PSPARQL Query Containment

Melisachew Wudage Chekol, Jérôme Euzenat, Pierre Genevès, Nabil Layaïda

INSTITUT NATIONAL DE RECHERCHE EN INFORMATIQUE ET EN AUTOMATIQUE

PSPARQL Query Containment

Melisachew Wudage Chekol — Jérôme Euzenat — Pierre Genevès — Nabil Layaïda

N° 7641

June 8, 2011

Thème COG

ISSN 0249-6399 ISRN INRIA/RR--7641--FR+ENG

INSTITUT NATIONAL
DE RECHERCHE
EN INFORMATIQUE
ET EN AUTOMATIQUE

INRIA

centre de recherche
GRENOBLE - RHÔNE-ALPES

PSPARQL Query Containment

Melisachew Wudage Chekol , Jérôme Euzenat , Pierre Genevès ,
Nabil Layaïda

Thème COG — Systèmes cognitifs
Équipes-Projets EXMO et WAM

Rapport de recherche n° 7641 — June 8, 2011 — 19 pages

Abstract: Querying the semantic web is mainly done through SPARQL. This language has been studied from different perspectives such as optimization and extension. One of its extensions, PSPARQL (Path SPARQL) provides queries with paths of arbitrary length. We study the static analysis of queries written in this language, in particular, containment of queries: determining whether, for any graph, the answers to a query are contained in those of another query. Our approach consists in encoding RDF graphs as transition systems and queries as μ-calculus formulas and then reducing the containment problem to testing satisfiability in the logic. We establish complexity bounds and report experimental results.

Key-words: Query Containment, PSPARQL, Semantic Web, RDF, Regular path queries

Centre de recherche INRIA Grenoble – Rhône-Alpes
655, avenue de l'Europe, 38334 Montbonnot Saint Ismier
Téléphone : +33 4 76 61 52 00 — Télécopie +33 4 76 61 52 52

Inclusion de Requêtes PSPARQL

Résumé : Pas de résumé

Mots-clés : Inclusion de requêtes, PSPARQL, web sémantique, RDF, requêtes

1 Introduction

Access to semantic web data expressed in Resource Description Framework (RDF) can be achieved through querying. Currently, querying RDF graphs is done mainly with the SPARQL query language. It has been a source of research from various perspectives mainly extending the language with new features and optimizing queries automatically. Querying RDF graphs with SPARQL amounts to matching graph patterns that are sets of triples of subjects, predicates and objects. These triples are usually connected to form graphs by means of joins expressed using several occurrences of the same variable. On the other hand, PSPARQL (Path SPARQL) allows querying of arbitrary length paths by using regular expression patterns. Regular path queries (RPQs) are useful for expressing complex navigations in a graph. In particular, union and transitive closure are crucial when one does not have a complete knowledge of the structure of the knowledge base. SPARQL 1.0 lacks recursion mechanism and supports a simple form of RPQs however its extensions such as PSPARQL [1] and its successor SPARQL1.1 support this feature.

Query optimization aims at improving the performance of query evaluation. Since queries in the semantic web are evaluated over huge RDF graphs, optimizations are crucial. Many studies contributed to query optimization using in particular the relational algebra from the database community [12]. This is very often achieved by using rules for rewriting queries into equivalent but faster ones. All these works, however, need at some point to prove the correctness of query optimization, i.e., the semantics of the optimized query remains the same as the original one. In other terms, the results of a given query are exactly the same as the optimized one regardless of the considered database. This can be reduced to query containment. Thus query containment plays a central role in database and knowledge base query optimization [12, 6, 3]. In addition, query containment can be of independent interest for performing other optimizations. For example, if a query q_1 is contained in q_2, then q_1 can be evaluated on the materialized view of q_2 rather then on the whole data graph.

Such approaches have also been applied to SPARQL [16], but not yet for PSPARQL.

We address the problem of static analysis of PSPARQL queries, encompassing satisfiability, containment and equivalence of queries. We introduce an approach which has already been successfully applied for XPath [8]. PSPARQL is interpreted over graphs, hence we encode it in a graph logic, specifically the alternation-free fragment of the μ-calculus [13] with converse and nominals [19] interpreted over labeled transition systems. We show that this logic is powerful enough to deal with query containment where queries are made of regular expression patterns which allow navigation through the graph. One benefit of using a μ-calculus encoding is to take advantage of fixpoints and modalities for encoding recursion. Furthermore, this logic admits exponential time decision procedures that can be implemented efficiently in practice [20, 8].

After presenting RDF, PSPARQL and the μ-calculus (§2), we show how to translate RDF graphs into transition systems (§3.1) and PSPARQL queries into μ-calculus formulas (§3.2). Therefore, query containment in PSPARQL can be reduced to unsatisfiability test in μ-calculus (§4).

2 Preliminaries

This section introduces the basics of RDF and PSPARQL.

2.1 RDF: Resource Description Framework

RDF is a language used to express structured information on the Web as graphs. Here we present a compact formalization of RDF [11]. Let U, B, and L be three disjoint infinite sets denoting the set of URIs (identify a resource), blank nodes (denote an unidentified resource) and literals (a character string or some other type of data) respectively. We abbreviate any union of these sets as for instance, UBL $= U \cup B \cup L$. A triple of the form $<s, p, o> \in$ UB $\times U \times$ UBL is called an *RDF triple*. s is the *subject*, p is the *predicate*, and o is the *object* of the triple. Each triple may be thought of as an edge between the subject and the object labelled by the predicate, hence a set of RDF triples is often referred to as an *RDF graph*.

Example 1 (RDF Graph)**.** *Here are 8 triples of an RDF graph about writers and their works: (all identifiers correspond to URIs, _:b is a blank node):*

Poe wrote thegoldbug .	*Baudelaire translated thegoldbug .*
Poe wrote theraven .	*Mallarmé translated theraven .*
theraven type Poem .	*Mallarmé wrote _:b .*
_:b type Poem .	*thegoldbug type Novel .*

RDF has a model theoretic semantics [11].

2.2 PSPARQL

PSPARQL (short for Path SPARQL) extends SPARQL with regular expression patterns. SPARQL [14] is a W3C recommended query language for RDF. PSPARQL overcomes the limitation of the current version of SPARQL which is the inability to express path queries. Before presenting the syntax and semantics of PSPARQL, let us briefly introduce the notion of regular expression patterns (cf. [1] for detailed discussion).

2.2.1 Regular Expressions

Regular expressions are patterns used to describe languages (i.e., sets of strings) from a given alphabet. Let $\Sigma = \{a_1, ..., a_n\}$ be an alphabet. A *string/word* is a finite sequence of symbols from the alphabet Σ. A *language* $\mathcal{L}$ is a set of words over Σ which is a subset of Σ^*, i.e, $\mathcal{L}(\Sigma) \subseteq \Sigma^*$. A word can be either empty ϵ or a sequence of alphabet symbols $a_1...a_n$. If $A = a_1...a_n$ and $B = b_1...b_m$ are two words over some alphabet Σ, then $A.B$ is a word over the same alphabet defined as: $A.B = a_1...a_n b_1...b_m$.

