

On termination of the direct sum of term rewriting systems

Michaël Rusinowitch

► **To cite this version:**

Michaël Rusinowitch. On termination of the direct sum of term rewriting systems. [Research Report] RR-0679, INRIA. 1987, pp.10. <inria-00075874>

HAL Id: inria-00075874

<https://hal.inria.fr/inria-00075874>

Submitted on 24 May 2006

HAL is a multi-disciplinary open access archive for the deposit and dissemination of scientific research documents, whether they are published or not. The documents may come from teaching and research institutions in France or abroad, or from public or private research centers.

L'archive ouverte pluridisciplinaire **HAL**, est destinée au dépôt et à la diffusion de documents scientifiques de niveau recherche, publiés ou non, émanant des établissements d'enseignement et de recherche français ou étrangers, des laboratoires publics ou privés.

INRIA

UNITÉ DE RECHERCHE
INRIA-LORRAINE

Institut National
de Recherche
en Informatique
et en Automatique

Domaine de Voluceau
Rocquencourt
B.P. 105
78153 Le Chesnay Cedex
France

Tel. (1) 39 63 55 11

Rapports de Recherche

N° 679

**ON TERMINATION OF THE
DIRECT SUM OF TERM
REWRITING
SYSTEMS**

Michael RUSINOWITCH

JUIN 1987

**ON TERMINATION OF THE DIRECT SUM OF TERM REWRITING
SYSTEMS**

PAR

MICHAEL RUSINOWITCH

Résumé : Pour que la somme directe de deux systèmes de réécriture termine, nous montrons qu'il suffit qu'aucun membre droit de règle ne soit réduit à une variable, ou bien qu'aucun membre droit ne contienne plus d'occurrence d'une variable que le membre gauche correspondant.

Abstract : Some sufficient conditions are given for the termination of the direct sum of two term rewriting systems : either no right-hand-side of a rule is a variable, or no right-hand-side contains more occurrences of a variable than the corresponding left-hand-side.

ON TERMINATION OF THE DIRECT SUM OF TERM REWRITING SYSTEMS.

Michael Rusinowitch.
CRIN Campus Scientifique BP 239
54506 Vandoeuvre-les-Nancy Cedex France.

ABSTRACT

Some sufficient conditions are given for the termination of the direct sum of two term rewriting systems: either no right-hand-side of a rule is a variable, or no right-hand side contains more occurrences of a variable than the corresponding left-hand-side.

Keywords: term-rewriting system, termination.

INTRODUCTION.

It was recently proved in [Toyama 86a] that the direct sum of two term-rewriting systems R_0 and R_1 is confluent iff R_0 and R_1 are confluent. However, the same theorem is false if we replace the word confluent by terminating; this was shown in [Toyama 86b].

Let F_0 and F_1 be two disjoint sets of symbols and X an infinite set of variables. We suppose that R_0 (resp. R_1) is a noetherian term-rewriting system on $T(F_0, X)$ (resp. $T(F_1, X)$). Then the direct sum of R_0 and R_1 , which is denoted by $R_0 \cup R_1$, is noetherian on $T(F_0 \cup F_1, X)$ if one of the following conditions is satisfied:

Condition A: for every rule $l \rightarrow r$ in $R_0 \cup R_1$ we have $V(r) \subseteq V(l)$, where $V(t)$ denotes the multiset of variables of t .

Condition B: for every rule $l \rightarrow r$ in $R_0 \cup R_1$, r is not a variable.

EXAMPLES:

* R_0 : $f(0,1,x,x,x) \rightarrow f(x,x,x,0,0)$

R_1 : $or(x,y) \rightarrow x$

$or(x,y) \rightarrow y$

* The previous result can be trivially applied to two closed TRS.

