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Isomorphisms of the De Bruijn Digraph and Free-Space Optical Networks

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Abstract

The de Bruijn digraph $B(d, D)$ has degree d , diameter D , d^D vertices and d^{D+1} arcs. It is usually defined by words of size D on an alphabet of cardinality d , through a cyclic left shift permutation on the words, after which the rightmost symbol is changed. In this paper, we show that any digraph defined on words of a given size, through an *arbitrary* permutation on the alphabet **and** an *arbitrary* permutation on the word indices, is isomorphic to the de Bruijn digraph, provided that this latter permutation is *cyclic*. We use this result to improve from $O(d^{D+1})$ to $\Theta(\sqrt{d^{D+1}})$ the number of lenses required for the implementation of $B(d, D)$ by the Optical Transpose Interconnection System proposed by Marsden *et al.* (Optics Letters 18(13):1083-1085, July 1993).

Keywords: Distributed systems, Optical networks, Free space of optical interconnections, Optical Transpose Interconnection System (OTIS), de Bruijn digraph, Graph isomorphisms.

1 Introduction

Many results describing technological and theoretical advances in optical network topologies and design exist in the literature. In this paper, we give optimal optical layouts of several de Bruijn-like topologies, using existing optical technologies, namely the Optical Transpose Interconnecting System (OTIS) architecture proposed by Marsden *et al.* in [17]. These results are obtained using graph isomorphisms developed here.

Our motivation stems from the fact that it has been shown that the break-even line length is less than 1cm, when optical communication lines become more effective than their electrical counterparts, in terms of speed and power consumption [13]. Therefore, the use of optical interconnections between processors on the same board is justified, and some studies even suggest that on-chip optical interconnects will soon be cost effective [22]. Moreover, the emergence of cutting-edge technologies such as Vertical Cavity Surface-Emitting Lasers (VCSELs), high sensibility optical transimpedance receivers, beam splitters, micro-lenses and holograms, makes possible the fabrication of complex optical communication networks [18, 22].

The OTIS architecture is a simple means of implementing very dense one-to-one interconnections between processors in a free space of optical interconnections [17]. In a nutshell, it consists of two lenslet arrays allowing a large number of optical interconnections from a set of transmitters to a set of receivers as shown in Figure 1. The OTIS architecture is used by the 3-D OptoElectronic Stacked Processors consortium to connect a 16-by-16 switch (for more details, see <http://soliton.ucsd.edu/3doesp/public/>).

It has been shown [23] that the OTIS architecture can be used to build all-optical complete

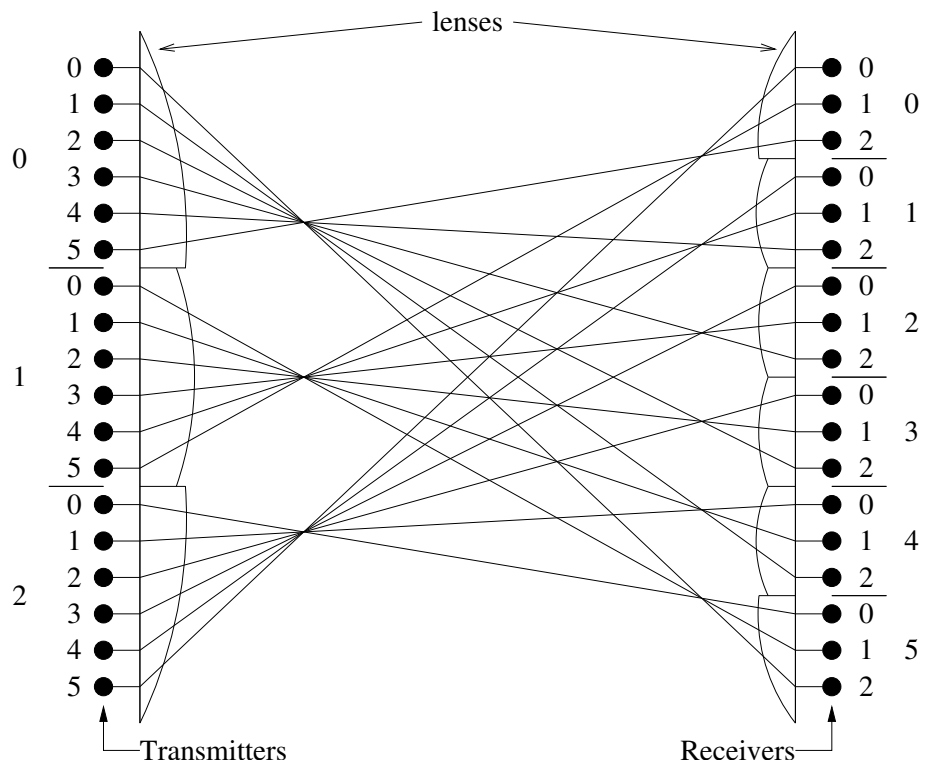


Figure 1: $OTIS(3,6)$.

networks based on a symmetric complete digraph with loops (K_n^+). The OTIS architecture and recent progress made in optoelectronics allow the implementation of K_{64}^+ by OTIS, in which each processor has 64 transceivers (one per neighbor in K_{64}^+). It has also been shown [11] how to implement the single-hop multi-OPS POPS network [7], and the multi-hop multi-OPS stack-Kautz network [10], with the OTIS architecture. However, these layouts do not scale well, since they require a large number of transceivers per processor, as well as a large number of optical lenses. Indeed, the efficiency of an OTIS layout is often measured by the number of lenses required. Consequently, an important research direction is to study the set of network topologies for which there are lens-efficient layouts with the OTIS architecture.

In our work, we focus on de Bruijn-like digraphs, since the de Bruijn digraph [1] is well known for its applications in parallel and distributed computing [6] and it has been extensively studied in graph theory [12]. It has also been used to build various multi-stage interconnection networks [18], and, more interestingly, the Galileo space probe of NASA uses a de Bruijn network with 8192 vertices to implement a signal decoder [8]. A de Bruijn digraph $B(d, D)$ is an alphabet digraph [1] of constant degree d and diameter D . It has for vertices the d^D words of length D on the alphabet $Z_d = \{0, 1, \dots, d-1\}$, and a vertex $x = x_{D-1} \dots x_1 x_0$, $x_i \in Z_d$, is adjacent to the set of vertices $\Gamma^+(x) = \{x_{D-2} \dots x_1 x_0 \beta, \beta \in Z_d\}$.

