

Automatic Region-based Memory Management for Real-time Embedded Systems

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Automatic Region-Based Memory Management for Real-Time Embedded Systems



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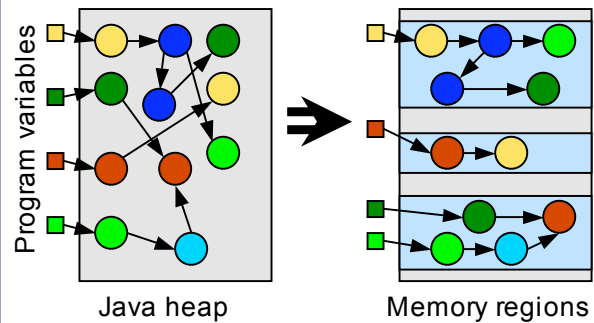
Motivation :

- Java means automatic memory management
- Garbage Collector means problems in a real-time context
 - Unpredictable pause times
 - Fragmentation of the heap
- *How can we provide automatic memory management without using a GC ?*

Our approach :

- Use region-based memory management
- Group data structures in regions
- Use a compile-time analysis to place objects in regions

Memory management with regions



- Objects allocated side by side : no more fragmentation
- Regions destroyed as a whole : predictable times
- Drawback: each object must be placed when allocated

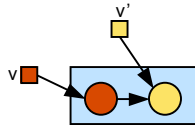
Pointer Interference Analysis :

- Build a partition of local variables
- $v \sim v'$ means they belong to the same data structure
- Simple algorithm :

$$\begin{aligned}
 v &:= u && \implies v \sim u \\
 v &:= u.f && \implies v \sim u \\
 v.f &:= u && \\
 \\
 m() & && \\
 \text{call } \sqrt{p1 \sim v1} & && \implies v1 \sim v2 \\
 \sqrt{p2 \sim v2} & && \\
 m'() & \text{ if } p1 \sim p2 &&
 \end{aligned}$$

Allocation Policy :

Simple allocation policy :
if two variables verify $v \sim v'$,
place their objects in the same region
Data structures will be automatically grouped by region



Other kinds of pointer analysis :

Escape analysis :

Does my object live longer than its method of origin ?

Points-to analysis :

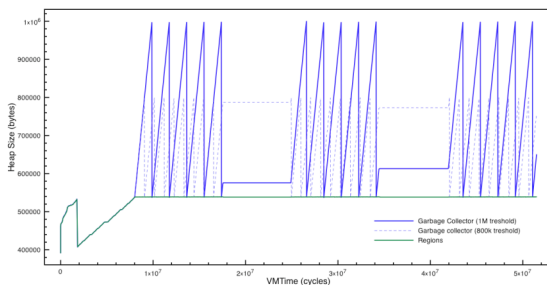
Where do the objects of my variable come from ?

Purity analysis :

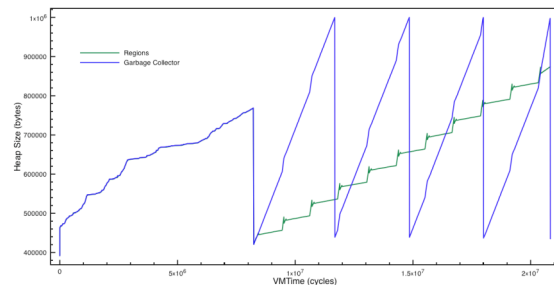
Does my method mutate the heap ?

Results :

Memory occupancy for two programs using a **Garbage Collector** or using **Regions**



- In this program, most regions are short-lived
- The program runs in nearly constant space
- No more need for a Garbage Collector



- This program uses a large mutating data structure
- Some of the generated garbage stays forever in the long-lived region
- Running this program without a GC may cause a memory leak

How to predict the runtime behaviour ?

Achieved :

- We propose a simple static analysis and allocation policy that groups data structures in regions
- Automatic region-based memory management can allow programs to run without a Garbage Collector

Perspectives :

- This approach has a tendency to place too many objects in the same region
- We need to find an algorithm to predict at compile time the behaviour of the region allocator