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# Optimal Control of Discrete Event Systems under Partial Observation

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## Abstract

We are interested in a new class of optimal control problems for Discrete Event Systems (DES). We adopt the formalism of supervisory control theory [7] and model the system as a finite state machine (FSM). Our control problem is characterized by the presence of uncontrollable as well as unobservable events, the notion of occurrence and control costs for events and a worst-case objective function. We first derive an observer for the partially unobservable FSM, which allows us to construct an approximation of the unobservable trajectory costs. We define the performance measure on this observer rather than on the original FSM itself. Further, we use the algorithm of [8] to synthesize an optimal submachine of the observer. This submachine leads to the desired supervisor for the system.

## 1 Introduction and Motivation

We are interested in a new class of optimal control problems for Discrete Event Systems (DES) [7]. The system to be controlled is modeled as a finite state machine (FSM). Our control problem follows the theory in [8] and is characterized by the presence of uncontrollable events, the notion of occurrence and control costs for events and a worst-case objective function. However, compared to the work in [8] and compared to [3, 6], we wish to take into account partial observability. Several concepts and properties of the supervisory control problem under partial observation were studied in [1, 4] among others. However, they only propose a qualitative theory for the control of DESs.

The starting point of our solution is a FSM which represents the global behavior of a given system, including its unobservable dynamics. The first step is the derivation of an observer for the partially unobservable FSM, called a C-observer. This step is necessary since unobservable events alone cannot trigger a specific behavior of a controller. We define the performance measure on the C-observer rather than on the original FSM itself. However, we will make the necessary efforts to keep track of the information that has disappeared with the initial structure. This observer allows us to remember an approximation of the unobservable costs between two observable events. This approximation corresponds to the worst, i.e., the highest, cost of the different unobservable trajec-

ries than can occur between two observable events. In the second step, we use the theory in [8] to synthesize an optimal controller corresponding to the optimal restricted behavior, insofar as it is achievable by an admissible (i.e., physically constructible) supervisor. We use back-propagation from the goal state to generate the supervisor, based on event cost functions. The supervisor is synthesized in a manner that gives them optimal sub-structure, consistent with the notion of DP-Optimality of [8].

## 2 Preliminaries

The system to be controlled is modeled as a FSM defined by a 5-tuple  $G = \langle \Sigma, Q, q_0, q_m, \delta \rangle$ , where  $\Sigma$  is the set of events,  $Q$  is the (finite) set of states,  $q_0$  is the initial state,  $q_m$  is the unique marked state, and  $\delta$  is the partial transition function defined on  $\Sigma^* \times Q$ . The behavior of the system is described by the prefix-closed language  $\mathcal{L}(G)$  [2], generated by  $G$ . Similarly, the language  $\mathcal{L}_m(G)$  corresponds to the marked behavior of the FSM  $G$ , i.e., the set of trajectories of the system ending in  $q_m$ . Some of the events in  $\Sigma$  are uncontrollable, i.e., their occurrence cannot be prevented by a controller, while the others are controllable. Likewise, control will be applied on a plant that is partially observable, i.e. the supervisor will observe only a subset of the events generated by plant  $G$ . Hence some of the events in  $\Sigma$  are observable whereas the others will be unobservable. In this regard,  $\Sigma$  can be partitioned as  $\Sigma = \Sigma_c \cup \Sigma_{uc}$  with  $\Sigma_c \cap \Sigma_{uc} = \emptyset$  and  $\Sigma = \Sigma_o \cup \Sigma_{uo}$  with  $\Sigma_o \cap \Sigma_{uo} = \emptyset$ , where  $\Sigma_c$ ,  $\Sigma_{uc}$ ,  $\Sigma_o$  and  $\Sigma_{uo}$  represent the set of controllable, uncontrollable, observable and unobservable events, respectively. Moreover, unobservable events are assumed to be uncontrollable, i.e.,  $\Sigma_{uo} \subseteq \Sigma_{uc}$ . In the sequel, we will only be interested in *trim* FSMs, i.e., FSMs for which all states of  $Q$  are accessible from  $q_0$  and coaccessible to  $q_m$  [2]. We say that FSM  $A = \langle \Sigma_A, Q_A, q_{0A}, q_m, \delta_A \rangle$  is a submachine of  $G$ , denoted  $A \subseteq G$ , if  $\Sigma_A \subseteq \Sigma$ ,  $Q_A \subseteq Q$ ,  $\forall \sigma \in \Sigma_A, q \in Q_A \delta_A(\sigma, q)! \Rightarrow (\delta_A(\sigma, q) = \delta(\sigma, q))$ . Notation  $\delta_A(\sigma, q)!$  means that  $\delta_A(\sigma, q)$  is defined, i.e., there is a transition labeled by event  $\sigma$  out of state  $q$  in  $A$ . We say that  $A$  is a submachine of  $G$  at  $q$  whenever  $q_{0A} = q \in Q$  and  $A \subseteq G$ . For any  $q \in Q$ , we will use  $\mathcal{M}(G, q) = \{A : A$

