

# Dynamic visibility in polygonal scenes with the visibility complex

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► **To cite this version:**

Stéphane Rivière. Dynamic visibility in polygonal scenes with the visibility complex. 13th Annual ACM Symposium on Computational Geometry, 1997, Nice, France. 1997. <inria-00510094>

**HAL Id: inria-00510094**

**<https://hal.inria.fr/inria-00510094>**

Submitted on 17 Aug 2010

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In the dual space, the set of rays passing through  $p_v$  is a dual curve  $p_v^*$  which passes through a set of faces of the visibility complex. The dual curve passes from one face  $f$  to another face  $f'$  by crossing an edge  $e$  separating the two faces; the point  $p$  associated with  $e$  is then a point of the view. The algorithm for computing the view computes the list of edges of the complex crossed by the dual line of the point of view. Depending on the method used to find the crossed edges, the view can be computed in  $O(v \log n)$  time, or in  $\Omega(v)$  and  $O(vn)$  time ( $v$  size of the view).

With this list of crossed edges, we can maintain the view when the point of view is moving in the scene.

#### 4 Maintaining a view around a moving point

Let  $p_v$  be a point moving in the scene. The view around  $p_v$  changes when  $p_v$  crosses an extended edge of the visibility graph (i.e., the ray containing the edge): A point of the view becomes invisible, or a new point becomes visible, or  $p_v$  crosses the non-extended edge. In the last case, although the view does not change, we still classify this event as a visibility change: It corresponds to a visibility event in the algorithms that we describe.

By extending each edge of the visibility graph in free space, we obtain a partition of free space into regions of constant visibility: The view around a point is the same for all points inside the same *visibility region*. To maintain the view around  $p_v$ , we do not compute all the  $O(n^4)$  visibility regions, but only the limits of the visibility region that contains  $p_v$ .

When  $p_v$  is moving, we check whether  $p_v$  has left its visibility region; if so, we search the side of the visibility region crossed by  $p_v$  and compute the limits of the region incident to that side. We repeat the process until finding the new visibility region containing  $p_v$ .

We compute the limits of the visibility region of  $p_v$  by using the list of edges of the complex crossed by the dual curve  $p_v^*$ . Let us consider a crossed edge  $e = (v_l, v_r)$  whose extremities are the vertices  $v_l$  and  $v_r$ . If  $p_v^*$  passes through one of these vertices, then the list of crossed edges — and therefore the view around  $p_v$  — changes. Let  $p$  be the point associated with  $e$ . In the scene, the two rays  $v_l^*$  and  $v_r^*$  associated with  $v_l$  and  $v_r$  respectively, form a wedge that contains  $p_v$  and no other (extended) edges of the visibility graph incident to  $p$ : The wedge contains the visibility region of  $p_v$  (figure 2).

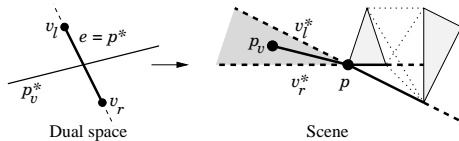


Figure 2: Wedge of visibility given by a point of the view.

More precisely, each wedge is made of an upper line and a lower line. The limits of the visibility region of  $p_v$  is then composed of the lower envelope of the set of upper lines and of the upper envelope of the set of lower lines (figure 3).

Once the view around  $p_v$  is computed with the visibility complex, the limits of the visibility region of  $p_v$  can be computed in  $O(v \log v)$  time ( $v$  size of the view).

When the point of view is moving from  $p_v$  to  $p'_v$ , we can check in logarithmic time whether the point has left its region and, if it is the case, find which line of the upper (resp. lower) has been crossed by  $(p_v, p'_v)$ . In the visibility

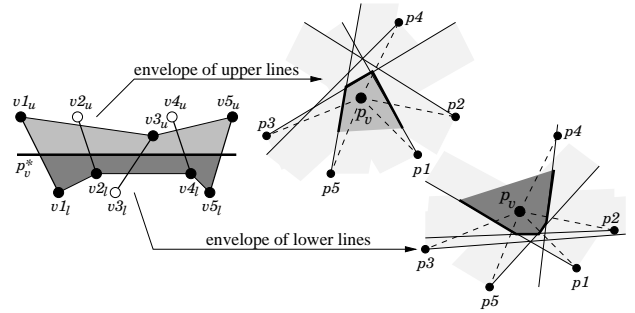


Figure 3: Computing the limits of the visibility region.

complex, the vertex  $v$  corresponding to the crossed line is swept by the dual curve  $p_v^*$ . Changes in the list of crossed edges (and therefore in the view) are local: They concern only the edges around the swept vertex and are done in constant time.

We must then update the limits of the visibility region. The changes are local too: We must remove (resp. insert) at most two lines in the set of lines, and the envelopes can be maintained by adapting the algorithms of Overmars and van Leeuwen [OvL81] or of Dobrindt and Yvinec [DY93] for dynamically maintaining a convex envelope.

**Proposition 1** *After computing the initial view around a point  $p_v$  and a preprocessing step in time  $O(v_0 \log v_0)$  ( $v_0$  size of view at  $t = 0$ ), the view around a moving point can be maintained in  $O(\log^2 v)$  time (resp. randomized  $O(\log v)$  time).*

We suppose now that the point of view is moving along a known trajectory, given for example by its position  $p_v(t)$ ,  $t \in [0, 1]$ . We do not need to compute all the limits of each visibility region crossed by the trajectory, but only the limits cut by the trajectory.

We suppose that we can compute in constant time the point of intersection of the trajectory with a line, and that we can compare in constant time two points on the trajectory (i.e., establish which one is before the other). It allows us to use a priority queue in which lines are ordered according to the position of their intersection point with the trajectory, so that we can pass the limits of visibility in order.

