
- *$\vec{\mathcal{B}}(2^iq,2)$ contains $(2^i-1)q$ Hamilton circuits.*

Proof. The result holds for $p = 1$, now for $p > 1$, by a result of D. Barth, J. Bond and A. Raspaud [6] we know that $\vec{\mathcal{B}}(p, 2)$ contains $p-1$ arc-disjoint Hamilton circuits for p a prime, and have at least 4 vertices. Hence theorem (**4.3**) implies that, as $\vec{\mathcal{B}}(d, 2) = \vec{\mathcal{B}}(pd_1, 2) = \vec{\mathcal{B}}(d_1, 2) \cdot \vec{\mathcal{B}}(p, 2)$, $\vec{\mathcal{B}}(d, 2)$ contains $(p-1)d_1$ arc-disjoint Hamilton circuits.

Similary, the second result follows from a result of R. Rowley and B. Bose [14] stating that $\vec{\mathcal{B}}(2^i, 2)$ contains $2^i - 1$ Hamilton circuits. □

4.3 Exhaustive Search for Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p, 2)$

As seen above the problem has been reduced to find a Hamilton decomposition of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p, 2) = L(\vec{\mathcal{K}}_{p,p})$ for a prime $p \geq 5$. In order to provides ideas and to strengthen our conjecture we have performed some exhaustive searches. The complexity of an exhautive search being exponential, we have restricted the set of solutions to the one such that *the i-th circuit H_i is obtained from H_0 by applying the automorphism ϕ_i of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(p, 2)$ which sends vertex (ab, l) on vertex $(a(b+i), l)$*. Furthermore we want solutions such that H_0 is Hamiltonian and the Hamilton circuits $H_i = \phi_i(H_0)$, $0 \leq i \leq p-1$ are arc-disjoint. However the search space is still exponentional in p and a computer search (with normal computation ressources) cannot be successful for p greater than 7. So we restricted again the search space to *"nearly linear"* solutions.

This restriction gave us solutions for small primes < 29. For example for $p = 5$ we found the cycle H_0 given by the following set of arcs:

$$\begin{array}{lcll}
(ab, 0) & \rightarrow & (a(2b), 1) & \text{for } a \notin \{0, 1\}, \\
(0b, 0) & \rightarrow & (0(2b+1), 1), & \\
(1b, 0) & \rightarrow & (1(2b+2), 1), & \\
(ab, 1) & \rightarrow & ((2a+b)b, 0). &
\end{array}$$

It induce the next cyclic permutation on level 0.

$$(00, 11, 14, 20, 40, 30, 10, 42, 24, 23, 01, 33, 21, 12, 31, 32, 04, 44, 13, 03, 22, 34, 43, 41, 02)$$

Finally we looked for very special Hamilton circuits H_0. That enables us to find a solution for every prime p between 7 and 12000. More precisely, we searched for parameters α and β in $\mathbb{Z}_p$ such that H_0 is given by the following set of arcs:

$$\begin{array}{lcll}
(ab, 0) & \rightarrow & (a(\alpha b), 1) & \text{for } a \neq 0, \\
(0b, 0) & \rightarrow & (0(\alpha b + \beta), 1), & \\
(ab, 1) & \rightarrow & ((a+b+1)b, 0). &
\end{array}$$

One can easily check that if $\alpha \neq 1$ the H_i's are arc-disjoint. So we have only to find α and β such that H_0 is a Hamilton circuit. In particular we should have $\alpha \neq 0$ (condition to obtain a one difactor) and $\beta \neq 0$ (otherwise we obtain a circuit of length p starting at vertex $(0, 0)$). We conjecture that:

Conjecture 7 *For any prime $p > 5$ there exist $\alpha \notin \{0,1\}$ and $\beta \neq 0$ such that the permutation π of $\mathbb{Z}_p^2$ defined by:*

$$\left.\begin{array}{rcll} \pi(ab) & = & (a+\alpha b+1)(\alpha b) & \text{if } a \neq 0, \\ \pi(0b) & = & (\alpha b+\beta+1)(\alpha b+\beta). & \end{array}\right\} \text{ is cyclic.}$$