From a theoretical point of view, we show in this paper that there are many ways to build digraphs which are isomorphic to the de Bruijn digraph. We first show, in Proposition 3.2, that the alphabet digraph $B_\pi(d, D)$, built by using a permutation π on Z_d and the adjacency relation $\Gamma_{B_\pi(d, D)}^+(x) = \{\pi(x_{D-2}) \dots \pi(x_1) \pi(x_0) \pi(\beta), \beta \in Z_d\}$, is isomorphic to $B(d, D)$. This result yields a new and simple proof of the isomorphism between the de Bruijn digraph and

the Imase-Itoh digraph [14]. Then, we show, in Theorem 3.8, that the alphabet digraph $A(f, \pi, j)$, built by using a permutation f on Z_D , a permutation π on Z_d and the adjacency relation $\Gamma_{A(f, \pi, j)}^+(x) = \{\pi(x_{f(D-1)}) \dots \pi(x_{f(j+1)})\pi(\beta)\pi(x_{f(j-1)}) \dots \pi(x_{f(1)})\pi(x_{f(0)})\}$, $\beta \in Z_d$, is isomorphic to $B(d, D)$ if and only if f is cyclic. Notice that the permutation π is applied on the alphabet and f is applied on the word indices (see Section 3.2). Finally, we show that the permutations π and f may be applied together to characterize a class of digraphs isomorphic to the de Bruijn digraph.

It has been shown in [11] that the OTIS architecture can implement the Imase and Itoh digraph $II(d, n)$ with n vertices of degree d . Our result on the isomorphism of $II(d, n)$ and $B(d, D)$ implies that OTIS can also implement the de Bruijn. Unfortunately, for a fixed degree d , such a layout uses $O(dn)$ lenses to interconnect n processors, and it is not very efficient. Our subsequent results help us to show how OTIS can implement the de Bruijn digraph with only $\Theta(\sqrt{dn})$ lenses, which is not only optimal, but also very efficient for practical values, as the number of lenses determines also their size.

This paper is organized as follows. In Section 2, we give the graph theoretical machinery needed in the remainder. Then, in Section 3, we show that there are many isomorphisms of the de Bruijn digraph, using first a permutation π on the alphabet Z_d and then a permutation f on the indices Z_D . In fact, there are exactly $d!(D-1)!$ such isomorphic digraphs. The application to optical networks is given in Section 4, where we present OTIS layouts of de Bruijn-like digraphs and we show how to obtain such layouts where the minimum number of lenses is used. We close the paper with some concluding remarks and directions for further research.

2 Preliminaries

2.1 Notation

Let $Z_n = \{0, 1, 2, \dots, n-1\}$ denote the ring of integers modulo n , Z_n^k the k dimensional vector space over Z_n and e_0, e_1, \dots, e_{k-1} the canonical basis. Addition of elements of Z_n will always be performed in Z_n (i.e., modulo n).

For any $\alpha \in Z_n$, let $\alpha.e_j$ denote the vector $(\overbrace{0, \dots, 0}^j, \alpha, 0, \dots, 0)$. Furthermore, let $Z_n.e_j$ denote the set of vectors $\{\alpha.e_j, \alpha \in Z_n\}$. We define the sum of two sets of vectors as $A + B = \{a + b, a \in A, b \in B\}$. Note that $Z_a.e_i + Z_a.e_j = \{\alpha.e_i + \beta.e_j, \alpha, \beta \in Z_a\} \neq Z_a.(e_i + e_j) = \{\alpha.(e_i + e_j), \alpha \in Z_a\}$.

For a permutation f and an integer i , we define f^i inductively by taking f^0 as the identity permutation and $f^{i+1} = f \circ f^i$.

The next permutation will be used frequently.

Definition 2.1 *Let \mathcal{C} denote the complement permutation in Z_n defined by $\mathcal{C}(u) = n - u - 1$. $\mathcal{C}(u)$ will also be denoted \bar{u} .*

2.2 Classical Digraphs

Given a digraph $G = (V, A)$, with vertex set V and arc set A , and a vertex $x \in V$, we will denote by $\Gamma_G^+(x)$ the set of the out-neighbors of x , i.e. $\{y \in V, (x, y) \in A\}$. When the digraph G is clear from the context we simply denote it by $\Gamma^+(x)$.

Definition 2.2 ([1]) *The de Bruijn digraph $B(d, D)$ of degree d and diameter D is defined as a digraph on the set of vertices made of words of length D on Z_d , and such that for*

$$x = x_{D-1}x_{D-2}\cdots x_1x_0, \Gamma^+(x) = \{x_{D-2}\cdots x_1x_0\beta, \beta \in Z_d\}.$$

Definition 2.3 *The conjunction $G_1 \otimes G_2$ of two digraphs $G_1 = (V_1, A_1)$ and $G_2 = (V_2, A_2)$ is the digraph with vertex set $V_1 \times V_2$, where there is an arc from (u_1, u_2) to (v_1, v_2) if and only if there is both an arc from u_1 to v_1 in G_1 and an arc from u_2 to v_2 in G_2 .*

Alternatively we have $\Gamma_{G_1 \otimes G_2}^+((u_1, u_2)) = \Gamma_{G_1}^+(u_1) \times \Gamma_{G_2}^+(u_2)$. Also, note that $B(d, k) \otimes B(d', k) = B(dd', k)$; one can see [3, 21] for more details about conjunction, de Bruijn digraphs and line digraphs.

Definition 2.4 ([19, 20]) *The Reddy-Pradhan-Kuhl digraph $RPK(d, n)$ of degree d with n vertices, is the digraph with vertex set Z_n and such that for $u \in Z_n$:*

$$\Gamma^+(u) = \{du + \alpha \pmod{n}, 0 \leq \alpha < d\}.$$

The digraphs $RPK(d, d^D)$ and $B(d, D)$ are isomorphic (see [19]). To see this, we associate with vertex $x = x_{D-1}x_{D-2}\cdots x_1x_0$ of $B(d, D)$ a number $u \in Z_n$ with $u = \sum_{i=0}^{D-1} x_i d^i$; then the out-neighbors of x in $B(d, D)$ are associated with numbers of the form $du + \alpha, 0 \leq \alpha < d$. Note also that this congruence definition is sometimes considered as the standard one for the de Bruijn digraph.

Definition 2.5 ([16]) *The Kautz digraph $K(d, D)$ of degree d and diameter D is defined as a digraph on the set of vertices made of words of length D on the alphabet Z_{d+1} , such that $x = x_{D-1}\cdots x_1x_0$, with $x_i \in Z_{d+1}$ and $x_i \neq x_{i+1}$. $\Gamma^+(x) = \{x_{D-2}\cdots x_1x_0\beta, \beta \neq x_0, \beta \in Z_{d+1}\}$.*

Definition 2.6 ([14]) *The Imase and Itoh digraph $II(d, n)$ of degree d with n vertices is the digraph with vertex set Z_n and such that for $u \in Z_n$: $\Gamma^+(u) = \{-du - \alpha \pmod{n}, 1 \leq \alpha \leq d\}$.*

It has been shown [15] that $II(d, d^{D-1}(d+1))$ is isomorphic to $K(d, D)$.

3 $B(d, D)$ as a digraph on an alphabet

In this section we define generalizations of the de Bruijn digraph. In order to obtain the out-neighbors of a vertex labeled by some word of length D over Z_d we proceed as follows:

- instead of performing a left-shift on the word letters we permute the letter positions in some arbitrary way;
- then each letter is replaced by another one according to some fixed permutation on the alphabet Z_d ;
- finally the letter of a given fixed position is replaced by an arbitrary letter of Z_d .