a trim submachine of  $G$  with respect to  $q_m$  and  $q_{0_A} = q$  to represent the set of trim submachines of  $G$  at  $q$  with respect to  $q_m$ . This set has a maximal element in the sense that the maximal element contains all other elements as submachines. It is denoted by  $M(G, q)$ . In order to consider the control problem under partial observation, we need to make sure that the initial FSM  $G$  has no unobservable cycle. Otherwise it would be impossible to alleviate the fact that it could make the system run indefinitely in that cycle, without the supervisor noticing. We then assume that  $G$  has no unobservable cycles.

Finally, to take into account the numerical aspect of the optimal control problem, two cost values are associated to each event of  $\Sigma$ . We introduce an occurrence cost function  $c_e : \Sigma \rightarrow \mathbb{R}^+$  and a control cost function  $c_c : \Sigma \rightarrow \mathbb{R}^+ \cup \{0, \infty\}$ . Control costs are used to represent the fact that disabling a transition possibly incurs a cost. The control cost function is infinity for events in  $\Sigma_{uc}$ . The cost functions are then used to introduce a cost on the trajectories of a submachine of  $G$ .

### 3 The C-Observer with respect to $\Sigma_{uo}$

The framework in which we develop our control theory is that of partially observable FSMs. The supervisor that will be generated should be able to take decisions based on the states and/or events that it observes. Consequently, we base our model upon a partially observed system, seen through an observer. However, in order to take into account unobservable events in the optimality under which we apply our control, we must keep track of their costs. The idea is to collect an approximation of the costs between two observable events in the states of the observer we want to build. For example, consider two states  $p$  and  $q$  of  $G$ , connected by (at least) a trace of the form  $\sigma s \in \Sigma_o \Sigma_{uo}^*$ . As we only observe the first event, it is not possible to know which trajectory has been taken between these two states. Hence, from an optimal control point of view, we have to consider that the plant evolves through the trajectory with the highest cost (there is no way to control the system in such a way that this trajectory is not taken). In order to collect these costs, we build a deterministic observer, named C-observer (Observer with Costs), and define the notion of a macro-state, allowing both to mask the underlying nondeterminism by abstracting away from the nondeterministic submachine and to keep track of the unobservable event costs of trajectories between two states. The C-observer constitutes the basic model on which the optimal control will be applied.

Before giving formally the definition of the C-observer, denoted by  $G_c$ , we need to check the original FSM  $G$  in order to account for unobservable events that may lead to  $q_m$  in  $G$ . Indeed, if an unobservable event leads to  $q_m$  in  $G$ , it may be impossible to determine whether or not the system has actually reached  $q_m$ . We therefore update  $G$  by adding a self-loop at  $q_m$ , labeled  $\varphi$  with  $\delta(\varphi, q_m) = q_m$ . The  $\varphi$  event is just an (observable) indicator event (e.g. a sensor) that signals that

$q_m$  has been reached. Without loss of generality, we can assume it is controllable and has zero occurrence and control costs.

#### 3.1 The C-Observer definition

The new structure that we define is called a C-observer. It is denoted by  $G_c = \langle \Sigma_o, X, x_0, x_m, f \rangle$ , where  $\Sigma_o$  is the set of observable events,  $X$  is the set of macro-states,  $x_0$  is the initial macro-state,  $x_m$  is the marked macro-state, and  $f$  is the partial transition function defined over  $\Sigma_o^* \times X \rightarrow X$ .