COUNTEREXAMPLES:

* [Toyama 86b] R_0 : $f(0,1,x) \rightarrow f(x,x,x)$

R_1 : $or(x,y) \rightarrow x$

$or(x,y) \rightarrow y$

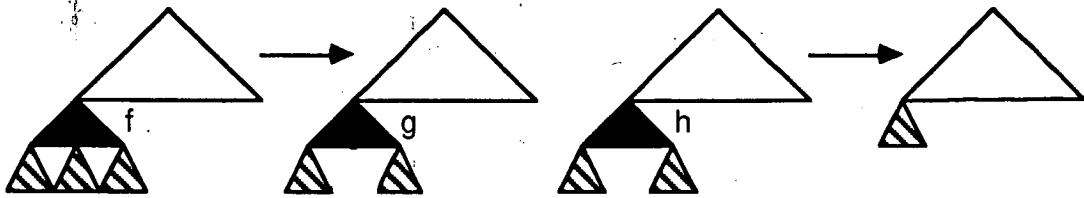
* If we replace $f(x,x,x)$ above by $f(f(x,x,x),f(x,x,x),f(x,x,x))$, we get a system that is not even quasi-terminating.

* There is in [Barendregt & Klop 86] a direct sum of canonical term-rewriting systems which is not canonical.

DEFINITION: a rewrite rule is **collapsing** when its right-hand side is a variable, and **duplicant** when its right-hand side has strictly more occurrences of one variable than its left-hand side.

A technique to prove termination is to embed the rewriting relation in a well-founded ordering on terms [Dershowitz 85] such as simplification orderings [Dershowitz 82][Plaisted 78][Jouannaud et al. 82] which are usually recursively defined. Another idea is to associate with each term a quantity which decreases when the term is rewritten, but cannot decrease for ever. The quantity may be a polynomial interpretation of the terms (see, for instance, [Huet-Oppen 80]).

Our goal now is to find such a quantity. In order to prove termination under hypothesis A, we divide terms into connected domains, such that all the symbols within a domain belong to the same F_i ($i=0$ or 1). During a rewriting sequence, the number of frontier occurrences between two domains cannot increase, since there is no duplicant rule in the system (figure a). It may decrease, due to the presence of collapsing rules (figure b). When it is not decreasing, then the multiset of connected domains of the term is rewritten by $R_0 \cup R_1$, which is known to terminate on homogeneous terms.



using the rule $f(x,x,y) \rightarrow g(x,y)$
figure a

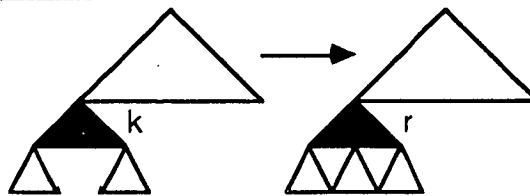
using the rule $h(x,x) \rightarrow x$
figure b

To deal with hypothesis B, we first note that the maximal number of connected domains traversed by a branch of a term is constant on every derivation starting from this term, due to the absence of collapsing rules. Then, we can show that the multiset of connected domains figuring at some level is decreasing, without the lower levels increasing (figure c). The termination is achieved as in the proof that nested multiset orderings [Dershowitz Manna 79] are well-founded.

level 1 is unchanged

level 2 decreases

level 3 increases



with the rule $k(x,x) \rightarrow r(x,x,x)$

figure c

The main difficulty is that a rewriting step by R_0 at an occurrence may generate new rewriting opportunities by R_1 above that occurrence as in the following example:

$$R_0: a \rightarrow b$$

$$R_1: f(x,x) \rightarrow g$$

$$f(a,b) \rightarrow f(b,b) \rightarrow g.$$

Let us prove now more formally our statements.

