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Derivation Schemes from OCL Expressions to B

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Abstract

In the continuity of our research on integration of UML and B, we address in this paper the transformation from OCL (Object Constraint Language), which is part and parcel of UML, into B. Our derivation schemes allow to automatically derive in B not only the complementary class invariants, the guard conditions in state-charts (in OCL) but also OCL specifications OCL class operations, events or use cases.

Keywords: *UML, OCL, OCL expression, B expression, B substitution.*

1 Introduction

The Unified Modelling Language (UML)[17] has become a de-facto standard notation for describing analysis and design models of object-oriented software systems. The graphical description of models is easily accessible. Developers and their customers intuitively grasp the general structure of a model and thus have a good basis for discussing system requirements and their possible implementation. However, since the UML concepts have English-based informal semantics, it is difficult even impossible to design tools for verifying or analysing formally UML specifications. This point is considered as a serious drawback of UML-based techniques.

To remedy such a drawback, one approach is to develop UML as a precise (i.e well defined) modelling language. The pUML (precise UML) group has been created to achieve this goal. However the main challenge [4] of pUML is to define a new formal notation that has been up to now an open issue. Furthermore, the support tool for such a new formalism is perhaps another challenge.

In waiting for a precise version of UML and its support tool, the necessity to analyse inconsistencies within UML

specifications should be solved in a pragmatic approach (cf. [3]) : formalising UML specifications by existing formal languages and then analysing UML specifications via the derived formal specifications. In this perspective, using the B language [1] to formalise UML specifications has been considered as a promising approach [13, 10]. By formalising UML specifications in B, one can use B powerful support tools like AtelierB [18], B-Toolkit [2] to analyse and detect inconsistencies within UML specifications [9]. On the other hand, we can also use UML specifications as the starting point to develop B specifications which can be then refined automatically to an executable code [6].

Meyer and Souquières [14] and Nguyen [15], based on the previous work of Lano [7], have proposed the derivation schemes from UML structural concepts into B. Each class, attribute, association and state is modelled as a B variable. The properties of those concepts are modelled as B invariants. The inheritance relationship between classes is also modelled as B invariant between B variables for the classes in question.

In [8, 11, 12] we have proposed approaches for modelling UML behavioural concepts. Each UML behavioral concept - *use case, class operation, event* - is firstly modelled by a B abstract operation in which the expected effects of such a concept on related data is specified directly on the derived data. The B operation for use cases, class operations and events may be refined afterward.

The UML-B derivation schemes for UML structural and behavioural concepts are used in three derivation procedures based on use cases [8], events [12] and class operation [11], which allow to integrate several kind of UML diagrams into the same B specification. At this stage, only the architecture, data and the operations' signature of the derived B specification are generated automatically. For the invariant within B specification, only the part that reflects the properties of UML structural concepts expressed graphically in the UML diagrams is generated. Therefore, the specification B should be completed with invariants for sup-

plementary class invariant, supplementary attribute properties as well as B operations' body.

As cited in the UML literature [16], OCL (Object Constraint Language) is often used to specify supplementary class invariant, supplementary attribute properties as well as pre- and post-conditions of behavioural concepts within UML specifications. In the continuity of our research on integration UML and B, we address in this paper the transformation from OCL expressions into B. This OCL-B translation is applied for generating supplementary invariant and the abstract operations' body of the derived B specification.

In Section 2 we outline what does look like the transformation from OCL expressions into B. The derivation schemes for OCL types and their operations are presented Section 3. The derivation schemes specific for postconditions are presented in Section 4. Discussions in Section 5 conclude our presentation.

2 From OCL expressions into B : an overview

2.1 The OCL language

The Object Constraint Language (OCL) is now part and parcel of the UML standard [16]. One can use OCL to write constraints that contain extra information about, or restrictions to, UML diagrams. OCL is intended to be simple to read and write. Its syntax is similar to object-oriented programming languages. Most OCL expressions can be read left-to-right where the left part usually represents - in object-oriented terminology - the receiver of a message. Frequently used language features are attribute access of objects, navigation to objects that are connected via association links, and operation calls. OCL expressions are not only used to define invariants on classes and types, they also allow specification of guard conditions in UML state-charts and pre- and postconditions on class operations, use cases or events. Figure 1 shows the OCL specification of the class operation `Pump::enable_Pump` according to its informal specification in [5].