Starting from  $G$ , the set  $X$  of macro-states of  $G_c$  will be constituted of pairs in  $Q \times \mathbb{R}^+$ . More specifically, the admissible states that are considered are states that can be reached by a trace of events constituted by an observable first event followed by a sequence of unobservable events. In language formalism, the latter trace should be in  $\Sigma_o \Sigma_{uo}^*$ . We present more formally the way the states of the system  $G_c$  are built. First, we introduce the set of triples  $\mathcal{D}$  defined by :

$$\mathcal{D} = \{(p, q, \sigma) \in Q \times Q \times \Sigma_o / \exists s \in \Sigma_{uo}^*, \delta(\sigma s, p) = q\}. \quad (1)$$

A triple  $(p, q, \sigma)$  belongs to set  $\mathcal{D}$  if there is a trace between  $p$  and  $q$  whose first event is  $\sigma$  and whose following events are all unobservable. Note that more than one trace  $s$  could verify this condition. We now define the set of traces that verify the above conditions, for a given triple  $(p, q, \sigma)$  :

$$\forall (p, q, \sigma) \in \mathcal{D}, \mathcal{S}(p, q, \sigma) = \{s \in \Sigma_{uo}^* / \delta(\sigma s, p) = q\}. \quad (2)$$

Using (2), we can easily deduce the following property:

**Property 1**  $\forall (p, q, \sigma) \in \mathcal{D}, |\mathcal{S}(p, q, \sigma)| < \infty$ .

Finally, we do not want to lose the cost of the unobservable events that have been projected. To this effect we introduce the notion of *locally computed cost* associated with a triple  $(p, q, \sigma)$  of  $\mathcal{D}$ . Formally, it is given by a function, denoted by  $c_o$ , over  $\mathcal{D} \rightarrow \mathbb{R}^+$ , and defined by:

$$\forall (p, q, \sigma) \in \mathcal{D}, c_o(p, q, \sigma) = \max_{s \in \mathcal{S}(p, q, \sigma)} c_e(s) \quad (3)$$

This way, we keep track of the worst unobservable trace that could lead from  $p$  to  $q$ . Using the previous notations,  $G_c$  is a FSM, defined as follows:

**Definition 1** Given an FSM  $G$ , the associated C-observer  $G_c$  is given by a tuple  $\langle \Sigma_o, X, x_0, x_m, f \rangle$ . It is an FSM whose elements are defined as follows:

1.  $X$  is the set of macro-states.  $x \in X$  is defined by a set of pairs  $(q, c) \in Q \times \mathbb{R}^+$ , called micro-states;
2. The final macro-state is defined by  $x_m = \{(q_m, 0)\}$  and the initial macro-state  $x_0$  as :

$$x_0 = \{(q, c_q), \exists s \in \Sigma_{uo}^*, \delta(s, q_0) = q \text{ and } c_q = \max_{t \in \Sigma_{uo}^*, \delta(t, q_0) = q} c_e(t)\}$$

3.  $\forall x \in X$  and  $\forall \sigma \in \Sigma_o$ , define

$$\forall (p, c_p) \in x, A_\sigma^x(p) = \{(q, c_o(p, q, \sigma)) \mid (p, q, \sigma) \in \mathcal{D}\}.$$

$A_\sigma^x(p)$  basically constitutes the set of states of  $G$  that can be reached via a trace  $\sigma \Sigma_{u_o}^*$  (from a micro-state of  $x$ ), together with the associated approximation of the unobservable trace cost.

4. The transition function  $f$  is recursively defined by:

$$\forall x \in X \forall \sigma \in \Sigma_o, f(\sigma, x) = \{(q, c_q) \in \bigcup_{(p, c_p) \in x} A_\sigma^x(p), c_q = \max_{s \in \mathcal{S}(p, q, \sigma)} c_e(s)\}$$

Hence, if there exists different micro-states of the form  $(q, \cdot)$  in  $f(\sigma, x)$ , then we only consider the pair with the maximal cost.

5. We only build the accessible part of the system (i.e. the states  $x \in X$  that are reachable from  $x_0$  by  $f$ ).

The way  $G_c$  is built masks the nondeterministic nature of the projected FSM.  $x_0$  is computed from the unobservable reach of  $(q_0, 0)$ .  $x_m$  is a single marked state, namely,  $\{(q_m, 0)\}$ . Finally,  $f$  can be constructed recursively from the initial state. Indeed, we can construct the set of states of  $G_c$  using point (2) and then point (3) and (4) of Definition 1 recursively. Note that due to Property 1, the recursion terminates. The structure that we obtain is another deterministic FSM, whose events are taken in  $\Sigma_o$ . States of  $G_c$  are macro-states with respect to  $G$ . However, we have computed and kept a local cost to avoid losing track of the costs of the unobservable events that have disappeared from the structure.