3.1 Permutation on the alphabet

We start by considering only the case of a permutation on the alphabet. We will show that there is an isomorphism between $B(d, D)$ and digraphs that we call $B_\pi(d, D)$, i.e., which are similar to $B(d, D)$ up to a permutation on the alphabet Z_d . Then we will prove that $II(d, d^D)$ and $B(d, D)$ are isomorphic, which, according to previous remarks will imply that $II(d, d^D), B(d, D), RPK(d, d^D)$ are isomorphic.

Definition 3.1 Given a permutation π on Z_d , $B_\pi(d, D)$ is defined as the digraph on the set of vertices made of words of length D on Z_d , such that for each vertex $x = x_{D-1}x_{D-2} \cdots x_1x_0$, $x_i \in Z_d$, $\Gamma^+(x) = \{\pi(x_{D-2}) \cdots \pi(x_1)\pi(x_0)\beta, \beta \in Z_d\}$.

Proposition 3.2 $B(d, D)$ and $B_\pi(d, D)$ are isomorphic.

Proof: We just have to define an isomorphism W from $B_\pi(d, D)$ to $B(d, D)$; W is the following one-to-one mapping from the words of length D on the alphabet Z_d onto itself:

$$W(x_{D-1}x_{D-2} \cdots x_1x_0) = \pi^0(x_{D-1})\pi^1(x_{D-2}) \cdots \pi^{D-2}(x_1)\pi^{D-1}(x_0).$$

Given a vertex x of $B_\pi(d, D)$ we have

$$\begin{aligned} W^{-1}\Gamma_{B(d,D)}^+(W(x_{D-1}x_{D-2} \cdots x_0)) &= W^{-1}\Gamma_{B(d,D)}^+(\pi^0(x_{D-1})\pi^1(x_{D-2}) \cdots \pi^{D-1}(x_0)) \\ &= W^{-1}(\pi^1(x_{D-2}) \cdots \pi^{D-2}(x_1)\pi^{D-1}(x_0)\beta), \beta \in Z_d \\ &= \pi(x_{D-2}) \cdots \pi(x_0)\pi(\beta) \\ &= \pi(x_{D-2}) \cdots \pi(x_0)\alpha, \alpha \in Z_d. \end{aligned}$$

As the last line corresponds to the adjacency relation in $B_\pi(d, D)$ we have defined the required isomorphism. □

Notice that more generally, the digraph of degree d , diameter D and $n = d^D$ vertices such that $\Gamma^+(x_{D-1}x_{D-2} \cdots x_1x_0) = \{\pi_0(x_{D-2})\pi_1(x_{D-3}) \cdots \pi_{D-2}(x_0)\pi_{D-1}(\alpha), \alpha \in Z_d\}$, where $\pi_i, i \in Z_D$, is a permutation on Z_d , is isomorphic to $B(d, D)$.

Proposition 3.3 The digraphs $B(d, D)$ and $II(d, d^D)$ are isomorphic.

Proof: Let $\overline{B}(d, D)$ denote $B_{\mathcal{C}}(d, D)$, where \mathcal{C} is the complement permutation over Z_d . As in the case of $B(d, D)$, $\overline{B}(d, D)$ can be alternatively defined as a digraph with vertex set Z_n

in which for $u \in Z_n$, $\Gamma^+(u) = \{-(du + \beta) - 1, 0 \leq \beta < d\}$, that is $\Gamma^+(u) = \{-du - \alpha, 1 \leq \alpha \leq d\}$. This is the definition of $II(d, d^D)$. So $B_C(d, D)$ and $II(d, d^D)$ are isomorphic. From Proposition 3.2 these digraphs are hence isomorphic to $B(d, D)$. \square

We thus obtain

Corollary 3.4 *The digraphs $B(d, D)$, $RPK(d, d^D)$ and $II(d, d^D)$ are isomorphic.*

To illustrate, Figure 2 shows the digraph $B(2, 3)$ which is isomorphic to $RPK(2, 8)$ of Figure 3 and to $II(2, 8)$ represented in Figure 4.

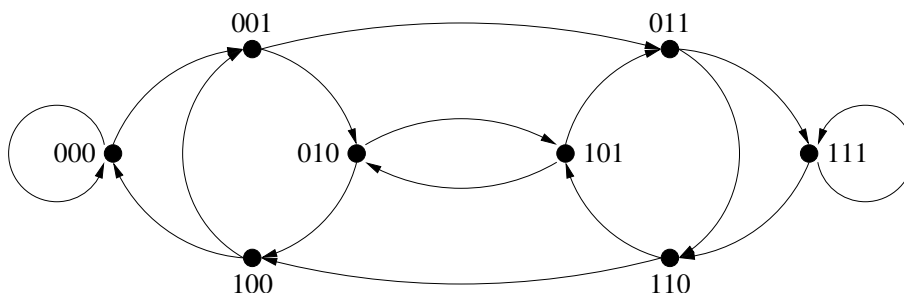


Figure 2: de Bruijn digraph of degree 2, diameter 3 and 8 vertices.

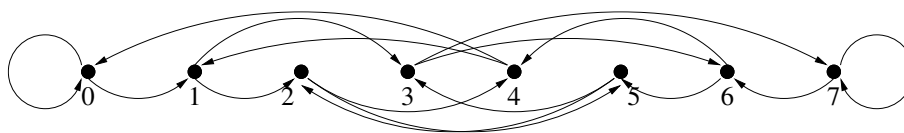


Figure 3: RPK digraph of degree 2, diameter 3 and 8 vertices.

3.2 Permutation on the indices, digraphs $A(f, \pi, j)$

Now we consider more general digraphs than $B_\pi(d, D)$ for which the left shift is replaced by any permutation on Z_D . First we need some definitions.

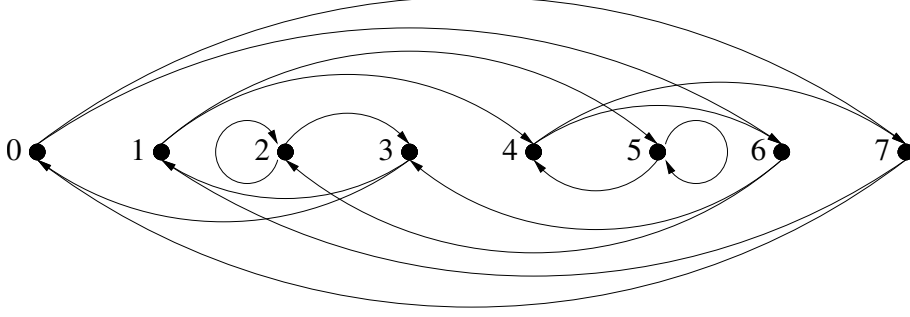


Figure 4: II digraph of degree 2, diameter 3 and 8 vertices.

Definition 3.5 To a permutation f on Z_D , we associate the linear map \vec{f} of Z_d^D defined on the basis by $\vec{f}(e_i) = e_{f(i)}$ (this linear map is associated with a permutation matrix). Note that $\vec{f}g = \vec{f}\vec{g}$.