2.2 The B language and method

B [1] is a formal software development method that covers the software process from specifications to implementations. The B notation is based on set theory, the language of generalised substitutions and first order logic. Specifications are composed of abstract machines similar to modules or classes; they consist of variables, invariance properties relating to those variables and operations. The state of the system, i.e. the set of variable values, is only modifiable by operations. The abstract machine can be composed in various ways. Thus, large systems can be specified in a modular

```

CONTEXT Pump::enable_Pump(pi : PUMPID, gg : GRADE, vi :
VEH_ID) : void
PRE
  Pump.allInstances→collect(pump_Id)→includes(pi)
POST
  let pp : Set(Pump) = Pump.allInstances → select(
    pump_Id@pre=pi and status@pre=disabled)
  in
  if pp → notEmpty then
    pp → forall(p | p.status = enabled) and
    pp → forall(p | p.display.grade = gg) and
    pp → forall(p | p.display.cost = costOfGrade(gg)) and
    pp → forall(p | p.display.volume = 0) and
    pp → forall(p | p.display.veh_Id = vi) and
    pp → forall(p | p.motor.status = on) and
    pp → forall(p | p.clutch.status = freed)
  else true endif

```

Figure 1. The operation `enable_Pump` in OCL

way, possibly reusing parts of other specifications. B refinement can be seen as an implementation technique but also as a specification technique to progressively augment a specification with more details until an implementation that can then be translated into a programming language like ADA, C or C++. At every stage of the specification, proof obligations ensure that operations preserve the system invariant. A set of proof obligations that is sufficient for correctness must be discharged when a refinement is postulated between two B components.

2.3 Principles to translate OCL expressions to B

The core of OCL is given by an expression language. OCL expressions can be used in various contexts, for example, to define constraints such as class invariants and pre- and postconditions on behavioural concepts. Our derivation schemes from OCL to B are therefore defined for concepts related to OCL expressions : (i) the OCL types and the associated operations and (ii) the postconditions on behavioural concepts.

It is natural to model an OCL type by a B type, which would be a B predefined type such as \mathcal{Z} , *BOOL* etc, or a B user-defined type such as sets or relations. In addition, the formalisation in B of OCL types is guided and motivated by the wish to facilitate the formalisation in B of operations on OCL types. Intuitively, an OCL expression for class invariants, for guard conditions or for preconditions on behavioural concepts should be modelled by a B expression; meaning that every OCL operation (except `oclIsNew`, which is used in postconditions on behavioural concepts) should be represented by a B expression.

The derivation schemes from OCL to B for the types and the associated operations are sufficient to derive a B expression from an OCL expression of class invariants on class diagrams, guard conditions on state-charts or preconditions on behavioural concepts. To model postconditions of be-

havioural concepts, the use of B generalised substitutions is necessary. The OCL expressions involving values after executing the behavioural concepts are translated into B substitutions.

3 Derivation schemes for OCL types and their operations

3.1 Types OCL

The types in OCL can be classified as follows. The group of predefined basic types includes Integer, Real, Boolean and String. Enumeration types are user-defined. An object type corresponds to a classifier in an object model.

Collections of values can be described by the collection types Set(T), Sequence(T) and Bag(T). These are the classical types for bulk data, namely sets, lists and multi-sets respectively. The parameter T denotes the type of the elements. Notice that types at the meta-level such as OclExpression are not considered in the translation from OCL expressions into B.

3.2 Predefined basic types

Derivation 1 (Integer) In B there are two predefined types corresponding to the OCL type Integer : \mathcal{Z} and *INT*. \mathcal{Z} is chosen as the formalisation of Integer since \mathcal{Z} is more abstract than *INT*. The OCL operations defined on Integer can be mapped to operations defined on \mathcal{Z} as shown in Table 1, where a, b are two integers and a, b denote their B formalisation.

Operations OCL	Semantics in B
a:Integer	$a \in \mathcal{Z}$
a=b	$a = b$
a<>b	$\neg(a = b)$
a+b	$a + b$
a-b	$a - b$
-a	$-a$
a*b	$a \times b$
a div b	a / b
a mod b	$a \text{ mod } b$
a<b	$a < b$
a<=b	$a \leq b$
a>b	$b < a$
a>=b	$b \leq a$
a.min(b)	$\min(\{a, b\})$
a.max(b)	$\max(\{a, b\})$
a.abs	$\text{map}(-a, a)$
a/b	$a \mapsto b$

Table 1. Modelling Integer OCL operations in B

Remark 1 (The operation “/”)

1. In OCL, the operation a/b, where a, b are two integers, gives as result a real value. Since B does not define the data type for real values, we propose to model a/b by a pair $a \mapsto b$, where a and b denote respectively the B formalisation of a and b.
2. The fact of using a rate to express the division between two integers implies to define the formalisation in B for operations between an integer and a rate.

Derivation 2 (Boolean) The OCL type Boolean is modelled in B by its correspondence *BOOL*. The Boolean OCL operations are modelled in B by expressions on *BOOL* as shown in Table 2, where a, b are two booleans and a, b denote their B formalisation.

Operations OCL	Semantics in B
a:Boolean	$a \in \text{BOOL}$
a=b	$a = b$
a<>b	$\neg(a = b)$
a or b	$a \vee b$
a xor b	$\neg(a = b)$
a and b	$a \wedge b$
not a	$\neg a$
a implies b	$\neg a \vee b$
if a then b else c endif	$\neg((a \wedge b) = (\neg a \wedge c))$

Table 2. Modelling Boolean OCL operations in B

Derivation 3 (String) The B predefined type *STRING* cannot be used to model the OCL type String due to restrictions of operations on *STRING* (only “=” and “<>” are defined for *STRING*). We propose therefore to model String by *seq(0..255)*. Hence we can use B expressions on sequences to define String OCL operations (except two operations toUpper and toLower as shown in Table 3).