Definition 1 (Regular expression pattern)**.** *Given an alphabet Σ and a set of variables V, a regular expression $\mathcal{R}(\Sigma, V)$ can be constructed inductively as follows:*

$$e := uri \mid x \mid e_1 \shortmid e_2 \mid e_1.e_2 \mid e^+ \mid e^*$$

Where $e \in \mathcal{R}(\Sigma, V)$ and x denotes a variable, $e_1 \shortmid e_2$ denotes disjunction, $e_1.e_2$ denotes concatenation, e^+ denotes positive closure, and e^* denotes Kleene closure. Let U be a set of URIs and V a set of variables, a regular expression over $\mathcal{R}(U, V)$ can be used to define a language over the alphabet $U \cup V$.

2.2.2 PSPARQL Syntax

The only difference between the syntax of SPARQL and PSPARQL is on triple patterns. Triple patterns in PSPARQL contain regular expressions in property positions instead of only URIs or variables as it is the case of SPARQL. Queries are formed based on the notion of query patterns defined inductively from triple patterns: a tuple $t \in \text{UBV} \times \mathcal{R}(U, V) \times \text{UBLV}$, with V a set of variables disjoint from UBL, is called a triple pattern. Triple patterns grouped together using connectives ($\mathsf{AND}, \mathsf{UNION}, \mathsf{OPT}$) form *graph patterns* (a.k.a query patterns). We use an abstract syntax that can be easily translated into μ-calculus.

Definition 2 (Query Pattern)**.** *A PSPARQL query pattern q is inductively defined as follows :*

$$\begin{aligned} q = \;& t \in \text{UBV} \times \mathcal{R}(U, V) \times \text{UBLV} \\ & \mid q_1 \; \mathit{AND} \; q_2 \mid q_1 \; \mathit{UNION} \; q_2 \mid q_1 \; \mathit{OPT} \; q_2 \mid q_1 \; \mathsf{MINUS} \; Q_2 \end{aligned}$$

Where $\mathcal{R}(U, V)$ is a regular expression pattern defined over URIs U and query variables V.

Definition 3. *A PSPARQL SELECT query is a query of the form $q\{\overrightarrow{w}\}$ where $\overrightarrow{w}$ is a tuple of variables in V which are called distinguished variables, and q is a query pattern.*

Example 2 (PSPARQL queries)**.** *Consider the following queries $q_1\{?x\}$ and $q_2\{?x\}$ on the graph of Example 1:*

q_1

```
SELECT ?x
  WHERE {
      ?x    (translated | wrote) . type   Poem.
  }
```

q_2

```
SELECT ?x
   WHERE {
       { ?x  (translated  .  type)    Poem }
          UNION
       { ?x wrote  ?l .}
   }
```

2.2.3 PSPARQL Semantics

The semantics of PSPARQL queries is given by a partial mapping function $\rho : V \mapsto UBL$. The domain of ρ, $dom(\rho)$, is the subset of V on which ρ is defined. Two mappings ρ_1 and ρ_2 are said to be *compatible* if $\forall x \in dom(\rho_1) \cap dom(\rho_2)$, $\rho_1(x) = \rho_2(x)$. Hence, $\rho_1 \cup \rho_2$ is also a mapping. This allows for defining the

join, union, and difference operations between two sets of mappings M_1, and M_2 as shown below:

$$\begin{array}{lcl} M_1 \bowtie M_2 & = & \{\rho_1 \cup \rho_2 \mid \rho_1 \in M_1, \rho_2 \in M_2 \\ & & \quad \text{are compatible mappings } \} \\ M_1 \cup M_2 & = & \{\rho \mid \rho \in M_1 \text{ or } \rho \in M_2\} \\ M_1 \setminus M_2 & = & \{\rho \in M_1 \mid \forall \rho_1 \in M_2,\ \rho \text{ and } \rho_1 \\ & & \quad \text{are not compatible } \} \end{array}$$

Now, we are ready to define the evaluation of PSPARQL triple patterns recursively as follows:

$$\begin{aligned} [\![<x, uri, y>]\!]_G &= \{\rho \mid\ <\rho(x), \rho(uri), \rho(y)> \in G\} \\ [\![<x, z, y>]\!]_G &= \{\rho \mid\ <\rho(x), \rho(z), \rho(y)> \in G\} \\ [\![<x, e + e', y>]\!]_G &= [\![<x, e, y>]\!]_G \cup [\![<x, e', y>]\!]_G \\ [\![<x, e.e', y>]\!]_G &= [\![<x, e, n>]\!]_G \bowtie [\![<n, e', y>]\!]_G \\ [\![<x, e^+, y>]\!]_G &= \{\rho \mid \exists <n_0, e, n_1>, <n_1, e, n_2>, ..., \\ &\qquad <n_{k-1}, e, n_k> \in G \text{ such that } n_0 = \rho(x), \\ &\qquad n_k = \rho(y) \text{ and } e...e \in \mathcal{L}(e^+)\} \\ [\![<x, e^*, y>]\!]_G &= \{\rho \mid \rho(x) = \rho(y)\} \cup [\![<x, e^+, y>]\!]_G \end{aligned}$$

The evaluation of query patterns over an RDF graph G is inductively defined by:

$$\begin{aligned} [\![.]\!]_G &: q \rightarrow 2^{\mathrm{V} \times \mathrm{UBL}} \\ [\![q_1 \text{ AND } q_2]\!]_G &= [\![q_1]\!]_G \bowtie [\![q_2]\!]_G \\ [\![q_1 \text{ UNION } q_2]\!]_G &= [\![q_1]\!]_G \cup [\![q_2]\!]_G \\ [\![q_1 \text{ OPT } q_2]\!]_G &= ([\![q_1]\!]_G \bowtie [\![q_2]\!]_G) \cup ([\![q_1]\!]_G \setminus [\![q_2]\!]_G) \\ [\![q_1 \text{ MINUS } q_2]\!]_G &= [\![q_1]\!]_G \setminus [\![q_2]\!]_G \\ [\![q\{\overrightarrow{w}\}]\!]_G &= \pi_{\overrightarrow{w}}([\![q]\!]_G) \end{aligned}$$

Where the projection operator $\pi_{\overrightarrow{w}}$ selects only those part of the mappings relevant to variables in $\overrightarrow{w}$.

Example 3 (Answers to SPARQL queries)**.** *The answers to query q_1 and q_2 of Example 2 on graph G of Example 1 are respectively $\{Poe, Mallarme\}$ and $\{Baudelaire, Poe, Mallarme\}$. Hence, $[\![q_1]\!]_G \subseteq [\![q_2]\!]_G$.*

Beyond this particular example, the goal of query containment is to determine whether this holds for any graph.

Definition 4 (Containment)**.** *Given queries q_1 and q_2 with the same arity, q_1 is contained in q_2, denoted $q_1 \sqsubseteq q_2$, iff for any graph G, $[\![q_1]\!]_G \subseteq [\![q_2]\!]_G$.*

Definition 5 (Equivalence)**.** *Two queries q_1 and q_2 are equivalent, $q_1 \equiv q_2$, iff $q_1 \sqsubseteq q_2$ and $q_2 \sqsubseteq q_1$.*

3 Encodings

In this section, encodings of RDF graphs as transition systems, and regular expressions and PSPARQL queries as μ-calculus formulas are explained.

3.1 Encoding RDF graphs as Transition Systems

Before presenting the encoding of RDF graphs as transition systems over which the μ-calculus is interpreted, we introduce the syntax and semantics of the μ-calculus.

3.1.1 μ-calculus

The modal μ-calculus [13] is an expressive logic which adds recursive features to modal logic using fixpoint operators.