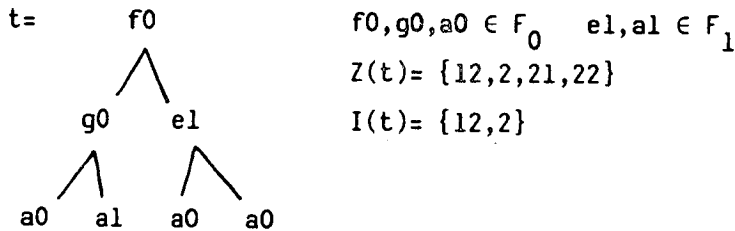
1. DEFINITIONS

Terms are considered as trees. Positions in trees are defined by occurrences which are sequences of integers (see [Huet 80] for rigorous definitions). We call $\text{occ}(t)$ (resp. $J(t)$) the set of non-variable (resp. variable) occurrences of a term t . Given an occurrence i of a term t , $p(i)$ denotes the maximal proper prefix of i . If ϵ is the root occurrence of t , $p(\epsilon)$ is undefined.

Let $Z(t)$ be the set of occurrences i of t , such that the symbol at i comes from a different theory than the symbol at $p(i)$. To put it more formally $Z(t) = \{i \in \text{occ}(t); t(i) \in F_j \text{ and } t(p(i)) \in F_{1-j}, \text{ for } j=0 \text{ or } 1\}$.

Let $I(t)$ the subset of elements of $Z(t)$ such that none of their strict left factors is in $Z(t)$. Hence, every symbol occurring at $i \in I(t)$ (and before i) is in the same theory than the root symbol. In the following, function symbols whose index is j ($j=0$ or 1) are implicitly supposed to belong to F_j .

example:



Let $SK(t)$ be the pure part of t . To be more precise, let Ω be a new variable, then we have:

$$SK(t) = t[i \leftarrow \Omega]_{i \in I(t) \cup J(t)}$$

It is easy to check that $J(SK(t)) = I(t) \cup J(t)$. In particular, when x is a variable, $SK(x) = \Omega$.

For any substitution $\theta: X \rightarrow T(F, X)$, and for $i=1,0$ we define $SK_i(\theta)$ to be the following substitution: for all $x \in X$,

$$SK_i(\theta)(x) = \begin{cases} \Omega & \text{if } \theta(x)(\epsilon) \in F_{1-i} \\ SK(\theta(x)) & \text{otherwise} \end{cases}$$

We just state the next easy proposition:

Proposition 1: if $t \in T(F_1, X)$ then $SK(\theta(t)) = SK_1(\theta)(t)$.

example :

$$\begin{array}{lcl}
 t = & f_0 & \theta(x) = a_1 \qquad SK_0(\theta)(x) = \Omega \\
 & \wedge & \theta(y) = f_0(a_1, a_1) \quad SK_0(\theta)(y) = f_0(\Omega, \Omega) \\
 & \begin{array}{cc} x & y \end{array} &
 \end{array}$$

Let $P(t)$ be the multiset of the maximal homogeneous terms we can find in t . To be more precise, P is the unique solution of the recursive equation: $P(t) = \{SK(t)\} \cup (\cup_{i \in I(t)} P(t/i))$. We shall call $SP(t)$ the sub-multiset $(\cup_{i \in I(t)} SK(t/i))$ of $P(t)$.

example :

$$P \left(\begin{array}{c} f_0 \\ \wedge \\ g_0 \quad e_1 \\ \wedge \quad \wedge \\ a_0 \quad a_1 \quad a_0 \quad a_0 \end{array} \right) = \left\{ \begin{array}{c} f_0 \\ \wedge \\ g_0 \quad \Omega \\ \wedge \\ a_0 \quad \Omega \end{array}, \begin{array}{c} e_1 \\ \wedge \\ \Omega \quad \Omega \end{array}, a_1, a_0, a_0 \right\}$$

The term-rewriting system $R_0 \cup R_1$ induces a rewriting relation " \rightarrow " on $T(F_0, \Omega) \cup T(F_1, \Omega)$. Since every " \rightarrow chain" is either a " \rightarrow_{R_0} chain" or a " \rightarrow_{R_1} chain", it is a finite chain:

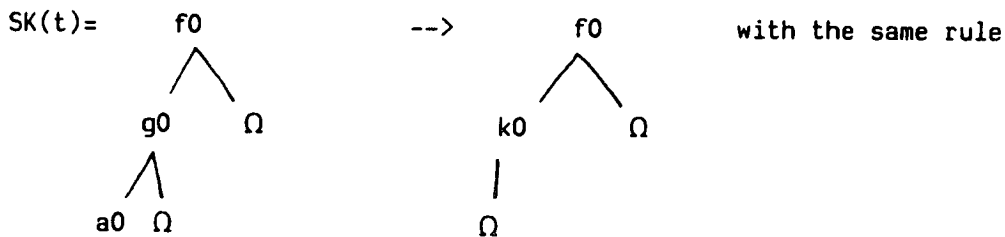
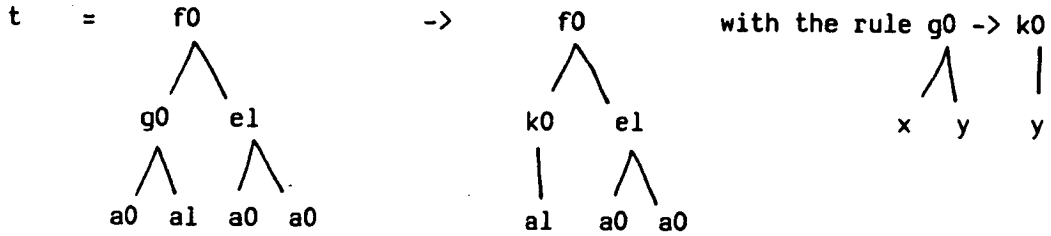
Proposition 2: \rightarrow is noetherian.

Every rewriting on t which is performed at an occurrence of the domain of $SK(t)$ can be transposed into a rewriting of $SK(t)$ (at the same occurrence), as it is showed by the next lemma.

Lifting Lemma: suppose that $t \rightarrow t'$ by application of the rule $l \rightarrow r$ at occurrence i of t with the substitution θ . We suppose that i is also an occurrence of $occ(SK(t))$ and $t(\varepsilon) \in F_0$ (the case F_1 is similar). Then

$$SK(t) \dashrightarrow \begin{cases} \Omega & \text{if } i=\varepsilon \text{ and } r \in X \text{ and } \theta(r)(\varepsilon) \in F_1 \\ SK(t') & \text{otherwise} \end{cases}$$

example :



proof: we have $t = t[i \leftarrow \theta(1)]$ and $t' = t[i \leftarrow \theta(r)]$. Since $i \in \text{occ}(SK(t))$ and $t(\varepsilon) \in F_0$, we also have $t(i) \in F_0$. Therefore $l(\varepsilon) = t(i) \in F_0$, $SK(t) = SK(t)[i \leftarrow SK(\theta(1))]$. The proposition above can be applied to $\theta(1)$: $SK(\theta(1)) = SK_0(\theta(1))$. Now, if $\theta(r)(\varepsilon) \in F_0$ we also have $SK(t') = SK(t)[i \leftarrow SK(\theta(r))]$. If $\theta(r)(\varepsilon) \in F_1$ and $i \neq \varepsilon$ then $SK(t') = SK(t)[i \leftarrow \Omega]$. In any of these two previous cases: $SK(t') = SK(t)[i \leftarrow SK_0(\theta(r))]$ and we proved that $SK(t) \dashrightarrow SK(t')$ by the same rule. Let us suppose now that $i = \varepsilon$ and $\theta(r)(\varepsilon) \in F_1$. This understands that r is a variable which obviously occurs in l and such that $t = \theta(1)$ and $t' = \theta(r)$. The definition of $SK(t)$ shows that $SK(t) = \theta'(1)$ for the substitution θ' defined by: $\theta'(x) = SK_0(\theta(x))$ for every $x \in X$. In particular, $\theta'(r) = SK_0(\theta(r)) = \Omega$. This last result takes care of the first case in the lemma.

Let MP be the set of finite multisets of elements in $T(F_0, \Omega) \cup T(F_1, \Omega)$. We extend \dashrightarrow to MP in a standard way [Jouannaud Lescanne 82]: we still get a noetherian relation on MP . This relation will be denoted by \Rightarrow . Let \Rightarrow^+ be the transitive closure of \Rightarrow .