Remark 2 Two operations toUpper and toLower involve a repetitive computation which is very sophisticated such that they cannot be expressed by an expression B at the level of an abstract machine.

Derivation 4 (Real) There is no B predefined type for real values, however Remark 1 suggests us a solution to approximate a real value by a rate. Hence the type Real can be modelled in B by relation $\mathcal{Z} \leftrightarrow \mathcal{Z}$. It remains to define the conversion from a real value to its corresponding rate as well as the formalisation of Real OCL operations using B expressions on $\mathcal{Z} \leftrightarrow \mathcal{Z}$, which need some further investigation and therefore is beyond the scope of the current paper.

Operations OCL	Semantics in B
a:String	$a \in seq(0..255)$
a=b	$a = b$
a<>b	$\neg(a = b)$
a.size	$size(a)$
a.concat(b)	$a \frown b$
a.subString(lower,upper)	$(a \uparrow upper) \downarrow lower$
a.toUpper	<i>no definition</i>
a.toLower	<i>no definition</i>

Table 3. Modelling String OCL operations in B

Derivation 5 (Enumeration types) Each enumeration type $Enum = \{val1, \dots, valn\}$ is modelled in B by a enumerated set serving as a user-defined type $Enum = \{val1, \dots, valn\}$. Each element $val_i\#$ in $Enum$ is modelled by an element val_i in $Enum$. The modelling in B of operations on an enumeration type is shown in Table 4, where $a\#$, $b\#$ are two values of type $Enum$ and a , b denote their B formalisation.

Operations OCL	Semantics in B
$a\#:Enum$	$a \in Enum$
$a\#=b\#$	$a = b$
$a\#\<b\#$	$\neg(a = b)$

Table 4. Modelling Enumeration OCL operations in B

Derivation 6 (Object types) According to Meyer and Souquières [14], for each class $class$, the B constant $CLASS$ models the possible instance set and the B variable $class$ model the effective instance set of $class$. Therefore, the object type $class$ is modelled in B as $CLASS$, whereas the operation $class.allInstances$ is modelled as $class$.

Derivation 7 (Collection types) Given T an OCL type. Let's call T the B formalisation of T, the B formalisation of collection types on T is as follows :

- $Set(T)$, which denotes all subsets of T, is modelled in B by $\mathcal{P}(T)$,
- $Bag(T)$, which denotes all multi-sets on T, is modelled as $T \rightarrow \mathcal{N}$. An element bag of $Bag(T)$ is therefore modelled as $bag \in T \rightarrow \mathcal{N}$ and for each element $tt : T$ of bag , $bag(tt)$ denotes the occurrence number of tt in bag ,
- $Sequence(T)$ is directly modelled by $Seq(T)$.

The formalisation of OCL operations on collection types is shown in Table 5, Table 6 and Table 7, where :

- T is an OCL type on which the collection types are defined, and $tt : T$; $ss, ss2 : Set(T)$; $bb, bb2 : Bag(T)$; $se, se2 : Sequence(T)$;
- $ss, ss2, bb, bb2, se, se2$ and tt are respectively the B formalisation of $ss, ss2, bb, bb2, se, se2$ and tt ;
- for the operation sum , T must be of type Integer.

Remark 3 (Operations on collection types) The semantics of the operation $asSequence$ on a set or a bag has not been defined in OCL therefore we cannot model it in B. It is the same for the operation $excluding$ on a sequence.

Operations OCL	Semantics in B
$ss:Set(T)$	$ss \subseteq T$
$ss=ss2$	$ss = ss2$
$ss\<ss2$	$\neg(ss = ss2)$
$ss \rightarrow union(ss2)$	$ss \cup ss2$
$ss \rightarrow union(bb)$	<i>cf. $ss \rightarrow asBag \rightarrow union(bb)$</i>
$ss \rightarrow intersection(ss2)$	$ss \cap ss2$
$ss \rightarrow intersection(bb)$	$ss \cap dom(bb)$
$ss-ss2$	$ss - ss2$
$ss \rightarrow symmetricDifference(ss2)$	$(ss \cup ss2) - (ss \cap ss2)$
$ss \rightarrow including(tt)$	$ss \cup \{tt\}$
$ss \rightarrow excluding(tt)$	$ss - \{tt\}$
$ss \rightarrow asBag$	$ss \times \{1\}$
$ss \rightarrow asSequence$	<i>no definition</i>
$ss \rightarrow size$	$card(ss)$
$ss \rightarrow count(tt)$	$card(ss \cap \{tt\})$
$ss \rightarrow includes(tt)$	$tt \in ss$
$ss \rightarrow includesAll(ss2)$	$ss2 \subseteq ss$
$ss \rightarrow includesAll(bb)$	$dom(bb) \subseteq ss$
$ss \rightarrow includesAll(se)$	$ran(se) \subseteq ss$
$ss \rightarrow excludes(tt)$	$\neg(tt \in ss)$
$ss \rightarrow excludesAll(ss2)$	$ss \cap ss2 = \phi$
$ss \rightarrow excludesAll(bb)$	$ss \cap dom(bb) = \phi$
$ss \rightarrow excludesAll(se)$	$ss \cap ran(se) = \phi$
$ss \rightarrow isEmpty$	$ss = \phi$
$ss \rightarrow notEmpty$	$\neg(ss = \phi)$
$ss \rightarrow sum$	$\Sigma(xx).(xx \in ss xx)$