The syntax of the μ-calculus is composed of countable sets of *atomic propositions* AP, a set of *nominals* Nom, a set of *variables* Var, a set of *programs* $Prog$ for navigating in graphs. A μ-calculus formula, φ, can be defined inductively as follows:

$$\begin{aligned}\varphi \quad ::= \quad & \top \mid \bot \mid p \mid X \mid \neg\varphi \mid \varphi \vee \psi \mid \\ & \varphi \wedge \psi \mid \langle a \rangle \varphi \mid [a]\varphi \mid \mu X \varphi \mid \nu X \varphi\end{aligned}$$

where $p \in AP \cup Nom, X \in Var$ and $a \in Prog$ is either an atomic program or its converse $\bar{a}$. The greatest and least fixpoint operators (ν and μ) respectively introduce general and finite recursion in graphs [13].

The semantics of the μ-calculus is given over a transition system, $K = (S, R, L)$ where S is a non-empty set of nodes, $R : Prog \to 2^{S \times S}$ is the transition function, and $L : AP \to 2^S$ assigns a set of nodes to each atomic proposition or nominal where it holds, such that $L(p)$ is a *singleton* for each nominal p. For converse programs, R can be extended as $R(\bar{a}) = \{(s', s) \mid (s, s') \in R(a)\}$. Besides, a valuation function $V : \text{Var} \to S$ is used to assign a set of nodes to each variable. For a valuation V, variable X, and a set of nodes $S' \subseteq S$, $V[X/S']$ is the valuation that is obtained from V by assigning S' to X. The semantics of a formula in terms of a transition system (a.k.a. Kripke structure) and a valuation function is represented by $[\![\varphi]\!]^K_V$. The semantics of basic μ-calculus formula is defined as follows:

$$\begin{aligned}
[\![\mathtt{p}]\!]^K_V &= L(p), p \in AP \cup Nom \\
[\![\mathtt{X}]\!]^K_V &= V(X), X \in Var \\
[\![\neg\varphi]\!]^K_V &= S \backslash [\![\varphi]\!]^K_V \\
[\![\varphi \wedge \psi]\!]^K_V &= [\![\varphi]\!]^K_V \cap [\![\psi]\!]^K_V \\
[\![\varphi \vee \psi]\!]^K_V &= [\![\varphi]\!]^K_V \cup [\![\psi]\!]^K_V \\
[\![\langle \mathtt{a} \rangle \varphi]\!]^K_V &= \{s \in S | \exists s' \in S . (s, s') \in R(a) \ \wedge s' \in [\![\varphi]\!]^K_V\} \\
[\![[\mathtt{a}]\varphi]\!]^K_V &= \{s \in S | \forall s' \in S . (s, s') \in R(a) \Rightarrow s' \in [\![\varphi]\!]^K_V\} \\
[\![\mu \mathtt{X} \varphi]\!]^K_V &= \bigcap \{S' \subseteq S | [\![\varphi]\!]^K_{V[X/S']} \subseteq S'\} \\
[\![\nu \mathtt{X} \varphi]\!]^K_V &= \bigcup \{S' \subseteq S | S' \subseteq [\![\varphi]\!]^K_{V[X/S']}\}
\end{aligned}$$

Once providing the syntax and semantics of μ-calculus, the next subsections introduce representation of RDF graphs as transition systems and queries as formulas.

3.1.2 Encoding of RDF graphs

An RDF graph is encoded as a transition system in which nodes correspond to RDF entities and RDF triples. Edges relate entities to the triples they occur in. Different edges are used for distinguishing the functions (subject, object, predicate). expressing predicates as nodes, instead of atomic programs, makes it possible to deal with full RDF expressiveness in which a predicate may also be the subject or object of a statement.

Definition 6 (Transition system for RDF graph)**.** *Given an RDF graph,* $G \subseteq \text{UB} \times U \times \text{UBL}$*, the transition system associated* G*,* $\sigma(G) = (S, R, L)$ *over* $AP = \text{UBL} \cup \{s', s''\}$*, is such that:*

- $S = S' \cup S''$ *with* S' *and* S'' *the smallest sets such that* $\forall u \in U_G, \exists n^u \in S'$*,* $\forall b \in B_G, \exists n^b \in S'$*,* $\forall l \in L_G, \exists n^l \in S'$ *and* $\forall t \in G, \exists n^t \in S''$*,*
- $\forall t =<s, p, o>\in G$*,* $\langle n^s, n^t\rangle \in R(s)$*,* $\langle n^t, n^p\rangle \in R(p)$*, and* $\langle n^t, n^o\rangle \in R(o)$*,*
- $L : \text{UBL} \to 2^S$*;* $\forall u \in U_G, L(u) = \{n^u\}$*,* $\forall b \in B_G, L(b) = S'$*,* $L(s') = S'$*,* $\forall l \in L_G, L(l) = \{n^l\}$ *and* $L(s'') = S''$*,*
- $\forall n^t, n^{t'} \in S''$*,* $\langle n^t, n^{t'}\rangle \in R(d)$*.*

The program d is introduced to render each triple accessible to the others and thus facilitate the encoding of queries. The function σ associates what we call a *restricted transition system* to any RDF graph. Formally, we say that a transition system K is a *restricted transition system* iff there exists an RDF graph G such that $K = \sigma(G)$.

A restricted transition system is thus a bipartite graph composed of two sets of nodes: S', those corresponding to RDF entities, and S'', those corresponding to RDF triples. For example, Figure 1 shows the restricted transition system associated with the graph of Example 1.

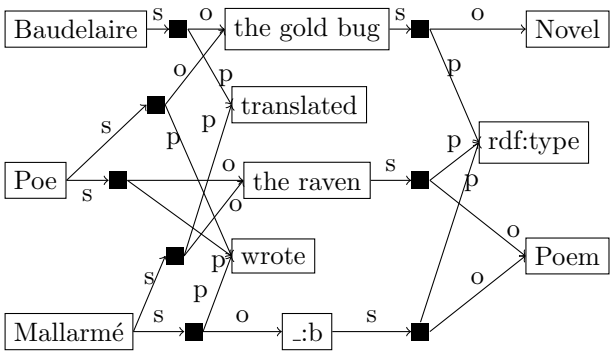

Figure 1: Transition system encoding the RDF graph of Example 1. Nodes in S'' are black anonymous nodes; nodes in S' are the other nodes (d-transitions are not displayed).

When checking for query containment, we consider the following restrictions:

- The set of programs is fixed: $Prog = \{s, p, o, d, \bar{s}, \bar{p}, \bar{o}, \bar{d}\}$ (note that $d = \bar{d}$).

- A model must be a restricted transition system.

This last constraint can be expressed in the μ-calculus as follows:

Proposition 1 (RDF restriction on transition systems)**.** *A formula φ is satisfied by some restricted transition system if and only if $\varphi \wedge \varphi_r$ is satisfiable by some transition system, i.e. $\exists K_r [\![\varphi]\!]^{K_r} \neq \emptyset \iff \exists K [\![\varphi \wedge \varphi_r]\!]^{K} \neq \emptyset$, where:*

$$\varphi_r = \quad \nu X.\theta \wedge \kappa \ \wedge (\neg\langle d\rangle\top \vee \langle d\rangle X)$$

in which $\theta = \langle\bar{s}\rangle s' \wedge \langle p\rangle s' \wedge \langle o\rangle s' \wedge \neg\langle s\rangle\top \wedge \neg\langle\bar{p}\rangle\top \wedge \neg\langle\bar{o}\rangle\top$ and $\kappa = [\bar{s}]\xi \wedge [p]\xi \wedge [o]\xi$ with

$$\xi = \begin{cases} \neg\langle\bar{s}\rangle\top \wedge \neg\langle o\rangle\top \wedge \neg\langle p\rangle\top \wedge \neg\langle d\rangle\top \wedge \neg\langle\bar{d}\rangle\top \\ \wedge\neg\langle s\rangle s' \wedge \neg\langle\bar{o}\rangle s' \wedge \neg\langle\bar{p}\rangle s'. \end{cases}$$

The formula φ_r ensures that θ and κ hold in every node reachable by a d edge, i.e. in every S'' node. The formula θ forces each S'' node to have a subject, predicate and object. The formula κ navigates from a s'' node to every reachable s' node, and forces the latter not to be directly connected to other subject, predicate or object nodes.