**Lemma 1 [5]** Let  $x \in X - \{x_m\}$  be a state of  $G_c$ , and let  $(q, c_q) \in x$  be a micro-state of  $x$ . We can state that

- (1) either  $\exists \sigma \in \Sigma_o, \exists q' \in Q, \delta(\sigma, q) = q'$  and, in this case,  $\exists x' \in X$ , s.t.  $f(\sigma, x) = x'$ , and  $(q', \cdot) \in x'$
- (2) or  $\exists \sigma \in \Sigma_{u_o}$  and  $\exists q' \in Q$  s.t.  $\delta(\sigma, q) = q'$  and, in this case,  $\exists (q', \cdot) \in x$ .

Moreover,  $\forall (q, c_q) \in x, \exists s \in \Sigma_{u_o}^* \Sigma_o, \delta(s, q)!$ .

What the above lemma states is that whatever the state  $x$  that can be reached during the execution of the plant, there eventually exists a way out of this state (either directly via an observable event or by an unobservable trajectory which reaches a micro-state of  $x$  having the previous property. Next, we state that the C-observer realized from  $G$  inherits properties of  $G$ .

**Proposition 1 [5]**  $G_c$  is non-blocking.

### 3.2 Extended notion of Controllability

In this section, we formalize the method used (by a supervisor) to generate a submachine from a C-observer.

**Submachines of a C-observer.** We wish to apply some control to the original system in order to verify a certain performance criterion. In other words, we wish to reduce the system  $G_c$ , and therefore  $G$ , to a particular behavior. This

leads us to define the notion of a submachine of  $G_c$ . In fact, even if the worlds in which they are defined (for  $G$  and  $G_c$ ) are different, the notion of submachine is the same as the one given in Section 2 (i.e. a submachine of  $G_c$  is any structure that has its states in those of  $G_c$ , the same initial state and final state and its events and transitions in those of  $G_c$ ).

Moreover, we are only interested in complete behavior, i.e. we wish to obtain a controlled system that reaches the state  $x_m$  and therefore the state  $q_m$ . Hence, we wish to consider the submachine of  $G_c$  that have this property. Hence the notion of  $G$ -live submachines.

**Definition 2** Let  $G_c = \langle \Sigma_o, X, x_0, x_m, f \rangle$  be the C-observer associated with  $G = \langle \Sigma, Q, q_0, q_m, \delta \rangle$ . A submachine  $H = \langle \Sigma_o, X_H, x_0, x_m, f_H \rangle$  of  $G_c$  is said to be  $G$ -live if the following condition holds:

$$\forall x_H \in X_H \setminus \{x_m\}, \forall (q, c_q) \in x_H, \exists (q', c_{q'}) \in x_H \text{ s.t.} \\ \{[\exists s \in \Sigma_{u_o}^*, \delta(s, q) = q'] \wedge [\exists \sigma \in f_H(x_H), \delta(\sigma, q')!]\}.$$

A submachine  $H$  of  $G_c$  is  $G$ -live whenever any micro-state of  $x_H$  has a transition that is either an observable transition for the initial FSM  $G$ , or an unobservable transition that leads to another micro-state of  $x_H$  from which there is a possibility of exiting the macro-state (except for the marked state). Quite naturally, using Lemma 1, we can state that :

**Proposition 2 [5]** If  $G_c$  is the C-observer associated with  $G$ , then  $G_c$  is  $G$ -live.

**Controllability in this framework.** The structure on which control will be applied is FSM  $G_c$ . We first have to adapt the classical definition of controllability introduced by [7]. Indeed, even if the control policy remains the same (we do not want to disable uncontrollable events), we have to take care of the fact that, by removing controllable transitions, the obtained submachine of  $G$  inherits some properties of the initial FSM  $= G_c$ . Hence the new definition of controllability:

**Definition 3** Let  $G_c = \langle \Sigma_o, X, x_0, x_m, f \rangle$  be the C-observer associated with  $G = \langle \Sigma, Q, q_0, q_m, \delta \rangle$ .  $H = \langle \Sigma_o, X_H, x_0, x_m, f_H \rangle$  is said to be a controllable submachine of  $G_c$  if the following conditions hold:

1.  $\forall x_H \in X_H$  that can be reached via a trace of  $\mathcal{L}(H)$ ,  $\forall \sigma \in \Sigma_{uc} \cap \Sigma_o, f(\sigma, x_H)! \Rightarrow f_H(\sigma, x_H)!$ ,
2.  $H$  is  $G$ -live.