Table 5. Modelling Set(T) OCL operations in B

3.3 Operations select, reject, collect, forAll, exists

Derivation 8 (select, reject, collect, forAll, exists) Given an OCL type T, let's call T the B formalisation of T. The B formalisation of OCL operations $select$, $collect$, $forAll$, $exists$ on collection types on T is shown in Table 8, where :

- $ss : Set(T)$, $bb : Bag(T)$, $se : Sequence(T)$, $tt : T$;
- $boolexprrt$ is a boolean expression on tt and $exprrt$ is an expression on tt ;

Operations OCL	Semantics in B
bb:Bag(T)	$bb \in T \rightarrow \mathcal{N}$
bb=bb2	$bb = bb2$
bb<>bb2	$\neg(bb = bb2)$
bb->union(bb2)	$\{vv, nn \mid vv \in T \wedge$ $vv \in dom(bb) \cup dom(bb2) \wedge$ $nn \in \mathcal{N} \wedge nn =$ $max(bb\{vv\} \cup bb2\{vv\})\}$
bb->union(ss)	cf. bb->union(ss->asBag)
bb->intersection(bb2)	$\{vv, nn \mid vv \in T \wedge$ $vv \in dom(bb) \cap dom(bb2) \wedge$ $nn \in \mathcal{N} \wedge nn =$ $min(bb\{vv\} \cup bb2\{vv\})\}$
bb->intersection(ss)	$dom(bb) \cap ss$
bb->including(tt)	$bb \triangleleft \{tt \mapsto (\Sigma(x). (xx \in$ $bb\{tt\} \cup \{0\} xx) + 1)\}$
bb->excluding(tt)	$bb \triangleleft (\{tt \mapsto max(\Sigma(x). (xx \in$ $bb\{tt\} \cup \{0\} xx) - 1, 0)\} \triangleright \{0\})$
bb->asSequence	no definition
bb->asSet	$dom(bb)$
bb->size	$\Sigma(vv). (vv \in dom(bb) bb(vv))$
bb->count(tt)	$\Sigma(xx). (xx \in bb\{tt\} \cup \{0\} xx)$
bb->includes(tt)	$tt \in dom(bb)$
bb->includesAll(bb2)	$dom(bb2) \subseteq dom(bb)$
bb->includesAll(ss)	$ss \subseteq dom(bb)$
bb->includesAll(se)	$ran(se) \subseteq dom(bb)$
bb->excludes(tt)	$\neg(tt \in dom(bb))$
bb->excludesAll(bb2)	$dom(bb) \cap dom(bb2) = \phi$
bb->excludesAll(ss)	$dom(bb) \cap ss = \phi$
bb->excludesAll(se)	$dom(bb) \cap ran(se) = \phi$
bb->isEmpty	$bb = \phi$
bb->notEmpty	$\neg(bb = \phi)$
bb->sum	$\Sigma(xx). (xx \in dom(bb) xx \times bb(xx))$

Table 6. Modelling Bag(T) OCL operations in B

- $ss, bb, se, tt, boolexprrt$ et $exprtt$ are respectively the B formalisation of $ss, bb, se, tt, boolexprrt$ and $exprtt$.

3.4 Attribute and navigation operations

An attribute or navigation operation on an object might return as a single value/object, a set of values/objects, a multi-set of values/objects or a sequence of values/objects. It is also possible to apply an attribute or a navigation operation on the result of another attribute or a navigation operation. Hence the target of an attribute or navigation operation can be an object, a set of objects, a multi-set of objects or even a sequence of objects. Our derivation schemes for attribute and navigation operations are based on the derivation schemes for UML structural concepts (cf. Derivation 9).