The number of possible solutions is then only p^2. So we have been able to verify the conjecture, by a computer search for large values of p (≤ 12000). Below we give some solutions for p less than 100.

p	7	11	13	17	19	23	29	31	37	41	43	47	53	59	61	67	71	73	79	83	89	97
α	2	3	4	2	6	2	7	2	3	2	4	3	4	3	4	6	2	2	2	3	2	2
β	3	7	4	14	4	13	28	11	19	25	22	18	29	1	25	14	28	27	51	37	25	16

For example for $p = 7$, $\alpha = 2$, $\beta = 3$ we obtain the following cyclic permutation on level 0:

$$\begin{array}{l}(00,43,46,35,03,32,14,31,62,44,61,22,04,54,01,65,33,36,25,63,66,55,23,26,15,\\ \quad 53,56,45,13,16,05,06,21,52,34,51,12,64,11,42,24,41,02,10,20,30,40,50,60)\end{array}$$

So using corollary (**4.3**) we have:

Theorem 4.5 *If d is divisible by any number q, $4 \leq q \leq 12000$ then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,2)$ and consequently $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ has a Hamilton decomposition.*

This result can be strengthened in the case of $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,4)$. Indeed we know that $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,4)$ and $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(3,4)$ have a Hamilton decomposition and we have been able to generalize lemma (**4.2**) for $\vec{\mathcal{B}}(p,4) \cdot \vec{C}_n$ with p odd prime and $n \geq 5$.

Theorem 4.6 *If d is divisible by any number q, $2 \leq q \leq 12000$ then $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,4)$ and consequently $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(d,n)$ for $n \geq 4$ has a Hamilton decomposition.*

As a consequence the butterflies $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2p,n)$ have an Hamilton decomposition for $n \geq 4$.

5 Conclusion

In this paper, we have shown that in a lot of cases Butterfly digraphs have a Hamilton decomposition and give strong evidence that the only exceptions should be $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,2)$, $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(2,3)$ and $\mathcal{W}\vec{\mathcal{B}}\mathcal{F}(3,2)$. We have furthermore reduced the problem to check if $L(\vec{\mathcal{K}}_{p,p})$ has a Hamilton decomposition for p prime (or equivalently that $\vec{\mathcal{K}}_{p,p}$ has an eulerian compatible decomposition). We have also shown that such a decomposition will follow from the solution of a problem (conjecture (**7**)) in number theory. Our interest came from a conjecture of D. Barth and A. Raspaud [7] concerning the decomposition of Butterfly networks into undirected Hamilton cycles. This conjecture is solved in [8] by generalizing the technics of section (**3.2**).

Finally we have seen in proposition (**4.4**) that the technics can be applied to obtain results on the Hamilton decomposition of de Bruijn digraphs. In the spirit it will be interesting to solve the following problem:

Problem: Determine the smallest integer $f_d(n)$ such that $\vec{\mathcal{B}}(d,n) \cdot \vec{C}_{f_d(n)}$ has a Hamilton decomposition.

A proof similar to that of lemma (**4.2**) should lead to $f_d(n) \leq n+1$. Conjecture (**6**) is, for a given d, equivalent to $f_d(n) \leq n$.

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Table of Contents

Unité de recherche INRIA Lorraine, Technopôle de Nancy-Brabois, Campus scientifique,
615 rue du Jardin Botanique, BP 101, 54600 VILLERS LÈS NANCY
Unité de recherche INRIA Rennes, Irisa, Campus universitaire de Beaulieu, 35042 RENNES Cedex
Unité de recherche INRIA Rhône-Alpes, 46 avenue Félix Viallet, 38031 GRENOBLE Cedex 1
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