Condition (1) imposes that any transition that needs to be disabled in  $G_c$  to generate  $H$  needs to be controllable. Condition (2) imposes that no submachine of a C-observer presents any deadlocks or livelocks. This condition imposes that any micro-state of a state  $x_H$  must have an active outgoing trace (in the original FSM from which  $G_c$  was derived) that is either unobservable (thereby leading to another micro-state of  $x_H$  and eventually leading to a state from which there is an observable outgoing event) or observable (thereby leading to another macro-state of  $H$ ).

**The supervisor.** Now that we have the definition of a controllable submachine of a C-observer, it is interesting to determine how such a submachine can be obtained via a supervisor acting upon  $G_c$ . However, control cannot be blindly performed. Disabling an event that was admissible in a state  $x$  of  $G_c$  can induce a deadlock in the initial FSM  $G$ . Hence, we introduce the notion of *Admissible Control Actions (ACA)*.

**Definition 4** Let  $G_c = \langle \Sigma_o, X, x_0, x_m, f \rangle$  be the C-observer associated with  $G = \langle \Sigma, Q, q_0, q_m, \delta \rangle$ . We define the set of Admissible Control Actions (ACA) at state  $x \in X$  as a function :

$$\Gamma_x = \{ \gamma \subseteq \Sigma_c, \forall (q, c_q) \in x, \exists (q', c_{q'}) \in x \text{ s.t.} \\ \{ [\exists s \in \Sigma_{uo}^*, \delta(s, q) = q'] \wedge [\exists \sigma \in f(x) \setminus \gamma, \delta(\sigma, q')!] \} \}$$

More precisely,  $\Gamma_x$  gives, for a state  $x$  of  $G_c$  all the possible sets of controllable events that can be disabled without risk of deadlock. In other words, given a state  $x$  of  $G_c$  and a given  $\gamma$  in  $\Gamma_x$ , if  $\sigma$  belongs to  $\gamma$ , it means that  $\sigma$  can be disabled because there actually exists at least one trajectory  $s \in \Sigma_{uo}^*$  that leads the system in another micro-state of  $x'$  for which there exists an observable event  $\sigma'$  that makes the system leave the macro-state  $x$  and eventually reach a state  $x' = f(x, \sigma')$  of  $G_c$ .

Using Definition 4, a supervisor of  $G_c$  is defined by:

**Definition 5** Let  $G_c$  be the C-observer associated with  $G$  and  $\Pi = (\Gamma_x)_{x \in X}$  be the set of admissible control actions, then a supervisor  $S$  is a function given by :

$$S : X \rightarrow 2^{\Sigma_c} \\ x \mapsto \gamma \in \Gamma_x \quad (4)$$

In other words, a supervisor of  $G_c$  is obtained by choosing a particular  $\gamma$  in a state  $x$ . By definition, the control action will always belongs to  $\Sigma_c$ , which ensures that  $S$  never disables an uncontrollable event.

Conceptually, the supervisor controlling the plant  $G$  is placed in feedback with  $G$  and  $G_c$ . Only the observable events can be seen by  $S$ . Therefore  $G_c$  plays the role of an observer that will somehow rebuild a part of the state in which the system has evolved. According to this information, the supervisor determines whether the observation corresponds to a (conditionally) controllable event and if it has to enable/disable this event in order to keep the closed loop system behaving "desirably".

To conclude this section, let us remark that Definition 5 is consistent with the definition of a controllable submachine of the C-observer  $G_c$ . This is summarized by the following proposition:

**Proposition 3**  $H \subseteq G_c$  is a controllable submachine of  $G_c$  if and only if there exists a supervisor  $S$ , such that  $\forall x_H \in X_H, f_H(x_H) = f(x_H) \setminus S(x_H)$ .

#### 4 Optimal Supervisory Control Problem

The aim of optimal control is to study the behavioral properties of a system, to take advantage of a particular structure,

and to generate a controller which constrains the system to a desired behavior according to quantitative and qualitative aspects [3, 6, 8]. This is performed by the addition of quantitative measures in the form of occurrence and control cost functions, to capture the fact that some legal behaviors are better than others.