Derivation 9 (Structural concepts (extracted from [14]))

- An attribute $attr$ of type $typeAttr$ in a class $Class$ is modelled by a B variable $attr$ defined as : $attr \in class \leftrightarrow typeAttr$, where the variable $class$

Operations OCL	Semantics in B
se:Sequence(T)	$se \in seq(T)$
se=se2	$se = se2$
se<>se2	$\neg(se = se2)$
se->union(se2)	$se \sim se2$
se->append(tt)	$se \leftarrow tt$
se->prepend(tt)	$tt \rightarrow se$
se->subSequence(i,j)	$(se \uparrow j) \downarrow i$
se->at(i)	$se(i)$
se->first	$first(se)$
se->last	$last(se)$
se->asSet	$ran(se)$
se->asBag	$\{vv, nn \mid vv \in ran(se) \wedge nn \in \mathcal{N} \wedge$ $nn = card(se^{-1}\{vv\})\}$
se->including(tt)	$se \leftarrow tt$
se->excluding(tt)	no definition
se->size	$size(se)$
se->count(tt)	$card(se^{-1}\{tt\})$
se->includes(tt)	$tt \in ran(se)$
se->includesAll(se2)	$ran(se2) \subseteq ran(se)$
se->includesAll(ss)	$ss \subseteq ran(se)$
se->includesAll(bb)	$dom(bb) \subseteq ran(se)$
se->excludes(tt)	$not(tt \in ran(se))$
se->excludesAll(se2)	$ran(se2) \cap ran(se) = \phi$
se->excludesAll(ss)	$ss \cap ran(se) = \phi$
se->excludesAll(bb)	$dom(bb) \cap ran(se) = \phi$
se->isEmpty	$se = \phi$
se->notEmpty	$\neg(se = \phi)$
se->sum	$\Sigma(xx). (xx \in dom(se) se(xx))$

Table 7. Modelling Sequence(T) OCL operations in B

models the effective instance set of Class and $typeAttr$ is defined as a B set to model $typeAttr$. The relation defining $attr$ might be further refined according to additional properties of $attr$.

- A binary association $assos$ between two classes $Class$ and $Class2$ is modelled by a B variable $assos$ defined as $assos \in class \leftrightarrow class2$. If there are eventual qualifiers $q1 : Q1, \dots, qn : Qn$ at the role end of $Class$, they are modelled in a similar manner to an attribute : $q1 \in class2 \leftrightarrow Q1, \dots, q1 \in class2 \leftrightarrow Q1$. We add also a B invariant linking $assos$ and $q1, \dots, qn$ as follows : $(assos^{-1} \otimes q1 \otimes \dots \otimes qn)^{-1} \in class \times Q1 \times \dots \times Qn \leftrightarrow class2$. As for attributes, the relation defining $assos$ could be further refined according to additional properties of $assos$.

Derivation 10 (Attribute operations) Given $attr, cc, sc, bc, seqc$ an attribute, an object, a set of objects, a bag of objects and a sequence of objects of a class $Class$. Let's call $attr, cc, sc, bc$ and $seqc$ their B formalisation according to Derivation 6, Derivation 7 and Derivation 9 :

1. the expression $cc.attr$, which denotes the value(s) of the

Operations OCL	Semantics in B
ss->select(tt boolexprrt)	$\{tt tt \in ss \wedge boolexprrt\}$
bb->select(tt boolexprrt)	$\{tt, nn tt \in dom(bb) \wedge boolexprrt \wedge nn \in \mathcal{N} \wedge nn = bb(tt)\}$
ss->reject(tt boolexprrt)	$\{tt tt \in ss \wedge \neg boolexprrt\}$
bb->reject(tt boolexprrt)	$\{tt, nn tt \in dom(bb) \wedge \neg boolexprrt \wedge nn \in \mathcal{N} \wedge nn = bb(tt)\}$
ss->collect(tt boolexprrt)	$\{tt, nn tt \in expr_tt[ss] \wedge nn \in \mathcal{N} \wedge nn = card(\{xx x \in ss \wedge exprrt(xx) = tt\})\}$
bb->collect(tt exprrt)	$\{tt, nn tt \in exprrt[dom(bb)] \wedge nn \in \mathcal{N} \wedge nn = \Sigma(xx).(xx \in dom(bb) \wedge exprrt(xx) = tt bb(xx))\}$
se->collect(tt exprrt)	$\lambda(ii).(ii \in dom(se) exprrt(se(ii)))$
ss->forAll(tt boolexprrt)	$\forall(tt).(tt \in ss \Rightarrow boolexprrt)$
bb->forAll(tt boolexprrt)	$\forall(tt).(tt \in dom(bb) \Rightarrow boolexprrt)$
se->forAll(tt boolexprrt)	$\forall(tt).(tt \in ran(se) \Rightarrow boolexprrt)$
ss->exists(tt boolexprrt)	$\exists(tt).(tt \in ss \wedge boolexprrt)$
bb->exists(tt boolexprrt)	$\exists(tt).(tt \in dom(bb) \wedge boolexprrt)$
se->exists(tt boolexprrt)	$\exists(tt).(tt \in ran(se) \wedge boolexprrt)$