Definition 3.6 Any permutation π on Z_d is extended naturally to Z_d^D by setting for $x = x_{D-1}x_{D-2}\cdots x_1x_0$: $\pi(x) = \pi(x_{D-1})\pi(x_{D-2})\cdots\pi(x_1)\pi(x_0)$.

Definition 3.7 Let $d, D \in \mathbb{N}$. Given a permutation f on Z_D , a permutation π on Z_d and a position $j \in Z_D$, we define the digraph $G = A(f, \pi, j)$ by

$$\begin{aligned} V(G) &= Z_d^D \\ \Gamma_G^+(x) &= \pi(f(x)) + Z_d \cdot e_j \end{aligned}$$

$A(f, \pi, j)$ has degree d and we will say that it has dimension D .

Note that vertices of $B(d, D)$ can be considered as elements of Z_d^D (by associating $x_{D-1}\cdots x_1x_0$ with $\sum x_i \cdot e_i$). Let σ be the permutation on Z_D defined by $\sigma : i \mapsto i + 1 \pmod{D}$, and let Id be the identity permutation on Z_d . Then the de Bruijn digraph is exactly $A(\sigma, Id, 0)$. Indeed for $x = x_{D-1}x_{D-2}\cdots x_1x_0$, $\Gamma_{A(\sigma, Id, 0)}^+(x) = \{x_{D-2}\cdots x_1x_0(x_{D-1} + \gamma), \gamma \in Z_d\} = \{x_{D-2}\cdots x_1x_0\beta, \beta \in Z_d\}$. Similarly, $B_\pi(d, D)$ and $A(\sigma, \pi, 0)$ are isomorphic.

Theorem 3.8 *The digraph $G = A(f, \pi, j)$ with degree d and dimension D is isomorphic to $B(d, D)$ if and only if f is a cyclic permutation on Z_D .*

Proof: We will prove the theorem in two steps. First, we show that when f is a cyclic permutation on Z_D , then the digraphs $G = A(f, \pi, j)$ and $B(d, D)$ are isomorphic. Then, we show, by contradiction, that if the digraphs $G = A(f, \pi, j)$ and $B(d, D)$ are isomorphic then f is a cyclic permutation on Z_D .

Case 1: Suppose that f is a cyclic permutation on Z_D . We show that the digraphs $A(f, \pi, j)$ and $B(d, D)$ are isomorphic.

Note that since f is a cyclic permutation it induces a unique orbit on Z_D . We associate with f the unique permutation g on Z_D , such that $\forall i \in Z_D, g(i) = f^i(j)$. From this definition, we have

$$g^{-1}fg(i) = g^{-1}ff^i(j) = g^{-1}f^{i+1}(j) = g^{-1}g(i+1) = i+1 \quad (1)$$

$$g^{-1}(j) = 0 \quad (\text{hence } \overrightarrow{g^{-1}}(e_j) = e_0) \quad (2)$$

Let us consider the digraph isomorphism induced by \overrightarrow{g} . For $x \in Z_d^D$, we compute $\overrightarrow{g^{-1}}\Gamma_G^+(\overrightarrow{g}(x))$. From the definition of G , $\Gamma_G^+(x) = \pi(\vec{f}(x)) + Z_d \cdot e_j$, thus by linearity $\overrightarrow{g^{-1}}\Gamma_G^+(\overrightarrow{g}(x)) = \pi\left(\overrightarrow{g^{-1}}\left(\vec{f}\overrightarrow{g}(x)\right)\right) + Z_d \cdot \overrightarrow{g^{-1}}(e_j)$. But from (1), $\overrightarrow{g^{-1}}fg = \vec{\sigma}$ and from (2), $\overrightarrow{g^{-1}}(e_j) = e_0$. Finally, $\overrightarrow{g^{-1}}\Gamma_G^+(\overrightarrow{g}(x)) = \pi(\vec{\sigma}(x)) + Z_d \cdot e_0$, and this adjacency relation is the one of $B_\pi(d, D)$ which is isomorphic to $B(d, D)$.

Case 2: Suppose now that the digraphs $A(f, \pi, j)$ and $B(d, D)$ are isomorphic and let us show that f is a cyclic permutation on Z_D .

Suppose, by contradiction, that f is not a cyclic permutation on Z_D . Then it induces

at least two orbits on Z_D . Let $orb(j)$ be the set of indices sharing the orbit of j , $orb(j) = \{i \in Z_D, \exists k < D, f^k(j) = i\}$. Let $f = f_{orb(j)} \circ f_{r(j)}$, where $f_{orb(j)}$ is a permutation on Z_D such that $f_{orb(j)}$ is a cyclic permutation on $orb(j)$ and is equal to the identity permutation on $r(j) = Z_D \setminus orb(j)$. Let also $f_{r(j)}$ be a permutation on Z_D such that $f_{r(j)}$ is a permutation on $r(j)$ and is equal to the identity permutation on $orb(j)$.

Then, $f^{|orb(j)|}(j) = (f_{orb(j)} \circ f_{r(j)})^{|orb(j)|}(j) = f_{orb(j)}^{|orb(j)|} \circ f_{r(j)}^{|orb(j)|}(j) = f_{orb(j)}^{|orb(j)|}(j) = j$.

Now, we shall consider separately the case where $|orb(j)| \leq D - 2$, and the case where $|orb(j)| = D - 1$.

A. Let $|orb(j)| \leq D - 2$. We show that the digraph $A(f, \pi, j)$ is not connected.

Let x be a vertex of $A(f, \pi, j)$, $x = x_{D-1}x_{D-2} \cdots x_1x_0 = \sum_{i \in orb(j)} x_i \cdot e_i + \sum_{k \in r(j)} x_k \cdot e_k$,

and let $X = \left\{ \sum_{i \in orb(j)} Z_d \cdot e_i + \sum_{k \in r(j)} y \cdot e_k, y \in Z_d \right\} = \left\{ \sum_{i \in orb(j)} Z_d \cdot e_i + Z_d \cdot \sum_{k \in r(j)} e_k \right\}$.

Note that $|X| = d^{|orb(j)|}$, and since $|orb(j)| \leq D - 2$, we have $X \subsetneq Z_d^D$.

Computing the neighborhood of X yields

$$\begin{aligned}
\Gamma_G^+(X) &= \pi(\vec{f}(X)) + Z_d \cdot e_j \\
&= \pi \left(\vec{f} \left(\sum_{i \in orb(j)} Z_d \cdot e_i + Z_d \cdot \sum_{k \in r(j)} e_k \right) \right) + Z_d \cdot e_j \\
&= \sum_{i \in orb(j)} \pi(Z_d) \cdot \vec{f}_{orb(j)}(e_i) + \pi(Z_d) \cdot \sum_{k \in r(j)} \vec{f}_{r(j)}(e_k) + Z_d \cdot e_j \\
&= \sum_{i \in orb(j)} \pi(Z_d) \cdot e_i + Z_d \cdot e_j + Z_d \cdot \sum_{k \in r(j)} e_k \\
&= \sum_{i \in orb(j)} Z_d \cdot e_i + Z_d \cdot \sum_{k \in r(j)} e_k \\
&= X
\end{aligned}$$

Hence, X and $Z_d^D \setminus X$ are not connected, which implies that $A(f, \pi, j)$ is not connected

when $|orb(j)| \leq D - 2$, and thus that $A(f, \pi, j)$ cannot be isomorphic to $B(d, D)$.