Proof. ($\Rightarrow$) Assume that $\exists K_r [\![\varphi]\!]^{K_r} \neq \emptyset$, since φ_r is satisfied by only transition systems, one gets $[\![\varphi_r]\!]^{k_r} \neq \emptyset$. Hence it follows that, $\exists K_r [\![\varphi]\!]^{K_r} \neq \emptyset$ and $[\![\varphi_r]\!]^{k_r} \neq \emptyset$ which implies $\exists K_r [\![\varphi]\!]^{k_r} \wedge [\![\varphi_r]\!]^{K_r} \neq \emptyset$. From this, using the semantics of μ-calculus formula, one obtains $\exists K_r [\![\varphi \wedge \varphi_r]\!]^{K_r} \neq \emptyset$. Since a restricted transition system is also a transition system, $K_r \subseteq K$, it follows that $\exists K.[\![\varphi \wedge \varphi_r]\!]^{K} \neq \emptyset$.
($\Leftarrow$) Assume that $\exists K [\![\varphi \wedge \varphi_r]\!]^{K} \neq \emptyset$. We construct a restricted transition system model $K_r = (S_r, R_r, L_r)$ and a function $f : K_r \to K$ from $K = (S, R, L)$. Add a node n'_0 to S_r with $f(n'_0) = n_0$ where $\varphi \wedge \varphi_r$ is satisfied in K. Supppose we have constructed a node n_r of S_r. For $j \in \{s, p, o\}$, if there is $n \in S$ with $(f(n_r), n) \in R(j)$, then pick one such n and add a node n'_r to S_r with $f(n'_r) = n$. Finally, for an atomic proposition p, $L_r(p) = \{n_r \in S_r \mid f(n_r) \in L(p)\}$.
The RDF triple structure is maintained in K_r i.e.
$\langle(s, s''), (s'', p), (s'', o)\rangle$ is valid through out the graph. If there were node pairs outside of this structure, then φ_r will not be satisfied. Throught out the graph, θ and κ ensure that for each triple node $s'' \in S_r$, there exists exactly one incoming subject edge, one outgoing property edge, and one outgoing object edge. Hence, $[\![\varphi_r]\!]^{K_r} \neq \emptyset$
To verify that $[\![\varphi]\!]^{K_r} \neq \emptyset$, it is enough to show $[\![\varphi]\!]^{K} \Rightarrow [\![\varphi]\!]^{K_r}$ by induction on the structure of φ. □

3.2 Encoding PSPARQL Queries as μ-calculus Formulas

Queries are translated to μ-calculus formulas. The principle of the translation is that each triple pattern is associated with a sub-formula stating the existence of the triple somewhere in the graph. Hence, they are quantified by μ so as to put them out of the context of a state. In this translation, variables are replaced by nominals which will be satisfied when they are matched in such triple relations. For that purpose, we use a function $\lambda : VUBL \to UBL$ such that:

$$\lambda(x) = \begin{cases} v_x & \text{if } x \in V \\ x & \text{if } x \in \text{UBL} \end{cases}$$

The function $\mathcal{A}$ encodes queries inductively on the structure of query patterns. AND and UNION are replaced by boolean connectives $\wedge$ and $\vee$ respectively. The MINUS operator is translated as $\wedge$ and $\neg$. OPT queries carry implicit negation in that they can be expressed as a logic formula in the following form: q_1 OPT $q_2 = (q_1 \wedge q_2) \vee (q_1 \wedge \neg q_2)$. Unfortunately, this formula can be reduced to just q_1 which is not the intended semantics of the operator. Hence, we need another approach in order to correctly encode this operator. To do so, we rely on the interpretation given below:

$$q_1 \text{ OPT } q_2 = \begin{cases} q_1 \text{ AND } q_2 & \text{if } \rho(q_2) \in G \\ q_1 & \text{if } \rho(q_2) \notin G \end{cases}$$

The above interpretation of OPT operator dictates that: q_1 OPT q_2 evaluates as q_1 AND q_2 if there exists a mapping ρ for q_2, otherwise it evaluates as q_1. Based on this, the μ-calculus encoding of the operator can be obtained. The formula $ew(q_1 \text{ AND } q_2)$ evaluates to all nodes S if q_2 exists in the graph, it evaluates to $\emptyset$ otherwise. Further, the function f translates query patterns into formulas recursively.

Definition 7. *The encoding of a PSPARQL query pattern q is $\mathcal{A}(q)$ such that:*

$$\begin{aligned}
\mathcal{A}(<x,e,z>) &= \mu X.(\langle\bar{s}\rangle\lambda(x) \wedge \mathcal{R}(\lambda(e),\lambda(z))) \\
&\qquad \vee \langle d\rangle X \vee \langle s\rangle X \vee \langle\bar{p}\rangle X \vee \langle\bar{o}\rangle X \\
\mathcal{A}(q_1 \textit{ AND } q_2) &= \mathcal{A}(q_1) \wedge \mathcal{A}(q_2) \\
\mathcal{A}(q_1 \textit{ UNION } q_2) &= \mathcal{A}(q_1) \vee \mathcal{A}(q_2) \\
\mathcal{A}(q_1 \textit{ MINUS } q_2) &= \mathcal{A}(q_1) \wedge \neg\mathcal{A}(q_2) \\
\mathcal{A}(q_1 \textit{ OPT } q_2) &= ew(f(q_1) \wedge f(q_2)) \wedge \mathcal{A}(q_1 \textit{ AND } q_2) \vee \\
&\qquad ew(f(q_1) \wedge \neg f(q_2)) \wedge \mathcal{A}(q_1) \\
ew(\varphi) &= \mu X.\varphi \vee \langle d\rangle X \vee \langle s\rangle X \vee \langle\bar{p}\rangle X \vee \langle\bar{o}\rangle X \\
f(<x,e,z>) &= \langle\bar{s}\rangle\lambda(x) \wedge \mathcal{R}(\lambda(e),\lambda(z)) \\
f(q_1 \textit{ AND } q_2) &= f(q_1) \wedge f(q_2) \\
f(q_1 \textit{ UNION } q_2) &= f(q_1) \vee f(q_2) \\
f(q_1 \textit{ OPT } q_2) &= f(q_1)
\end{aligned}$$

Definition 7 introduces regular expression encoding function $\mathcal{R}$ that takes two arguments (the predicate which is a regular expression pattern, and the object of a triple).This function is inductively defined as follows:

Definition 8. *Regular expressions are encoded recursively using the function $\mathcal{R}$, detailed below:*

$$\begin{aligned}
\mathcal{R}(uri,y) &= \langle p\rangle uri \ \wedge \langle o\rangle y \\
\mathcal{R}(x,y) &= \langle p\rangle x \wedge \langle o\rangle y \\
\mathcal{R}(e_1 \shortmid e_2,y) &= (\mathcal{R}(e_1,y) \vee \mathcal{R}(e_2,y)) \\
\mathcal{R}(e_1.e_2,y) &= \mathcal{R}(e_1,\langle s\rangle\mathcal{R}(e_2,y)) \\
\mathcal{R}(e^+,y) &= \mu X.\mathcal{R}(e,y) \vee \mathcal{R}(e,\langle s\rangle X) \\
\mathcal{R}(e^*,y) &= \mathcal{R}(e^+,y) \vee \langle\bar{s}\rangle y
\end{aligned}$$

Example 4. *This example shows an encoding of query 1 of Example 2 as a μ-calculus formula.*

$$\begin{aligned}\mathcal{A}(q_2) = \mu X.(&\langle\bar{s}\rangle\lambda(x) \wedge \langle p\rangle translated \wedge \langle o\rangle\langle s\rangle(\langle p\rangle type \\ &\wedge \langle o\rangle Poem)) \vee \langle d\rangle X \vee \langle s\rangle X \vee \langle\bar{p}\rangle X \vee \langle\bar{o}\rangle X \vee \\ \mu X.(&\langle\bar{s}\rangle\lambda(x) \wedge \langle p\rangle wrote \wedge \langle o\rangle l) \\ &\vee \langle d\rangle X \vee \langle s\rangle X \vee \langle\bar{p}\rangle X \vee \langle\bar{o}\rangle X\end{aligned}$$

4 Reducing Query Containment to Unsatisfiability

In this section, we address the problem of query containment, $q_1\{\overrightarrow{w}\} \sqsubseteq q_2\{\overrightarrow{w}\}$, by reducing it to the problem of unsatisfiability in the logic. The first theorem expressed the correctness and completeness of the encodings.

Theorem 2. *For any graph G and PSPARQL query $q\{\vec{w}\}$,*
$\forall\rho.(\rho \in [\![q\{\vec{w}\}]\!]_G \iff [\![\mathcal{A}(\rho(q\{\vec{w}\}))]\!]^{\sigma(G)} \neq \emptyset)$

Proof. This is proved inductively:
(Base case) The base case is proved for triple patterns containing regular expression patterns of the form: $y \mid e_1.e_2 \mid e^+$. First, when $q\{x,y,z\} = <x,y,z>$.
$\forall G.\forall\rho.(\rho \in [\![<x,y,z>]\!]_G \iff [\![\mathcal{A}(\rho(<x,y,z>))]\!]^{\sigma(G)} \neq \emptyset)$
($\Rightarrow$) If $\rho \in [\![<x,y,z>]\!]_G$, then $<\rho(x),\rho(y),\rho(z)> \in G$. Hence $\sigma(G) = (S,R,L)$ contains:

- $t \in S'', n^{\rho(x)}, n^{\rho(y)}, n^{\rho(z)} \in S'$,
- $(n^{\rho(x)}, t) \in R(s), (t, n^{\rho(y)}) \in R(p), (t, n^{\rho(z)}) \in R(o)$, and
- $L(\rho(x)) = n^{\rho(x)}, L(\rho(y)) = n^{\rho(y)}, L(\rho(z)) = n^{\rho(z)}$.

$<\rho(x),\rho(y),\rho(z)>$ can be encoded as a μ-calculus formula. This encoding when evaluated over the transition system $\sigma(G)$ is non empty because if the triple exists in G, it also exists in the transition system. Consequently,

$\Rightarrow$ $[\![\mu X.(\langle\bar{s}\rangle\lambda(\rho(x)) \wedge \langle p\rangle\lambda(\rho(y)) \wedge \langle o\rangle\lambda(\rho(z)))$
$\vee \langle d\rangle X \vee \langle s\rangle X \vee \langle\bar{p}\rangle X \vee \langle\bar{o}\rangle X]\!]^{\sigma(G)} \neq \emptyset$
$\Rightarrow$ $[\![\mathcal{A}(<\rho(x),\rho(y),\rho(z)>)]\!]^{\sigma(G)} \neq \emptyset$
$\Rightarrow$ $[\![\mathcal{A}(\rho(<x,y,z>))]\!]^{\sigma(G)} \neq \emptyset$

($\Leftarrow$) $[\![\mathcal{A}(\rho(<x,y,z>))]\!]^{\sigma(G)} \neq \emptyset$ entails that there is a state $n^t \in S''$ and $n^{\rho(x)}$, $n^{\rho(y)}$, $n^{\rho(z)} \in S'$, such that $\langle n^{\rho(x)}, n_t\rangle \in R(s)$, $\langle n_t, n^{\rho(y)}\rangle \in R(p)$, and $\langle n_t, n^{\rho(z)}\rangle \in R(o)$ and $n^{\rho(x)} \in L(\lambda(\rho(x)))$, $n^{\rho(y)} \in L(\lambda(\rho(y)))$ and $n^{\rho(z)} \in L(\lambda(\rho(z)))$. Since the transition system $\sigma(G)$ is the encoding of an RDF graph G, this means that $<\lambda(\rho(x)),\lambda(\rho(y)),\lambda(\rho(z))> \in G$. Subsequently, $[\![<x,y,z>]\!]_G \neq \emptyset$, thus there exists a mapping ρ such that $\rho \in [\![<x,y,z>]\!]_G$.

• When $q\{x,y\} = <x,e_1.e_2,y>$.
$\forall G.\forall\rho.(\rho \in [\![<x,e_1.e_2,y>]\!]_G \Leftrightarrow [\![\mathcal{A}(\rho(<x,e_1.e_2,y>))]\!]^{\sigma(G)} \neq \emptyset)$
Assume that, $\forall G.\forall\rho.(\rho \in [\![<x,e_1,z>]\!]_G \Leftrightarrow [\![\mathcal{A}(\rho(<x,e_1,z>))]\!]^{\sigma(G)} \neq \emptyset)$ and $\forall G.\forall\rho.(\rho \in [\![<z,e_2,y>]\!]_G \Leftrightarrow [\![\mathcal{A}(\rho(<z,e_2,y>))]\!]^{\sigma(G)} \neq \emptyset)$

$\rho \in [\![<x, e_1.e_2, y>]\!]_G$

$\Leftrightarrow \rho \in ([\![<x, e_1, z>]\!]_G \bowtie [\![<z, e_2, y>]\!]_G)$

$\Leftrightarrow \rho \in [\![<x, e_1, z>]\!]_G$ and $\rho \in [\![<z, e_2, y>]\!]_G$

$\Leftrightarrow [\![\mathcal{A}(\rho(<x, e_1, z>))]\!]^{\sigma(G)} \neq \emptyset$ and $[\![\mathcal{A}(\rho(<z, e_2, y>))]\!]^{\sigma(G)} \neq \emptyset$ by the assumption

$\Leftrightarrow [\![\mu X.(\langle \bar{s} \rangle \lambda(\rho(x)) \wedge \langle p \rangle \lambda(\rho(e_1)) \wedge \langle o \rangle \lambda(\rho(z))) \vee \langle d \rangle X \vee \langle s \rangle X \vee \langle \bar{p} \rangle X \vee \langle \bar{o} \rangle X]\!]^{\sigma(G)} \neq \emptyset$
and
$[\![\mu X.(\langle \bar{s} \rangle \lambda(\rho(z)) \wedge \langle p \rangle \lambda(\rho(e_2)) \wedge \langle o \rangle \lambda(\rho(y))) \vee \langle d \rangle X \vee \langle s \rangle X \vee \langle \bar{p} \rangle X \vee \langle \bar{o} \rangle X]\!]^{\sigma(G)} \neq \emptyset$