Table 8. Modelling in B of operations select, reject, collect, forAll and exists

attribute *attr* associated to the object *cc*, is generally modelled in B by $attr[\{cc\}]$. If the cardinality of *attr* is equal to 1, *cc.attr* can be modelled by $attr(cc)$; otherwise and if *attr* is ordered, $attr[\{cc\}]$ is interpreted as a sequence;

- the expression *sc.attr* denotes a collection of values for *attr* associated with elements in *sc*. If the cardinality of *attr* is equal to 1 then *sc.attr* denotes a set and is modelled by $attr[sc]$. If the cardinality of *attr* is multiple but *attr* is not ordered then *sc.attr* denotes a bag and is modelled by $\{vv, nn|vv \in attr[sc] \wedge nn \in \mathcal{N} \wedge nn = card(attr^{-1}[\{vv\}] \cap sc)\}$. Otherwise there is no semantics for *sc.attr* and there is no therefore corresponding B formalisation;
- the expression *bc.attr* has no semantics if *attr* is multiple and ordered; otherwise *bc.attr* denotes a bag of values of *attr* associated to the bag *bc* and is modelled by $\{vv, nn|vv \in attr[dom(bc)] \wedge nn \in \mathcal{N} \wedge nn = \Sigma(cc).(cc \in attr^{-1}[\{vv\}] \cap dom(bc)|bc(cc))\}$;
- the expression *seqc.attr* has only semantics if the cardinality of *attr* is equal to 1 and in that case it denotes a sequence of values for *attr* and is modelled by $\lambda(ii).(ii \in dom(seqc)|attr(seqc(ii)))$.

Remark 4 The navigation operations without qualifiers are modelled in a similar manner to the attribute operations. For reasons of space, we omit here those derivation schemes.

Derivation 11 (Navigation operations with qualifiers)

Given a binary association *assos* from the class *Class* to

Class2. The association *assos* is qualified by attributes $q1 : Q1, \dots, qn : Qn$ at the role end of *Class*. Given values and objects $cc : Class, v1 : Q1, \dots, vn : Qn$. Let's call *roleClass2* the role end attached to *Class2* in *assos*. The expression $cc.roleClass2[v1, \dots, vn]$ is modelled in B by $(assos^{-1} \otimes q1 \otimes \dots \otimes qn)^{-1}[\{cc \mapsto v1 \mapsto \dots \mapsto vn\}]$, where $cc, v1, \dots, vn$ are the B formalisation of *cc, v1, ..., vn* and $q1, \dots, qn$ are defined according to Derivation 9. Furthermore, if the multiplicity property of *Class2* in *assos* is equal to 1, $cc.roleClass2[v1, \dots, vn]$ can be expressed in B by $(assos^{-1} \otimes q1 \otimes \dots \otimes qn)^{-1}(cc \mapsto v1 \mapsto \dots \mapsto vn)$.

Remark 5 It is always possible to define the B semantics for navigation operations with qualifiers on a set or a bag of objects. However those situations are rarely encountered and the corresponding derivation schemes are omitted in the current paper for reasons of spaces.

Derivation 12 (Navigation to association classes)

Given *assos* a binary association class between two classes *Class* and *Class2*. Given *cc, sc, bc, seqc* an attribute, an object, a set of objects, a bag of objects and a sequence of objects of *Class*. Let's call *assos, cc, sc, bc* and *seqc* the B formalisation of *assos, cc, sc, bc* and *seqc* according to Derivation 9:

- the expression *cc.assos*, which denotes the instance(s) of *assos* associated to *cc*, is modelled by $\{cc\} \triangleleft assos$; if the cardinality of the role end of *Class2* in *assos* is equal to 1, *cc.assos* can also be modelled by $cc \mapsto assos(cc)$;
- the expression *sc.assos*, which denotes the instances of *assos* associated to the elements of *sc*, is modelled by $sc \triangleleft assos$;
- the expression *bc.assos*, which denotes a bag of instances of *assos* associated to the elements of *bc*, is modelled by $\{cc, nn|cc \in dom(bc) \triangleleft assos \wedge nn \in \mathcal{N} \wedge nn = bc(dom(\{cc\}))\}$;
- the expression *seqc.assos* has only semantics if the cardinality of the role end of *Class2* in *assos* is equal to 1 and in that case it denotes a sequence of instances of *assos* associated to the elements of *seqc* and is modelled by $\lambda(ii).(ii \in dom(seqc)|seqc(ii) \mapsto assos(seqc(ii)))$.

Remark 6 (Let expressions) All the variables declared by let expressions should be replaced by their values before the transformation.

4 Derivation schemes specific for postconditions

This section presents the modelling of OCL expressions on postconditions of behavioural concepts in B . In the sequel, the terms postconditions refers to a class operation. However the derivation schemes can also be applied for use cases and events. As said earlier (cf. Section 2.3), the postconditions OCL of an operation $oper$ are modelled in B by substitutions B in the body of the abstract operation B *oper* which correspond to $oper$. First of all are some definitions.