B. Let $|orb(j)| = D - 1$. We show that $A(f, \pi, j)$ cannot be isomorphic to $B(d, D)$ if π is a cyclic permutation on Z_d , and that, otherwise, $A(f, \pi, j)$ is not connected.

Let $Y_l = \left\{ \sum_{i \in orb(j)} Z_d \cdot e_i + l \cdot e_k, k = Z_d \setminus orb(j) \right\}$, $l \in Z_d$. We have $\Gamma_G^+(Y_l) = Y_{\pi(l)}$.

(i) When π is a cyclic permutation on Z_d , it follows that $\pi(l) \neq l$, and thus $\forall l \in Z_d$, $Y_l \cap Y_{\pi(l)} = \emptyset$. Therefore, $A(f, \pi, j)$ has no loops, implying that it cannot be isomorphic to $B(d, D)$.

(ii) When π is not a cyclic permutation on Z_d , we have $\cup_{1 \leq s \leq d} \pi^s(l) \subsetneq Z_d$. Hence, $\cup_{1 \leq s \leq d} \Gamma_G^{+s}(Y_l) = \cup_{1 \leq s \leq d} Y_{\pi^s(l)} \subsetneq Z_d^D$ and $A(f, \pi, j)$ is not connected.

□

Notice that since there are $(D - 1)!$ cyclic permutation functions f on Z_D , Theorem 3.8 leads to $(D - 1)!$ alternative ways of defining $B(d, D)$. Furthermore, as Proposition 3.2 leads to $d!$ other ways of defining $B(d, D)$, using permutations π on Z_d , we finally obtain $d!(D - 1)!$ alternative ways of defining $B(d, D)$.

If f is not a cyclic permutation on Z_D , it can be shown that each connected component of $A(f, \pi, j)$ is the conjunction of a de Bruijn digraph with a circuit. Furthermore, in case $A(f, \pi, j)$ is connected, and $d = D - 1$, then the digraph $A(f, \pi, j)$ is isomorphic to the wrapped Butterfly digraph $WBF(d, D - 1)$ of degree d and $D - 1$ levels with d^{D-1} vertices each. See [2, 9] for more details.

3.3 Examples

The following examples illustrate Theorem 3.8. First, we show an alphabet digraph $G = A(f, Id, 2)$ which is isomorphic to $B(d, 6)$. Then we show an alphabet digraph $H = A(f, Id, 1)$ where the permutation f is not cyclic and the orbit of 1 has size $D - 2$, and hence H is not connected.

3.3.1 $G = A(f, Id, 2)$

Let G be the digraph with degree d , dimension 6 and vertex set Z_d^6 , defined below:

$$\Gamma_G^+(x_5x_4x_3x_2x_1x_0) = x_2x_1x_0\beta x_5x_4, \text{ with } \beta \in Z_d.$$

G is indeed $A(f, Id, 2)$, where f is the following permutation on Z_6 :

$$f(i) = \begin{cases} i + 3 & \text{if } i < 3; \\ 2 & \text{if } i = 3; \\ i + 2 \pmod{6} & \text{otherwise.} \end{cases}$$

Then, \vec{f} is defined on Z_d^6 by

$$\vec{f}(x_5x_4x_3x_2x_1x_0) = x_2x_1x_0x_3x_5x_4$$

The permutation g associated with f , used in Theorem 3.8 (i.e., $g(i) = f^i(2)$), is such that $g(0) = 2$, $g(1) = 5$, $g(2) = 1$, $g(3) = 4$, $g(4) = 0$, $g(5) = 3$, and $\vec{g}(x_5x_4x_3x_2x_1x_0) = x_1x_3x_5x_0x_2x_4$, and hence $\vec{g}^{-1}(x_5x_4x_3x_2x_1x_0) = x_3x_0x_4x_1x_5x_2$. See Figure 5 for an illustration.

$$\begin{array}{ccccc}
g(0) = 2 & \xrightarrow{f} & g(1) = 5 & \xrightarrow{f} & g(2) = 1 \\
\uparrow f & & & & \downarrow f \\
g(5) = 3 & \xleftarrow{f} & g(4) = 0 & \xleftarrow{f} & g(3) = 4
\end{array}$$

Figure 5: Illustration of the permutation g defined as $g(i) = f^i(2)$.

Then

$$\begin{aligned}
\overrightarrow{g^{-1}}\Gamma_G^+\overrightarrow{g}(x_5x_4x_3x_2x_1x_0) &= \overrightarrow{g^{-1}}\Gamma_G^+(x_1x_3x_5x_0x_2x_4) \\
&= \overrightarrow{g^{-1}}(x_0x_2x_4\beta x_1x_3), \beta \in Z_d \\
&= x_4x_3x_2x_1x_0\beta \\
&= \Gamma_{B(d,6)}^+(x_5x_4x_3x_2x_1x_0).
\end{aligned}$$

Therefore, G is isomorphic to $B(d, 6)$.

3.3.2 $H = A(f, Id, 1)$

Let H be the digraph with degree d , dimension 3, and vertex set Z_d^3 defined below:

$$\Gamma_H^+(x_2x_1x_0) = x_0\beta x_2, \text{ with } \beta \in Z_d.$$

H is indeed $A(f, Id, 1)$, where f is the permutation on Z_3 such that $f(i) = \bar{i} = 2 - i$.

Then, \overrightarrow{f} is defined on Z_d^3 by $\overrightarrow{f}(x_2x_1x_0) = x_0x_1x_2$.

If we define the function g associated with f by $g(i) = f^i(1)$, we get $g(0) = g(1) = g(2) = 1$; hence g is not a cyclic permutation on Z_3 . Furthermore, the orbit induced by $j = 1$ is $orb(j) = \{1\}$ with $|orb(j)| = 1 \leq D - 2 = 1$ and thus H is not connected. Notice that the digraph H can be alternatively defined by

$$\Gamma_H^+([x_2x_0], [x_1]) = \{[x_0x_2], [\beta]\}$$

It is composed of $\frac{d^2-d}{2}$ digraphs $C_2 \otimes B(d, 1)$ plus d digraphs $C_1 \otimes B(d, 1)$, as shown in Figure 6 for $d = 2$.

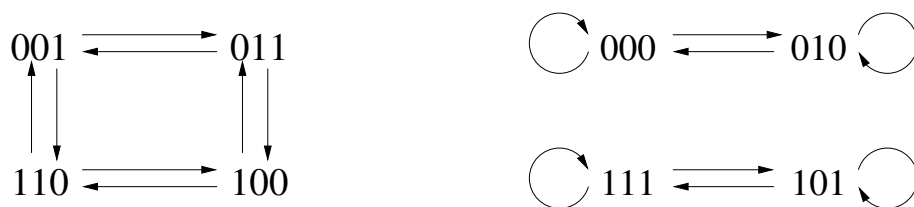


Figure 6: The digraph H of Example 3.3.2, when $d = 2$.