$\Leftrightarrow [\![\mu X.(\langle \bar{s} \rangle \lambda(\rho(x)) \wedge \langle p \rangle \lambda(\rho(e_1)) \wedge \langle o \rangle \langle s \rangle (\langle p \rangle \lambda(\rho(e_2)) \wedge \langle o \rangle \lambda(\rho(y)))) \vee$
$\langle d \rangle X \vee \langle s \rangle X \vee \langle \bar{p} \rangle X \vee \langle \bar{o} \rangle X]\!]^{\sigma(G)} \neq \emptyset$ *

$\Leftrightarrow [\![\mathcal{A}(<\rho(x), \rho(e_1).\rho(e_2), \rho(y)>)]\!]^{\sigma(G)} \neq \emptyset$

$\Leftrightarrow [\![\mathcal{A}(\rho(<x, e_1.e_2, y>))]\!]^{\sigma(G)} \neq \emptyset$

* the two conjunct formulas are joined and remain satisfiable because of the nominal $\lambda(\rho(z))$ that appears in each of the recursive formulas.

• When the triple pattern contains positive concatenation, $q\{x, y\} = <x, e^+, y>$.

$\rho \in [\![<x, e^+, y>]\!]_G$

$\Leftrightarrow \exists r_1, ..., r_k. <\rho(x), \rho(e), \rho(r_1)> \in G$ *and* $<\rho(r_1), \rho(e), \rho(r_2)> \in G$
and ... and $<\rho(r_k), \rho(e), \rho(y)> \in G$

$\Leftrightarrow \rho \in [\![<x, e, r_1>]\!]_G$ and $\rho \in [\![<r_1, e, r_2>]\!]_G$ and ... and $\rho \in [\![<r_k, e, y>]\!]_G$

$\Leftrightarrow [\![\mathcal{A}(\rho(<x, e, r_1>))]\!]^{\sigma(G)} \neq \emptyset$ and $[\![\mathcal{A}(\rho(<r_1, e, r_2>))]\!]^{\sigma(G)} \neq \emptyset$
and ... and $[\![\mathcal{A}(\rho(<r_k, e, y>))]\!]^{\sigma(G)} \neq \emptyset$

$\Leftrightarrow [\![\mathcal{A}(\rho(<x, e, r_1>)) \wedge \mathcal{A}(\rho(<r_1, e, r_2>)) \wedge ... \wedge \mathcal{A}(\rho(<r_k, e, y>))]\!]^{\sigma(G)} \neq \emptyset$

$\Leftrightarrow [\![\mathcal{A}(\rho(<x, e, r_1> \wedge <r_1, e, r_2> \wedge ... \wedge <r_k, e, y>))]\!]^{\sigma(G)} \neq \emptyset$ **

$\Leftrightarrow [\![\mathcal{A}(\rho(<x, e^+, y>)]\!]^{\sigma(G)} \neq \emptyset$

** the same argument as that of, the proof for $<x, e_1.e_2, y>$, is used for joining separate satisfiable formulas.

(Inductive case) Query patterns: q_1 AND q_2 | q_1 UNION q_2 | q_1 OPT q_2 | q_1 MINUS q_2. We provide the transcriptions of AND and OPT. The proof of UNION and MINUS follows similarly. First, consider when $q\{\vec{w}\} = q_1$ AND q_2.

$\rho \in [\![q_1 \text{ AND } q_2]\!]_G$

$\Leftrightarrow \rho \in [\![q_1]\!]_G$ and $\rho \in [\![q_2]\!]_G$

$\Leftrightarrow [\![\mathcal{A}(\rho(q_1))]\!]^{\sigma(G)} \neq \emptyset$ and $[\![\mathcal{A}(\rho(q_2))]\!]^{\sigma(G)} \neq \emptyset$ by induction hypothesis.

$\Leftrightarrow [\![\mathcal{A}(\rho(q_1)) \wedge \mathcal{A}(\rho(q_2))]\!]^{\sigma(G)} \neq \emptyset$ *

$\Leftrightarrow [\![\mathcal{A}(\rho(q_1) \text{ AND } \rho(q_2))]\!]^{\sigma(G)} \neq \emptyset$

$\Leftrightarrow [\![\mathcal{A}(\rho(q_1 \text{ AND } q_2))]\!]^{\sigma(G)} \neq \emptyset$

* this formula remains satisfiable because there exists ϕ a satisfiable subformula of $\mathcal{A}(\rho(q_1))$ and $\mathcal{A}(\rho(q_2))$ in $\sigma(G)$. In fact, ϕ is a nominal obtained by encoding a variable which is common to both q_1 and q_2.

Inductive case for OPT i.e., when $q\{\vec{w}\} = q_1 \text{ OPT } q_2$.
$\rho \in [\![q_1 \text{ OPT } q_2]\!]_G$

$\Leftrightarrow \rho \in ([\![q_1]\!]_G \bowtie [\![q_2]\!]_G)$ or $\rho \in ([\![q_1]\!]_G \setminus [\![q_2]\!]_G)$
$\Leftrightarrow \rho \in ([\![q_1]\!]_G \bowtie [\![q_2]\!]_G)$ if $\rho(q_2) \in G$ or
$\rho \in [\![q_1]\!]_G$ if $\rho(q_2) \notin G$
$\Leftrightarrow [\![\mathcal{A}(\rho(q_1 \text{ AND } q_2))]\!]^{\sigma(G)} \neq \emptyset$ if $\rho(q_2) \in G$ or
$[\![\mathcal{A}(\rho(q_1))]\!]^{\sigma(G)} \neq \emptyset$ if $\rho(q_2) \notin G$ by induction hypothesis.
$\Leftrightarrow [\![\mathcal{A}(\rho(q_1 \text{ AND } q_2))]\!]^{\sigma(G)} \neq \emptyset \cap [\![ew(q_1 \text{ AND } q_2)]\!]^{\sigma(G)} = S \cup$
$[\![\mathcal{A}(\rho(q_1))]\!]^{\sigma(G)} \neq \emptyset \cap [\![ew(q_1 \text{ AND } \neg q_2)]\!]^{\sigma(G)} = S$ *
$\Leftrightarrow [\![\mathcal{A}(\rho(q_1 \text{ AND } q_2)) \wedge ew(q_1 \text{ AND } q_2)]\!]^{\sigma(G)} \neq \emptyset \cup$
$[\![\mathcal{A}(\rho(q_1)) \wedge ew(q_1 \text{ AND } \neg q_2)]\!]^{\sigma(G)} \neq \emptyset$
$\Leftrightarrow [\![\mathcal{A}(\rho(q_1 \text{ OPT } q_2))]\!]^{\sigma(G)} \neq \emptyset$

* here we added a formula which evaluates to the entire set of states S if both $q_1$and q_2 are found in the transition system. Hence, the first part of the disjunction evaluates to a non empty result whereas if q_2 does not exist in the transition system, the second part of the disjunction is non empty. Thereby, retaining the semantics of the OPT operator. □

This theorem is the key to reduce query containment to satisfiability. For a proper translation of the query encodings, we use a variable renaming function $\phi_q^{\vec{w}}$ which renames all variables in q, but the distinguished variables in $\vec{w}$, to independent variables.