4.1 Definitions

Given an operation $oper$, the postconditions of $oper$ can be considered as a constraint $P(out_1, \dots, out_n, in_1, \dots, in_m)$ which links the potential “outputs” (cf. Definition 2) and potential “inputs” (cf. Definition 1) of $oper$.

Definition 1 (Operation potential inputs) The set $In_{pot} = \{in_1, \dots, in_m\}$ of potential inputs of an operation $oper$ consists of : (i) the eventual parameters stereotyped by “in” or “inout” whose value is provided upon every call to $oper$ and (ii) the objects, the attributes and the associations available upon the operation call.

Definition 2 (Operation potential output) The set $Out_{pot} = \{out_1, \dots, out_n\}$ of potential outputs of an operation $oper$ consists of : (i) the eventual return parameter, which is referenced by the name *result* in OCL, of $oper$; (ii) the eventual parameters stereotyped by “out” or “inout” of $oper$; (iii) the eventual newly created objects during the execution of $oper$ and (iv) the eventual updated attributes and associations.

Definition 3 presents a standard style of the constraint $P(out_1, \dots, out_n, in_1, \dots, in_m)$. In our opinion, the definition is enough generalised to be able to cover almost class operations. Our derivation schemes in the sequel are defined in reference to this definition.

Definition 3 (Well-formed postconditions)

1. Every potential output out_i is defined by an elementary constraint $P_i(out_i, [Input], [NewObject])$, which defines out_i according to elements of $Input$ as well as the newly created objects (the elements of $NewObject$, which is a subset of $Output$) :
 - (a) P_i states the creation of an object (out_i) by $oper$;
 - (b) P_i is a comparison between the value of out_i and the values of elements in $Input \cup NewObject$. Two cases should be distinguished : (i)

P_i is represented by $out_i = \text{expression-CL}(Input \cup NewObject)$, meaning that out_i is defined deterministically in terms of elements of $Input \cup NewObject$; (ii) P_i is represented by a boolean expression but not an equality on out_i and eventual elements of $Input \cup NewObject$, meaning that out_i is defined non deterministically in terms of elements of $Input \cup NewObject$;

- (c) P_i might represent the application of the operation *forall* on a set of objects/values to be updated by $oper$.

2. The constraint P is a combination between the elementary constraints P_1, \dots, P_n and the operations and and if ... then ... else ... endif :

- (a) the condition part in an expression if ... then ... else ... endif refers to potential inputs ;
- (b) the elementary constraints P_1, \dots, P_n are linked by and in order to compose the body of expressions if ... then ... else ... endif ;
- (c) the expressions if ... then ... else ... endif can be nested ;
- (d) two body expressions in an expression if ... then ... else ... endif contain either another expression if ... then ... else ... endif or an expression and elementary constraints ;

4.2 Modelling elementary constraints

Derivation 13 (Return parameter)

1. Given P_i in form $result = \text{expr}(Input \cup NewObject)$ to define deterministically the return value of $oper$. We add in the operation B *oper* the following substitution : $out = \text{expr}(Input \cup NewObject)$, where out represents the return parameter of $oper$ and expr is the B formalisation of expr .
2. Given P_i in form $\text{expr}(result, Input \cup NewObject)$ to define non-deterministically the return parameter of $oper$. We can rewrite the constraint P_i in the following manner :
 - we introduce a temporary variable res which takes the place of $result$ in the old P_i ;
 - we rewrite P_i in form : $\text{expr}(res, Input \cup NewObject)$ and $result = res$,

the new form of P_i enables us to update $oper$ as the following manner :

- we create a clause **any...where...** if it has not been created (cf. Derivation 17) ;
- we declare res , which models res :
any..., res where
 $...expr(res, Input[\cup NewObject])$
- we add the substitution $out: = res$ in the body of **any...where...**

Remark 7

1. The B formalisation of $expr$ is done using derivation schemes in Section 3. All the eventual occurrences of $@pre$ are omitted.
2. Derivation 13 can be extended to apply for eventual parameters stereotyped by out or $inout$ of $oper$.

Derivation 14 (Object creation) Given Pi a constraint specifying that an object cc of a class $Class$ is created by $oper$. We create in $oper$:

- a clause **any...where...** if this clause has not been created (cf. Derivation 17) ;
- a temporary variable cc , which models cc :
any ..., cc where
 $... \wedge cc \in CLASS - class$
- a B substitution, which models the updating of effective instance set of $Class$ by the object cc , in the body **any...where...**
 $\parallel class: = class \cup \{cc\}...$

Derivation 15 (Updating attribute value of an object)

Given an attribute $attr$ and an object cc of a class $Class$:

1. given Pi in form $cc.attr=expr(Input[\cup NewObject])$ to define deterministically the new value of $cc.attr$ in which the cardinality of $attr$ is equal to 1. We model the constraint Pi by the substitution B
 $attr := attr \leftarrow \{cc \mapsto expr\}...$;
2. given Pi in form $cc.attr=expr(Input[\cup NewObject])$ to define deterministically $cc.attr$ in which the cardinality of $attr$ is “*”. We model Pi by the substitution :
 $attr := attr \cup \{cc\} \times expr...$,
in the two cases above, $attr$, cc and $expr$ are the B formalisation of $attr$, cc and $expr$.