4 OTIS layout for digraphs

In this section, we first review the definition of the OTIS architecture, and the known digraphs having a layout with OTIS. Then, we study the particular case of the OTIS layout of de Bruijn digraphs, and, using the theoretical results developed in the previous section, we show how to obtain optimal OTIS layouts for the de Bruijn digraphs.

4.1 The OTIS architecture

$OTIS(p, q)$ is a free-space optical system which allows one-to-one communications from p groups of q transmitters to q groups of p receivers, using $p + q$ lenses. This architecture connects the transmitter (i, j) , $0 \leq i \leq p-1$, $0 \leq j \leq q-1$, to the receiver $(q-j-1, p-i-1)$. Optical interconnections in the OTIS architecture are realized with a pair of lenslet arrays [4] in a free space of optical interconnections as shown in Figure 1.

4.2 Known OTIS layouts

Let $m = pq$. The $OTIS(p, q)$ architecture connects m transmitter units of the processors to m receiver units of the same set of processors. The interconnection network can be defined as follows. Given a degree d such that d divides m , let $H(p, q, d)$ be the d -regular digraph with $n = \frac{m}{d}$ vertices and vertex set Z_n , in which the d transmitters ($\lfloor (du + \alpha)/q \rfloor, du + \alpha \pmod{q}$), $\alpha \in Z_d$, and the d receivers ($\lfloor (du + \alpha)/p \rfloor, du + \alpha \pmod{p}$), $\alpha \in Z_d$, belong to the vertex $u \in Z_n$ (i.e., a processing unit). There is a connection from vertex u to vertex v in $H(p, q, d)$ if one transmitter of u is connected to one receiver of v . For instance, Figure 7 represents the digraph $H(4, 8, 2)$. A formal definition of $H(p, q, d)$ is as follows.

Definition 4.1 *The digraph $H(p, q, d)$ is defined from the $OTIS(p, q)$ architecture. It has degree d , such that d divides pq , and $n = \frac{pq}{d}$ vertices and $m = pq$ arcs. It has vertex set $V(H) = Z_n$, and its neighborhood is defined as follows:*

$$\Gamma_H^+(u) = \left\{ v = \left\lfloor \frac{(pq - 1) \left(\left\lfloor \frac{du + \alpha}{q} \right\rfloor + 1 \right) - p(du + \alpha)}{d} \right\rfloor, \quad 0 \leq \alpha \leq d - 1 \right\}$$

Definition 4.2 *A digraph $G = (V, E)$, with $|V| = n$ vertices of constant degree d and $|E| = m = dn$ arcs, has an $OTIS(p, q)$ -layout if and only if there exist p and q , with $pq = m$, and an isomorphism from $G = (V, E)$ onto $H(p, q, d)$.*

Let G^- be the digraph obtained by reversing all the arcs of a digraph G . We remark that if G admits an $OTIS(p, q)$ -layout then G^- has an $OTIS(q, p)$ -layout.

It has been proved that the Imase and Itoh digraph $II(d, n)$ with n vertices of degree d has an $OTIS(d, n)$ -layout [11], and consequently, so do the de Bruijn and the Kautz digraphs by Proposition 3.3. For fixed degree d , such a layout uses $O(dn)$ lenses to interconnect

n processors, which is not very efficient. In the following we will show how to build an $OTIS(p, q)$ -layout for the de Bruijn digraph, such that the number $p + q$ of lenses used is $\Theta(\sqrt{dn})$, which is also very efficient for practical values. Note that the values of p and q determine the size of the lenses and, by technological considerations, it is preferable to work with lenses of almost equal size (i.e., $p \sim q$).

4.3 The degree-diameter problem for OTIS

For given degree d and diameter D , one can study the maximum number of vertices of the family of connected digraphs $H(p, q, d)$ (i.e., the d -regular digraph built from an $OTIS(p, q)$). We implemented a simple and exhaustive search, based on a computer program which returns the diameter of any input digraph $H(p, q, d)$ (in case H is connected). We executed this program for all pairs (p, q) , such that $\frac{pq}{d} \leq \sum_{i=0}^{D_m} d^i = \frac{d^{D_m+1}-1}{d-1}$, which is the Moore bound for degree d digraphs of maximum diameter D_m [5]. The results obtained are reported in Table 1, for degree 2 and diameters 8, 9, and 10. Note that the table contains only the largest digraphs found for each diameter.

We remark that the Kautz digraph appears to be the largest digraph of degree d and diameter D which has an $OTIS(p, q)$ -layout. This would not be surprising, since it is among the best known constructions for the classic degree-diameter problem [5].

On the other hand, for fixed degree d , diameter D , and number of vertices $n = d^D$, we found several digraphs $H(p, q, d)$, with $pq = d^{D+1}$ for different values of p and q such that both p and q were powers of d . All these digraphs happened to be isomorphic to $B(d, D)$. As an example, $H(2, 256, 2)$, $H(4, 128, 2)$ and $H(16, 32, 2)$ are isomorphic to $B(2, 8)$. In the next

section, we will prove that whenever D is even and $H(p, q, d)$, as defined above, is connected, then it is isomorphic to $B(d, D)$. We will also address the case where D is odd.

n	p	q		n	p	q		n	p	q	
⋮	⋮	⋮		⋮	⋮	⋮		⋮	⋮	⋮	
253	2	253		509	2	509		1022	2	1022	
254	2	254		510	2	510		1023	2	1023	
255	2	255		511	2	511		1024	2	1024	B(2,10)
256	2	256	B(2,8)	512	2	512	B(2,9)		4	512	
	4	128			8	128			8	256	
	16	32		513	2	513			16	128	
									32	64	
258	2	258		516	2	516		1026	2	1026	
264	2	264		528	2	528		1032	2	1032	
288	2	288		576	2	576		1056	2	1056	
384	2	384	K(2,8)	768	2	768	K(2,9)	1152	2	1152	
								1536	2	1536	K(2,10)
$D = 8$				$D = 9$				$D = 10$			

Table 1: $H(p, q, 2)$ with diameters 8, 9 and 10.

4.4 Efficient $OTIS(p, q)$ -layout

In this section we characterize $OTIS(p, q)$ -layouts of de Bruijn digraphs having a minimum number of lenses. For this, we will study cases where $p = d^{p'}$ and $q = d^{q'}$, as indicated in the discussion above.

Proposition 4.3 *The digraph $H(p, q, d)$ is isomorphic to $A(f, \mathcal{C}, p' - 1)$, where \mathcal{C} is the complement permutation defined in Definition 2.1, and where the permutation f is defined on Z_D , $D = p' + q' - 1$, as follows:*

$$f : i \mapsto \begin{cases} i + p' & \text{if } i < q' - 1; \\ p' - 1 & \text{if } i = q' - 1; \\ i + p' - 1 \pmod{D} & \text{otherwise.} \end{cases}$$

Proof: Notice that in this proof, we associate with a number $x \in Z_{d^k}$, its d -ary representation, denoted $\omega(x)$ (i.e., $\omega(x) = x_{k-1} \dots x_1 x_0$, $x_i \in Z_d$, such that $x = \sum_{i=0}^{k-1} x_i d^i$).