#### 4.1 Transformation of $G_c$

We first need to transform the C-observer, in order to exactly fit within the framework developed by [8]. Indeed, unlike in the case of total observability where costs are defined in events only, we have incorporated cost information in the macro-states of the  $G_c$ . These costs were attached to the states in order to keep track of the unobservable cost of the trajectory between two macro-states (see Section 3.1). Basically, the transformation we will perform on  $G_c$ , consists in "shifting" the cost of the macro-state to the events that can be executed in this macro-state. For a given  $x$ , and a given  $\sigma$  admissible in  $x$ , we consider the worst cost of the pairs  $(q, c_q) \in x$  such that  $\sigma$  belongs to the active event set of  $q$  in  $G$ . The transformation is performed as follows: let  $x \in X$  and let  $f(x)$  be the set of events that  $G_c$  can execute in  $x$ . For each  $\sigma \in f(x)$ , we rename  $\sigma$  as  $\sigma_x$  and we attach to this new event the cost  $c_e(\sigma_x)$  defined by :

$$c_e(\sigma_x) = \max_{(q, c_q) \in x, \delta(\sigma, q)!} \{c_q\} + c_e(\sigma) \quad (5)$$

The controllability status of the event as well as the control cost of the events do not change (namely, we have  $c_c(\sigma_x) = c_c(\sigma)$ ). Call  $\Sigma'_o$  the new set of event. The transition function  $f$  remains the same (i.e.  $f(x, \sigma_x)$  is defined and equal to  $x'$  whenever  $x' = f(x, \sigma)$ ).

The new C-observer  $G'_c$  we obtain is still a FSM. It is defined by  $\langle \Sigma'_o, X, x_0, x_m, f \rangle$ . Compared to  $G_c$ , the global structure of  $G'_c$  does not change. The only difference is that we change the original alphabet of  $G_c$  in such a way that costs are now defined on events only, as carried out in [8]. From now on,  $G'_c$  is a deterministic and trim FSM. To each event is attached two values, which respectively correspond to its event and control costs. The only difference with [8] lies in the notion of controllability that, in our framework, takes into account the notion of liveness of the underlying system  $G$ . However, this does not affect the use of the theory of [8] to compute the optimal supervisor of  $G_c$ , and therefore the optimal supervisor of  $G$ . Indeed, as in our case, the theory is based on the notion of acceptable control actions that have to be computed at first. In [8], a control action in a state  $x$  is admissible whenever it does not disable uncontrollable events and it does not produce local deadlock (i.e. no output event.)

#### 4.2 Trajectory costs of a submachine of $G'_c$

In order to be able to discuss optimality, we now explain how to compute the cost of a trajectory of  $G'_c$ .

**Control cost function over the states.** In order to model this particular aspect, let us define the control cost of an event

according to a state. We first introduce  $\Sigma_d(x, H) = f_{G'_c}(x) \setminus f_H(x)$  as the set of disabled events at state  $x$  for the system to remain in submachine  $H$  of  $G'_c$ . Whereas in [8] the control cost function was defined on an event, in the case of partial observation, it is defined on a state as follows: considering a submachine  $H$  of  $G'_c$ , we have

$$C_c(x, H) = \begin{cases} \infty & \text{if } \Sigma_d(H, x) \notin \Gamma_x \\ \sum_{\sigma' \in \Sigma_d(H, x)} c_c(\sigma') & \text{otherwise} \end{cases} \quad (6)$$

The cost of a state  $x$  is equal to  $\infty$  whenever there does not exist a particular control policy  $\gamma \in \Gamma_x$  that restricts the behavior of  $G'_c$  to  $H$  (i.e. when an uncontrollable event has been removed or when a controllable event has been removed, then inducing a deadlock).

**Cost of a trajectory and of a submachine of  $G'_c$ .** We are now ready to define the cost of a trajectory  $s$  of a submachine  $H$  as well as the objective cost function of a submachine  $H$  of  $G'_c$ .

**Definition 6** Let  $H = \langle \Sigma'_o, X_H, x_{0,H}, x_{m,H}, f_H \rangle$  be a sub-machine of  $G'_c$  derived from  $G$  and  $\mathcal{L}_m(H)$  be the marked language generated by  $H$ , then

1. for all  $y$  in  $H$  and trajectory  $s = \sigma'_1 \dots \sigma'_n, \forall i, 1 \leq i \leq n, \sigma'_i \in \Sigma'_o$  such that  $f_H(y, s)$  exists, the cost of  $s$  is given by :

$$C_O(y, H, s) = \sum_{i=1}^n c_e(\sigma'_i) + \sum_{i=0}^n C_c(f_H(y, \|s\|_i), H), \quad (7)$$

where  $\|s\|_i$  denotes the prefix of  $s$  of length  $i$ ,

2. the objective cost function denoted by  $C_{Sup}(H)$  is given by:

$$C_{Sup}(H) = \sup_{s \in \mathcal{L}_m(H)} (C_O(x_0, H, s)) \quad (8)$$

The cost of a trajectory is the sum of the occurrence costs of the events composing it, to which is added the cost of controlling events on the way to remain in machine  $H$ . If an uncontrollable event is disabled, the cost of a trajectory becomes infinite because of the second term of (7). Finally,  $C_{Sup}(H)$  represents the worst case behavior that is possible in submachine  $H$ . The next lemma characterizes the interaction of event and control costs:

**Lemma 2 [5]** Let  $H_1 \subseteq H_2 \in \mathcal{M}(G'_c, x)$  and  $s \in \mathcal{L}_m(H_1)$ , then  $C_O(x, H_1, s) \geq C_O(x, H_2, s)$ .