2. Duplication-free systems

Let us prove now that condition A is enough to ensure the termination property for the term-rewriting system $R_0 \cup R_1$, denoted by R. Our goal will be reached by defining a noetherian ordering on $T(F,X)$ which is compatible with R [Manna & Ness theorem]:

Let us define the ordering $>$ on $T(F,X)$ by:

$t > t'$ if $|Z(t)| > |Z(t')|$ or
 $|Z(t)| = |Z(t')|$ and $P(t) \ni P(t')$

Since $>$ is the lexicographical composition of the noetherian orderings $(>, N)$ and (\ni, MP) , it is also noetherian. Now let us show that: $t \rightarrow t'$ implies $t > t'$.

Suppose $t \rightarrow t'$ by applying the rule $l \rightarrow r$ at occurrence i with the substitution θ . Let ι be the maximal prefix occurrence of i such that $\iota \in Z(t)$; otherwise ι is ϵ . We also suppose that $t(\iota) \in F_0$ (the case F_1 will be similar).

CASE 1: We suppose that i is equal to ι and $r \in X$. We also suppose that $\theta(r)(\epsilon) \in F_1$ (the root symbol of $\theta(r)$ is in F_1).

We can get t' by the replacement of t/i at occurrence i within t , by $\theta(r)$. Let h be an occurrence of t/i such that $(t/i)/h = \theta(r)$. (it suffices to take h such that $l/h = r$). We can notice that $Z((t/i)/h) \subseteq Z(t/i)$. Then, we have $(t/i)/h = t/i.h$. The hypothesis implies that i does not belong any more to $Z(t')$. Hence we have:

$$Z(t') = ((Z(t) \setminus \{i.k; k \in Z(t/i)\}) \setminus \{i\}) \cup \{i.k; k \in Z(t/i.h)\}.$$

Since $|Z(t/i.h)| \leq |Z(t/i)|$, we have $|Z(t')| < |Z(t)|$.

CASE 2: Otherwise.

In this case $|Z(t')| \leq |Z(t)|$, but one can no more ensure that the inequality is strict. However, we have:

$$P(t') = (P(t) \setminus P(t/i)) \cup P(t'/i).$$

$$P(t/i) = \{SK(t/i)\} \cup SP(t/i)$$

$$P(t'/i) = \{SK(t'/i)\} \cup SP(t'/i)$$

Firstly, the Lifting lemma implies that $SK(t/i) \rightarrow SK(t'/i)$. Now, if we can prove $SP(t'/i) \subseteq SP(t/i)$ we are done, since it is easy to derive from this inclusion : $P(t'/i) \Rightarrow P(t/i)$ and then $P(t) \Rightarrow P(t')$.

We have $t'/i = t/i[i \leftarrow \theta(r)]$. According to condition A, $V(r) \subseteq V(1)$. As a consequence, we have (*) $\{\theta(x); x \in V(r)\} \subseteq \{\theta(x); x \in V(1)\}$ (multiset inclusion). As $r \in T(F_0, X)$ and $SK(t/i) \in T(F_0, \Omega)$ we can remark that every element of $SP(t'/i)$ appears within a $\theta(x)$ for some $x \in V(r)$. Therefore (*) implies $SP(t'/i) \subseteq SP(t/i)$.

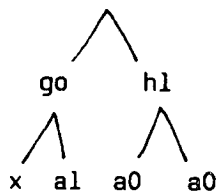
3. Collapse-free systems.

Let us prove now that Condition B ensures that $R_0 \cup R_1$ is noetherian. The technique used above fails, since, in the actual case, we may have $t' \rightarrow t$ with $Z(t) > Z(t')$. Let us introduce the notion of **level** of a term, which is the maximal number of homogeneous layers of symbols that can be encountered on a branch of the term.