The $OTIS(p, q)$ architecture connects p groups of q transmitters to q groups of p receivers, such that the transmitter (i, j) , $0 \leq i < p$, $0 \leq j < q$, is connected to the receiver $(q - j - 1, p - i - 1)$. Note the $\omega(i)$ has length p' and $\omega(j)$ has length q' . We also label transceiver (i, j) by $\omega(i)\omega(j)$, that has length $p' + q'$. As $\omega(q - j - 1) = \overline{\omega(j)}$ and $\omega(p - i - 1) = \overline{\omega(i)}$, we have that the transmitter $\omega(i)\omega(j)$ is connected to the receiver $\overline{\omega(j)\omega(i)}$.

$H(d^{p'}, d^{q'}, d)$ has constant degree d and n vertices. Note that since d divides both p and q , each group of $d^{q'}$ transmitters contains $d^{q'-1}$ vertices. So we can label the k -th vertex ($k \in Z_{d^{q'-1}}$) of the i -th group of transmitters ($i \in Z_p$) by $\omega(i)\omega(k)$, of length $p' + q' - 1$. This vertex has the transmitters $\{(i, dk + \beta), \beta \in Z_d\}$, which are labeled by the set of strings $\{\omega(i)\omega(k)\omega(\beta), \beta \in Z_d\}$. Thus, the transmitters of the vertex $\omega(i)\omega(k)$ are connected to

the receivers $\{\overline{\omega(k)\omega(\beta)\omega(i)}, \beta \in Z_d\}$. Let $\omega(i) = \omega(l)\omega(\alpha)$, where $\omega(l)$ is a string of length $p' - 1$ and $\alpha \in Z_d$. We have $\{\overline{\omega(k)\omega(\beta)\omega(i)}, \beta \in Z_d\} = \{\overline{\omega(k)\omega(\beta)\omega(l)\omega(\alpha)}, \beta \in Z_d\}$ and these receivers belong to the vertices $\{\overline{\omega(k)\omega(\beta)\omega(l)}, \beta \in Z_d\}$.

Finally, a vertex x represented by the string $\omega(x) = \omega(l)\omega(\alpha)\omega(k)$, with $\omega(l) \in Z_d^{p'-1}$, $\omega(\alpha) \in Z_d$ and $\omega(k) \in Z_d^{q'-1}$, is connected through the *OTIS*(p, q) architecture, to the vertices represented by the strings $\{\overline{\omega(k)\omega(\beta)\omega(l)}, \beta \in Z_d\}$. As $\omega(x) = x_{D-1} \dots x_1 x_0$, $x_i \in Z_d$, we have $\omega(l) = x_{D-1} x_{D-2} \dots x_{q'}$, $\omega(\alpha) = x_{q'-1}$, and $\omega(k) = x_{q'-2} \dots x_1 x_0$. Thus, the digraph $H(d^{p'}, d^{q'}, d)$ can be defined on the vertex set Z_d^D with

$$\Gamma_H^+(x_{D-1} \dots x_1 x_0) = \{\overline{x_{q'-2} \dots x_0 \beta x_{D-1} \dots x_{q'}}\}.$$

According to the definition of f this is exactly $A(f, \mathcal{C}, p' - 1)$.

□

The following lemma allows us to characterize all *OTIS*(p, q)-layouts for de Bruijn digraphs $B(d, D)$.

Lemma 4.4 *Let $p' + q' - 1 = D$. For any degree d , the digraphs $B(d, D)$ and $H(d^{p'}, d^{q'}, d)$ are isomorphic if and only if the permutation f of Z_D defined by*

$$f : i \mapsto \begin{cases} i + p' & \text{if } i < q' - 1; \\ p' - 1 & \text{if } i = q' - 1; \\ i + p' - 1 \pmod{D} & \text{otherwise,} \end{cases}$$

is cyclic.

Proof: We showed in Proposition 4.3 that $H(d^{p'}, d^{q'}, d)$ is isomorphic to the alphabet digraph $A(f, \mathcal{C}, p' - 1)$. We have also shown in Theorem 3.8 that the alphabet digraph

$A(f, \mathcal{C}, j)$ of degree d and dimension D is isomorphic to the de Bruijn digraph $B(d, D)$ if and only if f is cyclic. The lemma follows. \square

Now, we show that there exist p' and q' such that $B(d, D)$ and $H(d^{p'}, d^{q'}, d)$ are isomorphic and $d^{p'} + d^{q'} = \Theta(\sqrt{d^{D+1}})$ (i.e., $p + q = \Theta(\sqrt{dn})$, since $dn = d^{D+1}$).

Lemma 4.5 *Any OTIS(p, q)-layout for $B(d, D)$ has $\Omega(\sqrt{dn})$ lenses.*

Proof: Since $pq = d^{D+1}$, we have $p + q \geq 2\sqrt{d^{D+1}} = 2\sqrt{dn} = \Omega(\sqrt{dn})$. \square

Proposition 4.6 *Let $p' = q' = \frac{D+1}{2}$. Then $B(d, D)$ and $H(d^{\frac{D+1}{2}}, d^{\frac{D+1}{2}}, d)$ are isomorphic if and only if $D = 1$.*

Proof: If $p' = q' > 1$, $f(p' - 1) = q' - 1 = p' - 1$ and f is not cyclic. If $p' = q' = 1$, $D = 1$ and $B(d, 1)$ is isomorphic to $H(d, d, d)$. \square

Proposition 4.7 *Let D be even, $p' = \frac{D}{2}$ and $q' = \frac{D}{2} + 1$. Then $B(d, D)$ and $H(d^{\frac{D}{2}}, d^{\frac{D}{2}+1}, d)$ are isomorphic, and therefore $B(d, D)$ has an OTIS(p, q)-layout with $p+q = \Theta(\sqrt{dn})$ lenses.*

Proof: The permutation f on Z_D is defined by

$$f : i \mapsto \begin{cases} i + \frac{D}{2} & \text{if } i < \frac{D}{2}; \\ \frac{D}{2} - 1 & \text{if } i = \frac{D}{2}; \\ i + \frac{D}{2} - 1 \pmod{D} & \text{otherwise,} \end{cases}$$

and one can easily verify that f is cyclic. Thus, using Lemma 4.4, the isomorphism is verified.

Furthermore, the OTIS(p, q)-layout has $d^{\frac{D}{2}} + d^{\frac{D}{2}+1} = O(\sqrt{dn})$ lenses, and thus, by Lemma 4.5, the layout has $\Theta(\sqrt{dn})$ lenses. \square

Example: Figure 7 shows the digraph $H(4, 8, 2)$ which has 16 vertices of degree 2 and 32 arcs. The vertex set of $H(4, 8, 2)$ is Z_2^4 , and $\Gamma_H^+(x_3x_2x_1x_0) = \{\overline{x_1x_0\beta x_3}\}$. We represent each vertex twice (once at the left of the figure and once at the right), and the arcs going from the left to the right. Figure 8 represents the digraph $B(2, 4)$ with the vertices labeling of $H(4, 8, 2)$.