Theorem 3. *Given PSPARQL queries $q_1\{\vec{w}\}$ and $q_2\{\vec{w}\}$, $q_1\{\vec{w}\} \sqsubseteq q_2\{\vec{w}\}$ iff $\mathcal{A}(q_1) \wedge \neg\phi_{q_1}^{\vec{w}}(\mathcal{A}(q_2)) \wedge \varphi_r$ is unsatisfiable.*

Proof. It can be proved as follows:
$q_1\{\vec{w}\} \sqsubseteq q_2\{\vec{w}\}$

$\Leftrightarrow \forall G.[\![q_1\{\vec{w}\}]\!]_G \subseteq [\![q_2\{\vec{w}\}]\!]_G$
$\Leftrightarrow \forall G.\forall\rho.\ (\rho \in [\![q_1]\!]_G \Rightarrow \rho \in [\![q_2]\!]_G)$
$\Leftrightarrow \forall G.\forall\rho.([\![\mathcal{A}(\rho(q_1))]\!]^{\sigma(G)} \neq \emptyset) \Rightarrow ([\![\mathcal{A}(\rho(q_2))]\!]^{\sigma(G)} \neq \emptyset)$
by Theorem 2
$\Leftrightarrow \forall G.\forall\rho.([\![\mathcal{A}(\rho(q_1))]\!]^{\sigma(G)} \neq \emptyset) \Rightarrow ([\![\phi_{q_1}^w(\mathcal{A}(\rho(q_2)))]\!]^{\sigma(G)} \neq \emptyset)$
by transparent renaming
$\Leftrightarrow \forall G.\forall\rho[\![\neg\mathcal{A}(\rho(q_1)) \vee \phi_{q_1}^w(\mathcal{A}(\rho(q_2)))]\!]^{\sigma(G)} \neq \emptyset$
$\Leftrightarrow \forall G.\forall\rho[\![\mathcal{A}(\rho(q_1)) \wedge \neg\phi_{q_1}^w(\mathcal{A}(\rho(q_2)))]\!]^{\sigma(G)} = \emptyset$
$\Leftrightarrow \forall G.[\![\mathcal{A}(q_1) \wedge \neg\phi_{q_1}^w(\mathcal{A}(q_2))]\!]^{\sigma(G)} = \emptyset$ *
$\Leftrightarrow \forall K.[\![\mathcal{A}(q_1) \wedge \neg\phi_{q_1}^w(\mathcal{A}(q_2)) \wedge \varphi_r]\!]^K = \emptyset$ by Proposition 1
$\Leftrightarrow \mathcal{A}(q_1) \wedge \neg\phi_{q_1}^w(\mathcal{A}(q_2)) \wedge \varphi_r$ unsatisfiable

□

(*) From Theorem 2, it follows that if there exists a set of mappings for an evaluation of a query over a graph, then the encoding of the query over the transition system obtained from the graph is satisfiable.

4.1 Complexity

The complexity of testing query containment through this reduction depends on the *lean* of the reduced formula. Thus, duplicate formulas that appear in the encoding for OPT does not affect the complexity of the reduction. The *lean* of a formula φ is a subset of $cl(\varphi)$ defined by $\{\psi \in cl(\varphi) \mid \psi \in AP \cup Nom$ or ψ is of theform $\langle\alpha\rangle\phi$ or $[\alpha]\phi\}$. Where $cl(\varphi)$ is the *Fisher-Ladner closure* of a formula φ as a set of all subformulas of it where fixpoint formulas are unwound once [8, 20].

Proposition 4. *Query containment can be solved in a time of* $2^{\mathcal{O}(n^2 log\ n)}$ *where* $n = |Lean(|q_1| + |q_2|)|$ *is the size of the lean, and* $|q_1|$ *and* $|q_2|$ *denote the sizes of queries* q_1 *and* q_2.

Note that this EXPTIME complexity is the complexity of satisfiability test for a μ-calculus formula with backward modalities and nominals [15, 19].

4.2 Experimentation

All experiments were conducted under Linux Ubuntu v10.04 LTS Lucid Lynx. The machine has an Intel Core2 Duo T7250 2.00GHz CPU and 2GB physical memory. The Java programs were executed with JRE v1.6.0_22. Further, the μ-calculus satisfiability solver from [20] is used to test containment. Table 1 shows the running times of containment tests for queries that are listed below:

- q_1– select all nodes connected to s through p.

```
SELECT * WHERE {
    s p ?o .
}
```

- q_2 – select all nodes connected to s through a sequence of p paths.

q_2

```
SELECT * WHERE {
    s p+ ?o .
}
```

- q_3 – select all instances of o.

q_3

```
SELECT * WHERE {
    ?s type o .
}
```

- q_4 – select all instances of o by eloring all its subclasses.

q_4

```
SELECT * WHERE {
    ?s type.sc* ?o
}
```

- q_5 – select all instances of o by loiting all RDF Schema assertions (triples).

q_5
```
SELECT * WHERE
{
   { ?s type.sc* o . }
        UNION
   { ?s ?p1 ?y . ?p1 sp* ?p2 . ?p2 domain.sc* o . }
       UNION
   { ?y ?p1 ?s . ?p1 sp* ?p2 . ?p2 range.sc* o . }
}
```

- q_6 – select all those authors who translated or wrote a Poem.

q_6
```
 SELECT ?x
  WHERE {
       ?x    (translated | wrote) . type   Poem.
  }
```

- q_7 – select all those authors who translated Poem and wrote something else.

q_7
```
SELECT ?x
  WHERE {
       { ?x  (translated  .  type)    Poem }
          UNION
       { ?x wrote  ?l .}
   }
```

- q_8 – select all the triples (i.e., the entire graph).

q_8
```
 SELECT * WHERE {
    ?s ?p ?o
 }
```

- q_9 – select all those capital cities which can be reached by sequence of train or plane from Paris.

q_9
```
SELECT ?city
WHERE {
       Paris (train | plane)+ ?city .
       ?city capitalOf ?country .
}
```

- q_{10} – select all cities that can be reached from Paris by any transport means.

q_{10}
```
SELECT ?city
WHERE {
    ?p sp* transport .
     Paris (?p)+ ?city .
}
```

- q_{11} – selects cities connected to Paris by plane or train followed by a bus.

q_{11}

```
SELECT ?city
WHERE {
      Paris (train | plane).bus ?city .
}
```

- q_{12} – selects cities connected to Paris by plane or train followed by a bus.

q_{12}

```
SELECT ?city
WHERE {
   { {Paris train ?c .} UNION {Paris plane ?c .} }
   ?c bus ?city .
}
```

q	q'	$q \sqsubseteq q'$?	**Running time (in millis)**
q_1	q_2	Yes	202
q_2	q_1	No	217
q_4	q_3	No	210
q_3	q_4	Yes	202
q_6	q_7	Yes	438
q_7	q_6	No	513
q_8	q_2	No	191
q_2	q_8	No	211

Table 1: Running time of query containment test.

5 Related Work

Query optimization has been the subject of an important research effort for many types of query languages, with the common goal of speeding up query processing. To the best of our knowledge, so far the problem of SPARQL with path query optimization has not been addressed. However, the works found in [18, 9, 16] considered the problem of SPARQL query optimization. So, the present work can be used to prove the correctness of query rewriting techniques.