To initialize the priority queue, we first compute the view around  $p_v(0)$ , and for each edge of the complex crossed by  $p_v^*(0)$ , we check whether the upper line (resp. lower line) intersects the trajectory and — if it is the case — put it in the priority queue. Then, at each step we consider the first line of the priority queue, line which indicates the limit of visibility being crossed. We update (locally) the view as in the previous algorithm, consider the (at most two) new lines that limit the region, and put them in the priority queue if they cut the trajectory.

**Proposition 2** *After the computation of the initial view around a point  $p_v$  and a preprocessing step done in  $O(v_0 \log v_0)$  time, the view along a trajectory can be maintained in  $O(\log v)$  time at each change of visibility.*

#### 5 Dynamically maintaining the visibility complex

We now show how the visibility complex can be maintained when the polygon vertices are moving in the scene. More

precisely, each of the  $n$  polygon vertices is defined by its position  $p_i(t)$  for  $t \in [0, 1]$ , and we want to maintain the topological structure of the complex from  $t = 0$  to  $t = 1$ . Modifications in the complex arise at *elementary visibility changes*. They occur when three points  $p_1$ ,  $p_2$ , and  $p_3$  such that  $(p_1, p_2)$  and  $(p_2, p_3)$  are edges of the visibility graph, change their relative positions:  $p_2$  previously below (resp. above)  $(p_1, p_3)$  becomes above (resp. below)  $(p_1, p_3)$  (figure 4a).

Corresponding changes in the visibility complex are local: They occur around the vertices  $v_{12}$  and  $v_{23}$  corresponding to  $(p_1, p_2)$  and  $(p_2, p_3)$  respectively (figure 4b), and can be made in constant time (there are 50 cases to consider).

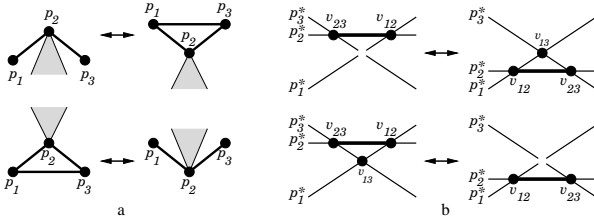


Figure 4: Elementary visibility changes. **a.** In the scene. **b.** In the visibility complex.

To maintain the complex from  $t = 0$  to  $t = 1$ , we must handle all the visibility changes in order. We compute for each edge  $e_j = (v_{ij}, v_{jk})$  ( $v_{ij} = p_i^* \cap p_j^*$ ) of the complex a “date of death”  $t_d$  that is the time when the three corresponding points  $p_i(t_d)$ ,  $p_j(t_d)$ , and  $p_k(t_d)$  in the scene will be collinear. Then we put the edges in a priority queue and handle the visibility changes in order. At each visibility change, the set of dates of death can be updated in constant time and the priority queue can be updated in  $O(\log k) = O(\log n)$  time ( $k$ , size of the visibility graph, verifies  $k = O(n^2)$ ). If dates of death can be computed in constant time, then:

**Proposition 3** *After a preprocessing step in  $O(k_0 \log n)$  time ( $k_0$  size of the visibility complex at  $t = 0$ ), the visibility complex of a moving polygonal scene can be maintained in  $O(\log n)$  time at each visibility change.*

## 6 Maintaining a view around a moving point among moving polygons

We finally show how to maintain a view around a point when the point of view and the obstacles are moving. In this context, using the data structures shown before for maintaining a view (dynamic convex envelope or priority queue) is not well-suited: At each change in the visibility complex, these structures must be updated and the modifications may not be local. Thus we integrate the visibility changes around the point of view with the visibility changes of the visibility complex.

For each vertex  $v_{ij}$  of the complex extremity of an edge crossed by the dual curve  $p_v^*$  we also compute a “date of death”, that is the date  $t_d$  when  $p_v^*(t_d)$  passes through  $v_{ij}(t_d)$ , that is when in the scene  $p_v(t_d)$ ,  $p_i(t_d)$ , and  $p_j(t_d)$  are collinear, and put all these events in the priority queue.

When an elementary visibility change occurs in the complex, some dates of death related to the view around  $p_v$  must also be updated. However, these changes are local, and if for each edge put in the priority queue we add a flag indicating whether the edge is crossed by the dual curve of the point of

view, all modified dates of death can be updated in constant time.

**Proposition 4** *After a preprocessing step in  $O(k_0 \log n)$  time ( $k_0$  size of the visibility complex at  $t = 0$ ), the view around a moving point among moving polygons can be maintained in  $O(\log n)$  time each time a visibility change occurs in the scene or in the view.*

## 7 Conclusion

We have shown two output-sensitive algorithms for maintaining the view around a moving point in a fixed scene. These algorithms use the visibility complex so that local changes of visibility in the scene yield also “local” recomputations. They are particularly efficient in large occluded environments, where the size of the visibility complex is rather linear than quadratic (in the number of polygon vertices) and the size of the view is much smaller than the size of the scene. They are relatively simple and have been implemented.

Then we have shown an algorithm for maintaining the visibility complex for dynamic scenes which can be used in global visibility computations. The algorithm handles the changes of visibility according to their date of occurrence. We are currently studying whether visibility changes happening in different hidden parts of the scene could be handled independently in any order. We are also studying algorithms for maintaining the complex upon insertion or deletion of polygons in a scene, algorithms which could reuse (for a first study) our algorithms for dynamic visibility.

Finally, we have shown an algorithm for maintaining the view around a point moving in a dynamic scene, algorithm more suited to the case when most of the objects of the scene are seen at some point in time from the point of view.

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