This lemma states that the cost associated with a trajectory admissible in a machine is lower than the cost of the same trajectory generated by one of its submachines. The purpose of “contracting a submachine” is to remove trajectories with high event costs. However this process is accompanied by rising control costs, hence the optimization problem we now define.

### 4.3 The optimization problem

We are only interested in machines that achieve a task (we only consider plants having a behavior which terminates at a marked state). Among all the trim and controllable submachines of  $G'_c$ , since we want to deal with optimal solutions, we want to extract the submachines that have a minimal objective cost function.

**The optimal submachines of  $G'_c$ .** Considering the trim hypothesis, we denote  $\mathcal{M}(G'_c, x)$  as the set of trim submachines of  $G'_c$  starting at state  $x$  with respect to the unique final state  $x_m$  and denote by  $M(G'_c, x)$  its maximal element (see Section 2). We now define the optimization problem.

**Definition 7**  $\forall x \in X, H_o \in \mathcal{M}(G'_c, x)$  is an optimal sub-machine of the FSM  $G'_c$  if

$$C_{Sup}(H_o) = \min_{H \in \mathcal{M}(G'_c, x)} C_{Sup}(H) < \infty.$$

The cost  $C_{Sup}(H_o)$  of  $H_o$  represents the minimum worst case cost incurred to reach  $x_m$  from  $x_0$  when the behavior of  $G'_c$  is restricted to a submachine of it. As some events in some states are not controllable (which induces an infinite cost), optimality is met when there is no other control policy with lower worst-case cost that allows to reach the marked state  $x_m$  certainly. At a lower level (in the world of  $G$ ), the control policy induced by submachine  $H_o$  corresponds to the one with lower worst-case cost, knowing that  $G$  could evolve through unobservable trajectories with the worst possible cost. In general, there will exist several optimal submachines for an FSM.

As in the case of total observation [8], the following lemma is stated to note that optimal solutions lie within the class of controllable submachines.

**Lemma 3 [5]** Let  $H \in \mathcal{M}(G, x)$ . If  $C_{Sup}(H) < \infty$  then  $H$  is controllable.

From Lemma 3, uncontrollable submachines are not candidates for optimality since the cost for restricting the system to those submachines is infinite. The following theorem gives necessary and sufficient conditions for the existence of optimal submachines:

**Theorem 1 [5]** An optimal submachine of  $G'_c$  exists if and only if there exists a submachine  $H$  of  $G'_c$  such that  $H$  is trim, controllable, with no cycles.

Intuitively, this theorem states that an optimal solution exists when there are controllable submachines of  $G'_c$  in which there does not exist cycles. The controllability assumption ensures that the cycles can be broken using controllable events alone. The submachine that includes all the other optimal submachines will be called the maximal optimal submachine and will be denoted by  $H^\uparrow_o$ .

**The DP-optimal submachines of  $G'_c$ .** In general, the solution to the Optimal Supervisory Control Problem is not unique. Moreover, all the optimal solutions do not structurally have optimal sub-solutions, which means that they do

not satisfy the principle of Dynamic Programming. In fact, in the previous section, optimality is obtained only regarding the paths between the initial and final state and never the post-fix paths between any state of the corresponding FSM and the final state. In this section, we will show that whenever an optimal solution exists, a solution having optimal sub-structure also exists. We call this latter type a DP-optimal solution (DP stands for Dynamical Programming) and define it as follows :

**Definition 8** A submachine  $H_{DO}$  of  $G'_c$  is DP-Optimal if it is optimal and  $\forall x' \in X_{H_{DO}}$ ,  $M(H_{DO}, x')$  is an optimal submachine in  $\mathcal{M}(G'_c, x')$ .

We have already seen that optimality actually exists when the worst-case cost from the initial state  $x_0$  to  $x_m$  is finite once minimized. DP-Optimality is obtained when any terminal path from any state of a submachine to the goal state  $x_m$  is optimal in the previous sense.