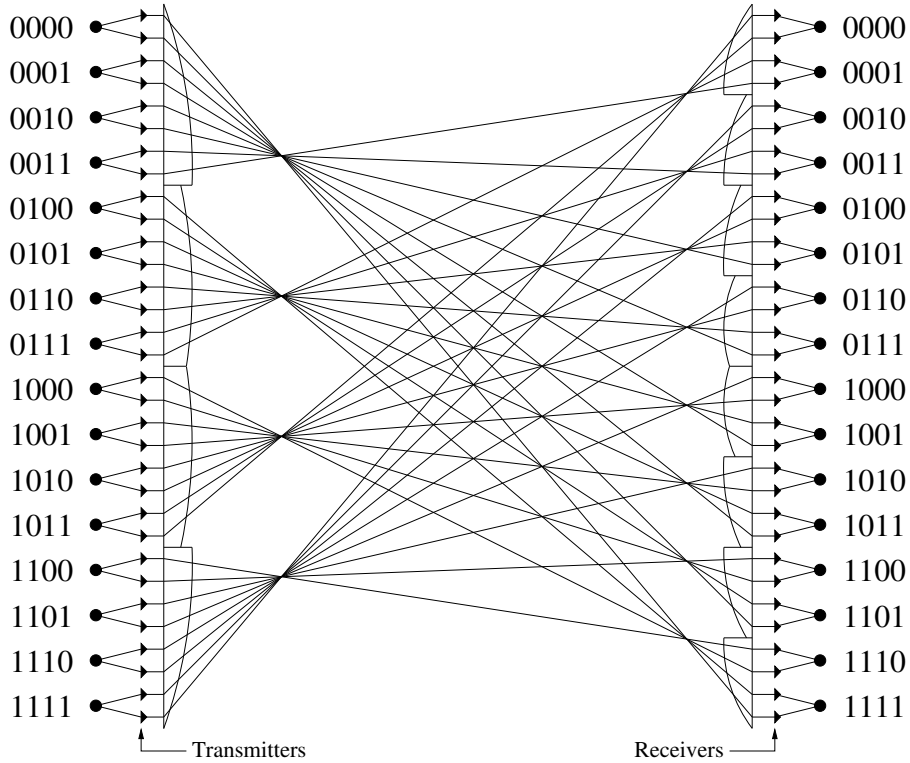


Figure 7: $H(4, 8, 2)$.

When the diameter $D > 1$ is odd, by Proposition 4.6, $B(d, D)$ and $H(p, q, d)$ cannot be isomorphic with $p = q$. The case where p is closest to q is hence when, w.l.o.g., $p' = \frac{D-1}{2}$ and $q' = \frac{D-1}{2} + 2$. However, it is not always the case that $B(d, D)$ and $H(d^{\frac{D-1}{2}}, d^{\frac{D-1}{2}+2}, d)$ are isomorphic. For instance, $H(2^5, 2^7, 2)$ and $B(2, 11)$ are isomorphic, while $H(d^6, d^8, d)$ and $B(d, 13)$ are not.

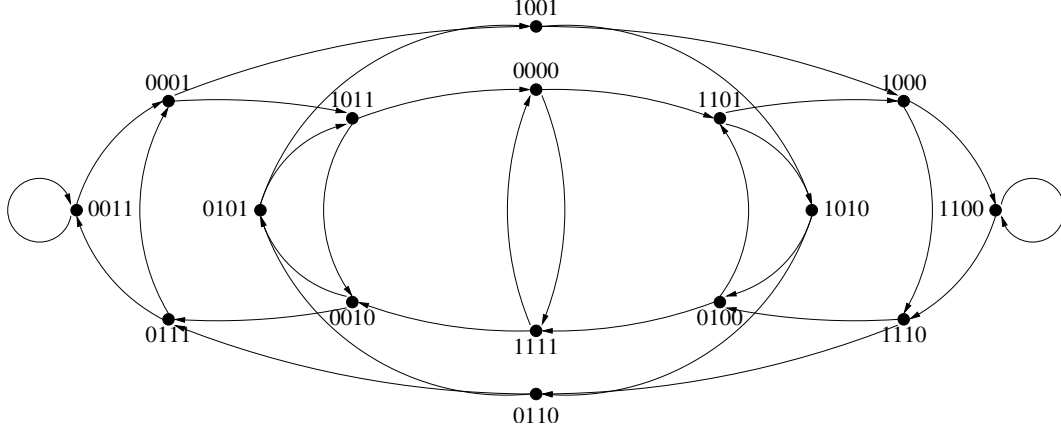


Figure 8: $B(2, 4)$ with the adjacency relation of $H(4, 8, 2)$, i.e.: $\Gamma_H^+(x_3x_2x_1x_0) = \{\overline{x_1x_0\beta x_3}\}$.

Corollary 4.8 *Given $B(d, D)$ and $H(p, q, d)$, their isomorphism can be verified in $O(D)$ time.*

Proof: Using Lemma 4.4 it is sufficient to verify that f is a cyclic permutation on Z_D . This requires $O(D)$ steps. □

Corollary 4.9 *Minimizing $d^{p'} + d^{q'}$, such that $B(d, D)$ and $H(d^{p'}, d^{q'}, d)$ are isomorphic, requires $O(D^2)$ time steps.*

Proof: If D is even, using Proposition 4.7, the minimum is obtained with $p' = \frac{D}{2}$ and $q' = \frac{D}{2} + 1$. Otherwise, there are D pairs (p', q') , with $p' < q'$, such that $D = p' + q' - 1$. Thus, using Corollary 4.8, we obtain the minimum value in $O(D^2)$ time. □

5 Conclusion

In this paper, we have shown that any alphabet digraph defined on words of a given size, through an arbitrary permutation on the alphabet and an arbitrary permutation on the

word indices, is isomorphic to the de Bruijn digraph, provided that this latter permutation is cyclic. In other words, we have shown that the alphabet digraph $A(f, \pi, j)$, built by using a permutation f on Z_D , a permutation π on Z_d , an index $0 \leq j \leq D - 1$, and the adjacency relation $\Gamma_{A(f, \pi, j)}^+(x) = \{y_{D-1} \dots y_{j+1} \beta y_{j-1} \dots y_1 y_0 \mid \beta \in Z_d, y_i = \pi(x_{f(i)})\}$ is isomorphic to $B(d, D)$ if and only if f is cyclic.

As an application of this result, we characterized the *OTIS*(p, q)-layout of the de Bruijn digraph $B(d, D)$ when p and q are powers of d ($p = d^{p'}$, $q = d^{q'}$ and $p' + q' - 1 = D$). In particular, we have shown that when the diameter D is even, the minimum number of lenses is obtained with $p = d^{\frac{D}{2}}$ and $q = d^{\frac{D}{2}+1}$, and that for odd diameter D , the minimum can be obtained in time $O(D^2)$. For the cases where $p, q \neq d^i$, our exhaustive search on all de Bruijn digraphs with less than 2^{32} arcs leads us to conjecture that such layouts do not exist.

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