The level of any term in $T(F_0, X) \cup T(F_1, X)$ is 1. The level of a term in $T(F_0 \cup F_1, X)$ is recursively defined: $level(t) = 1 + \max \{level(t/i); i \in I(t)\}$ (with $\max\{\emptyset\} = 0$).

example:

level (f0) = 3



Let us notice that Condition B ensures that the level of terms is preserved under rewriting. As usual, we start defining an ordering $>$ on $T(F_0 \cup F_1, X)$:

$t > t'$ if t and t' have the same level and (recursively)

$SK(t') \rightarrow SK(t)$ or

$SK(t') = SK(t)$ and $\{t'/i\}_{i \in I(t')} \gg \{t/i\}_{i \in I(t)}$

where \gg is the multiset extension of $>$.

This ordering presents some similarity with the "nested multiset ordering" in [Dershowitz Manna 79]. Let us recall that \rightarrow is noetherian. By induction on the level of terms, we obtain :

Proposition 4: $>$ is noetherian.

We are now in position to apply Manna & Ness method, contingent on proving:

Proposition 5: $t \rightarrow_R t'$ implies $t > t'$.

proof: if t is reduced to t' at occurrence i and i belongs to $\text{occ}(\text{SK}(t))$ then the Lifting Lemma yields: $\text{SK}(t) \rightarrow \text{SK}(t')$. Otherwise, since there is no collapsing rule, we have $\text{level}(t) = \text{level}(t')$, $\text{SK}(t) = \text{SK}(t')$ and the rewriting takes place in a subterm of t , say t/i with $i \in I(t)$. Now, by induction on the level of terms, $\{t/i\}_{i \in I(t)} >> \{t'/i\}_{i \in I(t')}$. From this we can derive $t > t'$.

CONCLUDING REMARKS.

1. We conjecture that a system $R_0 \cup R_1$ is non-terminating only if R_0 contains a collapsing rule and R_1 contains a duplicant rule or vice-versa. However, we did not succeed to build a noetherian extension of the rewriting relation, in that case.

2. Attention has been recently focused on unification algorithms in combination of theories (see for instance [Yellick 1985]). We think our results may be useful in a term-rewriting approach of these problems, as it was developed in [Kirchner 1985], especially in the presence of collapsing axioms, which are difficult to handle in unification theory.

I would like to thank J.P. Jouannaud, P. Lescanne and J.L. Remy for their comments.

REFERENCES

Dershowitz N. (1985) Termination. Proceedings of RTA. J.P. Jouannaud. LNCS 220. pp180-224.

Dershowitz N. (1982) Orderings for term-rewriting systems. Theoretical Computer Science 17(1982) pp279-301

Dershowitz N. and Manna Z. (1979). Proving termination with multiset orderings Comm. ACM 22 (8) pp465-476

Huet G. and Oppen D. C. (1980). Equations and rewrite rules : a survey. Formal Language Theory : Perspectives and Open Problems, Book R. Academic Press pp349-405

Jouannaud J-P, Lescanne P., Reinig F. (1982). Recursive decomposition ordering Formal description of programming concepts Bjorner D. North Holland pp331-348

Jouannaud J-P., Lescanne P. (1982). On multiset orderings. Information Processing Letters No 15_1982 pp57-63

Kirchner C. (1985). Methodes et outils de conception systematique d'algorithmes d'unification dans les theories equationnelles. These de Doctorat d'Etat, Universite Nancy I (1985).

Klop J.W. and Barendregt H.P. (1986). Unpublished note.

Plaisted D. (1978). A recursively defined ordering for proving termination of term rewriting systems. University of Illinois sept.78 UIUCDCS-R-78-943

Toyama Y. (1986a). On the Church-Rosser property for the direct sum of term rewriting systems, to appear in J.ACM.

Toyama Y. (1986b). Counterexamples to terminating for the direct sum of term rewriting systems. submitted .

Yelick K. (1985). Combining unification algorithms for confined regular equational theories. Proceedings of RTA. J.P.Jouannaud. LNCS 220. pp365-

Imprimé en France
par

l'Institut National de Recherche en Informatique et en Automatique