If a particular DP-Optimal FSM includes all other DP-Optimal FSMs as submachines of itself, then we call it the *maximal DP-Optimal submachine*. The maximal DP-Optimal submachine of a machine  $G'_c$  at  $q$  w.r.t.  $x_m$  will be denoted by  $M_D^o(G'_c, x)$ . Note that all DP-Optimal submachines are acyclic. The existence of a DP-Optimal submachine of  $G'_c$  is given by the following theorem (the proof can be found in [8]).

**Theorem 2** If an optimal submachine of  $G'_c$  exists, then the unique maximal DP-Optimal submachine  $M_D^o(G'_c, x_0)$  of  $G$  w.r.t.  $x_m$  also exists.

**The DP-Optimal algorithm.** Consider a FSM  $G = \langle \Sigma, Q, q_0, q_m, \delta \rangle$  with a unique initial state  $q_0$ , and a unique marked state  $q_m$  and its corresponding transformed C-observer  $G'_c = \langle \Sigma'_o, X, x_0, x_m, f \rangle$ . Then there exists an algorithm [8], named **DP-Opt**, with a worst-case complexity  $\mathcal{O}(|X|^2|\Sigma_o| \log(|\Sigma_o|) + |X|^3|\Sigma_o|)$  (Theorem 6.10 of [8]), that constructs the desired maximal DP-Optimal submachine  $M_D^o(G'_c, x_0)$  of the FSM  $G'_c$  w.r.t.  $x_0$  and  $x_m$ . The algorithm also returns the worst inevitable cost  $c_{sup}^g(M_D^o(G'_c, x_0))$ . We refer the reader to [8] for a complete description of **DP-Opt**.

#### 4.4 The supervisor

The supervisor computation consists of different steps. Once the C-observer  $G_c$  derived from the initial FSM  $G$  is computed, we first have to transform it into  $G'_c$  by attaching the cost induced by the unobservable trajectories to the events in order to fit within the framework of [8] (see Section 4.1). From this machine, using the algorithm of [8], we compute (if it exists) the DP-Optimal solution  $M_D^o(G'_c, x_0)$  of  $G'_c$ . At this point, we disable in  $G_c$  the corresponding sets of events in  $\Sigma'_o$  and for all  $x \in X$ , we retrieve  $\Sigma_d(G_c, x)$ , the set of disabled event at state  $x$  for the system to remain in submachine  $M_D^o(G_c, x_0)$  of  $G_c$ . Call  $f_c$  the new transition function. It is formally given by :

$$f_c : X \times \Sigma_o \rightarrow X$$

$$(x, \sigma) \mapsto \begin{cases} f(x, \sigma) & \text{if it is defined and } \sigma \notin \Sigma_d(G_c, x) \\ \text{undefined} & \text{otherwise} \end{cases}$$

Now, a supervisor  $S$  of  $G_c$  can be derived from  $M_D^o(G_c, x_0)$  by attaching to this FSM an output function  $O$  that for a given states  $x$  delivers the set of disabled events  $\Sigma_d(G_c, x)$ . The supervisor  $S = \langle \Sigma_o, X, x_0, x_m, f_c, O \rangle$  will in fact be used for two purposes. It first plays the role of an observer that is able to rebuild part of the state in which the system has evolved. Based on this information,  $S$  sends back to the system the set of events that have to be disabled in order to force the closed loop system to eventually reach the marked state  $q_m$  by minimizing the global cost of the trajectory.

## 5 Conclusion

In this paper, we have introduced a new type of optimal control for DESs by adding the notion of partial observation. The system to be controlled is represented by an FSM  $G$  with a unique marked state and some unobservable events. The first step was the derivation of a C-observer  $G_c$  from the partially unobservable FSM, which allows us to remember an approximation of the unobservable trajectory costs. We then presented a new definition of controllability derived from the classical one introduced by [7], that allows us to avoid the blocking of  $G$  without observing it. We then define the performance measure on this observer rather than on the FSM itself. In the second step, we first transform  $G_c$  into  $G'_c$  by shifting the cost of the macro-state to the events that can be executed in this macro-state. We then use the algorithm presented in [8] to synthesize an optimal submachine of the C-observer, which leads to the desired supervisor for the system. The behavior of the obtained controlled system is optimal w.r.t.  $\Sigma_o$ , in the sense that  $G_c$  carries on the best approximation of the unobservable trajectories. Moreover it is optimal for  $G'_c$  and therefore for  $G_c$ . This optimality status is due to [8].

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