An early formalization of RDF(S) graphs has been presented in [10], in which the complexity of query evaluation and containment is also studied. The authors investigate a Datalog-style, rule-based query language for RDF(S) graphs. In particular, they establish an NP-completeness result for query containment over simple RDF graphs. The work found in [17] provides algorithms for the containment and minimization of RDF(S) query patterns utilizing concept and property hierarchies for the query language RQL (RDF Query Language). The NP-completeness is established for query containment concerning conjunctive and union of conjunctive queries. Query containment is studied in [3] based on an encoding in propositional dynamic logic with converse (CPDL). They establish an EXPTIME complexity bound for containment of union of conjunctive queries under description logic constraints. Our work is similar in spirit, in the sense that the μ-calculus is a logic that subsumes CPDL and may open the

way for extensions of the query language. In particular, we consider the OPT operator, previously overlooked, and regular graph patterns including paths of arbitrary length.

Most notably and closely related results on query containment come from the study of regular path queries (RPQs) [2]. The difference between [2] and our work lies in the features supported by the languages. While RPQs in [2] support backward navigation and conjunction, PSPARQL supports variables in paths, union, and negation of queries (as implicit negation is carried by the query operator OPT).

Conjunctive RPQs have been studied in [7, 5] where an EXPSPACE algorithm for query containment is proved. On the other hand, containment of conjunctive RPQs with inverse have an EXPSPACE worst case complexity [2]. Most recently, containment of RPQs under description logic constraints have been studied in [4], and it has been show that the problem is 2EXPTIME complete.

6 Conclusions

In this paper we addressed query containment of SPARQL queries with paths. We took a similar approach to [8] that established the optimal complexity for XPath query containment. The problem of PSPARQL query containment has been reduced to satisfiability test in μ-calculus. For that purpose, we encoded RDF graphs as transition systems and PSPARQL queries as formulas. The reduction is proved to be sound and complete and the problem is shown to be EXPTIME-complete. In addition, we implemented the proposed approach via an encoding using the μ-calculus solver of [20] and this demonstrated the effectiveness of the encoding.

Paths are included in the new version of SPARQL[1] which is currently under standardization by W3C hence our results are a step towards query containment for SPARQL 1.1. The proposed encodings are not specific to PSPARQL. The same RDF encoding can be used for SPARQL query containment.

Similarly, the PSPARQL query encoding can be extended to richer types of queries, e.g., query modulo RDFS constraints or OWL ontologies, only Theorem 2 would change. So, we plan to extend this work towards different query types.

References

[1] F. Alkhateeb, J. Baget, and J. Euzenat. Extending SPARQL with regular expression patterns (for querying RDF). *Journal of web semantics*, 7(2):57–73, 2009.

[2] D. Calvanese, G. De Giacomo, M. Lenzerini, and M. Y. Vardi. Containment of conjunctive regular path queries with inverse. In *Proc. of the 7th Int. Conf. on the Principles of Knowledge Representation and Reasoning (KR 2000)*, pages 176–185, 2000.

[1] `http://www.w3.org/TR/sparql11-query`

[3] D. Calvanese, G. D. Giacomo, and M. Lenzerini. Conjunctive query containment and answering under description logics constraints. *ACM Trans. on Computational Logic*, 9(3):22.1–31, 2008.

[4] D. Calvanese, M. Ortiz, and M. Simkus. Containment of regular path queries under description logic constraints. In *Proc. of the 22nd Int. Joint Conf. on Artificial Intelligence (IJCAI 2011)*, 2011. To appear.

[5] D. Calvanese and R. Rosati. Answering recursive queries under keys and foreign keys is undecidable. In *Proc. of the 10th Int. Workshop on Knowledge Representation meets Databases (KRDB 2003)*.

[6] A. Chandra and P. Merlin. Optimal implementation of conjunctive queries in relational databases. In *Proc. STOC*, pages 77–90, 1977.

[7] D. Florescu, A. Levy, and D. Suciu. Query containment for conjunctive queries with regular expressions. In *Proceedings of the seventeenth ACM SIGACT-SIGMOD-SIGART symposium on Principles of database systems*, pages 139–148. ACM, 1998.

[8] P. Genevès, N. Layaïda, and A. Schmitt. Efficient static analysis of XML paths and types. In *Proc. PLDI*, pages 342–351, 2007.

[9] J. Groppe, S. Groppe, and J. Kolbaum. Optimization of SPARQL by using coreSPARQL. In *Proc. ICEIS (1)*, pages 107–112, 2009.

[10] C. Gutiérrez, C. Hurtado, and A. Mendelzon. Foundations of semantic web databases. In *Proc. PODS*, pages 95–106, 2004.

[11] P. Hayes. RDF semantics. W3C Rec., 2004.

[12] Y. Ioannidis. Query optimization. *ACM Comput. Surv.*, 28(1):121–123, 1996.

[13] D. Kozen. Results on the propositional μ-calculus. *Theor. Comp. Sci.*, 27:333–354, 1983.

[14] E. Prud'hommeaux and A. Seaborne. SPARQL Query Language for RDF. W3C Rec., 2008.

[15] U. Sattler and M. Y. Vardi. The hybrid μ-calculus. In *IJCAR*, pages 76–91, 2001.

[16] M. Schmidt, M. Meier, and G. Lausen. Foundations of SPARQL query optimization. In *Proc. ICDT*, pages 4–33, 2010.

[17] G. Serfiotis, I. Koffina, V. Christophides, and V. Tannen. Containment and minimization of RDF/S query patterns. In *Proc. ISWC*, pages 607–623, 2005.

[18] M. Stocker, A. Seaborne, A. Bernstein, C. Kiefer, and D. Reynolds. SPARQL basic graph pattern optimization using selectivity estimation. Proc. WWW, pages 595–604, 2008.

[19] Y. Tanabe, K. Takahashi, and M. Hagiya. A decision procedure for alternation-free modal μ-calculi. In *Advances in Modal Logic*, pages 341–362, 2008.

[20] Y. Tanabe, K. Takahashi, M. Yamamoto, A. Tozawa, and M. Hagiya. A decision procedure for the alternation-free two-way modal mu-calculus. In *Proc. TABLEAUX*, pages 277–291, 2005.

Centre de recherche INRIA Grenoble – Rhône-Alpes
655, avenue de l'Europe - 38334 Montbonnot Saint-Ismier (France)

Centre de recherche INRIA Bordeaux – Sud Ouest : Domaine Universitaire - 351, cours de la Libération - 33405 Talence Cedex
Centre de recherche INRIA Lille – Nord Europe : Parc Scientifique de la Haute Borne - 40, avenue Halley - 59650 Villeneuve d'Ascq
Centre de recherche INRIA Nancy – Grand Est : LORIA, Technopôle de Nancy-Brabois - Campus scientifique
615, rue du Jardin Botanique - BP 101 - 54602 Villers-lès-Nancy Cedex
Centre de recherche INRIA Paris – Rocquencourt : Domaine de Voluceau - Rocquencourt - BP 105 - 78153 Le Chesnay Cedex
Centre de recherche INRIA Rennes – Bretagne Atlantique : IRISA, Campus universitaire de Beaulieu - 35042 Rennes Cedex
Centre de recherche INRIA Saclay – Île-de-France : Parc Orsay Université - ZAC des Vignes : 4, rue Jacques Monod - 91893 Orsay Cedex
Centre de recherche INRIA Sophia Antipolis – Méditerranée : 2004, route des Lucioles - BP 93 - 06902 Sophia Antipolis Cedex

Éditeur
INRIA - Domaine de Voluceau - Rocquencourt, BP 105 - 78153 Le Chesnay Cedex (France)
http://www.inria.fr
ISSN 0249-6399