Remark 8 Derivation 15 can be extended for updating associations.

Derivation 16 (The operation forAll on a set of objects)

Given a set sc of objects and an attribute $attr$ of the cardinality 1 of a class $Class$:

1. given Pi in form $sc \rightarrow \text{forall}(p | p.attr = expr(Input[\cup NewObject]))$ to define deterministically the value of attribute $attr$ of all elements in sc . We model Pi by the following substitution B : $attr := attr \leftarrow sc \times \{expr\}...$,
2. given Pi in form $sc \rightarrow \text{forall}(p | expr(p.attr, Input[\cup NewObject]))$ to define non deterministically the value of attribute $attr$ of all elements in sc . We model Pi in creating in $oper$:

- a clause **any...where...** if this clause has not been created (cf. Derivation 17) ;
- a temporary variable at defined as :
any..., at where
 $at \in class \mapsto typeAttr \wedge$
 $dom(at) = sc \wedge$
 $\forall(tt).(tt \in ran(at) \Rightarrow$
 $expr(at, Input[\cup NewObject]))...$
- the substitution :
 $attr: = attr \leftarrow at...$ ||
in the body of the clause **any...where...**,

in the two cases above, $attr$ and sc are the B formalisation of $attr$ and sc ; and $expr$ is the B formalisation of $expr$ in which tt replaces $p.attr$.

Remark 9

1. Derivation 16 did not consider the case where several attributes of the same set of objects have changed their values. However this situation can be solved by applying Derivation 16 several times.
2. We did not consider the case where forall is applied on a sequence or a bag of objects since those situation has no semantics in the context of postconditions.
3. Derivation 16 can be extended for the case where the body of forall is the navigation operation.
4. Derivation 16 can also be extended for the case where the cardinality of $attr$ is greater than 1. However this situation is rarely encountered.

4.3 Substitution unification

In the previous section, we have presented the principles for modelling an elementary constraint. In general, each elementary constraint gives rise to a B substitution. This section discusses the way to unify the derived substitutions.

Derivation 17 (Substitution unification) Given $P1$ and ... and Pk an expression OCL with the operation and and the elementary constraint $P1, ..., Pk$ in the postconditions of

oper. The substitutions B derived from this expression are unified by the following manner :

- all the eventual temporary variables as well as the eventual created objects are declared in the same clause **any...where...** ;
- all the substitutions involving the same B variable are unified ;
- the substitutions for the different B variables are placed on parallel (“||”) and they constitutes the body of the clause **any...where...** if this clause is created.

Derivation 18 (The expression if ... then ... else ... endif)

Given if cond then expr1 else expr2 endif an expression of postconditions of an operation oper. The B formalisation of this expression gives rise to a clause :

- **if cond then subst(expr1) else subst(expr2) end**, if expr2 is not an expression if ... then ... else ... endif or
- **if cond then subst(expr1) elsif cond2 then subst(expr2) ... end**, if expr2 is also in form if cond2 then expr2 ... ;

where *cond* is the B formalisation of cond and *subst(expr)* denotes the substitutions derived from expr.

Remark 10 *skip* is the substitution of true.

Derivation 19 (Operation) Given {pre,post} the preconditions and the postconditions in OCL of an operation oper, where post is defined according to Definition 3. The B operation *oper* modelling oper is generated by the following manner :

- the signature and a part of precondition of *oper* for typing eventual parameters are generated according to derivation schemes described in [14, 11] ;
- if pre is not true then the precondition in *oper* is augmented by predicates derived from pre in using derivation schemes in Section 3 ;
- the substitution part of *oper* is generated from post using derivation schemes previously presented in this section.

In the context of postconditions Remark 6 may be applied. However, there is also another alternative for let expressions as described in Derivation 20.

Derivation 20 (let ... in postconditions) The expression let $v1 : type1 = val1, \dots, vn : typen = valn$ in expr in postconditions of an operation oper gives rise to a clause *Let... in oper* :

```

let  $v1, \dots, vn$  be
   $v1 = val1 \wedge$ 
  ..  $\wedge$ 
   $vn = valn$ 
in
   $subst(expr)$ 
end

```

where $v1, \dots, vn, val1, \dots, valn$ and *subst(expr)* denote the B formalisation of $v1, \dots, vn, val1, \dots, valn$ and expr.

5 Conclusion

This paper presents a systematic way for transforming OCL expressions into B, which can be applied to generate B supplementary invariants as well as B abstract operations in B specifications, which are generated according to derivation procedures in our previous works [8, 11, 12], from the OCL specifications for the supplementary class invariants and for behavioural concepts in UML specifications.

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